

Algorithms

- Classical algorithm (focus)
- FastDTW
- DTW with bands

Step 1: create a matrix D

DTW($s[1\dots n]$, $t[1\dots m]$) // s is reference sequence, t is query sequence

//Initialization

$D = \text{array}[0\dots n, 0\dots m]$

$D[0,0] = 0$

for $i=1$ to n $D[i,0] = \infty$

for $j=1$ to m $D[0,j] = \infty$

//Fill in the matrix

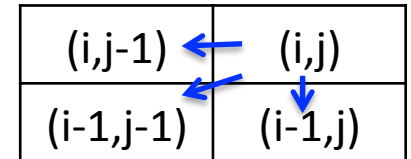
for $i=1$ to n

for $j=1$ to m

cost = $d(s[i], t[j])$

$D[i,j] = \text{cost} + \min(D[i-1,j], D[i,j-1], D[i-1,j-1])$

return $D[n,m]$



Example: Initialization

Query sequence
j: 0...m, m=7

[7]	1	∞						
[6]	2	∞						
[5]	3	∞						
[4]	2	∞						
[3]	1	∞						
[2]	1	∞						
[1]	0	∞						
[0]		0	∞	∞	∞	∞	∞	∞
t			1	1	2	3	2	0
	s	[0]	[1]	[2]	[3]	[4]	[5]	[6]

i: 0...n, n=6

Reference sequence: s

Initialization

$D = \text{array}[0\dots n, 0\dots m]$

$D[0,0] = 0$

for $i=1$ to n $D[i,0] = \infty$

for $j=1$ to m $D[0,j] = \infty$

Column i = 1

Query sequence
j: 0...m, m=7

[7]	1	∞	5					
[6]	2	∞	5					
[5]	3	∞	4					
[4]	2	∞	2					
[3]	1	∞	1					
[2]	1	∞	1					
[1]	0	∞	1					
[0]		0	∞	∞	∞	∞	∞	∞
t			1	1	2	3	2	0
	s	[0]	[1]	[2]	[3]	[4]	[5]	[6]

i: 0...n, n=6

Reference sequence: s

Loop

```

for i=1 to n
  for j=1 to m
    cost = d(s[i],t[j]) //  $d(s[i],t[j]) = |s[i]-t[j]|$ 
     $D[i,j] = \text{cost} + \min(D[i-1,j], D[i,j-1], D[i-1,j-1])$ 
  return D[n,m]
```

Column i = 2

Query sequence
j: 0...m, m=7

[7]	1	∞	5	5				
[6]	2	∞	5	5				
[5]	3	∞	4	4				
[4]	2	∞	2	2				
[3]	1	∞	1	1				
[2]	1	∞	1	1				
[1]	0	∞	1	2				
[0]		0	∞	∞	∞	∞	∞	∞
t			1	1	2	3	2	0
	s	[0]	[1]	[2]	[3]	[4]	[5]	[6]

i: 0...n, n=6

Reference sequence: s

Loop

```

for i=1 to n
  for j=1 to m
    cost = d(s[i],t[j]) //  $d(s[i],t[j]) = |s[i]-t[j]|$ 
    D[i,j] = cost + min(D[i-1,j], D[i,j-1], D[i-1,j-1])
return D[n,m]
```

Column i = 3

Query sequence
j: 0...m, m=7

[7]	1	∞	5	5	3			
[6]	2	∞	5	5	2			
[5]	3	∞	4	4	2			
[4]	2	∞	2	2	1			
[3]	1	∞	1	1	2			
[2]	1	∞	1	1	2			
[1]	0	∞	1	2	4			
[0]		0	∞	∞	∞	∞	∞	∞
t			1	1	2	3	2	0
	s	[0]	[1]	[2]	[3]	[4]	[5]	[6]

i: 0...n, n=6

Reference sequence: s

Loop

```

for i=1 to n
  for j=1 to m
    cost = d(s[i],t[j]) //  $d(s[i],t[j]) = |s[i]-t[j]|$ 
     $D[i,j] = \text{cost} + \min(D[i-1,j], D[i,j-1], D[i-1,j-1])$ 
return D[n,m]
```

