

10.3 Deterministic RBGs Based on Number Theoretic Problems

10.3.1 Discussion

A DRBG can be designed to take advantage of number theoretic problems (e.g., the discrete logarithm problem). If done correctly, such a generator's properties of randomness and/or unpredictability will be assured by the difficulty of finding a solution to that problem. Section 10.3.2 specifies a DRBG based on elliptic curves; Section 10.3.3 specifies a DRBG based on the RSA integer factorization problem.

10.3.2 Dual Elliptic Curve Deterministic RBG (Dual_EC_DRBG)

10.3.2.1 Discussion

Dual_EC_DRBG (...) is based on the following hard problem, sometimes known as the "elliptic curve logarithm problem": given points P and Q on an elliptic curve modulo n , find a such that $Q = aP$.

Dual_EC_DRBG (...) uses a seed m bits in length to initiate the generation of m -bit pseudorandom strings by performing scalar multiplications using two random points in an elliptic curve group, where the curve is defined over a field approximately 2^m in size, where $m \geq 192$. Figure 16 depicts the **Dual_EC_DRBG (...)**.

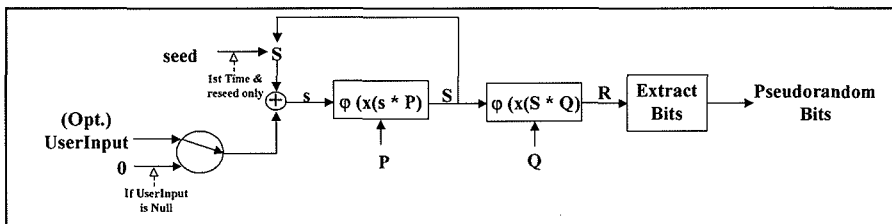


Figure 16: Dual_EC_DRBG (...)

The initialization of this DRBG requires the selection of an appropriate elliptic curve for the desired security strength. The curve **shall** be generated in accordance with ANSI X9.62 or X9.63, or **shall** be selected from the NIST Recommended curves. If the DRBG is reseeded, the same or a different curve may be used for the new DRBG instance.

During initialization and reseeding, two random points (P and Q) **shall** be generated for the curve, such that each point generates a large cyclic group on the curve. These points may be considered as the equivalent of keys and **shall** be handled as keys (see SP 800-57). The points may be generated externally and entered into the DRBG during initialization and reseeding, or may be generated internally.

The *seed* used to define the initial value (S) of the DRBG **shall** have entropy that is at least twice the desired security *strength*. The length of the *seed* **shall** be m bits in length. Further requirements for the *seed* are provided in Section 9.4.

When optional user input (*UserInput*) is used, the length and value of *UserInput* are arbitrary.

Comment [ebb1]: Page: 91
The revised version.

Comment [ebb2]: Page: 91
What reference should be provided? The revised X9.62? FIPS 186-2? Should there be other references?

Comment [ebb3]: Page: 91
How is this guaranteed, given our current philosophy? Presumably, when $m = 192$, 80 bits of security are obtained, requiring 160 bits of entropy. How does one get exactly 192 bits? Do we use a KDF and run until we get enough bits? Which KDF?

Figure 17 depicts the insertion of test input for the *seed* and the *UserInput*. The tests **shall** be run on the output of the generator. Validation and Operational testing are discussed in Section 11. Detected errors **shall** result in a transition to the error state.

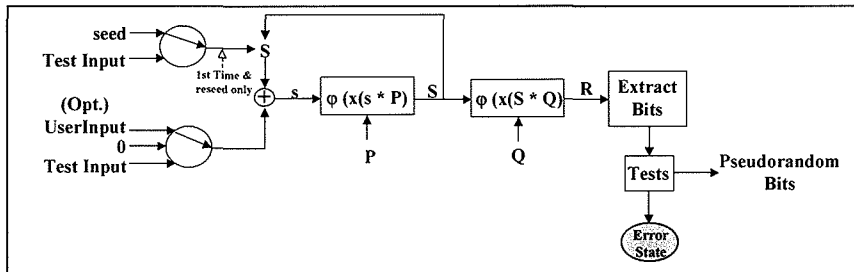


Figure 17: Dual_EC_DRBG (...) (with Tests)

