# Abstracting Control Structures to Control the Application Flow



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# Imperative Control Structure

```
List<Product> products = order.getProducts();
Double discount = 0.0;
for (Product p : products) {
    if (this.isGiftProduct(p)) {
        discount = p.getPrice();
        break;
```

# Using High-order Functions



### Using High-order Functions and Ternary Operator



```
if (isGiftProduct(p)) {
    return p.getPrice() * 0.5;
} else {
    return p.getPrice();
}
```

```
Function<Product, Double> price = p -> {
    if (isGiftProduct(p)) {
        return p.getPrice() * 0.5;
    } else {
        return p.getPrice();
    }
}
```

# The Optional Type





#### optionalElement

```
.ifPresent(val -> doSomethingWithValue(val)),
.orElseGet(() -> doSomethingElse());
```



#### You Have to Decide if It Is Worth the Effort



Does this make the code more readable?

Is this an over-engineered implementation?



# Using Recursion Instead of Loops



# Loops Iterate over Lists

```
This is a list from 1 to 10

for (int i = 1; i <= 10; i++) {

// ...
}
```



Whatever you can do with recursion, you can also do it with iteration, and vice versa.



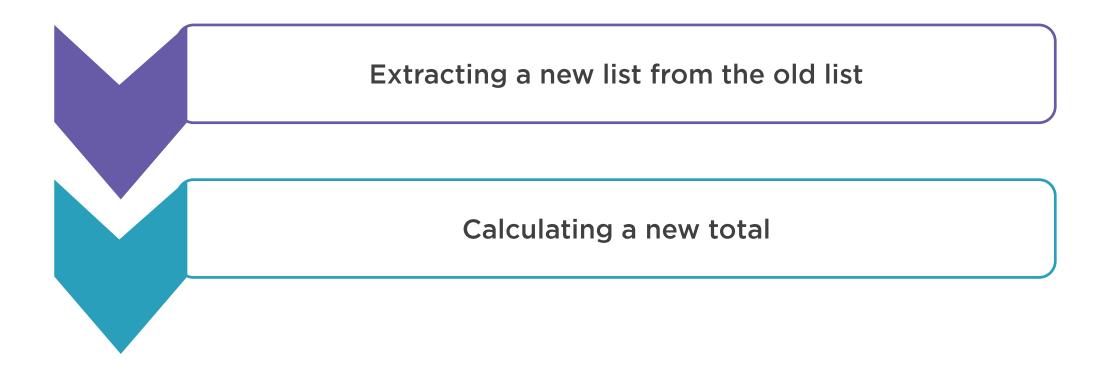
```
[1 2 3 4]
1 + [2 3 4]
```

```
head
        tail
[1 2 3 4]
 1 + [ 2 3 4 ]
      2+[34]
```

```
head
        tail
[1 2 3 4]
 1 + [ 2 3 4 ]
      2+[34]
          3+[4]
```

```
head
    tail
[1 2 3 4]
 1 + [ 2 3 4 ]
       2+[34]
          3 + [ 4 ] : Sum is zero
               4+[]
```

# Honoring Immutability

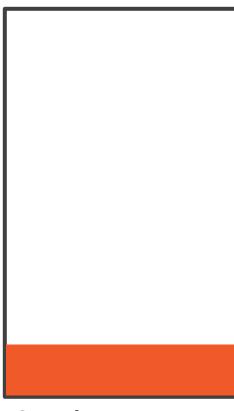




```
sum ([1, 2, 3, 4])
1 + sum ([2, 3, 4])
1+(2+sum([3, 4]))
1 + (2 + (3 + sum([4])))
1 + (2 + (3 + (4 + sum([]))))
1 + (2 + (3 + (4 + 0)))
1 + (2 + (3 + 4))
1 + (2 + 7)
1 + 9
10
```