Column $i = 4, 5, 6$

Query sequence
j: 0...m, m=7

[7]	1	∞	5	5	3	4	2	2
[6]	2	∞	5	5	2	2	1	3
[5]	3	∞	4	4	2	1	2	5
[4]	2	∞	2	2	1	2	2	4
[3]	1	∞	1	1	2	4	5	6
[2]	1	∞	1	1	2	4	5	6
[1]	0	∞	1	2	4	7	9	9
[0]		0	∞	∞	∞	∞	∞	∞
t			1	1	2	3	2	0
	s	[0]	[1]	[2]	[3]	[4]	[5]	[6]

i: 0...n, n=6

Reference sequence: s

Loop

```

for i=1 to n
  for j=1 to m
    cost = d(s[i],t[j]) //  $d(s[i],t[j]) = |s[i]-t[j]|$ 
     $D[i,j] = \text{cost} + \min(D[i-1,j], D[i,j-1], D[i-1,j-1])$ 
  return D[n,m]
```

Example in R

```
library("dtw")
```

```
s<-c(1,1,2,3,2,0)
```

```
t<-c(0,1,1,2,3,2,1)
```

```
d<-dtw(t,s,dist.method="Manhattan", keep.internals=TRUE,step.pattern=symmetric1)
```

//query is t , s is reference

```
> d$costMatrix
```

	[,1]	[,2]	[,3]	[,4]	[,5]	[,6]
[1,]	1	2	4	7	9	9
[2,]	1	1	2	4	5	6
[3,]	1	1	2	4	5	6
[4,]	2	2	1	2	2	4
[5,]	4	4	2	1	2	5
[6,]	5	5	2	2	1	3
[7,]	5	5	3	4	2	2

```
> d$stepPattern
```

Step pattern recursion:

$$g[i,j] = \min(\begin{aligned} &g[i-1,j-1] + d[i,j], \\ &g[i,j-1] + d[i,j], \\ &g[i-1,j] + d[i,j], \end{aligned})$$

Normalization hint: NA

Other parameters (dtw)

> **print(symmetric2) //this is the default step patterns**

Step pattern recursion:

```
g[i,j] = min(  
    g[i-1,j-1] + 2 * d[i ,j ] ,  
    g[i ,j-1] +    d[i ,j ] ,  
    g[i-1,j ] +    d[i ,j ] ,  
)
```

Normalization hint: N+M

> **d\$dist //the alignment distance**

[1] 2

Step 2: backtrack to get the warping path

OptimalWarpingPath(D) [1]

- Let $p_L = (n, m)$
- Suppose $p_i = (n_i, m_i)$ has been computed, compute p_{i-1} .
 - If $(n_i, m_i) = (1, 1)$, terminate.
 - $p_{i-1} = \begin{cases} (1, m_i - 1), & \text{if } n_i = 1 \\ (n_i - 1, 1), & \text{if } m_i = 1 \\ \operatorname{argmin}\{D(n_i - 1, m_i - 1), D(n_i - 1, m_i), D(n_i, m_i - 1)\}, & \text{otherwise} \end{cases}$
 - We take the lexicographically smallest pair in case “argmin” is not unique.
- Return the optimal path $p^* = (p_1, \dots, p_L)$

Query sequence
j: 0...m, m=7

[7]	1	∞	5	5	3	4	2	2
[6]	2	∞	5	5	2	2	1	3
[5]	3	∞	4	4	2	1	2	5
[4]	2	∞	2	2	1	2	2	4
[3]	1	∞	1	1	2	4	5	6
[2]	1	∞	1	1	2	4	5	6
[1]	0	∞	1	2	4	7	9	9
[0]		0	∞	∞	∞	∞	∞	∞
t			1	1	2	3	2	0
	s	[0]	[1]	[2]	[3]	[4]	[5]	[6]

i: 0...n, n=6

Reference sequence: s

Backtrack to get warping path

$$p_{i-1} = \operatorname{argmin}\{D(n_{i-1}, m_{i-1}), D(n_{i-1}, m_i), D(n_i, m_{i-1})\}$$

