

# CS/ECE/ISyE 524

## Final Exam Preparation



# Details!

## Final Exam

- 07:45AM-09:45AM, May 10, 2024.
- Van Vleck B102
- McBurney Students: 7:45AM-10:45AM, 3126 Mechanical Engineering Building

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- Closed Book, Closed Notes, Closed Everything!
  - Calculators are allowed, but not your phones.
  - I will provide the Slide of Trix (and other notes)—final standard “cheat sheet” posted by next Tuesday

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## Other Announcements

- Prof. L Office Hours: May 7, 9AM-10AM

# Today's Docket

- 1 Final Details
- 2 Review Course Topics
- 3 Questions
- 4 Practice Problems?

# Final Exam

- **Coverage:** Everything in Lecture Notes (Lecture 1-Lecture 18)
- Probably Six Problems
  - 1 True/False
  - 2 Short Answer
  - 3 Duality
  - 4 Modeling
  - 5 Modeling
  - 6 Modeling

# Final Exam

- No Note Sheet Allowed
  - No calculators or other electronic devices
  - Bring your ID to the exam and place it out. We will be checking IDs during the exam.
    - Also remember to take it with you
  - Spread out in the room! (We may ask you to move)
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# Final Exam

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## When Exam is Over...

- I will ask everyone to put their pencils down
- You are not to pick up your pencil again, even to write your name.
- If you pick up your pencil again, I will take off 10% of grade.

# Math Modeling—Being More Strict

- You *must* write proper mathematical notation and syntax to get credit. We will give less partial credit.
  - Writing JuMP code `@constraint(m, sum(x) == 1)` is *not* mathematical syntax
- You have to write specifically the set of entities over which you are doing a sum.
  - $\sum_i x_i$  is incorrect.  $i$  is an index, not a set
  - $\sum x_i \forall i \in S$  is incorrect
  - $\sum_{i=1}^4 x_i$  or  $\sum_{i \in S} x_i$  is proper notation
- $x_k \geq 42 \forall k$  is also incorrect notation.  $k$  is an index. For all  $k$  in *what set*?

## Lecture 2

- Geometric properties of LP Feasible Region
- Graphing and Solving Linear Programs

## Lecture 3

- Linear Algebra, Matrices, and Definitions
  - Linear Combinations, Convex combinations, Polyhedra
- Converting LP to Standard Form



## Lecture 4

- Input-Output Linear Programs
  - General Modeling:
    - Sets, Parameters, Vars
    - When to use  $\Sigma$ , when to use  $\forall$
  - Infeasible, Unbounded, Optimal Solution
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## Lecture 5

- Blending Constraints
- Multiperiod Planning
  - Backlogging

## Lecture 6

- Modeling epigraph of convex piecewise linear-functions
  - $|\cdot|$
  - Mini-max (Maxi-Min)
- CPM/Project Scheduling

## Lecture 7

- (More) Multiperiod
- Networks
- Min-Cost Network Flow
  - Transportation Problem, Assignment Problem
  - Shortest Path, Longest Path (DAG)
  - Max Flow
- Integer Properties of Solutions

## Lecture 8

- Duality
  - Taking dual
  - Weak duality
  - Strong Duality
  - Complementary Slackness
- Economic Interpretation: Shadow Prices

## Lecture 9

- Duality / Complementary Slackness
- Network Duality (and Interpretation)
- Max-Flow Min Cut

## Lecture 10

- Quadratic Functions/Quadratic Forms
- Spectral Theorem
- PSD Matrices
- Definition of Convex and Concave Functions
- Examples of convex and concave functions

## Lecture 11

- Least Squares is a (Convex) Optimization Problem
- Geometry of Least Squares
- Regression and Curve fitting
- Multiobjective and Tradeoffs
- Pareto Curves
- Regularization
- Norm balls
- Hierarchical Optimization

## Lecture 12

- Ellipsoids
- Degenerate ellipsoids
- Shifted Ellipdoids
- Quadraic Programs: Solution properties (Convex/Nonconvex)
- Where do Quadratic occur?
- Portfolio Optimization

## Lecture 13

- Definition of Cone
- Types of Cones
- 1 Second Order Cone
- 2 Rotated Second Order Cone
- 3 Modeling them with Julia/JuMP
- 4 Rational Powers
- 5 Semidefinite Cone



## Lectures 14 and 15

- Geometry of Feasible Region
- Relaxations and LP relaxation and Convex hull
- IP Modeling: Fixed Costs. Big M
- IP Modeling Variable Lower Bounds
- IP Modeling: "Simple" Logic

## Lecture 16

- Logic Constraints and “Slide of Trix”
- Modeling restricted set of values (like SOS1)
- Sudoku

## Lecture 17

- Set Cover Formulations
  - Enumerate/encode “feasible” patterns
- Cutting stock (and column generation)
- Traveling Salesman Problem
  - Subtour elimination and cutting planes
  - Miller-Tucker Zemlin Constraints

