

CISC304: Homework 6

Due on December 2, 11:00 am

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Problem 1

Use the Hilbert Deductive System for First Order Logic (aka Predicate Calc.) to prove the following. You may use the Deductive Theorem, and other theorems proved in class or in previous handouts.

- a. (10%) Prove that: $\forall_x \forall_y A \rightarrow \forall_y \forall_x A$
- b. (10%) Add to the Hilbert Proof System the generalization Rule version of Ax. 6:

$$\frac{U \vdash B(a)}{U \vdash \forall_x B(x)}$$

where U is a set of WFFs (can be empty), B is a WFF, and a does not appear in U . For a WFF A that does not contain x as a free variable, prove that:

$$\forall_x (A \rightarrow B(x)) \leftrightarrow (A \rightarrow \forall_x B(x)).$$

Solution 1a.

Prove that: $\forall_x \forall_y A \rightarrow \forall_y \forall_x A$

1. $A \rightarrow \forall_x A$ (Axiom 6)
2. $\forall_x A \rightarrow \forall_y \forall_x A$ (Axiom 6)
3. $\forall_x \forall_y A \rightarrow \forall_y A$ (Axiom 4 with $t = y$)
4. $(A \rightarrow \forall_x A) \rightarrow (\forall_y (A \rightarrow \forall_x A) \rightarrow (A \rightarrow \forall_x A))$ (Axiom 1)
5. $\forall_y (A \rightarrow \forall_x A) \rightarrow (A \rightarrow \forall_x A)$ (MP, 1, 4)
6. $\forall_y (A \rightarrow \forall_x A)^1$ (A, $(\forall_x A) \rightarrow C \vdash \forall_x C$ In Class)
7. $\forall_y A \rightarrow \forall_y \forall_x A$ (Axiom 5)
8. $\forall_x \forall_y A \rightarrow \forall_y \forall_x A$ (Syllogism Rule 3, 7)

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¹Using Axiom 6 seems more trivial here, but Axiom 6 will give us $\forall_y A \rightarrow \forall_x \forall_y A$. We want to get $\forall_y A \rightarrow \forall_y \forall_x A$.

Problem 2

Represent the following four premises and the conclusion in Propositional Calculus. *Employ in your representation only the four propositional variables indicated below.* Then use resolution to prove that the conclusion indeed soundly follows from the premises. Clearly show/explain all steps involved.

Premises:

Politicians attend the quarterly meetings.

Non-politicians attend the annual meetings.

If one attends the quarterly or the annual meetings one is invited to the symposium.

Conclusion: *One gets invited to the Symposium*

Propositional Variables:

Politicians: P

Attending Quarterly Meetings: Q

Attending Annual Meetings: A

Being invited to the Symposium: S

Solution

Premises (in respective order): $P \rightarrow Q$, $\neg P \rightarrow A$, $(Q \vee A) \rightarrow S$

Conclusion: S

To prove that the conclusion soundly follows from the premise, we must show that

$$(P \rightarrow Q) \wedge (\neg P \rightarrow A) \wedge ((Q \vee A) \rightarrow S) \wedge \neg S$$

is unsatisfiable.

1. $(\neg P \vee Q) \wedge (P \vee A) \wedge (\neg(Q \vee A) \vee S) \wedge \neg S$ (Break \rightarrow 's)
2. $(\neg P \vee Q) \wedge (P \vee A) \wedge ((\neg Q \wedge \neg A) \vee S) \wedge \neg S$ (Apply De-Morgan's Law)
3. $(\neg P \vee Q) \wedge (P \vee A) \wedge (S \vee \neg Q) \wedge (S \vee \neg A) \wedge \neg S$ (Distribute until CNF)
4. $\{\{\neg P, Q\}, \{P, A\}, \{S, \neg Q\}, \{S, \neg A\}, \{\neg S\}\}$ (Write in Clausal Normal Form)

Resolve using the Davis-Putnam procedure:

$\{\{\neg P, Q\}_1, \{P, A\}_2, \{S, \neg Q\}_3, \{S, \neg A\}_4, \{\neg S\}_5\}$

