

Reducing a feasible soln to a Basic Feasible soln

Theorem - If there is a fis to a set of m simultaneous linear equation Ax=b in on unknowns and r(A) = m, then there is a bit.s.

Example-

x=2, x=3, x=1 is a for of the LPP.

Max Z = x + 2x2 + 4x3

A.t. 2x+ x+ 4x = 11

4x = 6  $3x_1 + x_2 + 5x_3 = 14$ 

 $x_1, x_2, x_3 \ge 0$ 

Find a b.f.s.

 $A = \begin{bmatrix} a_1 & a_2 & a_3 \end{bmatrix} = \begin{bmatrix} 2 & 1 & 4 \\ 3 & 1 & 5 \end{bmatrix} \quad b = \begin{bmatrix} 11 \\ 14 \end{bmatrix}$ 

x1 a + 22 az + 23 az = 6

2a1+3a2+03=b-D

a, az, az see linearly dependent

2) as + dz az + dz az = 0

 $\lambda_1 \begin{pmatrix} 2 \\ 3 \end{pmatrix} + \lambda_2 \begin{pmatrix} 1 \\ 1 \end{pmatrix} + \lambda_3 \begin{pmatrix} 4 \\ 5 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \end{pmatrix}$ 

2 1 + 12 + 423 =0

321 + 2 + 523 = 0

 $\frac{\lambda_1}{5-4} = \frac{\lambda_2}{12-10} = \frac{\lambda_3}{2-3} = k (say)$ 

 $\lambda_1 = k$   $\lambda_2 = 2k$   $\lambda_3 = -k$   $\lambda_3 = -k$   $\lambda_4 = k$   $\lambda_5 = k$   $\lambda_6 = k$   $\lambda_6 = k$   $\lambda_6 = k$ 

a1+2a2-a3=0 -0

(2,3,D) to not a bits as the bits of this problem cannot have more than 2 positive variable. To determine the departing variable  $\min_{j} \left\{ \frac{x_{j}}{\lambda_{j}} \mid \lambda_{j} > 0 \right\}$ = min  $\left\{\frac{2}{7}, \frac{3}{2}\right\} = \frac{3}{2} \Rightarrow j = 2$ on is the depending vector and the to 20. 0 x2 - @ 0 X3 = 4ay + 6g/2 + 2as = 2b  $3m + 6a_1 - 3a_3 = 0$ ay 1+503 = 2b 1 a + 0.02 + 503 = p  $\left(\frac{1}{2},0,\frac{5}{2}\right)$ Instead we can calculate max votio as

mox  $\begin{cases} x_j & |\lambda_j| < 0 \end{cases} = \begin{cases} x_3 \\ x_3 \end{cases} \Rightarrow j=3$ as departing vector new  $x_j = dd x_j - \frac{x_3}{x_3} \lambda_j$ Men x = 2-(-1).1=3 New x= 3-(-1) 2 = 5 New x3 = 1-(-1)(-1)=0 Man potio tof. Another bf.s. (3,5,0)

P-10 N1

MMITTE Example xx+ 2x+ 4x3+ x4=+ 2x1 - x2 + 3x3 - 2x4=4 f.s. > x=1, z=1, x=1, xy=0 Reduce it to a before. x a + x a + x a + x a + x m = b at az + az = b May + hay + may = 0 Example -Show that x=5, x=0, x=-1 ip a b.f.s. of the system of ear x1+2x+x=4 2xx + xx +5x3=5 find the basic soln if there is any.  $A = (a_1, a_2, a_3) = (1 \ 2 \ 1 \ 5) b = (4 \ 5)$  $g_2 = 0$   $(a_1, a_2) = \begin{bmatrix} 1 & 1 \\ 2 & 5 \end{bmatrix}$  $det (11) = 5-2 = 3 \neq 0$ a, , as linearly independent So the given solution is basic Other basic soln.  $B = \begin{pmatrix} 1 & 1 \\ 2 & 5 \end{pmatrix} \qquad B^{-1} = \frac{1}{3} \begin{bmatrix} 5 & -1 \\ -2 & 1 \end{bmatrix}$  $x_8 = 8^{-1}b = \frac{1}{3} \left(\frac{5}{-2} - 1\right) \left(\frac{4}{5}\right) = \frac{1}{3} \left(\frac{15}{-3}\right) = \left(\frac{5}{-1}\right) \left(\frac{24}{3}\right)$ 

Marked No

For non-basic variable xz,

$$Y_2 = B^{-1}a_2$$
  
=  $\frac{1}{3}\begin{pmatrix} 5 & -1 \\ -2 & 1 \end{pmatrix}\begin{pmatrix} 2 \\ 1 \end{pmatrix} = \frac{1}{3}\begin{pmatrix} 9 \\ -3 \end{pmatrix} = \begin{pmatrix} 3 \\ -1 \end{pmatrix} \rightarrow t_{re}$ 

atleast one component of 1/2 is -ve as can enter the basis replacing as.

