



Optimization Theory and Methods

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- Multi-Stage Game With Perfect Information
- Two-Stage Quantity Competition
- Empty Threats
- Backward Induction
- Subgame Perfect PSNE
- Strategic Investment Example

↳ Multi-Stage Game with Perfect Information

- A multi-stage game with perfect information is one where:
 - 1) All players know the actions chosen at all previous stages $0, 1, 2, \dots, k - 1$, when choosing their actions at stage k .
 - 2) In each stage k , some players take an action.
 - 3) All players who take an action in a particular stage take actions simultaneously in that stage.
- Thus, when a player takes an action in stage k , it is aware of all the actions taken by all players in the previous stages, but is not aware of the actions taken by other players in stage k .
- **Important:** A pure strategy is no longer the same as an action.
- A pure strategy is now a full contingency plan for all possible combinations of actions by all players at all previous stages.

- ❑ Consider a 2 player game. Similar to the Cournot case, the actions by the players are the choices of the output levels, q_1 for player 1 and q_2 for player 2.
- We assume that at first player 1 chooses q_1 , then player 2 chooses q_2 .
- Thus player 2 knows the value of q_1 before deciding q_2 . So player 1's strategy is just a value of q_1 , while player 2's strategy is a mapping giving one value of q_2 for each value of q_1 .
- Let $c = 0$ (for simplicity) and $p(q) = 12 - q$. So player i 's payoff is $u_i(q_1, q_2) = [12 - (q_1 + q_2)] * q_i$.
- So a PSNE of this game should include one number, q_1 , and a mapping that gives **one value of q_2 for each value of q_1** .
- How should we find this PSNE?

- The easy part is to compute the best q_2 value for each q_1 value.
- First order conditions give: $0 = \frac{\partial}{\partial q_2} [u_2(q_1, q_2)] = \frac{\partial}{\partial q_2} [[12 - (q_1 + q_2)]q_2] = 12 - q_1 - 2q_2$.
- So, $q_2 = 6 - \frac{q_1}{2}$. Thus, for any value of q_1 chosen by player 1, player 2 will observe that and accordingly pick a q_2 value equal to $6 - \frac{q_1}{2}$.
- So the equilibrium payoff of player 1 is given by: $u_1(q_1, q_2) = [12 - (q_1 + q_2)] * q_1 = \left[12 - \left(q_1 + 6 - \frac{q_1}{2}\right)\right] * q_1 = \left[6 - \frac{q_1}{2}\right] * q_1$.
- So the first order conditions give: $0 = \frac{\partial}{\partial q_1} \left[\left[6 - \frac{q_1}{2}\right] * q_1\right] = 6 - q_1$.
- So the PSNE is given by: **$q_1^* = 6$ and $q_2^*(q_1) = 6 - \frac{q_1}{2}$** .
- This PSNE is called a *Stackelberg Equilibrium*.

↳ Other PSNEs of 2-Stage Quantity Competition Game

- However, there are other PSNEs of this game!!
- For example, what if player 2 decides to have a strategy where it will pick a constant value $q_2 = q_2^C$ for any value q_1 picked by player 1. In other words, player 2 does not change its strategy in response to the observed value of q_1 .
- In this case, $u_1(q_1, q_2) = [12 - (q_1 + q_2)] * q_1$. So 1st order conditions for player 1 give $0 = \frac{\partial}{\partial q_1} [u_1(q_1, q_2)] = \frac{\partial}{\partial q_1} [[12 - (q_1 + q_2)] * q_1] = 12 - 2q_1 - q_2$.
- Since q_2 does not depend on q_1 , we get, $0 = \frac{\partial}{\partial q_2} [u_2(q_1, q_2)] = \frac{\partial}{\partial q_2} [[12 - (q_1 + q_2)] * q_2] = 12 - q_1 - 2q_2$.
- So, we get $q_1^* = 6 - \frac{q_2^C}{2}$ and $q_2^* = 6 - \frac{q_1^C}{2}$. Solving these simultaneously, we get $q_1^* = q_2^* = 4$.

- In this new PSNE, player 1 assumes that player 2 will actually pick the exact same value of output q_2 even if player 1 picks a different value of q_1 (other than 4).
- Under this assumption, the best strategy for player 1 is to pick output $q_1 = 4$. But this is fully reliant on the assumption that player 1 actually believes that if it produces another level of output (e.g., $q_1 = 6$), player 2 will still play $q_2 = 4$ rather than playing $q_2 = 3$ which is the optimal response of player 2.
- In other words, player 2 threatens to player 1 that it will produce 4 units even though it is not optimal to player 2. This is what we call as empty threat. Note that $q_2 = 4$ is optimal for player 2 only when player 1 produces $q_1 = 4$. For all other levels of q_1 , it is suboptimal.
- The PSNE on the previous slide assumes that player 1 will produce $q_1 = 4$ because it believes in player 2's threat.