1. $\{\{\neg P, Q\}_1, \{P, A\}_2, \{\neg Q\}_{35}, \{\neg A\}_{45}\}$ (Resolve $S < 3, 5 >, < 4, 5 >$)
2. $\{\{\neg P, Q\}_1, \{P, A\}_2, \{\neg Q\}_{35}, \{\neg A\}_{45}\}$ (Resolve $Q < 1, 35 >$)
3. $\{\{\neg P\}_1, \{P, A\}_2, \{\neg A\}_3\}$ (Resolve $P < 1, 2 >$)
4. $\{\{A\}_{12}, \{\neg A\}_3\}$ (Resolve $A < 12, 3 >$)
5. $\{\square\}$

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Problem 3

Use *Resolution*, utilizing the **Davis-Putnam procedure**, to show that the following set of propositional clauses is unsatisfiable. Show every step of the procedure and clearly indicate which clauses and literals are being resolved, eliminated and/or added in each step.

$$S = \{\{p, \neg r\}_1; \{q, \neg r\}_2; \{q, \neg s\}_3; \{\neg p, o\}_4; \{\neg q, \neg o\}_5; \{\neg q, r, o\}_6; \{p, s, \neg o\}_7; \{\neg p, q, r\}_8; \{q, r, s, o\}_9\}$$

Solution

1. No trivial or literal clauses found.

2. Resolve for p by resolving clauses $\langle 1, 4 \rangle, \langle 1, 8 \rangle, \langle 7, 4 \rangle, \langle 7, 8 \rangle$:

$$\{\{\neg r, o\}_{14}; \{\neg r, q, r\}_{18}; \{s, \neg o, o\}_{74}; \{s, \neg o, q, r\}_{78}; \{q, \neg r\}_2; \{q, \neg s\}_3; \{\neg q, \neg o\}_5; \{\neg q, r, o\}_6; \{q, r, s, o\}_9\}$$

3. Delete trivial clauses 18, 74

$$\{\{\neg r, o\}_{14}; \{s, \neg o, q, r\}_{78}; \{q, \neg r\}_2; \{q, \neg s\}_3; \{\neg q, \neg o\}_5; \{\neg q, r, o\}_6; \{q, r, s, o\}_9\}$$

4. Resolve for r by resolving clauses $\langle 78, 14 \rangle, \langle 78, 2 \rangle, \langle 6, 14 \rangle, \langle 6, 2 \rangle, \langle 9, 14 \rangle, \langle 9, 2 \rangle$:

$$\{\{o, \neg o, s, q\}_{7814}; \{s, \neg o, q\}_{782}; \{\neg q, o\}_{614}; \{\neg q, o, q\}_{62}; \{q, s, o\}_{914}; \{q, s, o, \neg s\}_{92}; \{q, \neg s\}_3; \{\neg q, \neg o\}_5\}$$

5. Delete trivial clauses 7814, 62, 92:

$$\{\{s, \neg o, q\}_1; \{\neg q, o\}_2; \{q, s, o\}_3; \{q, \neg s\}_4; \{\neg q, \neg o\}_5\}$$

6. Resolve for o by resolving clauses $\langle 2, 1 \rangle, \langle 2, 5 \rangle, \langle 3, 1 \rangle, \langle 3, 5 \rangle$:

$$\{\{s, q, \neg q\}_{21}; \{\neg q\}_{25}; \{q, s\}_{31}; \{q, s, \neg q\}_{35}; \{q, \neg s\}_4\}$$

7. Delete trivial clauses 21 and 35:

$$\{\{\neg q\}_1; \{q, s\}_2; \{q, \neg s\}_3\}$$

8. Resolve for q by resolving clauses $\langle 2, 1 \rangle, \langle 3, 1 \rangle$

$$\{\{s\}_{21}; \{\neg s\}_{31}\}$$

9. Resolve for s by resolving clauses $\langle 21, 31 \rangle$

$$\{\square\}$$

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Problem 4

Transform each of the following formulae into **Clausal Form**. Clearly indicate the main stages of the transformation.

