

disk updation

28th march

Variable	Value		
p	43		
h	13		
m	29		
c	49		
h	26		

$$h = h * 2$$

$$C = C + 17$$

on disk write

$$h = 26, C = 66$$

if power fails in middle of execution

entire program is executed again.

If power fails in middle of variable updation.

start variable updation again.

on disk write

$h = 26$

done written:

$c = 66$

update variables

done.

done not written:

execute the operations
again.

t	19
d	94
m	36

Variable	Valid	Value	shadow
t	0	19	41
d	1	26	94
m	1	49	36

0=actual value valid
1=shadow valid

Variable	Valid	Value	Shadow
p	0	43	39
h	x _{0,1}	12/84	42/84
m	1	29	98
c	x _{0,1}	49	19/68/66
h	1	26	42

on disk write

$$h = 84$$

$$C = 66$$

done.

if power fails before "C" shadow updation

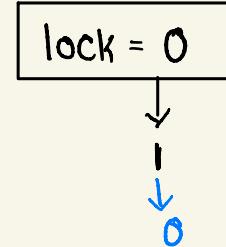
then update starts again,

shadow of h is also updated.

but during power cut, $h = 84, C = 49$ are

values.

Variable	Valid	Value	Shadow
p	0	43	39
h	0 10	28 46	31 46
m	0	29	98
c	0 10	49 66	38 66
h	0	36	42



On disk write

$$h = 46$$

$$c = 66$$

done.

if ($lock == valid == 1$)

(value) might be wrong.

after done is written:

update in shadow

make lock 1

make valid of h and c $\rightarrow 1$

transfer (shadow \rightarrow value) of h

make valid = 0

transfer (shadow \rightarrow value) of c

make valid = 0

lock = 0

3rd April

Buddy

Size: 2^k

int p, a, g, h, k, l, u, m, f, y

Virtual address.

000-	{	0000 p	
0001		0001 a	
0010		0010 g	
0011		0011 h	
0100		0100 k	
0101		0101 f	
0110		0110 u	
0111		0111 m	
1000		1000 f	
1001		1001 y	
1010		1010 e	

BUDDY

000 - -

0010 - -

0011 - -

01 - - -

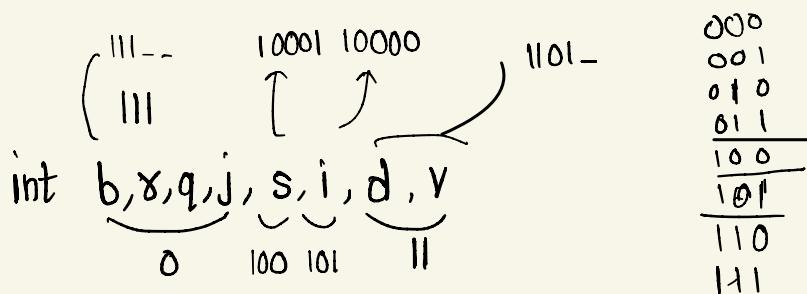
00000 | 00001 | 00010 | 00011 | 00100 | 00101 | 00110 | 00111 | 01000 | 01001 | 01010 | 01011 | 01100 | 01101 | 01110 | 01111 | 10000 | 10001 | 10010 | 10011 | 10100 | 10101 | 10110 | 10111

Physical
address.

000	001	01	100	← virtual buddy
1001	0011	101	0010	

Segment table :

virtual	01 <u>01</u>	00 <u>11</u>
physical	10101	00111



Segment table : 111, 10001, 10000, 1101

Virtual 011 110

Physical 11111 11010

000	001	01	10	virtual buddy
1001	0011	101	000	

virtual buddy

virtual size (45)

000, 001, 01, 10

(8) (8) (16) (16) δp_{q8}

abcde fghij klmn minimum

Physical size:

$$\text{wastage} = 3\%$$

Segment table :

a b c d e f g h i j k l m n o p q x

Segment table for 000 is x

01 is x-1

001101

$$2x + 2x - 2 = 18$$

fg hij 101

$$x = 5$$

011010

k l m n | o | o

Page table :

abcdefghijklmno

qrstuvwxyz

^(2^8)

find physical for

virtual

0101101

let pagesize = 8

0101101
page no

10 pages

9 to 16
pages.