10.3.2.2 Description

10.3.2.2.1 General

The initialization of **Dual_EC_DRBG (...)** consists of selecting an appropriate elliptic curve, selecting two random points on that curve and obtaining a *seed* that is used to determine an initial value (*S*) for the DRBG that is one element of the initial *state*. The state consists of:

1. (Optional) The *purpose* of the DRBG instantiation; if the DRBG is used for multiple *purposes*, requiring multiple instantiations, then the *purpose* **shall** be indicated, and the implementation **shall** accommodate multiple *states* simultaneously; if the DRBG will be used for only one *purpose*, then the *purpose* **may** be omitted),
2. A value (*S*) that is updated during each request for pseudorandom bits,
3. (Optional) The *m*-bit prime modulus *p* for curves over F_p ; if the DRBG will be using a curve over F_p , the value of *p* **shall** be present; otherwise, a value for *p* **shall not** be present,
5. The elliptic curve domain parameters $(q, a, b, G, n [, h])$, where *q* is the field size, *a* and *b* are two field elements that define the equation of the curve, *G* is a generating point of prime order on the curve, *n* is the order of the point *G*, and *h* is the optional cofactor,
6. The two random points on the curve (*P* and *Q*),
7. The maximum security *strength* provided by the instance of the DRBG, and
8. (Optional) A record of the seeding material in the form of a one-way function that is performed on the *seed* for later comparison with a new *seed* when the DRBG is reseeded; this value **shall** be present if the DRBG will potentially be reseeded; it **may** be omitted if the DRBG will not be reseeded.

Comment [ebb4]: Page: 92
Note that the basis indication and the SEED are not present. Do we need them ?

The *state* **shall** be retained within the DRBG.

For each request of pseudorandom bits, a *requested_strength* is provided and checked against the *strength* indicated in the *state*. If the *requested_strength* exceeds the *strength* afforded by the DRBG, an indication of failure is returned.

The variables used in the description of **Dual_EC_DRBG** (...) are:

a, b	Two field elements that define the equation of the curve.
E	An elliptic curve defined over F_2^m or over F_p .
<i>entropy</i>	The entropy of the <i>seed</i> .
G	A generating point of prime order on the curve.
h	The cofactor.
i	A temporary value that is used as a loop counter.
m	Length of the seed; $m \geq 192$.
n	The order of the point G on the curve.
<i>no_of_bits_per_block</i>	The number of pseudorandom bits to be used from each iteration of the elliptic curve process.
NRBG (<i>entropy</i>)	A function that acquires a string of bits from an Approved NRBG or an Approved DRBG (or chain of Approved DRBGs) that is seeded by an Approved NRBG. The parameter indicates the <i>entropy</i> to be provided in the returned bits.
<i>order_P</i>	The order of the point P .
<i>order_Q</i>	The order of the point Q .
p	The modulus; an m -bit prime, where $m \geq 192$.
P, Q	Random points on the elliptic curve E , such that each generates a large cyclic subgroup on E .
<i>pseudorandom_bits</i>	The pseudorandom_bits produced by the DRBG.
<i>purpose</i>	The <i>purpose</i> of a DRBG instance.
q	The field size of the curve.
R	A value from which pseudorandom bits are extracted.
<i>random_bits</i>	A temporary value.
<i>random_for_P</i>	A random value used to determine the point P .
<i>random_for_Q</i>	A random value used to determine the point Q .
<i>requested_no_of_bits</i>	The number of pseudorandom bits to be generated.
<i>requested_strength</i>	The security <i>strength</i> of the bits requested from the DRBG.
s	A temporary value.