- $\forall_x (P(x) \vee Q(x)) \rightarrow (\forall_x P(x) \vee \exists_x Q(x))$
- (from Ben Ari 9.1) $\forall_x \forall_y (\exists_z P(z) \wedge \exists_u (Q(x, u) \rightarrow \exists_v Q(y, v)))$
- (from Ben Ari 9.1) $\exists_x (\neg \exists_y P(y) \rightarrow \exists_z (Q(z) \rightarrow R(x)))$

Solution 4a.

$$\forall_x (P(x) \vee Q(x)) \rightarrow (\forall_x P(x) \vee \exists_x Q(x))$$

- Rename bound variables

$$\forall_x (P(x) \vee Q(x)) \rightarrow (\forall_y P(y) \vee \exists_z Q(z))$$

- Eliminate all binary boolean operators other than \vee and \wedge

$$\neg \forall_x (P(x) \vee Q(x)) \vee (\forall_y P(y) \vee \exists_z Q(z))$$

- Push negations inwards until they apply to atomic formulae only

$$\exists_x \neg (P(x) \vee Q(x)) \vee (\forall_y P(y) \vee \exists_z Q(z)) \quad (\neg \forall_x A(x) \equiv \exists_x \neg A(x))$$

$$\exists_x (\neg P(x) \wedge \neg Q(x)) \vee (\forall_y P(y) \vee \exists_z Q(z)) \quad (\text{De-Morgan's Law})$$

- Extract quantifiers from the matrix into the prefix

$$\exists_x \forall_y ((\neg P(x) \wedge \neg Q(x)) \vee P(y) \vee \exists_z Q(z))$$

$$\exists_x \forall_y \exists_z ((\neg P(x) \wedge \neg Q(x)) \vee P(y) \vee Q(z))$$

- Convert the matrix into conjunctive normal form.

$$\exists_x \forall_y \exists_z ((\neg P(x) \vee P(y) \vee Q(z)) \wedge (\neg Q(x) \vee P(y) \vee Q(z)))$$

- Remove existential quantifiers by *Skolemization*.

$$\forall_y \exists_z ((\neg P(a) \vee P(y) \vee Q(z)) \wedge (\neg Q(a) \vee P(y) \vee Q(z))) \quad (\text{Replaced } x \text{ with } a)$$

$$\forall_y ((\neg P(a) \vee P(y) \vee Q(f(y))) \wedge (\neg Q(a) \vee P(y) \vee Q(f(y)))) \quad (\text{Replaced } z \text{ with } f(y))$$

- Drop universal quantifiers, and write it in clausal normal form

$$\{\{\neg P(a), P(y), Q(f(y))\}, \{\neg Q(a), P(y), Q(f(y))\}\}$$

Solution 4b.

$$\forall_x \forall_y (\exists_z P(z) \wedge \exists_u (Q(x, u) \rightarrow \exists_v Q(y, v)))$$

1. Rename bound variables
(no change)
2. Eliminate all binary boolean operators other than \vee and \wedge
 $\forall_x \forall_y (\exists_z P(z) \wedge \exists_u (\neg Q(x, u) \vee \exists_v Q(y, v)))$
3. Push negations inwards until they apply to atomic formulae only
(no change)
4. Extract quantifiers from the matrix into the prefix
 $\forall_x \forall_y \exists_z (P(z) \wedge \exists_u (\neg Q(x, u) \vee \exists_v Q(y, v)))$
 $\forall_x \forall_y \exists_z \exists_u (P(z) \wedge (\neg Q(x, u) \vee \exists_v Q(y, v)))$
 $\forall_x \forall_y \exists_z \exists_u \exists_v (P(z) \wedge (\neg Q(x, u) \vee Q(y, v)))$
5. Convert the matrix into conjunctive normal form.
(already in conjunctive normal form)
5. Remove existential quantifiers by *Skolemization*.
 $\forall_x \forall_y \exists_u \exists_v (P(f(x, y)) \wedge (\neg Q(x, u) \vee Q(y, v)))$ (Replaced z with $f(x, y)$)
 $\forall_x \forall_y \exists_v (P(f(x, y)) \wedge (\neg Q(x, g(x, y)) \vee Q(y, v)))$ (Replaced u with $g(x, y)$)
 $\forall_x \forall_y (P(f(x, y)) \wedge (\neg Q(x, g(x, y)) \vee Q(y, h(x, y))))$ (Replaced u with $h(x, y)$)
6. Drop universal quantifiers, and write it in clausal normal form
 $\{\{P(f(x, y))\}, \{\neg Q(x, g(x, y)), Q(y, h(x, y))\}\}$

