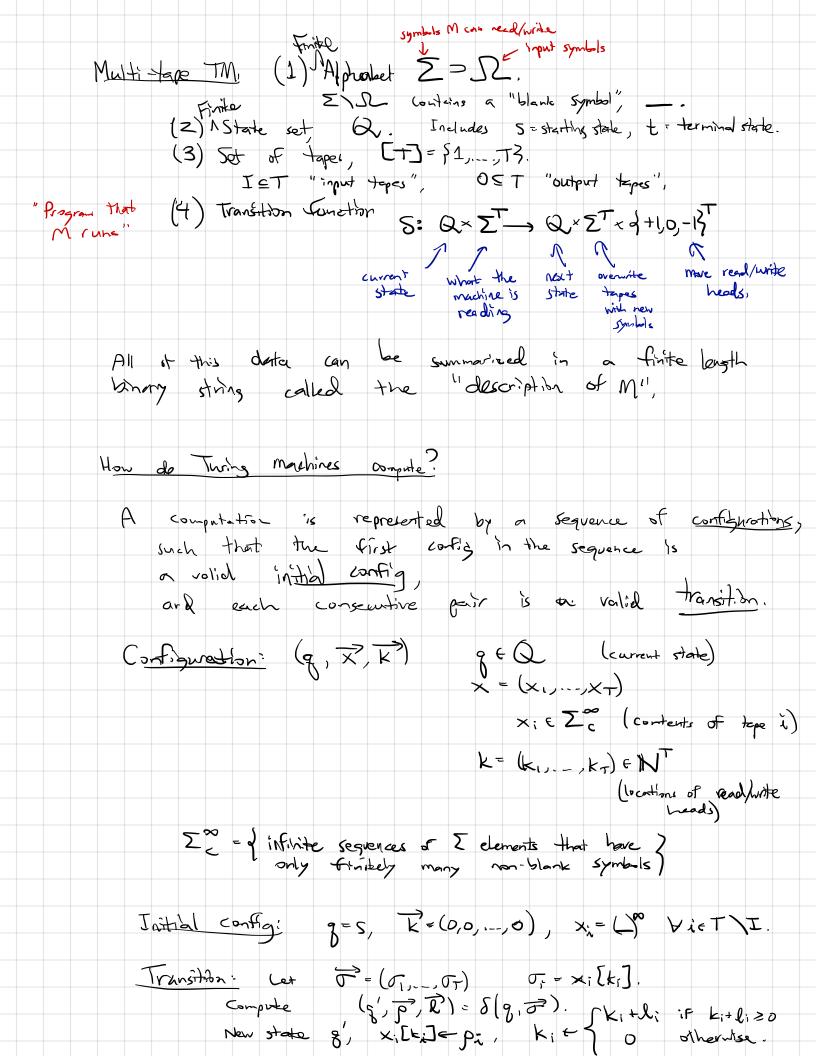
10 Apr 2019	TURING N	NACHINES AN	OMPUT	TABILITY
	2 tomorrow 101 and Plazza we'll			based on net ID.
Formalizing "al=		"computer."		
Alan luring These 2  To be That cou	: Turing formal: sms computationally	machine.  and many a equivalent.  ated in one	thurs, turne Any function of these	d out
જાજ જ	computational p	in all of the in these veces that se — Even d we can he with. Churc	can be co	urned
Turing machin	ne Basical DFA) augme		state mod	hae
Ho	ive you sees	Symbols M can really bet $Z = D^{2}$ L contains a	(A) Yes (	B) No
(3) 5	state set 6 St of taper, IST "input	[neludes] [+] = \$1,,T	5 = startly state, 3. T "output tag	t = terminal state.
"Program that (4) To	ransition Gunetion	S: Q×Σ <sup>T</sup> -	$\rightarrow Q^{\times} \Sigma^{T} \leftarrow 0$	\+1,0,-1\ <sup>T</sup>



An infinite computation on input & means IV naver halts.

Denoted M(x) = ? A finde computation is one where g = t at some time on. We say on is the running time.

And if z denotes the contents of the output types (discarding every symbol offiter initial blank on each type)

We write M(x) = zand usi z the output of M or input x.