

# Random problems with R

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## Abstract

R (Version 3.5.1 patched) has two issues with its random sampling functionality. First, it uses a version of the Mersenne Twister known to have a seeding problem, which was corrected by the authors of the Mersenne Twister in 2002. Updated C source code is available at <http://www.math.sci.hiroshima-u.ac.jp/~m-mat/MT/MT2002/CODES/mt19937ar.c>. Second, R generates random integers between 1 and  $m$  by multiplying random floats by  $m$ , taking the floor, and adding 1 to the result. Well-known quantization effects in this approach result in a non-uniform distribution on  $\{1, \dots, m\}$ . The difference, which depends on  $m$ , can be substantial. Because the `sample` function in R relies on generating random integers, random sampling in R is biased. There is an easy fix: construct random integers directly from random bits, rather than multiplying a random float by  $m$ . That is the strategy taken in Python's `numpy.random.randint()` function, and recommended by the authors of the Mersenne Twister algorithm, among others. Example source code in Python is available at <https://github.com/statlab/cryptorandom/blob/master/cryptorandom/cryptorandom.py> (see functions `getrandbits()` and `randbelow_from_randbits()`).

A textbook way to generate a random integer on  $\{1, \dots, m\}$  is to start with  $X \sim U[0, 1)$  and define  $Y \equiv 1 + \lfloor mX \rfloor$ . If  $X$  is truly uniform on  $[0, 1)$ ,  $Y$  is then uniform on  $\{1, \dots, m\}$ . But if  $X$  has a discrete distribution derived by scaling a pseudorandom  $w$ -bit integer (typically  $w = 32$ ) or floating-point number, the resulting distribution is, in general, not uniformly distributed on  $\{1, \dots, m\}$  even if the underlying pseudorandom number generator (PRNG) is perfect. Theorem ?? illustrates the problem.

**Theorem 1** (?). *Suppose  $X$  is uniformly distributed on  $w$ -bit binary fractions, and let  $Y_m \equiv 1 + \lfloor mX \rfloor$ . Let  $p_+(m) = \max_{1 \leq k \leq m} \Pr\{Y_m = k\}$  and  $p_-(m) = \min_{1 \leq k \leq m} \Pr\{Y_m = k\}$ . There exists  $m < 2^w$  such that, to first order,  $p_+(m)/p_-(m) = 1 + m2^{-w+1}$ .*

A better way to generate random elements of  $\{1, \dots, m\}$  is to use pseudorandom bits directly, avoiding floating-point representation, multiplication, and the floor operator. Integers between 0 and  $m - 1$  can be represented with  $\mu(m) \equiv \lceil \log_2(m - 1) \rceil$  bits. To generate a pseudorandom integer between 1 and  $m$ , first generate  $\mu(m)$  pseudorandom bits (for instance, by taking the most significant  $\mu(m)$  bits from the PRNG output, if  $w \geq \mu(m)$ , or by concatenating successive outputs of the PRNG and taking the first  $\mu(m)$  bits of the result, if  $w < \mu(m)$ ). Cast the result as a binary integer  $M$ . If  $M > m - 1$ , discard it and draw another  $\mu(m)$  bits; otherwise, return  $M + 1$ .<sup>1</sup> Unless  $m = 2^{\mu(m)}$ , this procedure is expected to discard some random draws—up to almost half the draws if  $m = 2^p + 1$  for some integer  $p$ . But if the input bits are IID Bernoulli(1/2), the output will be uniformly distributed on  $\{1, \dots, m\}$ . This is how the Python function `numpy.random.randint()` (Version 1.14) generates

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<sup>1</sup>See ?, p.114. This is also the approach recommended by the authors of the Mersenne Twister. See <http://www.math.sci.hiroshima-u.ac.jp/~m-mat/MT/efaq.html>, last accessed 18 September 2018.

pseudorandom integers.<sup>2</sup>

The algorithm that R (Version 3.5.1 patched) (?) uses to generate random integers in `R.unif.index()` (in `RNG.c`) has the issue pointed out in Theorem ?? in a more complicated form, because R uses a pseudorandom float at an intermediate step, rather than multiplying a binary fraction by  $m$ . The way the float is constructed depends on  $m$ . Because `sample` relies on random integers, it inherits the problem.

When  $m$  is small, R uses `unif.rand` to generate pseudorandom floating-point numbers  $X$  on  $[0, 1)$  starting from a 32-bit random integer generated from the (obsolete version of the) Mersenne Twister algorithm (?). The range of `unif.rand` contains (at most)  $2^{32}$  values, which are approximately equi-spaced (but for the vagaries of converting a binary number into a floating-point number (?), which R does using floating-point multiplication by 2.3283064365386963e-10).

When  $m > 2^{31}$ , `R.unif.index()` calls `ru` instead of `unif.rand`.<sup>3</sup> `ru` combines two floating-point numbers,  $R_1$  and  $R_2$ , each generated from a 32-bit integer, to produce the floating-point number  $X$ , as follows: the first float is multiplied by  $U = 2^{25}$ , added to the second float, and the result is divided by  $U$ :

$$X = \frac{\lfloor UR_1 \rfloor + R_2}{U}.$$

The relevant code is in `RNG.c`.

The cardinality of the range of `ru` is certainly not larger than  $2^{64}$ . The range of `ru` is unevenly spaced on  $[0, 1)$  because of how floating-point representation works.

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<sup>2</sup>However, Python's built-in `random.choice()` (Versions 2.7 through 3.6) does something else biased: it finds the closest integer to  $mX$ , where  $X$  is a binary fraction between 0 and 1.

<sup>3</sup>A different function, `sample2`, is called when  $m > 10^7$  and  $k < m/2$ . `sample2` uses the same method to generate pseudorandom integers.

The inhomogeneity can make the probability that  $X \in [x, x + \delta) \subset [0, 1)$  vary widely with  $x$ .

For the way `R_unif_index()` generates random integers, the non-uniformity of the probabilities of  $\{1, \dots, m\}$  is largest when  $m$  is just below  $2^{31}$ . The upper bound on the ratio of selection probabilities approaches 2 as  $m$  approaches  $2^{31}$ , about 2 billion. For  $m$  close to 1 million, the upper bound is about 1.0004.

We recommend that the R developers replace the obsolete (1997) version of the Mersenne Twister code with the current (2002) version, and replace the algorithm in `R_unif_index()` with the algorithm based on generating a random bit string large enough to represent  $m$  and discarding integers that are larger than  $m$ . The resulting code would be simpler and more accurate. Other routines that generate random integers using the multiply-and-floor method `(int) unif_rand() * n`, for instance, `walker_ProbSampleReplace()` in `random.c`, should also be updated to use an unbiased integer generator (e.g., to call the new version of `R_unif_index()`).

## References

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