Let page size = 32

0101101

3 or 4 pages

abcde, fghij, klmno, pqrst

→ fghij 01101

Physical memory = $2^{(5+5)}$
 $= 1024,$

Let page size = 16

0101101

number of pages = 5

abcd, efg, ijkl, mnop, qrst

0101101

ijkl 1101

Physical memory = 2^8

Segment table : a b c d e f (6)

3 buddies

ab, cd, ef - possible

a, bc, def ✓
 $2^k \quad 2^{k-1} \quad 2^{k-2}$
4 2 1

abc, de, f - ✗

a | bcd | cfg - ✓

abc, de, f

1 write smallest sequence so that work is done faster
in FCFS, composed SSTF

40 60 30 10

40 80 90 110

40 80 50 130



10

30

40

60

Job	Arrival	Serve
A	0	50
B	10	35

how to serve so that

- 1) Work over smallest
 - 2) Average wait is smallest.
-

Job Arrival service

A	0	50	for which 'K' preemption
B	10	K	reduces average wait.

let swap in S
swap out Z

$$k < 50$$

$$\frac{k}{2}$$

$$40/2$$

$$40/2 > k/2$$

10th April

RAM	operation	Disk
	Disk(45)	45
19	RAM(19)	
19, 24	RAM(24)	
	Disk(13)	45, 13
19, 24, 17	RAM(17)	
	Disk(42)	45, 13, 42
NULL	RAM(91)	45, 13, 42, 19, 24, 17, 91
71	RAM(71)	

(max capacity
of RAM = 3)

automatic transfer
from RAM → Disk

RAM	operation	Disk
	Disk(12)	12
19, 16	RAM(19)	
	Disk(13)	12, 13
17	RAM→Disk	12, 13, 19, 16
	Disk(49)	12, 13, 19, 16, 49
	OVER	12, 13, 19, 16, 49, 17

write operation

RAM	Operation	Disk	Work
		12, 17, 45, 19, 36, 84, 71, 93, 86	
	a = Disk	12, 17, 45, 19, 36, 84, 71, 93, 86	a = 12
	b = Disk	12, 17, 45, 19, 36, 84, 71, 93, 86	b = 17
19, 36	c = RAM	12, 17, 45, 19, 36, 84, 71, 93, 86	c = 45
same	d = Disk	12, 17, 45, 19, 36, 84, 71, 93, 86	d = 84
36	e = RAM	12, 17, 45, 19, 36, 84, 71, 93, 86	e = 19
nil	RAM \rightarrow Disk	12, 17, 45, 19, 36, 84, 71, 93, 86	
71, 93	f = RAM	12, 17, 45, 19, 36, 84, 71, 93, 86	f = 84
	RAM \rightarrow Disk	12, 17, 45, 19, 36, 84, 71, 93, 86	

read operation

RAM	operation	Disk
	Disk(45)	45
19	RAM(19)	
19,24	RAM(24)	
	Disk(13)	45, 13
111	RAM(17)	45, 13, 19, 24, 17
	Disk(42)	45, 13, 19, 24, 17, 42
91	RAM(91)	45, 13, 19, 24, 17, 42
91,71	RAM(71)	same

automatic transfer

write operation when RAM reaches
max capacity.

during read operation from RAM

if write RAM is not
empty then RAM \rightarrow D

while handling both read and
write simultaneously
during RAM \rightarrow DISK
(1st read, 2nd write)

RAM - write	RAM - read	operation	Disk
NULL	NULL		12, 17, 34, 81, 29, 38, 73, 86, 49
17, 34	a = RAM		12, 17, 34, 81, 29, 38, 73, 86, 49 a = 12
	DISK(39)		12, 17, 34, 81, 29, 38, 73, 86, 49
	RAM(18)		
16, 93	RAM(93)		
NULL	RAM → DISK	12, 17, 16, 93, 29, 38, 73, 86, 49	
38, 73	b = RAM	12, 17, 16, 93, 29, 38, 73, 86, 49 b = 29	
46	RAM(46)		
38, 73	C = . disk	12, 17, 16, 93, 29, 38, 73, 86, 49	C = 26
49	d = RAM	12, 17, 16, 93, 29, 38, 46, 86, 49	d = 26