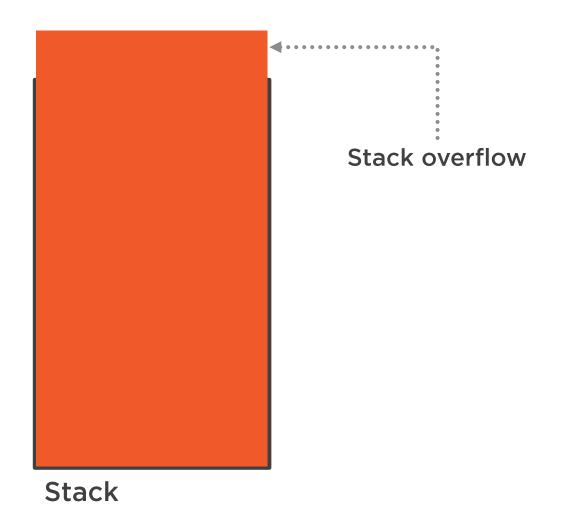
```
sum ([1, 2, 3, 4])
1 + sum ([2, 3, 4])
1+(2+sum([3, 4]))
1 + (2 + (3 + sum([4])))
1 + (2 + (3 + (4 + sum([]))))
1 + (2 + (3 + (4 + 0)))
1 + (2 + (3 + 4))
1 + (2 + 7)
1 + 9
10
```

sum([]) sum(4) sum(3, 4) sum(2, 3, 4) sum(1, 2, 3, 4) Stack



Stack







# Tail Recursion



```
int recursiveMethod(int i) {
    // ...
    return recursiveMethod(j);
}
```



#### Tail Recursion

```
public static void main(String args[]) {
   // ...
   int result = recursiveMethod(3);
  // ...
int recursiveMethod(int i) {
```



Stack



#### Tail-Recursive Function

```
sum ([1, 2, 3, 4], 0)
sum ([2, 3, 4], 1)
sum ([3, 4], 3)
sum ([4], 6)
sum ([], 10)
```



#### With and Without Tail Recursion

#### With tail recursion

```
sum ([1, 2, 3, 4], 0)
sum ([2, 3, 4], 1)
sum ([3, 4], 3)
sum ([4], 6)
sum ([], 10)
```

#### Without tail recursion

```
sum ([1, 2, 3, 4])
1 + sum ([2, 3, 4])
1+(2+sum([3, 4]))
1 + (2 + (3 + sum([4])))
1 + (2 + (3 + (4 + sum([]))))
1 + (2 + (3 + (4 + 0)))
1 + (2 + (3 + 4))
1 + (2 + 7)
```

#### To Create a Tail-recursive Function

Create a private recursive function with an additional accumulator parameter

The base case of the recursive function returns the accumulator

The recursive invocation provides an updated value for the accumulator

Create a public function that calls the tail-recursive function using the appropriate initial values



## In Summary

Non-tail recursive functions will use the stack to remember the state

Tail recursive functions use accumulators to remember state



# Tail Call Optimization (TCO)

When the compiler automatically make tail-recursive functions more efficient.



# Tail Call Optimization with Trampolines



# Thunk

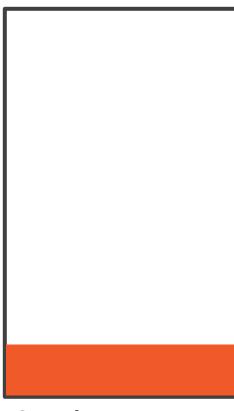
A function that is returned by another function to delay a computation.



# Implemented with the Supplier Interface

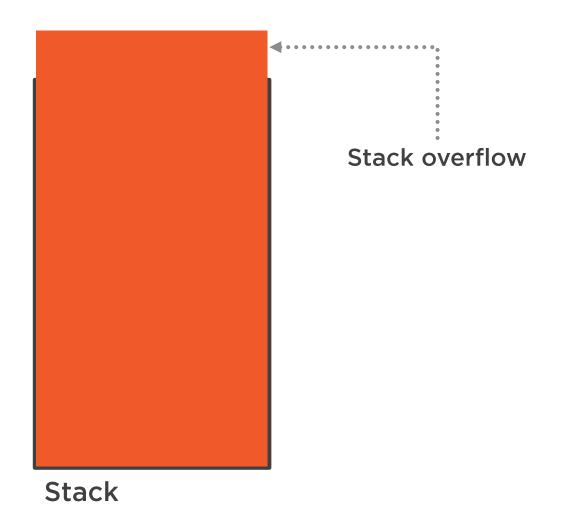
```
@FunctionalInterface
public interface Supplier<T> {
    T get();
}
```





