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Subset Sum. (NP Complete problems in arithmetic)

Remark about Euclidean TSP.

Input: list of points in  $\mathbb{R}^2$  with int coords  
 $(x_1, y_1), \dots, (x_n, y_n)$

integer  $k$   
Question:  $\exists?$  a TSP tour that visits all points and  
with combined (Euclidean) length  $\leq k$

This problem is NP-hard. NOT KNOWN TO BE NP-COMplete!  
We don't know a verifier.

Subset Sum Problem:

Input: list of integers  $w_1, \dots, w_n$  (in binary)  
target sum  $W$  (in binary)

Question:  $\exists? S \subseteq [n]$  s.t.  $\sum_{i \in S} w_i = W$

Remark: When we worked on dyn prog we saw an algorithm that solves this in time  $O(nW)$ .  
*exponential in input size, which is  $O(n \log W)$ .*

Boolean table:  $T[i, j] = \begin{cases} \text{TRUE} & \text{if a subset of } \{w_1, \dots, w_i\} \text{ sums up to } j \\ \text{FALSE} & \text{otherwise} \end{cases}$

$i = 0, \dots, n; \quad j = 0, \dots, W.$

$T[0, j] = \begin{cases} \text{TRUE} & \text{if } j = 0 \\ \text{FALSE} & \text{if } j > 0 \end{cases}$

$(i > 0)$   
 $T[i, j] = \begin{cases} T[i-1, j] & \text{if } w_i > j \\ T[i-1, j] \vee T[i-1, j-w_i] & \text{if } w_i \leq j \end{cases}$

### 3SAT $\leq_p$ SUBSET SUM:

Variables:  $x_1, \dots, x_n$   
Clauses:  $C_1, \dots, C_\ell$

(1) Convert each variable to a pair of numbers.

$$\begin{aligned} x_1 &\mapsto w_1 = 1, & w_2 = 2 \\ x_2 &\mapsto w_3 = 4, & w_4 = 8 \\ x_3 &\mapsto w_5 = 16, & w_6 = 32 \\ &\vdots \end{aligned}$$

$$x_n \mapsto w_{2n-1} = 2^{2n-2}, \quad w_{2n} = 2^{2n-1}$$

This is not good enough: nothing to prevent us from choosing both  $w_1$  &  $w_2$  which corresponds to  $x_1 = T$  &  $x_1 = F$ .

(2)

$$x_1 \mapsto \begin{array}{r} 101 \\ 110 \\ 001 \\ 010 \\ \hline 111 \end{array} \begin{array}{l} \text{F} \\ \text{T} \\ \text{T} \\ \text{F} \end{array} w$$

$$x_2 \mapsto \begin{array}{r} 101000 \\ 110000 \\ 001000 \\ 010000 \\ \hline 111111 \end{array} \begin{array}{l} \text{F} \\ \text{T} \\ \text{T} \\ \text{F} \end{array} w$$