

10.02.25

TEOP. 43.2

$$1.4. \quad T(n) = T(n-1) + 3n = T(n-2) + 3(n-1) + 3n \quad (\equiv)$$

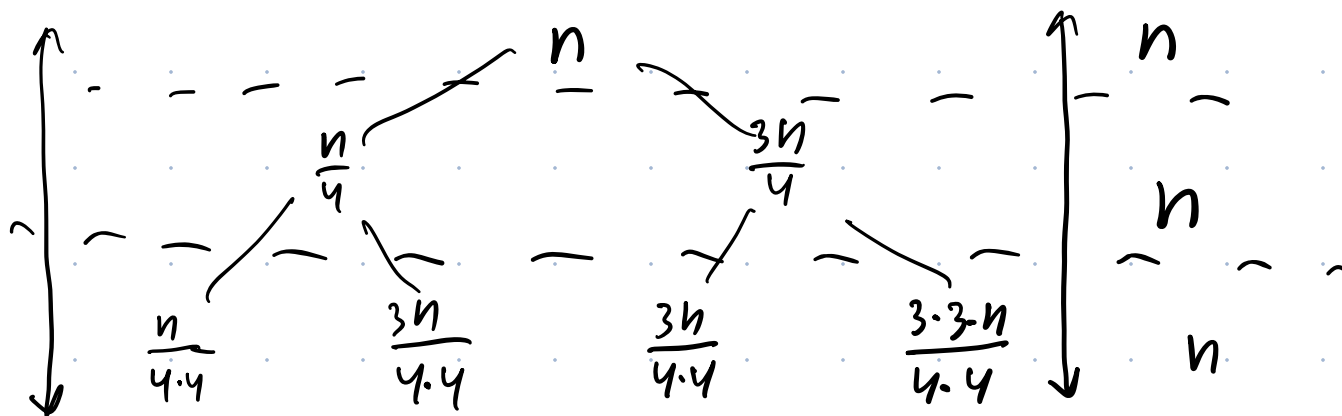
$$T(1) = C$$

$$(\equiv) \quad \sum_{i=0}^{n-2} 3(n-i) + C = 3 \sum_{i=2}^n i + C = 3 \left[\frac{n(n+1)}{2} - 3 \right] + C = \Theta(n^2)$$

$$1.5. \quad T(n) = T\left(\frac{n}{4}\right) + T\left(\frac{3n}{4}\right) + n$$

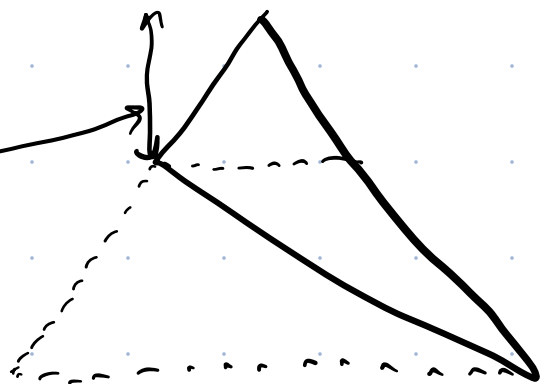
$$\geq 2T\left(\frac{n}{4}\right) + n; \quad c = \log_4 2 = \frac{1}{2} < 1; \quad T(n) = \Theta(n)$$

$$\leq 2T\left(\frac{3n}{4}\right) + n; \quad c = \log_{\frac{4}{3}} 2 > 1; \quad T(n) = \Theta(n^{\log_{\frac{4}{3}} 2})$$



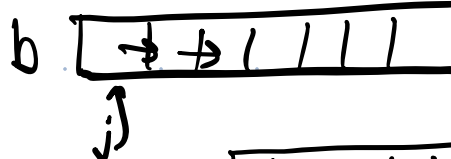
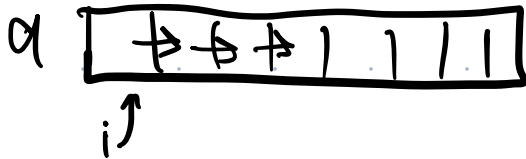
$$h_l = \log_4 n$$

$$h_r = \log_{\frac{4}{3}} n$$



$$n \log_{\frac{4}{3}} n \leq T(n) \leq n \log_{\frac{4}{3}} n$$

$$T(n) = \Theta(n \log n)$$



```
def merge(a, b):
```

```
    len_a = len(a)    # k
```

```
    len_b = len(b)    # m
```

```
    i, j = 0, 0
```

```
    merged = []
```