Stack







```
int recursiveMethod(int i) {
    // ...
return () -> recursiveMethod(j);
}
```

```
int recursiveMethod(int i) {
    // ...
    return () -> recursiveMethod(j);
}
```

recursiveMethod

Stack



# Trampoline

A loop that iteratively invokes a function that can return a thunk.



## Trampoline Pseudocode

```
while (f == 'function') {
    f = f();
}
return f;
```



## The Fold Operation



```
Integer sum(List<Integer> list, int acc) {
  return list.isEmpty()
      ? acc
      : sum( tail(list), acc + head(list) );
}
```

```
Integer sum(List<Integer> list, int acc) {
  return list.isEmpty()
      ? acc
      : sum( tail(list), acc + 1 );
}
```





### Higher-level of Abstraction



#### The Function Takes Two Values

The head of the list

The accumulator



## Type of the Function

$$(T,U)\rightarrow U$$



## Type of the Function (Curried Version)

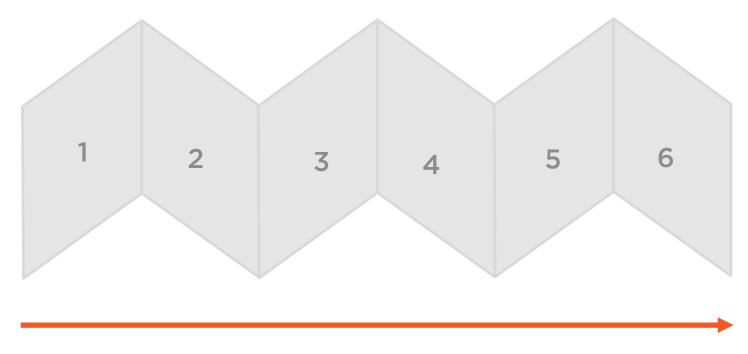
$$T \rightarrow U \rightarrow U$$



```
Function<Integer, Function<Integer, Integer>> sum = x \rightarrow y \rightarrow x + y;
Integer generic(List<Integer> list,
                 int acc,
                 Function<Integer, Function<Integer, Integer>> f) {
   return list.isEmpty()
            ? acc
            : generic(tail(list), f.apply(head(list)).apply(acc), f);
```

```
Function<Integer, Function<Integer, Integer>> sum = x \rightarrow y \rightarrow x + y;
Integer foldLeft(List<Integer> list,
                 int acc,
                 Function<Integer, Function<Integer, Integer>> f) {
   return list.isEmpty()
           ? acc
            : foldLeft(tail(list), f.apply(head(list)).apply(acc), f);
```

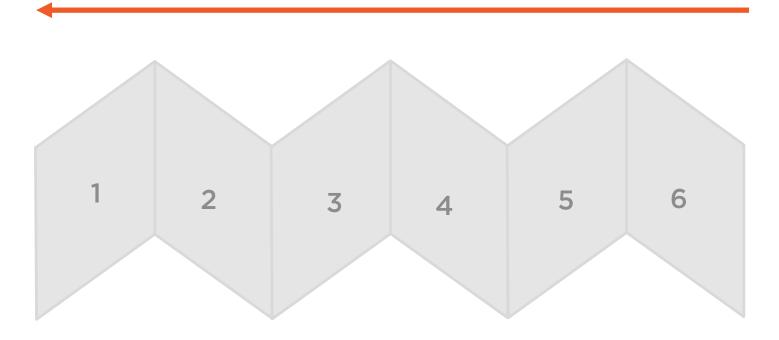




foldLeft



#### foldRight





## Folding in Two Directions

If the operation is commutative, both ways of folding are equivalent

If the operation is not commutative, the two ways of folding give different results





Is Stream.reduce() equivalent to foldLeft?



