

2 of 2

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9.6.3 Derivation Function Using a Block Cipher Algorithm

Let **CBC_MAC** be the function specified in Section 9.6.4. Let **ECB_Encrypt** be an encryption operation in the ECB mode using the selected block cipher algorithm. Let *outlen* be its output block length, and let *keylen* be the key length.

The following or an equivalent process **shall** be used to derive the requested number of bits.

Input:

1. *input_string*: The string to be operated on.
2. *no_of_bits_to_return*: The number of bits to be returned by **Block_Cipher_df**. The maximum length (*max_number_of_bits*) is implementation dependent, but **shall** be **2³¹** bits.

Output:

1. *status*: The status returned from **Block_Cipher_df**. The status will indicate **SUCCESS** or **ERROR**.
2. *requested_bits* : The result of performing the **Block_Cipher_df**.

Process:

1. If (*number_of_bits_to_return* > *max_number_of_bits*), then return an **ERROR**.
2. $L = \text{len}(\text{input_string})/8$. Comment: L is the bit string representation of the integer resulting from $\text{len}(\text{input_string})/8$.
3. $N = \text{number_of_bits_to_return}/8$. Comment : N is the bitstring representation of the integer resulting from $\text{number_of_bits_to_return}/8$.
Comment: Prepend the string length and the requested length of the output to the *input_string*.
3. $S = L \parallel N \parallel \text{input_string} \parallel 0x80$.
Comment : Pad S with zeros, if necessary.
4. While $(\text{len}(S) \bmod \text{outlen}) \neq 0$, $S = S \parallel 0x00$.
Comment : Compute the starting value.
5. *temp* = the Null string.
6. $i = 0$.
7. K = Leftmost *keylen* bits of 0x010203...1F.
8. While $\text{len}(\text{temp}) < \text{keylen} + \text{outlen}$, do
 - 8.1 $IV = i \parallel 0^{\text{outlen} - \text{len}(i)}$. Comment: The integer representation of i is padded with zeros to *outlen* bits.
 - 8.2 $\text{temp} = \text{temp} \parallel \text{CBC-MAC}(K, (IV \parallel S))$.

8.3 $i = i + 1$.

Comment: Compute the requested number of bits.

9. K = Leftmost *keylen* bits of *temp*.
10. X = Next *outlen* bits of *temp*.
11. *temp* = the Null string.
12. While $\text{len}(\text{temp}) < \text{number_of_bits_to_return}$, do
 - 12.1 $X = \text{ECB_Encrypt}(K, X)$.
 - 12.2 $\text{temp} = \text{temp} \parallel X$.
13. requested_bits = Leftmost *number_of_bits_to_return* of *temp*.
14. Return SUCCESS and *requested_bits*.

9.6.4 CBC-MAC Function

The CBC-MAC function was an Approved method for computing a message authentication code. Let **ECB_Encrypt** be an encryption operation in the ECB mode using the selected block cipher algorithm. Let *outlen* be the length of the output block of the block cipher algorithm to be used.

The following or an equivalent process shall be used to derive the requested number of bits.

Input:

1. *Key*: The key to be used for the block cipher operation.
2. *data_to_MAC*: The data to be operated upon.

Output:

1. *output_block*: The result to be returned from the CBC-MAC operation.

Process:

1. $\text{chaining_value} = 0^{\text{outlen}}$.
Comment: Set the first chaining value to *outlen* zeros.
2. $n = \text{len}(\text{data_to_MAC}) / \text{outlen}$.
3. Split the *data_to_MAC* into n blocks of *outlen* bits each forming block_1 to block_n .
4. For $i = 1$ to n do
 - 4.1 $\text{input_block} = \text{chaining_value} \oplus \text{block}_i$.
 - 4.2 $\text{chaining_value} = \text{ECB_Encrypt}(\text{Key}, \text{input_block})$.
5. $\text{output_block} = \text{chaining_value}$.
6. Return *output_block*.