S	A value that is initially determined by a <i>seed</i> , but assumes new values during each request of pseudorandom bits from the DRBG.
<i>save_seed</i>	A representation of the <i>seed</i> of this instance of the DRBG.
<i>seed</i>	The <i>seed</i> for this instance of the DRBG. The <i>seed</i> is used to derive the initial value of S .
<i>selected_random_bits</i>	A temporary value.
<i>state</i>	The state of the DRBG that is carried between calls to the generator. In the following specifications, the entire state is (<i>[purpose,] S, m, [p,] q, a, b, G, n, [h], P, order_P, Q, order_Q, strength [, save_initial_state]</i>). A particular element of the <i>state</i> is specified as <i>state.element</i> , e.g., <i>state.S</i> .
<i>status</i>	The <i>status</i> returned from the initialization, reseeding or Dual_EC_DRBG (..) processes, where <i>status</i> = Success or Failure.
<i>strength</i>	The maximum strength of an instance of the DRBG (see Table XX).
<i>temp</i>	A temporary value.
<i>UserInput</i>	Optional user input.
<i>UserInput_flag</i>	A flag that indicates whether or not user input may be used, with values as follows: <ul style="list-style-type: none"> 1 = Request <i>UserInput</i>, but return 0 if no input is available. 2 = Request <i>UserInput</i>; wait until <i>UserInput</i> is available before continuing. 3 = Obtain <i>UserInput</i> only during the first loop of the call. If no input is available, return 0. Subsequent loops should use the same value as the first loop. 4 = Obtain <i>UserInput</i> only during the first loop of the call; wait until <i>UserInput</i> is available before continuing. Subsequent loops should use the same value as the first loop. <p>If an implementation does not require this flexibility, the implementation may include only the appropriate capability (e.g., if <i>UserInput</i> is not desired, then the step requesting <i>UserInput</i> may be eliminated).</p>
$x(A)$	The x -coordinate of the point A on the curve E .
φ	A mapping from field elements to integers, which takes the bit vector representation of a field element and interprets it as the binary expansion of an integer.

Comment [ebb5]: Page: 95
Does this need to be defined here ?

* Scalar multiplication.

10.3.2.2.2 Initialization of Dual_EC_DRBG (...)

The following process or its equivalent **shall** be used to initialize the **Dual_EC_DRBG** (...) process. Let **One_Way** (...) be a one-way function.

Initialize_Dual_EC_DRBG (...):

Input: integer ([*purpose*,] m , *requested_strength*).

Output: integer *status*, where *status* = Success or Failure.

Process:

1. If ($m < 192$), then **Return** (Failure).

Comment : Determine the *strength* by the size of m .

2. If ($m < 224$), then *strength* = 80

Else if ($m < 256$), then *strength* = 112

Else if ($m < 384$), then *strength* = 128

Else if ($m < 512$), then *strength* = 192

Else *strength* = 256.

3. If (*requested_strength* > *strength*), then **Return** (Failure).

4. Choose a suitable elliptic curve E defined over F_{2^m} or over F_p , where p is an m -bit prime and $m \geq 192$. The curve may be entered from an external source, or may be generated internally. NIST Recommended elliptic curves are provided in FIPS 186-2/3. This process results in values for $q, a, b, G, n, [h]$ where q is the field size, a and b are two field elements that define the equation of the curve, G is a generating point of prime order on the curve, n is the order of the point G , and h is the cofactor. For curves over F_p , the value of p **shall** also be obtained.

Comment: Select random values for two points, P and Q .

5. If P and Q are to be obtained from an external source, then

5.1 Get P .

5.2 Determine the *order_of_P*.

5.3 If (*order_of_P* < ???), then go to step 5.1.

5.4 Get Q .

5.5 Determine the *order_of_Q*.

5.6 If (*order_of_Q* < ???), then go to step 5.4.

5.7 Go to step 8.

Comment [ebb6]: Page: 95
We probably want to specify something here, e.g., an Approved hash function of the appropriate strength.

Comment [ebb7]: Page: 95
Is this the entropy ?

Comment [ebb8]: Page: 96
Is this right ?

Comment: Get points P and Q using the random values.

6. If the curve is over F_p :

6.1 $entropy = \|p\| + 64$.

6.2 $random_for_P = (NRBG(entropy)) \bmod p$.

6.3 Using $random_for_P$, determine a point P .

6.4 Determine the $order_of_P$.

6.5 If $(order_of_P < ???)$, then go to step 6.2.

6.6 $random_for_Q = (NRBG(entropy)) \bmod p$.

