

# File System and Storage Performance

Operating Systems Lecture



Instructor: Tanaya Bowade & Dr Shabih Fatima

# Introduction to File Systems

## What is a File System?

A method and data structure that an operating system uses to control how data is stored and retrieved.

### Key Roles:



#### Data Management

Organizes and manages files and directories on storage devices



#### Abstraction Layer

Provides a logical view of data storage, hiding physical hardware details



#### Data Integrity

Ensures consistency and reliability through journaling and error checking



#### Access Control

Implements security measures like permissions and ownership

## Importance of I/O Performance

Storage and I/O performance directly impact system responsiveness and user experience.

### Why I/O Performance Matters:



Slow I/O operations cause CPU waiting time, reducing system throughput



I/O bottlenecks manifest as slow application loading and unresponsive interfaces



Efficient I/O management is critical for modern computing environments



Storage technology evolution (HDD to SSD) significantly impacts performance

# File System Concepts

## Core Components



### Files

Basic unit of data storage. A named collection of related information recorded on secondary storage. Can contain text documents, images, or executable programs.



### Directories

Special files that contain lists of other files and subdirectories. Provide hierarchical structure for organizing files, making them easier to locate and manage.



### Metadata

Data about data. Key information includes:

- Size
- Permissions
- Location
- Owner
- Timestamps
- Type

## Attributes & Naming

### File Attributes

#### Read-only

Prevents modification

#### Hidden

Invisible in listings

#### System

Critical OS file

#### Archive

Needs backup

### Naming Conventions

**Length Restrictions:** Maximum characters (e.g., 255 in many systems)

**Forbidden Characters:** Characters not allowed (e.g., /, \, :, \*, ? in Windows)

**Case Sensitivity:** Some systems are case-sensitive (UNIX/Linux)

**Extensions:** Suffixes indicating file type (e.g., .txt, .docx)

# File System Structure

## Layered Architecture

File systems are organized in layers, each providing services to the layer above.



### Application Interface

System calls: open(), read(), write(), close()



### Logical File System

File organization, directory management, access control



### File-Organization Module

Logical to physical block mapping, free space management



### Basic File System

I/O scheduling, device-independent I/O



### I/O Control

Device drivers, interrupt handling



### Devices

Disk drives, SSDs, other storage media

## Real-world Implementations

### UNIX File System (UFS)

- ✓ Uses inodes to store metadata separately from data blocks
- ✓ Case-sensitive file names
- ✓ Efficient for sequential access workloads

### NTFS (Windows)

- ✓ Case-insensitive but preserves case
- ✓ Advanced security features and permissions
- ✓ Better support for large volumes

# File Allocation and Access Methods

## File Allocation Methods

### Contiguous Allocation



File occupies consecutive blocks on disk

- + Simple implementation
- + Fast sequential access
- External fragmentation

### Linked Allocation



File stored as linked or chain list of blocks

- + No external fragmentation
- + Flexible file size
- Slow direct access

### Indexed Allocation



Index block points to data blocks

- + No external fragmentation
- + Fast direct access
- Index block overhead

## File Access Methods

### Sequential Access



Data accessed in order, from start to finish

- > Similar to reading a tape or log file
- > Processed one record at a time
- > **Example:** Reading text files line by line

### Direct Access (Random Access)



Access data at any arbitrary position

- > Jump to specific block or record
- > Essential for database applications
- > Requires efficient indexing mechanisms

# Directory Implementation and Management

## Directory Structures

### Single-Level Directory

All files contained in a single directory.

📁 Directory  
    📄 fileA 📄 fileB 📄 fileC

- + Simplicity
- Naming conflicts

### Two-Level Directory

Separate directory for each user.

📁 Master Dir  
    📁 User1\_Dir  
        📄 fileA 📄 fileB  
    📁 User2\_Dir  
        📄 fileA 📄 fileC

- + Resolves naming conflicts
- No user grouping

### Tree-Structured Directory

Hierarchical structure with multiple levels.