New Basis
$$\vec{B} = \begin{pmatrix} 2 & 1 \\ 1 & 5 \end{pmatrix}$$
 $(B)^{-1} = \begin{pmatrix} 1 \\ -1 & 2 \end{pmatrix}$ 

$$x_{\overline{8}} = (\overline{8})^{-1}b = 1 (5 - 1) (4) = (5/3) = (x_2)$$

$$x_1 = 0$$
,  $x_2 = 5/3$ ,  $x_3 = 2/3$ 

for Non-basic variable 
$$x_1$$

$$y_1 = (\overline{R})^{-1} \quad \text{ay} = \underbrace{1}_{9} \left( 5 - 1 \right) \left( \frac{1}{2} \right) = \begin{pmatrix} y_3 \\ y_3 \end{pmatrix}$$

ay enter replacing as in the booking

$$B' = (a_2, a_1) = \begin{pmatrix} 2 & 1 \\ 1 & 2 \end{pmatrix} = \begin{pmatrix} 8'5' = 1 \\ 3 & -1 \end{pmatrix}$$

$$x_{\beta}' = (\beta')^{-1} b$$

$$= \frac{1}{3} \left( \frac{2}{-1} - \frac{1}{2} \right) \left( \frac{4}{5} \right) = \left( \frac{1}{2} \right) = \left( \frac{x_2}{x_4} \right)$$

S- Apply the simplex process to solve the LPP (without using simplex table)

Max 
$$Z = 2x_1 - 3x_2$$

At  $2x_1 + 5x_2 \ge 10$ 
 $3x_1 + 8x_2 \le 2y$ 
 $x_1, x_2 \ge 0$ .

Max  $Z = 2x_1 - 3x_2$ 

At  $2x_1 + 5x_2 - x_3 = 10$ 

Ax  $1 + 8x_1 + x_2 = 2y$ 

Ax  $1 + 8x_1$ 

Book feasible Solution (BFS)

$$x_8 = B^{-1}b = \frac{1}{2}\begin{pmatrix} 1 & 0 \\ -3 & 2 \end{pmatrix}\begin{pmatrix} 10 \\ 24 \end{pmatrix} = \begin{pmatrix} 5 \\ 9 \end{pmatrix} = \begin{pmatrix} x_1 \\ x_{14} \end{pmatrix}$$
 $x_1 = 5$ ,  $x_2 = 0$ ,  $x_3 = 0$ ,  $x_{14} = 9$  —  $0$ .

$$C_8 = (2,0)$$
  $Z_8 = C_8 \times_B = (2,0) \begin{pmatrix} 5 \\ 9 \end{pmatrix} = 10$ 

= 25 x = 16 8=4

For non-book vector, compute  

$$Y_2 = B^{-1}a_2 = \frac{1}{2} \begin{pmatrix} 1 & 0 \\ -3 & 2 \end{pmatrix} \begin{pmatrix} 5 \\ 8 \end{pmatrix} = \begin{pmatrix} 5/2 \\ 1/2 \end{pmatrix}$$

$$\frac{1}{2} = \frac{1}{2} \begin{pmatrix} 1 & 0 \\ -3 & 2 \end{pmatrix} \begin{pmatrix} -1 \\ 0 \end{pmatrix} = \begin{pmatrix} -1/2 \\ 3/2 \end{pmatrix}$$