Solution 4c.

$$\exists_x (\neg \exists_y P(y) \rightarrow \exists_z (Q(z) \rightarrow R(x)))$$

1. Rename bound variables so that no variables appear in two quantifiers
(No changes)

2. Eliminate all binary boolean operators other than \vee and \wedge

$$\exists_x (\exists_y P(y) \vee \exists_z (Q(z) \rightarrow R(x))) \quad (\text{Elim } \rightarrow)$$

$$\exists_x (\exists_y P(y) \vee \exists_z (\neg Q(z) \vee R(x))) \quad (\text{Elim } \rightarrow)$$

3. Push negations inwards until they apply to atomic formulae only
(No changes)

4. Extract quantifiers from the matrix into the prefix

$$\exists_x \exists_y (P(y) \vee \exists_z (\neg Q(z) \vee R(x))) \quad (\text{Extracted } \exists_y)$$

$$\exists_x \exists_y \exists_z (P(y) \vee (\neg Q(z) \vee R(x))) \quad (\text{Extracted } \exists_z)$$

5. Convert the matrix into conjunctive normal form.
(No changes)

5. Remove existential quantifiers by *Skolemization*.

$$\exists_y \exists_z (P(y) \vee (\neg Q(z) \vee R(a))) \quad (\text{Replaced } x \text{ with } a)$$

$$\exists_z (P(b) \vee (\neg Q(z) \vee R(a))) \quad (\text{Replaced } y \text{ with } b)$$

$$P(b) \vee \neg Q(c) \vee R(a) \quad (\text{Replaced } z \text{ with } c)$$

6. Drop universal quantifiers, and write it in clausal normal form
 $\{\{P(b), \neg Q(c), R(a)\}\}$

Problem 6

Translate the three premises and conclusion into sentences in the Predicate Calculus, using the predicate symbols provided below. Then use resolution to prove that the conclusion indeed follows from the premises.

$W(x)$ x is a walrus; $C(x)$ x is a carpenter; $B(x)$ x lives in the bay; $M(x)$ x likes the mountains

Premises *No Walrus Likes the Mountains. All carpenters likes the Mountains. Everyone who lives in the bay is a walrus.*

Conclusion: *Those who live in the bay are not carpenters*

Solution

No Walrus Likes the Mountains.

$$\forall x \neg(W(x) \rightarrow M(x))$$

All carpenters likes the Mountains.

$$\forall x (C(x) \rightarrow M(x))$$

Everyone who lives in the bay is a walrus.

$$\forall x (B(x) \rightarrow W(x))$$

Problem 7

Use resolution to prove that the following formulae are valid. Show all stages of your reasoning.

- a. $\exists x \forall y (R(y, x, g(y, x))) \rightarrow \exists x \forall y \exists z (R(y, x, z))$
- b. $\forall x (P(x) \rightarrow Q(x)) \models \forall x (\exists y (P(y) \wedge R(x, y)) \rightarrow \exists y (Q(y) \wedge R(x, y)))$

Solution 7a.

We must show that $\neg(\exists x \forall y (R(y, x, g(y, x))) \rightarrow \exists x \forall y \exists z (R(y, x, z)))$ is unsatisfiable.

1. $\neg(\exists x \forall y (R(y, x, g(y, x))) \rightarrow \exists w \forall u \exists v (R(u, w, v)))$ (Rename variables)
2. $\exists x \forall y (R(y, x, g(y, x))) \wedge \neg \exists w \forall u \exists v (R(u, w, v))$ (Break $\neg(A \rightarrow B)$)
3. $\exists x \forall y (R(y, x, g(y, x))) \wedge \forall w \exists u \forall v (\neg R(u, w, v))$ (Push negations in)
4. $\exists x \forall y \forall w \exists u \forall v (R(y, x, g(y, x)) \wedge \neg R(u, w, v))$ (Pull quantifiers out to prefix)
5. $\forall y \forall w \forall v (R(y, a, g(y, a)) \wedge \neg R(f(y, w), w, v))$ (Skolemize: $x = a, u = f(y, w)$)
6. $\{\{R(y, a, g(y, a))\}_1, \{\neg R(f(y, w), w, v)\}_2\}$ (Delete \forall 's, write in clausal form)
7. Unify with $\theta = [y/f(y, w), a/w, v/g(y, a)]$:
 $\{\{R(f(y, w), w, g(y, a))\}_1, \{\neg R(f(y, w), w, g(y, a))\}_2\}$
8. $\{\square\}$ (Resolve 1, 2)

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Solution 7b.