$\Theta(k + m)$

```
    while (i != len_a or j != len_b):
```

```
        if (i == len_a):
```

```
            merged.append(b[j])
```

```
            j += 1
```

```
        elif (j == len_b):
```

```
            merged.append(a[i])
```

```
            i += 1
```

```
        else:
```

```
            if (a[i] > b[j]):
```

```
                merged.append(b[j])
```

```
                j += 1
```

```
            else:
```

```
                merged.append(a[i])
```

```
                i += 1
```

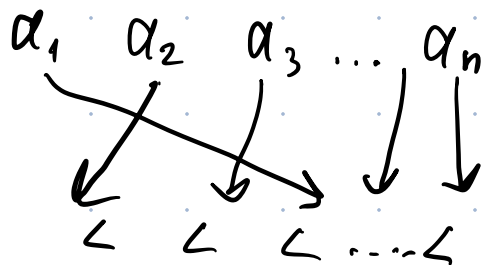
return merged

НИЖНЯЯ ОЦ. НА СОРТ. СРАВНЕНИЯМИ

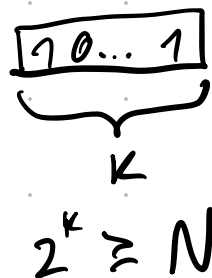
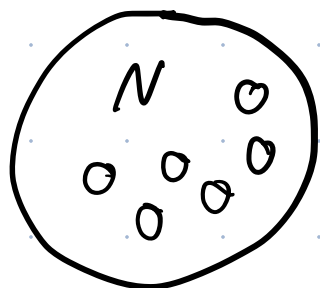
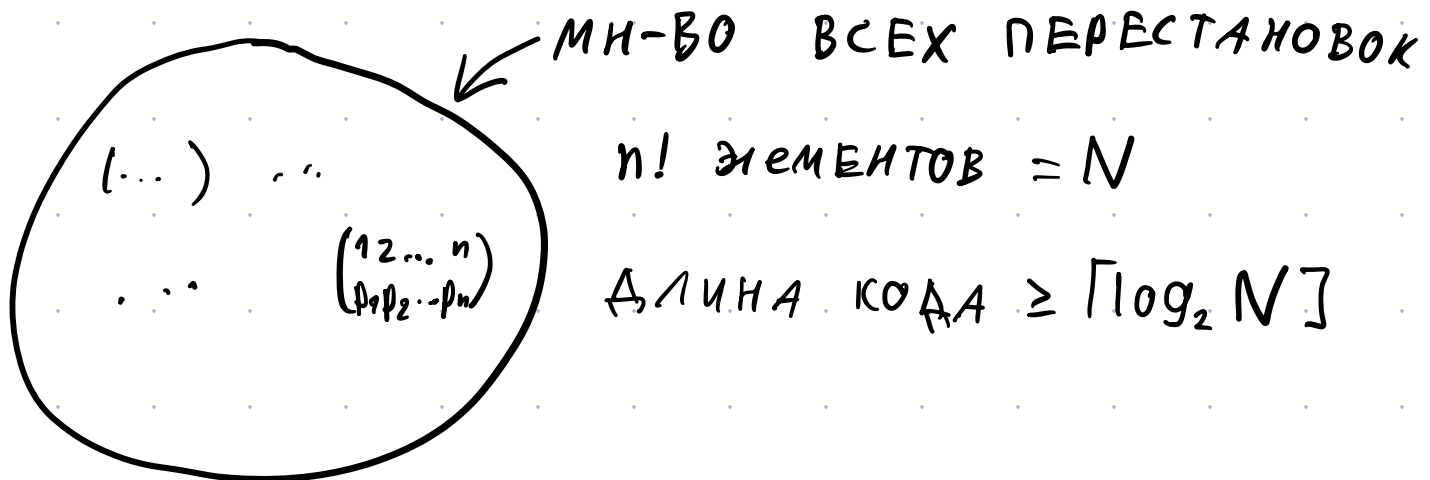
И ЧИСЛА: $1, 2, \dots, n$

$O(1)$ $a ? b$ (КАКОЙ ЭТО БОЛЬШЕ)