Compute 
$$Z_2 - C_2$$
,  $Z_3 - C_3$ 

$$Z_2 - C_2 = C_8 Y_2 - C_2 = (2, 0) \begin{pmatrix} 5/2 \\ \sqrt{2} \end{pmatrix} - (-3)$$

$$= 5 + 3 = 8$$

$$z_3 - c_3 = c_0 y_3 - c_3 = (2,0) (-y_2) - 0$$

$$= -1 < 0$$

O is not an optimal solution as is the entering vector in the basis

find the min rotio,

min 
$$\left\{\frac{z_{Bi}}{y_{i3}}\right\}$$

my is leaving

New basis

$$B = \begin{pmatrix} a_1 \cdot a_3 \end{pmatrix} = \begin{pmatrix} 2 & -1 \\ 3 & 0 \end{pmatrix}$$

$$8^{-1} = \frac{1}{3} \begin{pmatrix} 0 & 1 \\ -3 & 2 \end{pmatrix}$$

$$x_{8} = 8^{-1}b = \frac{1}{3}\begin{pmatrix} 0 & 1 \\ -3 & 2 \end{pmatrix}\begin{pmatrix} 1 & 0 \\ 2 & 4 \end{pmatrix} = \begin{pmatrix} 8 \\ 6 \end{pmatrix} = \begin{pmatrix} x_{1} \\ x_{3} \end{pmatrix}$$
 $x_{1} = 8$ ,  $x_{2} = 0$ ,  $x_{3} = 6$ ,  $x_{4} = 0$  —  $\infty$ 

CB = (2,0)  $Z = Z_B = C_B X_B = (2.0)(8) = 16$ 

$$\binom{8}{6} = 16$$

((saatht))

For non-basic variables 
$$x_1, x_4$$
 we compute  $Y_2 = 6^{-1}a_2 = \binom{8/3}{1/3}$  generally  $y_4 = 8^{-1}a_4 = \binom{1}{3}$ 

Compute -22-(2= CB Y2-(-3) = 25/3 >0

Zy-Cy = CB /4 - 0 = 2/3 70

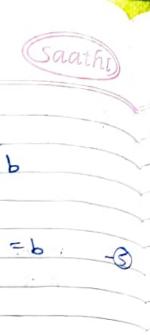
Optimality is reached a is the optimal solution

Date 05 / 09 / 2023



attains unbounded solv. 8- Man Z = 3x4+ 4x2 p. x. x. -x2 20 -x+3x2 <3 × 2 20 Q- Man Z = 6x1 + 10x2 attains on alternate optimal A.t. 324 + 522 =10 2j-G=0 for some 5x+3x =15 non-basic variable. ×1, ×20 Improving a BFS (Simplex Method) Man Z= CX Amxn man Ax = bRank of A = m 220 as az aj an 13-5 Bosis B = [B1, B2, B3, --, Bm] 811 812 813 Der der des den 0- Aj = 2 tij Bi = yij Bi + yzj Bz + ... + ymj Pm. Ymi ymz ymj ging yoj ≠0 ⇒ Br can be replaced by aj Br = 0; - 12 41; Bi - 0 XB => BXB=b E β; xε; = b - 3.

 $C = (C_B | C_R)$   $Z_B = C_B \times_B = \sum_{i=1}^{\infty} C_{Bi} \times_{Bi} - \Phi$ 



## New Basic solution-

saathi xor < zeri yij > a yri = min { xxi / yij >0 } (min ratio) a; is the enting rector = \( \frac{2}{1=1} \) \( \text{Ce} \) \( \text{Xe} \) \( \text{E} ZB - XBr (zj - cj)

Zx-Cx = Min {zj-G | zj-G < O}

Schoice of EV {Entering variable}



Alternative Optima-

Theorem - If there is an optimal BFS to a LPP and zj-cj=0 for some non-basic vector aj sud yij >0 for atleast one i then there exists an alternative Basic Optimal soln.

1.b.

Theorem = If there is an optimal BFS to a LPP and zj = g =0 for some non-basic vector of and yis < 0 4 i=1,2,then there exists an alternative non-basic optimal solution.

~ (xsi - yij 0) βi + θaj = b

- \ xxi - O yij i=1,2,--, m 1 0 maj bosic

( a new for nith (m+1) positive variables remaining zero

Z' = \( \text{CBi} \ (\text{xBi} - \text{Oyij}) + \text{GO}

= E Csi xs: - O E cai yij + GO

= ZB - 8 (zj - cj)

Zj-Cj=O for some non-basic variable for the existence of alternative option -> not a necessary condition



## Zj-cj70 for a non-book vector oj

By - leaving rector X 6x=0, 70,70 I degenerate cose.

× 88 = 0

2) = ZB - XBr (zj - cj)

= ZB

Two-Phase Method

Phose T-

Max Z\* = 0.x4 + 0.x5 + --- + 0.x4

-1. xay -1-xaz - ...- 1 xam.

→ no artificial variable → artificial variable at zero level.

Phase II-

Careful about artificial variable.

Should always remain at zero level

Case I -

ax EV ar LV

ZBi = ZBi - XBr yik