DESIGN DOCUMENT FOR ASGN2

```
My program has 5 modules:
Main()
readfromfile()
writetofile()
dispatch()
entry point()
DATA STRUCTURES USED:
struct args{
     bufferdata = used to save the header of the request
     filename = name of the file
     length = length of the file or O(initially) when it is a get request
     fd = file descriptor which is initialized to 0
     rgst socket = the file descriptor of the socket where the request is
coming from
     rqst flag = to check if it was a PUT(1) or GET(0)
vector<args> v :used to store the requests as structs
Main():
/* I will be skipping over the logistics of the server
     because they are the same as asgn1. I will focus on the multi-
threading.
*/
Creates one thread that is send to the dispatch() function because it will
be our dispatch thread
The main function creates the specified number of threads and sends them
to the entry point(), each thread has a struct pointer which will be used
once the requests come in
The main continues to where it waits for requests
Once request is recieved it:
First, parses the request (either PUT or GET)
     if(PUT){
                 //parse filename and content length
                 //push a struct to the vector
     v.push back({bufferdata,filename,length,0,rqst socket,1});
     else if(GET) {
                 //parse the filename
                 //push struct to the vector
                 v.push back({bufferdta,filename,0,0,rqst socket,0});
     else{
                 send back an error
      }
The main functions only job is to recieve requests and push them onto the
ofcourse it is always running inside of a while (1) loop so that the server
only
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void readfromfile(void *i)
This function performs the same details as the asgn1
The only difference is that instead of passing each varible to it such as:
filename
length etc....
the thread uses the struct pointer that was assigned to it to access those
variables
This allows for their to be minimal race conditions since each thread has
a different instances of its pointer
Some psuedocode:
 for(int i = 0; i<strlen(filename);i++){</pre>
     check if each char in filename comlies with requirements
     if it doesnt:
     send err msg to socket
     close socket
     return;
 if(length of filename != 27){
     send err msg to socket
     close socket
     return;
 }
mutex.down(nreader);
num readers +=1;
if(first reader){
     mutex.down(&writing);
mutex.up(nreader);
get file descriptor(fd) for file by using open()
if(fd > 0)
     mutex.down(read);
     //this is because i was having issues with sockets so i decided that
only one thread can read at a time
     if(fstat doesnt return error)
                 read data of file to a buffer
                 close file
                 send ok response
                 send content length
                 send buffer of the file content
                 close socket
      }
```

```
mutex.up(read);
     mutex.down(nreader);
     nreaders--;
     if(youre the last one){
     writing.up();
     mutex.up(nreader);
}
}
else{
     send appropriate err msg
     close socket
     return;
}
return;
void writetofile(void *i)
writetofile is similar to readfrom file
I will skip the filename syntax check because it is also done in here as
well
psuedocode:
/*syntax check for file name and length of filename*/
//open the file and retrieve the file descriptor(fd)
//the algorithm used here is the same as profs r/w problem in the lecture
mutex.down(&writing)
if(open was succesful){
     //valreadfile will hold number of bytes read
     valreadile = read(in from the socket pointed by the i ptr)
     write (write the valreadfile number of bytes to the file);
     close file
     send the creat resp
     close socket
     mutex.up(&writing);
     return;
}
else{
     send appropriate err msg
     close socket
     return;
}
void dispatch(void *i)
This fucntion is responsible for alerting the sleeping thread when there
is a new request
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```
available in the vector. count is a global variable which is incremented everytime a new request comes in from main() so if all threads are working it continues to signal it until one of them goes to sleep. This is because the signal function has no effect if all the threads are awake. psuedocode:
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while(1){
     while (count == 0) {
      //wait
      signal thread
void *entry point(void *i)
This is where the working threads are sent in order to wait for the signal
of the dispatch thread
psuedocode:
while(1){
     mutex.down(&threads)
     wait();
     mutex.down(count);
      count--;//since one request is being popped off the vector
      //ptr is the pointer pointing to the element in the vector
      set i = ptr;
      //move ptr of the vector to the next element in the vector
     ptr++;
     mutex.up(&thread);
      if(PUT){
                 go to writetofile
      }
      else{
                 go to readfromfile;
      }
}
```