

Networks II: Market Design—Lecture 8

Markets with Initial Endowments: The Core

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Recap: Last time

- Theorem 1: Serial dictatorship is strategyproof and therefore Pareto-efficient wrt true preferences
 - Note: Ordering of agents f is **independent** of preference profile \succ !
 - Serial dictatorship 2.0: Pareto-efficient, but **not** strategyproof!
- How general are serial dictatorships?
 - Theorem 2: Every Pareto-efficient matching in (A, H, \succ) is the output of **some** serial dictatorship.
 - Theorem 3: A mechanism is non-bossy, neutral, and strategyproof if and only if it is a serial dictatorship.

- Looking back: What we've seen so far
- Initial endowments: The 'Housing Market' problem
- The core
- Existence of a core allocation: Gale's TTC algorithm
- Properties of core allocations

So where are we now?

- Markets without money: Choosing a matching
 - One-sided preferences: Agents only on one side
 - Lots of real-world applications: Scheduling, college housing assignments, school choice, organ exchange, ...
 - Binary preferences: Easy to define 'ideal', 'best' allocations
 - Rank-order preferences: 'Best' difficult to define
 - Good matchings: Pareto-efficient (*i.e.*, not-bad) matchings
 - Recall: *Concept* of Pareto-efficiency applies more generally
 - *Not* restricted to strict rank-order preferences!

So where are we now?

- Mechanisms: Choosing a matching given market (A, H, \succ)
 - Pareto-efficiency of mechanism φ : Matching $M = \varphi(A, H, \succ)$ Pareto-efficient wrt \succ for all inputs
 - Strategyproofness of φ : Agents *never* have reason to lie about preferences to φ
 - Existence and characterization:
 - Serial dictatorship: A strategyproof, Pareto-efficient mechanism
 - Any strategyproof, non-bossy, neutral mechanism must be a serial dictatorship

Before we go on, though

- M is for: Models, matchings, mechanisms,... and mix-ups!
- Models:
 - How do agents express preferences over objects?
 - Binary versus rank-order preferences
 - Initial endowments or not
 - Agents on one versus both sides
 - Centralized allocation or decentralized, ...
- Matchings:
 - Allocation: Which agent is assigned which item
- Mechanisms (in our settings)
 - Choosing a matching
 - Takes as input a *particular instance* of preference profiles

- Pareto-efficiency is a property of
 - A The strict rank-order preference model
 - B Matchings
 - C Mechanisms
 - D Matchings and mechanisms
 - E All three: Model, matchings and mechanisms
- Strategyproofness is a property of
 - A The strict rank-order preference model
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 - E All three: Model, matchings and mechanisms

Moving on: One-sided markets with initial endowments

- Markets so far: No ‘initial endowments’, or ‘initial property rights’
 - No agent pre-owned (or had claims to) any item
 - So: No property rights issues
- What if agents own an item to start with?
 - Dorm room allocation
 - Organ donation exchanges
 - School assignments
 - Online used-good, service exchanges (Swapdom, Swapace, ...)

Swapdom FAQ x

www.swapdom.com/k/o.d?af=dha&laf=1

Swapping 101

Swapping Basics

Uploading Items

Creating a Swap

Completing a Swap

Shipping

Login / Register

Need more help?

Report an issue

Swapping Basics

- ✓ What is Swapdom?
- ✓ How does Swapdom work?
- ✓ How is Swapdom different than traditional swapping?
- ✓ How much does it cost to use Swapdom?
- ✓ Who can use Swapdom?
- ✓ What will better my chances of getting a swap?
- ✓ Do I need to price my items?
- ✓ How do you make sure swaps are fair?

How Swapdom Works

- 1 List like-new items.**
Show off the awesome items you want to swap.
- 2 Tell us what you want.**
Search the items posted by our community and request your favorites.
- 3 Offer items in exchange.**
Let us know which items you'd happily trade for your request.
- 4 Take it easy.**
We'll work behind the scenes to organize a multi-person swap and contact you when it's ready to roll.
- 5 Voilà! A successful swap.**
Swappers ship out their items and everyone in the swap ends up with the item they want!

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One sided matching markets with initial endowments

- General markets with initial endowments: Some agents have initial endowments while others do not
 - Undergraduate and graduate housing: First years versus remainder
 - Organ donation exchanges: Patients with live donors, and those without
- Study simpler version: *All* agents have an initial endowment
 - So far: Note *no* agents had initial endowments
- Central question: Allocation—who should get what?
 - Fairness, individual rights protection: ‘Individual rationality’, no ‘justified envy’

Initial endowments: What are good matchings?

