P4. "Spla in Spla"

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Spla

- P1..P3 → Splb: "Simple C-like stuff..."
 - i.e.
 - first-order functions at the toplevel
 - lexical scope, by block (function body, if, while, { ... })
 - follow the standard C-like stack discipline
 - lists and tuples by reference, on the heap
- P4 → Spla: whatever needed to do the language thing
 - Original goal: parse itself
 - Revised goal: ?

Language overview

- "whatever needed to do the language thing"
 - Parser combinators → (1) higher order functions
 - Somewhat misc. expressivity
 - **(2)** Nested function definitions, as expressions
 - Strings, chars:
 - (3) Type aliases, e.g. type @char = int; type @string = [@char]; (which really don't amount to anything interesting)
 - Slightly extend SSM to include "trapchr" instruction, then just (map print str)
 - Unneccessary semantic enhancements
 - Partial function application
 - Lazy evaluation, though this is a subtle point
 - Lists/tuples/functions by value

Variable binding

- Binding order etc?
 - "Just do something"
 - shadow new variables at block level, only allow access to initialized variables
 - Mutual recursion?
 - Nope; just let the programmer write:

```
(int → bool, int → bool) is_even_odd = (
  fun (n) { ... },
  fun (n) { ... }
);
(int → bool) is_even = is_even_odd.fst();
(int → bool) is_odd = is_even_odd.snd();
(Note that function definition and application can't occur in a single expression.)
```

- Hack:
 - Treat sequential declarations on functions and simple things (i.e. no applications) as such.

Blocks / lexical scoping

- Blocks / lexical scopes
 - Per let, { ... }, function body, and global
- Introduction of lexical contexts:
 - Global @ program execution
 - let, { ... } @ entry
 - Function body @ function application
- We have to keep all these contexts
 - Because variables can outlive execution of scopes
 - Instead of on the stack, we keep them on the heap

Blocks / lexical scoping

Context layout:

- 1. "Function context?" flag
- 2. Parent context
- 3. Return context (to be restored after exit)
- 4... locals

Blocks / lexical scoping

- Obviously wildly inefficient (in time and space)
- Is it a problem?
 - For "seriously" parsing Spla itself, somewhat
 - I had to hack SSM to have 10x as much heap to just parse simple things like int x = 10;
- What can we do?
 - Lambda-lift instead but I didn't
 - Remove contexts without variable declarations (at any depth) after execution
 - Put them on the stack I tried this, but such a hassle
 - Compile to C++ (with lambda lifting), x86 or LLVM instead of SSM
 - am doing this (x86) at the moment

Some more pragmatics

• Intermediate "intstructions", i.e. macros

```
- enter_ctxt n :: n → 0
```

 Create and enter new lexical context with top n stack cells as contents – cf. link n

```
- exit_ctxt :: 0 → 0
```

• Exit/restore lexical context – cf. unlink

```
- lex_load u i :: 0 → 1
```

Get variable #1, u contexts up

```
- lex_store u i :: 1 → 0
```

Store variable #1, u contexts up

Some more pragmatics

• Intermediate "intstructions", i.e. macros

```
- ret' :: 0 → 0
```

 Jumps to runtime code that traverses up the context tree until jumping out of function context

Returning to the original goal

- (Parsing Spla in Spla)
- Maybe I'll just:
 - parse expressions (same thing, but less)
 - or type-check (less computationally intensive, but the same HO functions and monads etc.)

```
type Qparser s a = s -> [(a, s)];
type @char
                 = int;
type @string
                 = [@char];
// List operations
// ========
([t] \rightarrow int) length = fun (list) {
    if (isEmpty(list)) {
      return 0;
    } else {
      return 1 + length(list.tl);
 };
([t] -> [t] -> [t]) concat = fun (a, b) {
    if (isEmpty(a)) {
      return b;
   } else {
      return a.hd : concat(a.tl, b);
 };
([[t]] \rightarrow [t]) concat_many = fun (lists) {
    if (isEmpty(lists)) {
      return [];
    } else {
      return concat(lists.hd,
        concat many(lists.tl));
 };
((a -> b) -> [a] -> [b]) map =
 fun (f, list) {
    if (isEmpty(list)) {
      return [];
    } else {
      return f(list.hd) : map(f, list.tl);
 };
((a -> b -> a) -> a -> [b] -> a) foldl =
 fun (f, e, list) {
    if (isEmpty(list)) {
      return e;
    } else {
      return foldl(f, f(e, list.hd), list.tl);
 };
([a] \rightarrow [a]) reverse = fun (list) {
    [a] rev = [];
    while (!isEmpty(list)) {
      rev = list.hd : rev;
      list = list.tl;
    return rev;
```

```
// The parser monad
// ========
(a -> @parser s a) mreturn = fun (e) {
    return fun (input) {
      return (e, input) : [];
  };
(@parser s a -> (a -> @parser s b) ->
 @parser s b) mbind = fun (p, f) {
    return fun (input) {
      return concat_many(map(fun (r) {
        (@parser s b) g = f(r.fst);
        return g(r.snd);
     }, p(input)));
   };
(@parser s a) mzero = fun (input) {
    return [];
(@parser s a -> @parser s a ->
 Qparser s a) mplus = fun (p, q) {
    return fun(input) {
      return concat(p(input), q(input));
   };
([@parser s a] -> @parser s a) mplus_list =
  fun (parsers) {
    return foldl(mplus, mzero, parsers);
// List input parsers
// =========
(@parser [t] t) next = fun (input) {
    return (input.hd, input.tl) : [];
((t \rightarrow bool) \rightarrow Qparser [t] t) sat =
  fun (P) {
    return fun (input) {
      if (P(input.hd)) {
        return next(input);
      } else {
        return [];
  };
(t \rightarrow 0parser [t] t) element = fun (e1) {
    return sat(fun (e2) {
      return (e1 == e2);
    });
  };
```

```
// String input parsers
// =========
(@parser (@string) (@char)) digit = sat(fun (c) {
   return c >= 48 && c <= 57;
 });
(@parser (@string) (@char)) lower = sat(fun (c) {
   return c >= 97 && c <= 122;
(@parser (@string) (@char)) upper = sat(fun (c) {
   return c >= 65 && c <= 90;
 });
(@parser (@string) (@char)) alpha =
 mplus(lower, upper);
(@parser (@string) (@char)) alphanum =
 mplus(alpha, digit);
(@string -> @parser (@string) (@string)) pstring =
 fun (match) {
   if (isEmpty(match)) {
     return mreturn([]);
   } else {
     return mbind(element(match.hd), fun (x) {
       return mbind(pstring(match.tl), fun (x) {
         return mreturn(match);
       });
     });
 };
([@string] -> @parser (@string) (@string)) pstring_any =
 fun (match_any) {
   return mplus_list(map(pstring, match_any));
 };
```

(Most of) the parser combinator library

(Des questions)