РАЗЛИЧНЫХ ПЕРЕСТАНОВОК $n!$ ШТУК



ДЛЯ КАЖДОЙ ПЕРЕСТАНОВКИ $\exists!$ ОБРАТНАЯ

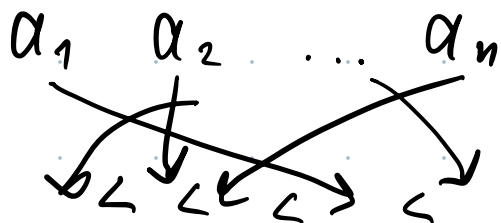
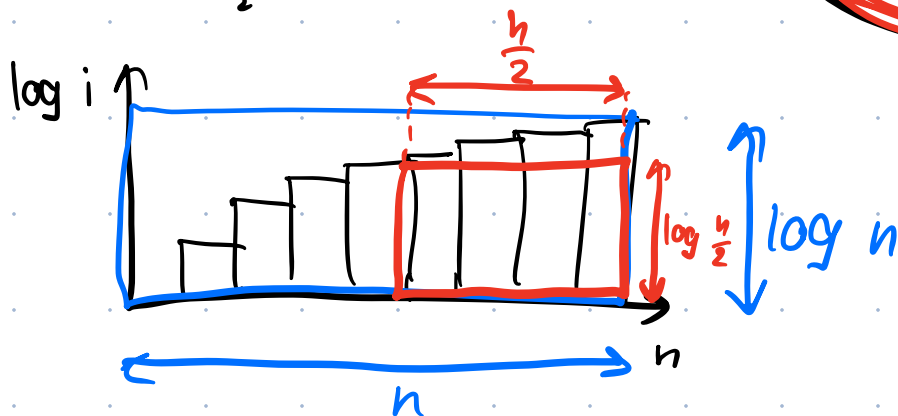


$$\log_2 2^k \geq \log_2 N$$
$$k \geq \lceil \log_2 N \rceil$$

$$\log_2 N = \log_2 n! = \sum_{i=1}^n \log_2 i \leq \sum_{i=1}^n \log_2 n = n \log_2 n = O(n \log n)$$

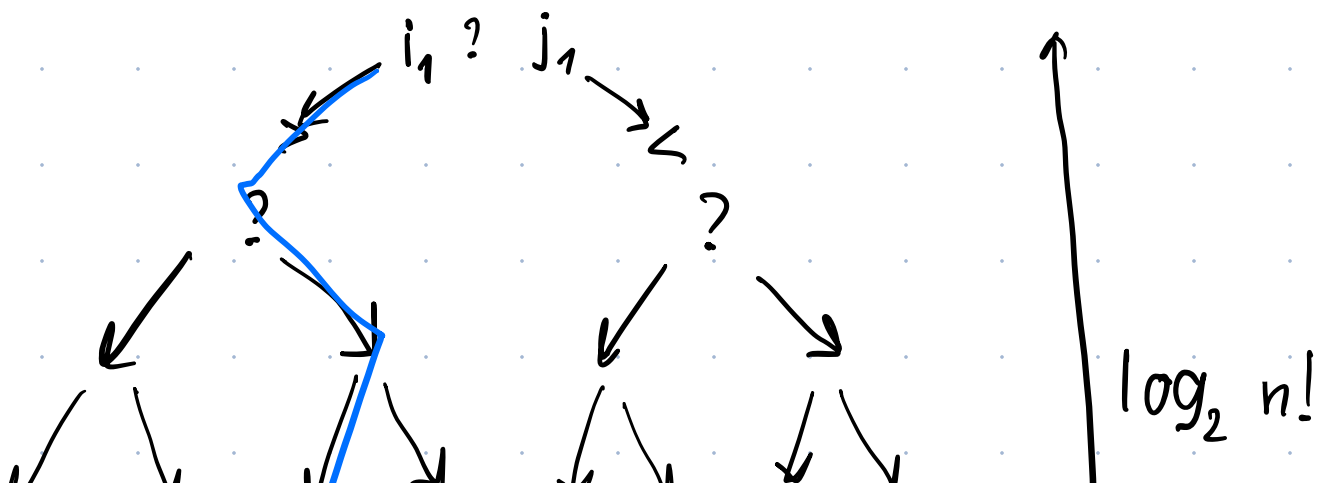
$$\log a \cdot b = \log a + \log b \quad \Rightarrow \sum_{i=\frac{n}{2}}^n \log_2 i \geq$$

