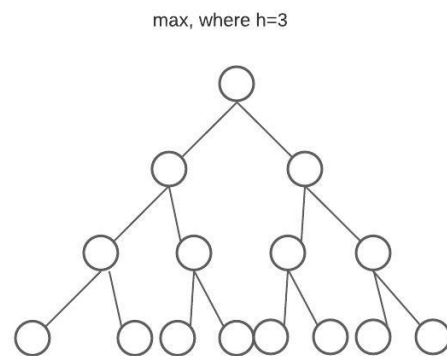
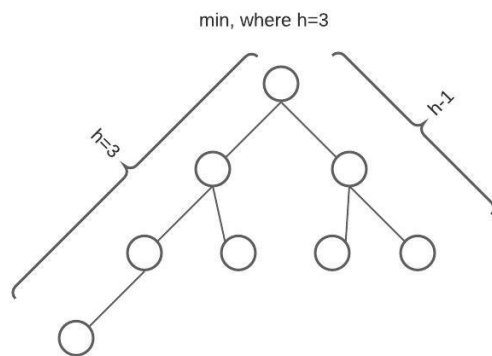


CSCI 4041 Assignment 1

🕒 Created	@Sep 24, 2020 9:47 AM
✉ Email	neis@umn.edu
👤 Name	Kelsey Neis
📅 Section	001

6-1

1) min and max numbers of elements in heap height h



Min: compute n for $h-1$, then add 1:

$$2^{(h-1+1)} - 1 = 2^h - 1$$

$$2^h - 1 + 1 \text{ (add the node on the last row)}$$

$$= 2^h$$

Max: follow formula

$$2^{(h+1)} - 1$$

Answer: $2^h \leq n \leq 2^{(h+1)} - 1$

2) Show that an n-element heap has height $\lg n$.

Given $n = 2^h$ is min, $h = \lfloor \log(n) \rfloor$

Given $n = 2^{(h+1)} - 1$ is max:

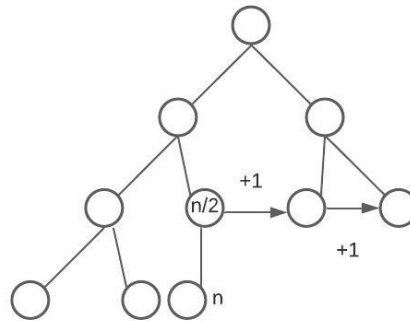
$$n+1 = 2^{(h+1)}$$

$$(n+1)/2 = 2^h$$

$$\log(n+1) - 1 = h$$

$$\lfloor \log(n+1) - 1 \rfloor = \lfloor \log(n) \rfloor$$

7) show that leaves are nodes indexed by $\lfloor n/2 \rfloor + 1, \lfloor n/2 \rfloor + 2, \dots, n$



to get the first leaf, calculate n 's parent $+ 1$, since that will be the first node without a child

$$\text{parent}(n) = \lfloor n/2 \rfloor$$

first leaf: $\lfloor n/2 \rfloor + 1$

second leaf: $\lfloor n/2 \rfloor + 2$

...and so on until we reach the last leaf, n

6-2

3)

The value at $A[i]$ remains in its position, and no further call is made to MaxHeapify, since the heap property is satisfied

4)

Since $i > A.\text{heap_size}/2$, the element $A[i]$ is a leaf, so it will remain in its position, as there are no children to compare it to.

5)

```
MAX-HEAPIFY(A, i)

l=LEFT(i)
r=RIGHT(i)
broken_heap=true
largest = i

while broken_heap == true:
    if l > A.heap_size and A[l] > A[i]
        largest = l
    if r <= A.heap_size and A[r] > A[largest]
        largest = r
    if largest == i
        broken_heap = false
    else
        exchange A[i] with A[largest]
        i = largest
```

6-5

8)

```
MAX-HEAP-DELETE(A, i)
    exchange A[i] with A[A.heap_size]
    A.heap_size = A.heap_size - 1
    MAX-HEAPIFY(A, i)
```

Since runtime for MAX-HEAPIFY is $O(\lg n)$, and the first two lines take constant time, the total runtime is $O(\lg n)$

9)

```
// k-arrays: array of k-arrays
// n: number of total elements

MERGE-SORT( k-arrays, n )
```

```

merged_array = []
for i in k-arrays:
    k-heap.push(original-index: i, value: k-arrays[i].pop)

BUILD-K-HEAP(k-heap) // change any reference to A[i] to A[i][value]
last-max = k-heap[1]

while(merged_array.length < n)
    merged_array.push(HEAP-EXTRACT-MAX-MODIFIED(k-heap, last-max.value))
    for j in merged_array.length down to 1 do:
        select-k = merged_array[j][original-index]
        if(k-arrays[select-k].length > 0)
            last-max = k-arrays[select-k].pop
            break

BUILD-K-HEAP(A)    // Time:  $O(k)$ 
    for i in [A.length/2] down to 1 do:
        MAX-HEAPIFY(A, i)

HEAP-EXTRACT-MAX-MODIFIED(A, next-elem)
    if A.heap-size < 1
        error "heap underflow"
    max = A[1]
    A[1] = A[next-elem]
    MAX-HEAPIFY(A, 1) // change any reference to A[i] to A[i][value]
    return max

```

Runtime: $O(n \lg k)$

BUILD-K-HEAP = $O(k)$ \rightarrow constant time

HEAP-EXTRACT-MAX-MODIFIED = $O(n \lg k)$