What properties must a 'good' matching have?

- Is Pareto efficiency 'enough'?
 - Example: Agents a, b, c ; Houses x, y, z
 - Preferences \succ : $a: x \succ z \succ y$; $b: x \succ z \succ y$; $c: x \succ y \succ z$
 - Initial endowments: a owns x , b owns y , c owns z
 - Matching $M_1 = (a-z, b-x, c-y)$
 - Is M_1 Pareto-efficient?
 - A Yes
 - B No
 - Is M_1 a 'good' matching?
- '*Individual rationality*': No agent is better off not participating in the market

Initial endowments: What are good matchings?

What properties must a 'good' matching have?

- Another example:
 - Agents a,b,c,d; Houses x,y,z,w
 - Preferences: a: $z \succ y \succ w \succ x$; b: $z \succ x \succ w \succ y$;
c: $y \succ w \succ x \succ z$; d: $z \succ x \succ y \succ w$
 - Initial endowments: a owns x, b owns y, c owns z, d owns w
 - Matching $M_1 = (a-y, b-x, c-w, d-z)$
 - Is M_1 Pareto-efficient?
 - A Yes
 - B No
 - What about individually rational? [A] Yes [B] No
 - Is M_1 a 'good' matching?
 - Hint: What if b and c met on their way to the market?
- No *group* of agents can do better by not participating in the market and trading *the items they own* amongst themselves

One sided matching markets with initial endowments

- Markets with initial endowments: From allocation to *trade*
 - The 'housing market' model (Shapley and Scarf 1974)
- **Housing market:** $((a_k, h_k)_{k \in \{1, \dots, n\}}, \succ)$
 - $\{a_1, \dots, a_n\}$ is a set of agents and $\{h_1, \dots, h_n\}$ is a set of houses, where agent a_k owns house h_k
 - Each agent a has strict preferences \succ_a over houses
 - Contrast market with no initial endowments: (A, H, \succ)
- Matching M : Function specifying who gets what
 - a receives $M(a)$ under M

Initial endowments: What are good matchings?

- What properties should a good matching have?
 - Pareto efficiency
 - '*Individual rationality*': No agent is better off not participating in the market

$$M(a_k) \succeq h_k$$

- No *group* of agents can do better by not participating in the market and trading their initial endowments amongst themselves

- **Core:** A central concept in game theory
- Matching M is in the **core** if there is **no** coalition of agents $C \subseteq A$, and a matching \hat{M} , such that
 - For any $a \in C$, $\hat{M}(a)$ is the *initial house* of some $a' \in C$
 - $\hat{M}(a) \succeq_a M(a)$ for all $a \in C$
 - $\hat{M}(b) \succ_b M(b)$ for some $b \in C$

Core: An Example

(A matching M is in the **core** if there is no coalition of agents $B \subseteq A$, and a matching \hat{M} , such that

- For any $a \in B$, $\hat{M}(a)$ is the *initial house* of some $b \in B$, and
- $\hat{M}(a) \succeq_a M(a)$ for all $a \in B$ and $\hat{M}(b) \succ_b M(b)$ for some $b \in B$)
- Agents a, b, c ; Houses x, y, z
- Preferences \succ are a : $y \succ z \succ x$; b : $z \succ y \succ x$; c : $y \succ x \succ z$
- Initial endowments: a owns x , b owns y , c owns z
- Matching $M_1 = (a-z, b-x, c-y)$, $M_2 = (a-x, b-z, c-y)$
[A] Yes [B] No
 - Is M_1 Pareto-efficient?
 - Is M_1 in the core?
 - Is M_2 in the core?

Core: Another Example

(A matching M is in the **core** if there is no coalition of agents $B \subseteq A$, and a matching \hat{M} , such that

- For any $a \in B$, $\hat{M}(a)$ is the *initial house* of some $b \in B$, and
- $\hat{M}(a) \succeq_a M(a)$ for all $a \in B$ and $\hat{M}(b) \succ_b M(b)$ for some $b \in B$)
- Agents a, b, c ; Houses x, y, z
- Preferences \succ are $a: y \succ z \succ x$; $b: x \succ y \succ z$; $c: z \succ y \succ x$
- Initial endowments: a owns z , b owns y , c owns x
- Matching $M = (a-z, b-x, c-y)$ [A] Yes [B] No
 - Is M individually rational?
 - Is M in the core?
 - Does coalition $\{a, c\}$ demonstrate M is not in the core?