- By assuming that player 1 will produce $q_1 = 4$, the response by player 2 to other levels of player 1 output does not matter.
- Yet, the very assumption that player 1 will produce $q_1 = 4$ is dependent on player 2 sticking to the strategy of always producing $q_2 = 4$ no matter what q_1 value is.
- As a result of the empty threat, neither player has a reason to deviate and hence this is a PSNE.
- However, this PSNE is highly suspect in terms of its portrayal of reality, because player 1 should be able to call player 2's "*bluff*" at some point.
- Hence, not all PSNEs are reasonable for a multi-stage game.
- We need a method to understand which equilibria are reasonable and which aren't.

- Consider how we calculated the Stackelberg equilibrium of the two-stage quantity competition game.
- Calculating the best strategy for player 2 in the second stage was relatively straightforward, because once q_1 was fixed, calculation of optimal q_2^* , as a function of q_1 , was easy.
- Next we calculated the optimal q_1 .
- So we started from the last stage of the game and worked backward to calculate the equilibrium of a multi-stage game.
- We will now extend the same idea to a general multi-stage game. This method of finding an equilibrium of a multi-stage game is called *backward induction*.
- For simplicity, we will only consider a situation where only one player takes an action at each stage of the game.

- Let h^k denote the set of actions actually taken by all players at all stages prior to stage k . Backward Induction algorithm works as follows:
 - 1) The algorithm begins by determining the optimal choices in the final stage K for each history h^K . This is determined by finding the action that maximizes the payoff of the player taking action at stage K contingent on a history h^K .
 - 2) Then it works back to stage $K - 1$ and determines the optimal action for the player taking action in stage $K - 1$, given that the player taking action in stage K with history h^K will take the action that we determined previously.
 - 3) We repeat this process one-by-one for stages $K - 2, K - 3, \dots, 2, 1$.

- This process yields a PSNE where each player's actions are optimal for each possible action history (even those not taken).
- In other words, the problem of empty threats does not occur.
- This pure strategy Nash equilibrium is called a subgame perfect pure strategy Nash equilibrium.
- The idea of subgame perfection is not merely restricted to the multi-stage games with only one player taking action at each stage.
- We will now look at an example of subgame perfect PSNE for a more general multi-stage game where one or more players can take actions at each stage.

- Problem: Firm 1 and firm 2 both have a production cost of 2 per unit. Firm 1 can install a new technology which will make the production cost 0 per unit. But installing the technology will cost f . Firm 2 will observe whether or not firm 1 has installed the new technology. Then the two firms will simultaneously decide their outputs (q_1 and q_2). Assume that the price as a function of demand is $p(q_1, q_2) = 14 - (q_1 + q_2)$. Find the subgame perfect PSNE of this game.
- Solution: The payoff functions of the two players are:

- $$u_1(q_1, q_2) = \begin{cases} (12 - (q_1 + q_2)) * q_1 & \dots \text{if it doesn't invest} \\ (14 - (q_1 + q_2)) * q_1 - f & \dots \text{if invests in new tech} \end{cases}$$
- and $u_2(q_1, q_2) = (12 - (q_1 + q_2)) * q_2$.

- Now we will work backward to find a subgame perfect PSNE.
- If firm 1 does not invest in new technology, then we have the basic Cournot game that we saw earlier. The PSNE is obtained by solving the following two FOCs simultaneously: $q_1 = 6 - \frac{q_2}{2}$ and $q_2 = 6 - \frac{q_1}{2}$. As found earlier, the unique PSNE is at $(q_1, q_2) = (4, 4)$. The payoff for each firm is 16 each.
- On the other hand, if firm 1 does invest in the new technology, then we will solve the following two FOCs simultaneously: $q_1 = 7 - \frac{q_2}{2}$ and $q_2 = 6 - \frac{q_1}{2}$. We get, $(q_1, q_2) = \left(\frac{16}{3}, \frac{10}{3}\right)$. Payoff of firm 1 is $\frac{256}{9} - f$. So firm 1 should make the investment if $\frac{256}{9} - f > 16$, i.e., if $f < \frac{112}{9}$.
- For example, if $f = 0$ then obviously firm 1 should invest in new technology, and if $f = 15$ then it shouldn't invest.

Objective :

Key Concepts :