RAM	Operation	Disk
Pq8	H(Pq8)	
	L(gh)	gh
cd	H(UKcd)	ghpq8uk
cedit	H(it)	
	L(hm)	ghpq8uk hm
<u>OVER</u>		ghpq8ukhm redit

RAM size = 5

H - high level operation

L - low level operation
(directly edits on disk) ↑

RAM	Operation	Disk
	H(Pq8)	Pq8
	L(gh)	pq8gh
d	H(UKcd)	pq8ghulk
	H(it)	pq8ghulk dit
	L(hm)	pq8ghulk dithm

RAM size = 3

write	read	operation	ukhcydmfzkcld
		Z(a)	ukhcydmfzkcld
		L(ts)	uts cydmfzkcld
		Place(7)	uts cydmfzkcld
		L(a,b)	uts cydmab kclrd
		Place(4)	uts ydmab kclrd
mab	Y(bc)	uts ydmab kclrd	b=y c=d
	L(gh)	uts cydmabgh lrd	
ab	X(d)		d=m
kch	H(kch)		
	Y(e)		

Z (a) = Disk read a=

L(a)=Disk write a

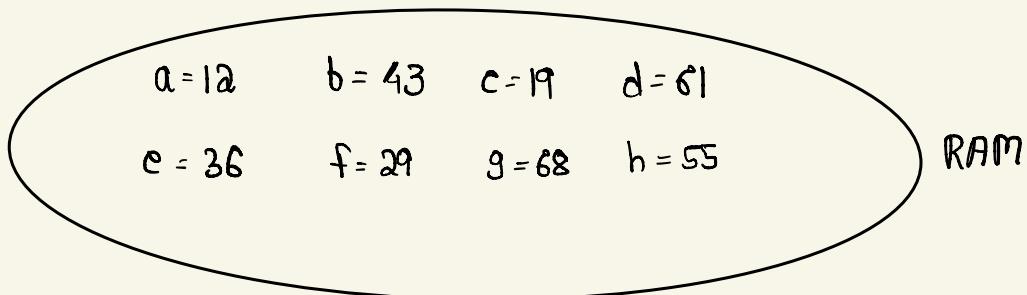
Y(a)=read from ram a=

H(a)=write on ram a

Place 4= take cursor to 4

11th April

Cache : Faster than RAM



{a,f,g} {b,c} {d,e,h}

Cache : U₀xy V₁pq W₂st

U=0 \Rightarrow a, f, g not available in cache.

U=2 \Rightarrow f in xy.

W=3 \Rightarrow h in st

Operation	Cache	
	042017036	
Print(c)	042219036	{19 R → C}
Print(g)	368219036	{68 R → C}
Print(b)	368143036	{43 R → C}
Print(e)	368143236	{36 R → C}
Print(g)	same	no transfer

Operation	Cache	
	029033067	
$b = 32$	029632067	
Point(e)	029632236	$\{e \rightarrow C\}$
$e = 41$	029632741	
Point(c)	029219741	$\begin{cases} C \rightarrow C \\ b(32) C \rightarrow R \end{cases}$
Point a	112219741	$12R \rightarrow C$
OVER		$\{e(41) C \rightarrow R\}$

+ 5 means
values to be
updated)

$a = 29$	$c = 73$	$e = 32$	$g = 86$
$b = 61$	$d = 84$	$f = 61$	$h = 19$
P	Q	R	S

group {PR} {QS}

Uxyzw Vijkl

U = 0

U = 1 1st variable updated

U = 2 second variable updated

U = 3 both updated

U = 4 neither updated

U = 6
U = 7
U = 8
U = 9

}

second page of group

Operation

0423906132

Print(d)

0423947384

7384 R → C

Print(e)

9326147384

3261 R → C

F = 69

8326947384

no transfer

h = 92

8326977392

no transfer

Print(a)

4296177392

{32 69} C → R

{29 61} R → C

Print(c)

4296147384

{73 84} R → C

{86 19} R → C

{86 92} C → R

c = 36

6366147384

d = 29

636C1

7329
3

OVER

{73 29} C → R
{32, 61} R → C
{36, 61} C → R

	04239 06132	
$f = 69$	74269	
$e = 11$	81169	
Print(a)		$\{1169\} C \rightarrow R$ $\{2961\} R \rightarrow C$