📁 Root Dir1  
    📁 SubDirA  
        📄 fileX  
        📄 fileY  
  
    📁 Dir2  
        📁 SubDirB  
            📄 fileZ

- + Hierarchical organization
- + No naming conflicts

# Storage Devices and Performance Basics

## HDD vs. SSD

Feature	HDD	SSD
Mechanism	Mechanical spinning platters	Flash memory (NAND-based)
Seek Time	Milliseconds (ms)	Microseconds (μs)
Transfer Rate	50-200 MB/s	500-7000 MB/s
Random I/O	Poor	Excellent
Durability	Susceptible to physical shock	More robust, no moving parts

## Key Performance Metrics



**Latency**  
Delay between I/O request and data transfer start. Measured in ms (HDD) or μs (SSD).

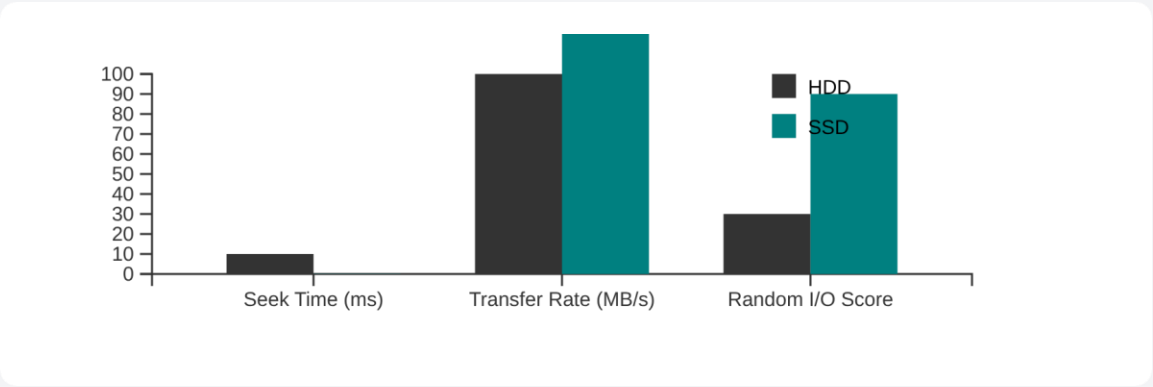


**Throughput**  
Amount of data transferred per second (MB/s or GB/s). Critical for sequential operations.



**IOPS**  
Input/output operations per second. Important for workloads with many small random I/Os.

## Performance Comparison



# Storage Performance Monitoring

## Key Metrics



### Latency

Delay between I/O request and data transfer, measured in ms (HDD) or  $\mu$ s (SSD)



### Throughput

Data transferred per unit time, measured in MB/s or GB/s



### IOPS

Input/output operations per second, critical for workloads with many small I/O requests



### Utilization

Percentage of time storage device is busy processing I/O requests



### Queue Depth

Number of I/O requests waiting to be processed

## Monitoring Tools

### Linux Tools

#### iostat - Reports I/O statistics

```
Linux 5.15.0-generic (hostname) 11/03/2025 _x86_64_ (8 CPU)
avg-cpu:  %user %nice %system %iowait %steal %idle
           0.50 0.00 0.50 0.00 0.00 99.00
Device tps kB_read/s kB_wrtn/s
sda 0.00 0.00 0.00
sdb 0.00 0.00 0.00
```

*tps: Transactions per second*

#### vmstat - Reports system statistics

Includes block I/O statistics (bi, bo) that indicate disk activity

### Windows Tools

#### Performance Monitor (perfmon)

Graphical tool for monitoring storage performance counters

Can track metrics like Disk Queue Length, % Disk Time, and Avg. Disk Response Time



# I/O Performance Analysis

## Analysis Techniques



### Bottleneck Identification

Pinpointing the component in the I/O path that limits overall performance.



### Workload Characterization

- Read-heavy vs. Write-heavy: Proportion of read vs. write operations
- Random vs. Sequential: Access patterns (impact on HDD performance)
- I/O Size: Average size of data blocks being read/written
- Concurrency: Number of simultaneous I/O requests

## Optimization Techniques



**Faster Storage Media**  
Upgrading from HDD to SSD



**RAID Configurations**  
Combining multiple disks for performance



**Load Balancing**  
Distributing I/O across multiple devices



**Application Improvements**  
Optimizing code to reduce I/O



**Caching & Buffering**  
Implementing effective caching



**File System Tuning**  
Adjusting parameters for workloads

# Caching and Buffering

## Role in Performance



### Caching

Stores copies of frequently accessed data in faster, temporary storage (e.g., RAM). Reduces latency for subsequent accesses.