To show that

$$\forall x(P(x) \rightarrow Q(x)) \models \forall x(\exists y(P(y) \wedge R(x, y)) \rightarrow \exists y(Q(y) \wedge R(x, y)))$$

is valid, we must show that

$$\forall x(P(x) \rightarrow Q(x)) \wedge \neg \forall x(\exists y(P(y) \wedge R(x, y)) \rightarrow \exists y(Q(y) \wedge R(x, y)))$$

is unsatisfiable.

1. Rename variables that appear in more than one quantifiers

$$\forall x(P(x) \rightarrow Q(x)) \wedge \neg \forall u(\exists y(P(y) \wedge R(u, y)) \rightarrow \exists v(Q(v) \wedge R(u, v)))$$

2. Break \rightarrow 's

$$\forall x(\neg P(x) \vee Q(x)) \wedge \neg \forall u(\neg \exists y(P(y) \wedge R(u, y)) \vee \exists v(Q(v) \wedge R(u, v)))$$

3. Push in negations

$$\begin{aligned} & \forall x(\neg P(x) \vee Q(x)) \wedge \exists u \neg(\neg \exists y(P(y) \wedge R(u, y)) \vee \exists v(Q(v) \wedge R(u, v))) \\ & \forall x(\neg P(x) \vee Q(x)) \wedge \exists u \neg(\forall y \neg(P(y) \wedge R(u, y)) \vee \exists v(Q(v) \wedge R(u, v))) \\ & \forall x(\neg P(x) \vee Q(x)) \wedge \exists u(\neg \forall y \neg(P(y) \wedge R(u, y)) \wedge \neg \exists v(Q(v) \wedge R(u, v))) \\ & \forall x(\neg P(x) \vee Q(x)) \wedge \exists u(\exists y(P(y) \wedge R(u, y)) \wedge \neg \exists v(Q(v) \wedge R(u, v))) \\ & \forall x(\neg P(x) \vee Q(x)) \wedge \exists u(\exists y(P(y) \wedge R(u, y)) \wedge \forall v \neg(Q(v) \wedge R(u, v))) \\ & \forall x(\neg P(x) \vee Q(x)) \wedge \exists u(\exists y(P(y) \wedge R(u, y)) \wedge \forall v(\neg Q(v) \vee \neg R(u, v))) \end{aligned}$$

4. Extract quantifiers out into prefix

$$\forall x \exists u \exists y \forall v ((\neg P(x) \vee Q(x)) \wedge (P(y) \wedge R(u, y)) \wedge (\neg Q(v) \vee \neg R(u, v)))$$

5. Skolemize to remove existential quantifiers: $u = f(x), y = g(x), v = a$

$$\forall x((\neg P(x) \vee Q(x)) \wedge (P(g(x)) \wedge R(f(x), g(x))) \wedge (\neg Q(a) \vee \neg R(f(x), a)))$$

6. Delete \forall 's, and write in clausal normal form

$$\{\{\neg P(x), Q(x)\}_1, \{P(g(x)), R(f(x), g(x))\}_2, \{\neg Q(a), \neg R(f(x), a)\}_3\}$$

7. Run Resolution to determine unsatisfiability

$$\begin{aligned} & \{P(g(x)), \neg Q(g(x))\}_4 && (\text{Resolve R} < 2, 3 > \text{sub. } [a/g(x)]) \\ & \{\neg P(g(x)), P(g(x))\}_5 && (\text{Resolve Q} < 1, 4 > \text{sub. } [a/g(x)]) \\ & \{\square\}_5 && (\text{Resolve P} < 1, 5 >) \end{aligned}$$

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