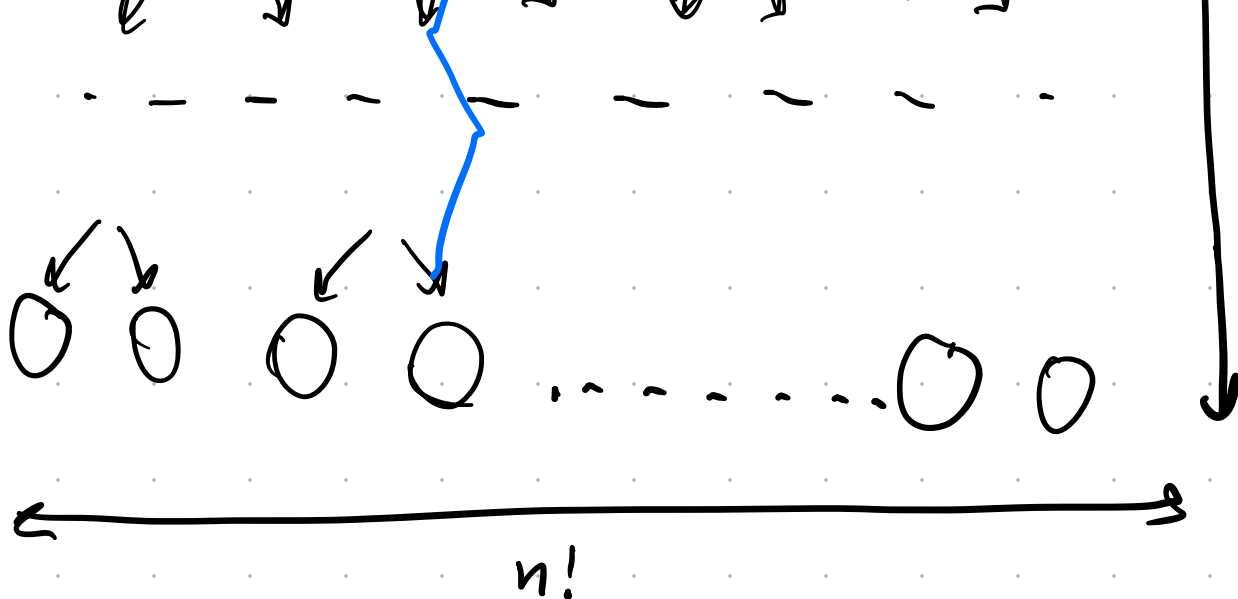
$$\geq \sum_{i=\frac{n}{2}}^n \log_2 \frac{n}{2} = \frac{n}{2} \log_2 \frac{n}{2} = \Omega(n \log n)$$



$\forall i \exists j$

$$\begin{pmatrix} 1 & 2 & 3 & 4 & \dots & n \\ p_1 & p_2 & p_3 & & & p_n \end{pmatrix}$$





`sorted_arr = sorted(arr, key = lambda i: i[0])`

$i[1]$

`arr = [(1, 3), (5, 8), (5, -14), ...]`

СОРТИРОВКИ НЕ ТОЛЬКО СРАВН

n
`a = [1, 1, 0, 1, ..., 1, 0]`

`zeros = 0`

`for i in range(len(a)):`
 `if (a[i] == 0):`
 `zeros += 1`

`for i in range(zeros):`
 `a[i] = 0`

`for i in range(zeros+1, len(a)):`
 `a[i] = 1`

0 0 0 1 0 1 1
`zeros = 5`

СОРТИРОВКА
ПОДСЧЁТОМ
(counting sort)

$\Theta(n)$ ПО ВРЕМЕНИ

$\Theta(1)$ ПО ПАМЯТИ

ОБОБЩИМ counting на int от 0 до M

```
f = [0 for _ in range(M+1)]
```

```
for i in range(len(a)):
```

```
if (a[i] == 0):  
    zeros += 1
```

```
    f[a[i]] += 1
```

```
ind = 0
```

```
for i in range(M+1):
```

```
    for j in range(f[i]):
```

```
        a[ind] = i
```

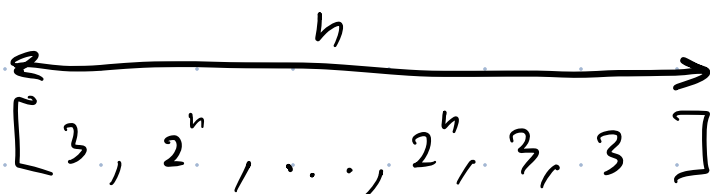
```
        ind += 1
```



ВРЕМЯ: $\Theta(n+M)$

ПАМЯТЬ: $\Theta(M)$

КОНТ Р ПРИМЕР:



$\Theta(n+2^n)$