6.7 Using $random_for_Q$, determine a point Q .

6.8 Determine the $order_of_Q$.

6.9 If $(order_of_Q < ???)$, then go to step 6.6.

6.10 Go to step 8.

Comment [ebb9]: Page: 96
In accordance with Section 11.3 of Part 1.

7. If the curve is over F_2^m :

7.1 $entropy = \|m\| + 64$.

7.2 $random_for_P = NRBG(entropy) \bmod m$.

7.3 Using $random_for_P$, determine a point P .

7.4 Determine the $order_of_P$.

7.5 If $(order_of_P < ???)$, then go to step 7.2.

7.6 $random_for_Q = NRBG(entropy) \bmod m$.

7.7 Using $random_for_Q$, determine a point Q .

7.8 Determine the $order_of_Q$.

7.9 If $(order_of_Q < ???)$, then go to step 7.6.

Comment [ebb10]: Page: 96
This will provide enough extra bits in steps 7.2 and 7.3 that a mod can be used to reduce to the desired number of bits (m).

8. $entropy = m + 64$.

Comment: Request the *seed* material.

9. $S = seed = NRBG(entropy) \bmod m$.

Comment : Perform a one-way function on the state values for later comparison.

11. (Optional) $save_seed = One_Way(seed)$.

12. $state = \{[purpose,] S, m, [p,] q, a, b, G, n, [h], P, Q, strength [, save_seed]\}$.

13. **Return** (Success).

10.3.2.2.3 Reseeding of Dual_EC_DRBG (...)

The following process or its equivalent **shall** be used to reseed the **Dual_EC_DRBG (...)** process. Let **One_Way (...)** be a one-way function.

Reseed_Dual_EC_DRBG (...):

Input: integer ($[purpose,] m, requested_strength$).

Output: integer *status*, where *status* = Success or Failure.

Process:

1. If ($m < 192$), then **Return** (Failure).
2. Get the appropriate *state* values for the indicated *purpose*, e.g., $q = state.q, a = state.a, b = state.b, save_seed = state.save_seed$. If a *state* is not available for the indicated *purpose*, **Return** (Failure).
3. If ($m < 224$), then *strength* = 80
Else if ($m < 256$), then *strength* = 112
Else if ($m < 384$), then *strength* = 128
Else if ($m < 512$), then *strength* = 192
Else *strength* = 256.
4. If ($requested_strength > strength$), then **Return** (Failure).

Comment : Select a new curve, if desired ; otherwise, use the old curve.
5. (Optional) Choose a suitable elliptic curve *E* defined over F_{2^m} or over F_p , where *p* is an *m*-bit prime and $m \geq 192$. The curve may be entered from an external source, or may be generated internally. NIST Recommended elliptic curves are provided in FIPS 186-2/3. This process results in values for $q, a, b, G, n, [h]$ where *q* is the field size, *a* and *b* are two field elements that define the equation of the curve, *G* is a generating point of prime order on the curve, *n* is the order of the point *G*, and *h* is the cofactor. For curves over F_p , the value of *p* **shall** also be obtained.
6. If *P* and *Q* are to be obtained from an external source, then
 - 6.1 Get *P*.
 - 6.2 Determine the *order_of_P*.
 - 6.3 If ($order_of_P < ???$), then go to step 6.1.
 - 6.4 Get *Q*.
 - 6.5 Determine the *order_of_Q*.
 - 6.6 If ($order_of_Q < ???$), then go to step 6.4.
 - 6.7 Go to step 9.
7. If the curve is over F_p .

Comment [ebb11]: Page: 97
There is an assumption that the strength of the new DRBG instance and the old DRBG instance need not be the same.

