

# CS70 Modular Arithmetic

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# Basic Definitions

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# Theorems

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$\gcd(x, m) = 1 \implies x$  has a multiplicative inverse modulo  $m$  and it is **unique**.

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Compute  $\gcd(16, 10)$ :

$$\begin{aligned}\gcd(16, 10) &= \gcd(10, 6) \\ &= \gcd(6, 4) \\ &= \gcd(4, 2) \\ &= \gcd(2, 0) \\ &= 2.\end{aligned}$$

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- When  $\gcd(x, y) = 1$ , we can deduce that  $b$  is an inverse of  $y \bmod x$ .
- This uses back substitutions repetitively so that the final expression is in terms of  $x$  and  $y$ .



# Functions

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- Pre-image is a **subset** of domain, and the image/range is the **subset** of co-domain.
  - If  $f(a) = b$ , where  $a \in A$  and  $b \in B$ , then we say that  $b$  is the image of  $a$  and  $a$  is the pre-image of  $b$ .

# Bijection



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- To show that a function is *onto*, choose  $a = f^{-1}(b)$  and so  $f(f^{-1}(b)) = b$ .

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- If  $f : A \rightarrow B$  is a bijection, it will have an **inverse** function (a lemma from notes), and  $|A| = |B|$ .

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*where  $N = \prod_{i=1}^k n_i$  and  $b_i = \frac{N}{n_i} \left( \frac{N}{n_i} \right)^{-1}_{n_i}$  where  $\left( \frac{N}{n_i} \right)^{-1}_{n_i}$  denotes the multiplicative inverse  $\pmod{n_i}$  of the integer  $\frac{N}{n_i}$ .*



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## ***Proof:***

To see why  $x$  is a solution, notice that for each  $i = 1, 2, \dots, k$ , we have

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- Then  $n_1 \mid (x - y)$ ,  $n_2 \mid (x - y)$ ,  $\dots$ ,  $n_k \mid (x - y)$ .
- Since  $n_1, n_2, \dots, n_k$  are relatively prime, we have that  $n_1 n_2 \dots n_k$  divides  $x - y$ , or

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Now, to prove uniqueness, suppose there are two solutions  $x$  and  $y$ .

- Then  $n_1 \mid (x - y)$ ,  $n_2 \mid (x - y)$ ,  $\dots$ ,  $n_k \mid (x - y)$ .
- Since  $n_1, n_2, \dots, n_k$  are relatively prime, we have that  $n_1 n_2 \dots n_k$  divides  $x - y$ , or

$$x \equiv y \pmod{N}.$$

# Chinese Remainder Theorem

## **Proof:**

To see why  $x$  is a solution, notice that for each  $i = 1, 2, \dots, k$ , we have

$$\begin{aligned}x &\equiv a_1 y_1 z_1 + a_2 y_2 z_2 + \cdots + a_k y_k z_k \pmod{n_i} \\&\equiv a_i y_i z_i \pmod{n_i} \\&\equiv a_i \pmod{n_i}.\end{aligned}$$

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Thus, the solution is unique modulo  $N$ . □

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④ Compute

$$x = \sum_{i=1}^k a_i y_i z_i$$

and  $x \bmod N$  is the unique solution modulo  $N$ .



# Chinese Remainder Theorem

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- 4 Solve this congruence for  $j_k$ .
- 5 Write the solved congruence as an equation, and then substitute this expression for  $j_k$  into the equation for  $x$ .
- 6 Continue substituting and solving congruences until the equation for  $x$  implies the solution to the system of congruences.

# Chinese Remainder Theorem

**Example:**



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- Now we have  $x = 7k + 6 = 7(5j + 4) + 6 = 35j + 34$ .

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- Then  $35j + 34 \equiv 1 \pmod{3} \implies j \equiv 0 \pmod{3} \implies j = 3t$ .

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- Now we have  $x = 7k + 6 = 7(5j + 4) + 6 = 35j + 34$ .
- Then  $35j + 34 \equiv 1 \pmod{3} \implies j \equiv 0 \pmod{3} \implies j = 3t$ .
- Finally, we have  $x = 35(3t) + 34 = 105t + 34 \implies x \equiv \boxed{34} \pmod{105}$ .



# Problem Time!