## lecture 18

- Quadratic Assignment Problem
- Linearize (0,1) quadratic term:

$$x, y \in \{0, 1\}, z = xy \Leftrightarrow z \geq x + y - 1, z \leq x, z \leq y$$

- Piecewise Linear Functions
  - SOS2
  - Convex Combination Model
  - Multiple Choice Model

# Problem 1

Consider the following LP:

$$\begin{array}{ll}\text{maximize} & Z = x_1 + 2x_2 \\ \text{s.t.} & -x_1 + x_2 \leq 4 \\ & x_1 - 2x_2 \leq 2 \\ & 4x_1 + \textcolor{red}{a}x_2 \leq 16 \\ & x_1, x_2 \geq 0\end{array}$$

## Problem 1, cont'd

- (a) Find a value of  $a$  such that there is a single optimal solution.
- (b) Find a value of  $a$  such that there are an infinite number of optimal solutions.
- (c) Find a value of  $a$  such that the problem is feasible but there are no optimal solutions.

## Problem 2

Consider the following LP:

$$\begin{array}{ll}\text{maximize} & Z = 2x_1 + 5x_2 + 3x_3 \\ \text{s.t.} & x_1 - 2x_2 + x_3 \geq 20 \\ & 2x_1 + 4x_2 + x_3 = 50 \\ & x_1, x_2, x_3 \geq 0\end{array}$$

- (a) Write the dual LP.
- (b) Find the complementary dual solution to the primal solution  $x_1 = 25, x_2 = 0, x_3 = 0$
- (c) Is  $x_1 = 25, x_2 = 0, x_3 = 0$  an optimal solution to the problem?

# Problem 1

(a)  $a = 0$

(b)  $a = 8$

(c)  $a = -8$

There are many possible answers for each. To see why these answers are valid, plot the lines.

# Example

- Use a 0-1 variable  $\delta$  to indicate whether *or not* the constraint  $2x_1 + 3x_2 \leq 1$  is satisfied.
  - $x_1, x_2$  are nonnegative continuous variables that cannot exceed 1.
- $\delta = 1 \Leftrightarrow 2x_1 + 3x_2 \leq 1$
- $M$  : Upper Bound on  $2x_1 + 3x_2 - 1$ . 4 works
- $m$  : Lower Bound on  $2x_1 + 3x_2 - 1$ . -1 works.
- $\varepsilon$  : 0.1



## Example, Cont.

- $(\Rightarrow)$  Recall the trick.
  - $z = 1 \Rightarrow \sum_{j \in N} a_j x_j \leq b \Leftrightarrow \sum_{j \in N} a_j x_j \leq b + M(1 - z)$
- $2x_1 + 3x_2 + 4z \leq 5$
- $(\Leftarrow)$ . Recall the trick.
  - $\sum_{j \in N} a_j x_j \leq b \Rightarrow z = 1 \Leftrightarrow \sum_{j \in N} a_j x_j \geq b + mz + \varepsilon(1 - z)$
- $2x_1 + 3x_2 + 1.1z \geq 1.1$

$$\begin{aligned} 2x_1 + 3x_2 + 4z &\leq 5 \\ 2x_1 + 3x_2 + 1.1z &\geq 1.1 \end{aligned}$$

# Slide o' Trix: Drink Production

- Facilities  $F = B \cup V \cup G$ 
  - $B \cap V = B \cap G = V \cap G = \emptyset$
- $B \subset F$ : Beer production
- $V \subset F$ : Vodka production
- $G \subset F$ : Gin production
- Drinks  
 $D = \{\text{Beer, Vodka, Gin}\}$
- Colleges  $C = T \cup I$
- $T \subset C$ : Big 10 Colleges
- $I \subset C$ : Ivy league college
- $a_{cd}$ : Amount of drink  $d \in D$  required by college  $c \in C$
- $b_f$ : Maximum amount of production at facility  $F$
- $\alpha_f$ : Per unit production cost at  $f \in F$
- $\beta_{fc}$ : Per unit transportation cost from  $f \in F$  to  $c \in C$

## A Linear Program

Meet college requirements for each drink at minimum cost

# Some Problems of Varying Degrees of Difficulty

- You can open at most 4 Vodka Production Facilities
- If you open a vodka production facility, you must ship at least  $\Delta$  units out of the facility
- If you open at least 3 vodka facilities, then you must open at least 6 Gin facilities
- You must open at least 2 beer facilities or 2 vodka facilities or 2 gin facilities
- If you ship more Gin to Big Ten College than to Ivy colleges, then you must ship less Beer to Ivy Colleges than Big 10 Colleges
- If Ivy League colleges drink more vodka and more Gin than Big Ten students, then the Big Ten students GET NO BEER!