## Folding Versus Mapping, Reducing, and Collecting





Is Stream.reduce() equivalent to foldLeft?

NO



# Reducing

Folding to a result that is the same type as the list elements.



#### FoldLeft



#### First Version of Reduce

Optional<T> reduce(BinaryOperator<T> accumulator)



#### First Version of Reduce

```
Optional<T> reduce(BinaryOperator<T> accumulator) {
     boolean foundAny = false;
     T result = null;
     for (T element : this stream) {
         if (!foundAny) {
             foundAny = true;
             result = element;
         else
             result = accumulator.apply(result, element);
     return foundAny ? Optional.of(result) : Optional.empty();
```

#### Second Version of Reduce

```
T reduce(T identity, BinaryOperator<T> accumulator)
```



#### Second Version of Reduce

```
T reduce(T identity, BinaryOperator<T> accumulator) {
    T result = identity;
    for (T element : this stream)
        result = accumulator.apply(result, element)
    return result;
}
```



#### Third Version of Reduce

```
U reduce(U identity, BiFunction<U, ? super T, U> accumulator, BinaryOperator<U> combiner)
```



#### Third Version of Reduce

```
U reduce(U identity, BiFunction<U, ? super T, U> accumulator, BinaryOperator<U> combiner) {
    U result = identity;
    for (T element : this stream)
        result = accumulator.apply(result, element)
    return result;
}
```



#### Collect

```
R collect(Supplier<R> supplier, BiConsumer<R,? super T> accumulator, BiConsumer<R,R> combiner)
{
    R result = supplier.get();
    for (T element : this stream)
        accumulator.accept(result, element);
    return result;
}
```



Reduce is designed to work with immutable objects, while collect can only work with mutable objects.



## Memoization



## Memoization

A technique that caches the result of an operation so it can be returned immediately if the same computation is performed in the future.



## Memoization Optimizes Execution Time

$$O(n) \rightarrow O(1)$$



#### Memoization Is a Trade





Reduces execution time

Increases memory



### Fibonacci Numbers

$$F_n = F_{n-1} + F_{n-2}$$

Except for the first two numbers that are always 0 and 1.



### Fib(4)

```
fib(4)

fib(3) + fib(2)

fib(2) + fib(1) + fib(1) + fib(0)
```



n	Fib(n)	Number of calls
0	1	1
1	1	1
2	2	3
3	3	5
4	5	9
5	8	15
6	13	25
7	21	41
8	34	67
9	55	109
10	89	177
11	144	287
12	233	465
13	377	753
14	610	1219
15	987	1973



Memoization is often an alternative to tail call optimization to improve performance.



### Memoization Is Compatible with Functional Programming

A memoized function always returns the same value for the same argument

If the side effect of storing the results is not visible from outside the function



```
public Double average(int number) {
    return IntStream.rangeClosed(1, number).average().orElse(0.0);
}

// ...

Function<Integer, Double> avg = this::average;
```

```
public Double average(int number) {
   return IntStream.rangeClosed(1, number).average().orElse(0.0);
// ...
Map<Integer, Double> cache = new ConcurrentHashMap<>();
Function<Integer, Double> avg =
                    x -> cache.computeIfAbsent(x, this::average);
```



For simple use cases, you can use the Stream API

For more advanced use cases, you must implement your own types





We can replace loops with recursive functions

Be careful with stack overflow exceptions

Write tail-recursive functions, where the recursive call is the last line of the function

- Tail-recursive functions use an accumulator to carry intermediate state
- Tail call optimization (TCO)
- Implement it with thunks and trampolines





#### Fold function

- Tail-recursive function for performing operations over collections
- Two directions, from left to right or from right to left
- If the operation is commutative, both ways of folding are equivalent
- Java doesn't have an equivalent of this function





#### Memoization

- Caching the result of an operation so it can be returned immediately if the same computation is performed in the future

Using imperative structures inside a pure function is not bad if the function remains pure



### In the Next Module

### Dealing with nulls functionally

