Optimizing Sparse Lower Triangular Solve Kei Imada

Optimizations. I developed and tested two potential optimizations to the given starter implementation of sparse lower triangular solve. The code for the optimizations are given in Figure 1. The first (Figure 1a) is a naive parallelization of the innermost loop. The second (Figure 1b) is an implementation of the level set partitioning described in [2] that finds blocks of parallelizable sections of the algorithm using a symbolic analysis of the input matrix.

```
template <typename T>
int parallel_lsolve(CCSMatrix<T> *matrix,
                                                                                int partitioned parallel lsolve(CCSMatrix<T> *matrix,
                                                                                                                      DenseVector<T> *input_vector,
                       DenseVector<T> *input vector) {
  if (!matrix || !input_vector) {
                                                                                                                      Partition *partition) {
    return 1:
                                                                                  if (!matrix || !input_vector || !partition) {
  int n = matrix->num_col_get();
                                                                                  int *Lp = matrix->column_pointer_get();
  int *Lp = matrix->column_pointer_get();
int *Li = matrix->row_index_get();
                                                                                  int *Li = matrix->row_index_get();
                                                                                  T *Lx = matrix->values_get();
T *x = input_vector->values_get();
  T *Lx = matrix->values_get();
  T *x = input_vector->values_get();
  int p, j;
if (!Lp || !Li || !x) {
   return 1; /* check inputs */
                                                                                   vector<vector<int>> partitioning = partition->partitioning_get();
                                                                                  if (!Lp || !Li || !x) {
  return 1; /* check inputs */
                                                                                  for (unsigned int i = 0; i < partitioning.size(); i++) {
  for (j = 0; j < n; j++) \{ // \text{ for every col in } L, \text{ row in } x \\ x[j] /= Lx[Lp[j]];
                                                                                vectorsint> partition = partitioning[i];

#pragma omp parallel default(shared) num_threads(8)
#pragma omp parallel default(shared) private(p) num_threads(8)
                                                                                #pragma omp for
#pragma omp for
    for (p = Lp[j] + 1; p < Lp[j + 1]; p++) {
    x[Li[p]] -= Lx[p] * x[j];
                                                                                     for (unsigned int part_idx = 0; part_idx < partition.size();</pre>
                                                                                          part_idx++) {
                                                                                       int j = partition[part_idx];
x[j] /= Lx[Lp[j]];
                                                                                       for (int p = Lp[j] + 1; p < Lp[j + 1]; p++) {
  return 0;
                                                                                        double tmp = Lx[p] * x[j];
int x_idx = Li[p];
                                                                                }
                                                                                    }
                                                                                  return 0;
                         (a) parallel
                                                                                                     (b) level partition
```

Figure 1: Two optimizations

Results. I ran the original and the two optimized algorithms on an 8-core Linux system. The full specification of the machine is listed in Figure 2. The code is available on my GitHub. As seen from Figure 3, the two optimizations did not produce a speedup from the original, but rather, exhibited slowdowns of least one magnitude. For both optimizations, it is possible that the given sparse matrices were too small to produce a speedup that surpassed OpenMP's thread bookkeeping overhead. Furthermore, C++ constructs such as vector were used in the level_partition optimization in order to store the partitioning. The overhead of using C++ libraries instead of faster C arrays may have contributed to the slowdown of the two optimizations.

```
OS: Ubuntu 18.04.5 LTS x86_64
Host: XPS 13 9370
Kernel: 5.4.0-62-generic
Shell: zsh 5.4.2
WM: i3
Theme: Adwaita [GTK3]
Icons: Adwaita [GTK3]
Terminal: alacritty
CPU: Intel i7-8550U (8 cores) @ 4.000GHz
GPU: Intel UHD Graphics 620
Total Memory: 15729MiB (16GB)
```

Figure 2: System specification

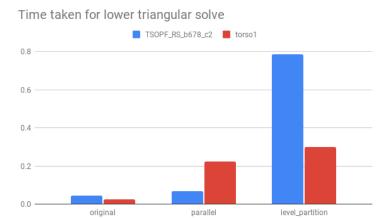


Figure 3: Time taken for sparse lower triangular solve

Development. I spent a nontrivial amount of time on code development. My unfamiliarity of available C++ libraries induced a longer development time. For example, I was inexperienced with C++ graph libraries such as those present in boost as well as unordered_map, and instead used unoptimized implementations from my college courses. I only used the libraries during the symbolic analysis of the level_partition optimization, hence this does not impact the solve runtime. Furthermore, my unfamiliarity with the lower triangular solve algorithm and the limited one week deadline contributed to the overall difficulty of this task. Nevertheless, I was able to fully understand the mechanism of this algorithm through this task and therefore more confident to produce further optimizations using parallel and distributed computing.

Future work. There are many improvements to this one week project. I can replace the C++ constructs in the optimizations with faster C arrays to further optimize the algorithm. Furthermore, implementing the h-level partitioning algorithm described in ParSy [1] may produce more reliable speedups compared to the naive optimizations. The symbolic analysis for the level_partition optimization is currently slow since the dependency graph is not pruned into a tree. Reading [2] further to understand the pruning algorithm should allow for a faster symbolic analysis of the input matrices.

- [1] K. Cheshmi, S. Kamil, M. M. Strout, and M. M. Dehnavi. *ParSy: Inspection and Transformation of Sparse Matrix Computations for Parallelism.* IEEE Press, 2018.
- [2] T. A. Davis, S. Rajamanickam, and W. M. Sid-Lakhdar. A survey of direct methods for sparse linear systems. *Acta Numerica*, 25:383–566, 2016.