- 7.1 $entropy = \|p\| + 64$.
- 7.2 $random_for_P = (NRBG(entropy)) \bmod p$.
- 7.3 Using $random_for_P$, determine a point P .
- 7.4 Determine the $order_of_P$.
- 7.5 If $(order_of_P < ???)$, then go to step 7.2.
- 7.6 $random_for_Q = (NRBG(entropy)) \bmod p$.
- 7.7 Using $random_for_Q$, determine a point Q .
- 7.8 Determine the $order_of_Q$.
- 7.9 If $(order_of_Q < ???)$, then go to step 7.6.
- 7.10 Go to step 9.
8. If the curve is over F_2^m :
 - 8.1 $entropy = m + 64$.
 - 8.2 $random_for_P = NRBG(entropy) \bmod m$.
 - 8.3 Using $random_for_P$, determine a point P .
 - 8.4 Determine the $order_of_P$.
 - 8.5 If $(order_of_P < ???)$, then go to step 8.2.
 - 8.6 $random_for_Q = NRBG(entropy) \bmod m$.
 - 8.7 Using $random_for_Q$, determine a point Q .
 - 8.8 Determine the $order_of_Q$.
 - 8.9 If $(order_of_Q < ???)$, then go to step 8.6.

Comment: Get points P and Q using the random values.
9. $entropy = m + 64$.

Comment: Request the *seed* material.
10. $S = seed = NRBG(entropy) \bmod m$.

Comment : Perform a one-way function on the state values for later comparison.
11. $temp = One_Way(seed)$.
12. If $(temp = save_seed)$, then go to step 5.
13. $state = \{[purpose,] S, m, [p,] q, a, b, G, n, [h], P, Q, strength [, save_seed]\}$.
14. **Return** (Success).

10.3.2.2.4 Generating Pseudorandom Bits Using Dual_EC_DRBG (...)

The following process or its equivalent **shall** be used to generate pseudorandom bits.

Dual_EC_DRBG (...):

Input: integer ([*purpose*,] *requested_no_of_bits*, *no_of_bits_per_block*, *requested_strength*, *UserInput_flag*).

Output: integer (*status*, *pseudorandom_bits*), where *status* = Success or Failure.

Process:

1. Set up the *state* in accordance with the indicated purpose, e.g., $S = \text{state}.S$, $m = \text{state}.m$, $p = \text{state}.p$, $P = \text{state}.P$, $Q = \text{state}.Q$, $\text{strength} = \text{state}.strength$, etc..
2. If $((\text{no_of_bits_per_block} < 0) \text{ or } (\text{no_of_bits_per_block} > m))$, then **Return** (Failure, 0).
3. If $(\text{requested_strength} > \text{strength})$, then **Return** (Failure, 0).
4. If $((\text{UserInput_flag} < 0) \text{ or } (\text{UserInput_flag} > 4))$, then **Return** (Failure, 0).
5. *temp* = then Null string.
6. If $(\text{UserInput_flag} = 0)$, then *UserInput* = 0
Else If $((\text{UserInput_flag} = 3) \text{ and } (i \neq 0))$, then go to step 7
Else *UserInput* = **Get_userInput** (*UserInput_flag*).
7. $s = (S \oplus \text{UserInput})$. Comment: If *UserInput* is shorter than *S*, then XOR *UserInput* to the rightmost bits of *S*.
8. $S = \phi(x(s * P))$
9. $R = \phi(x(S * Q))$. Comment: *R* is an *m*-bit number.
10. If the curve is over F_p :
 - 10.1 Do $i = m$ to 1 by -1
If $((\text{Bit } i \text{ of } p = 1) \text{ and } (\text{Bit } i \text{ of } R = 0))$, then go to step 10.2.
 - 10.2 *random_bits* = Rightmost ($i - 1$) bits of *R*.
 - 10.3 Go to step 12.
11. If the curve is over F_2^m : *random_bits* = Rightmost ($m-8$) bits of *R*.
12. If $(\text{no_of_bits_per_block} > \|\text{random_bits}\|)$, then *selected_random_bits* = *random_bits*
Else *selected_random_bits* = Leftmost *no_of_bits_per_block* of *random_bits*.
13. *temp* = *temp* || *selected_random_bits*.
14. If $(\|\text{temp}\| < \text{requested_no_of_bits})$, then go to step 6.
15. *pseudorandom_bits* = Leftmost *requested_no_of_bits* of *temp*.

Comment [ebb12]: Page: 99
Do these require substeps ?

16. Update the changed values in the *state*, e.g., $state.S = S$.

17. **Return** (Success, *pseudorandom_bits*).

10.3.2.3 Generator Strength and Attributes

[To be determined]

10.3.2.4 Reseeding and Rekeying

[To be determined]