Example in R

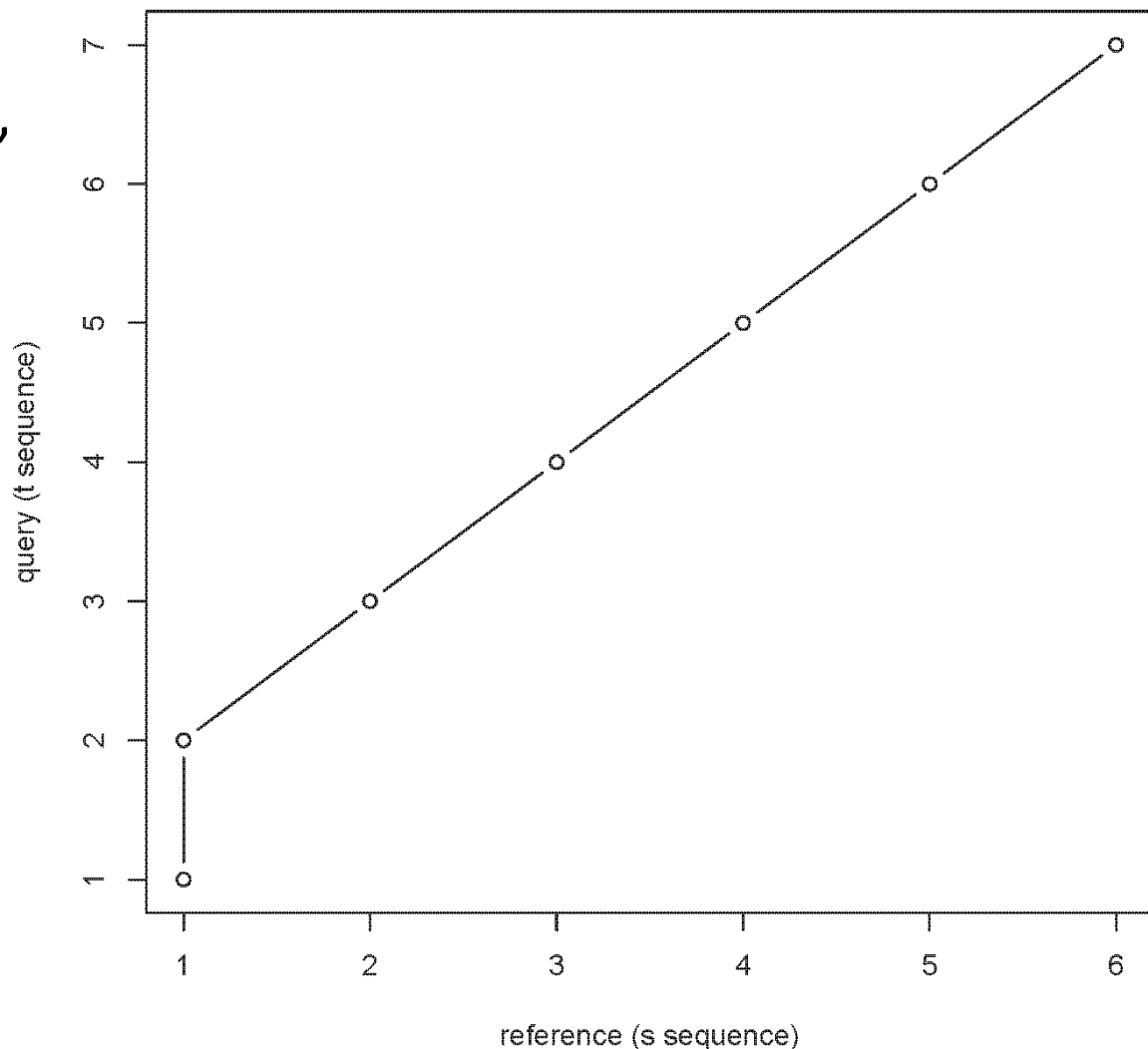
$p^* = ((1,1), (1,2), (2,3), (3,4),$
 $(4,5), (5,6), (6,7))$

> d\$index2

[1] 1 1 2 3 4 5 6

> d\$index1

[1] 1 2 3 4 5 6 7



`plot(d$index2,d$index1,xlab="reference (s sequence)",ylab="query (t sequence)",type="b")`

Operations (Edit distance?)

- Insertion // $D[i-1,j] + d(s[i],t[j])$
 - $s_1 \dots s_{i-1} s_i$
 - $t_1 \dots t_j$
- Deletion // $D[i,j-1] + d(s[i],t[j])$
 - $s_1 \dots s_i$
 - $t_1 \dots t_{j-1} t_j$
- Match // $D[i-1,j-1] + d(s[i],t[j])$
 - $s_1 \dots s_{i-1} s_i$
 - $t_1 \dots t_{j-1} t_j$