### Buffering

Temporarily holds data during transfer between devices or between a device and an application. Smooths out differences in data transfer rates.

### Cache Operation:

Cache Miss



Cache Hit



Cache Hit



Cache improves read performance by storing frequently accessed data in faster memory

## Modern OS Examples

### Implementation in Modern OS:



#### Linux

Implementing sophisticated caching through the unified buffer cache/page cache, which dynamically manages memory for both file data and metadata.



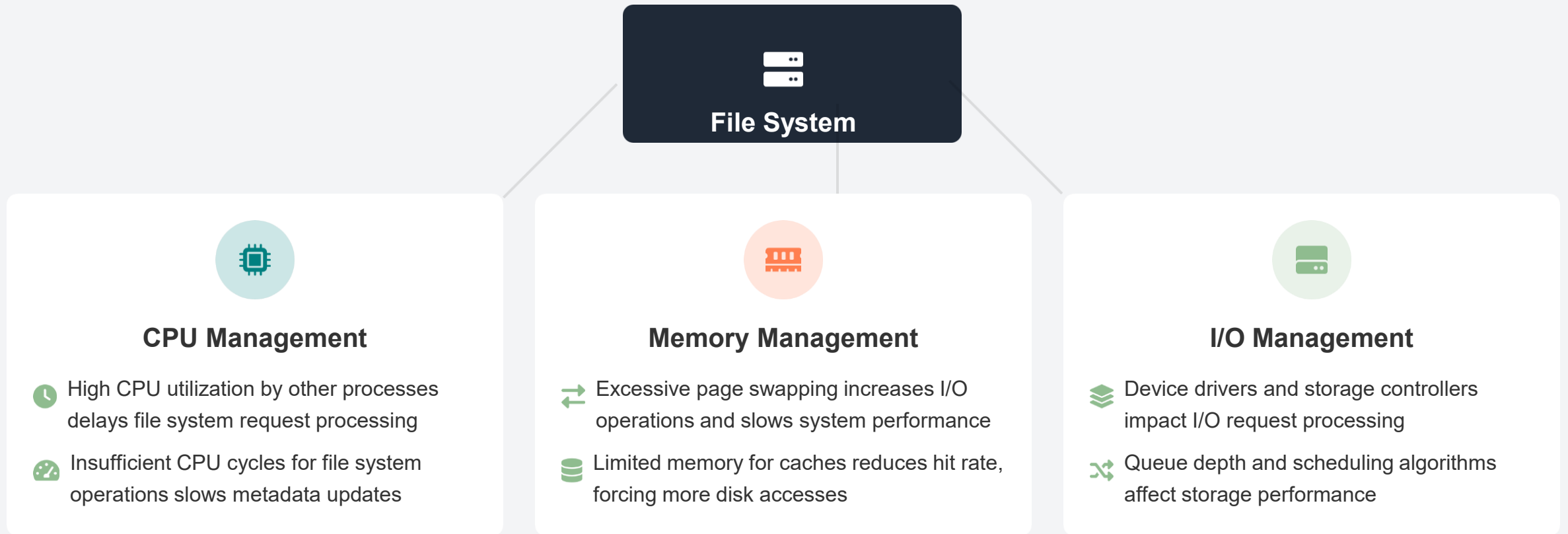
#### Windows

Uses memory-mapped files and I/O buffering to optimize file system performance and reduce physical disk accesses.

**Key Benefit:** Modern operating systems dynamically allocate and deallocate memory for caching based on system load and application demands, ensuring optimal performance.

# Resource Management Context

File system performance is deeply intertwined with the broader concept of operating system resource management. The OS orchestrates the efficient use of CPU, memory, and I/O devices, which are interconnected and affect each other's performance.



💡 **Key Insight:** Effective resource management requires a holistic approach, as performance issues in one area can impact others.

# Review Questions

Key questions for test preparation



1. Define and explain the main components of a file system.



2. What are the key differences between HDD and SSD performance?



3. How can caching improve I/O performance?



4. Describe one method used to monitor storage performance.



5. How does resource management influence file system efficiency?

💡 These questions cover the key topics. Review your notes and slides for detailed answers.