

Assignment 2 - Placement
EECE 583
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Info:

Developed in Linux.

Usage: ./Assignment2.exe [filename] [display mode] [anneal speed]

Display mode

0 = normal

1 = slower animations

2 = Final result only

3 = no graphics, console only.

Anneal Speed:

0 = slow anneal

1 = fast anneal

Ex. ./Assignment2.exe alu2.txt 0 0

Makefile and test script provided.

1. Simulated Annealing:

After reading the netfile and placing the cells on the board, the function ScheduleAnneal() is called. This function constructs the annealing schedule. The number of iterations per temperature is determined by $(\# \text{ of cells})^{4/3}$. Initial temperature is determined by taking the standard deviation of 50 random placements and multiplying it by 20 ($20 * \text{STD}$). Next the annealing is started within a while loop.

After each iteration the Temperature is updated by multiplying it by a constant that is 1.0 (Temperature = Temperature * beta). Through testing, I decided upon a few different values for Beta. When the acceptance rate is around the 0.44 range (determined by checking if it is within the range of 0.3 to 0.5 to allow for some fluctuations between individual iterations) I use the value of Beta = 0.997. Outside of that range of acceptance rates Beta = 0.991 is used. Using a larger value allows more iterations to occur when the acceptance rate is near the ideal range of 44%. To try and keep the acceptance rate around 44%, the window is also shrunk when the rate drops below 44%. Having a smaller window allows swaps to happen with cells that are closer by, which increases the chance that the swap will be accepted.

I initially tried to update the Temperature based on the overall STD on every iteration. Doing this would require storing an array of all costs and then calculating STD for every temperature iteration which would be an expensive operation which would increase the overall time.

The annealing ends when the Acceptance rate = 0%, meaning there was a single Temperature level in which no swaps were accepted. At that point two final Anneals are performed at 0 temperature and the loop ends.

Through many tests, it is clear that performing a slower/longer anneal with more iterations results in lower overall cost. To exemplify this, I added another option = Fast Anneal. For Fast anneal I used beta

values of: Beta = 0.985 (if acceptance rate is < 0.5 & > 0.3) else 0.94. This results in an anneal that is about 9X faster than the standard mode, but the cost is only around 1.5% higher.

2. Cost Function

This is a pretty simple calculation. Given a net, I simply grab the smallest and largest X and Y values contained within the net. So $\text{deltaX} = \text{maxX} - \text{minX}$ and $\text{deltaY} = \text{maxY} - \text{minY}$. The cost is then $\text{deltaX} + \text{deltaY}$, which is equal to the half perimeter of the cells contained within the net. The cost function takes into account the height of the routing in between rows. So the Y distance between a Cell in the first row and a cell in the 2nd row would be 2 and not 1 because of the routing in between. Hopefully that is the intended way to compute the cost for this assignment.

For each individual swap, however, cells are being swapped, not nets. So to determine whether to take the swap, we need to calculate the cost of all the nets that may have changed. The function `CostOfCell(x)` finds all the nets that cell X belongs to and adds their individual costs together.

3. Efficiency

The `DoAnneal()` function is kept as simple as possible so that it can be performed quickly. After choosing two cells and random, it determines the cost before the swap, and the cost after the swap. If the cost has decreased, we use the annealing formula $e^{(-1 * \text{costDelta} / \text{Temperature})}$ to determine if we keep the change anyways. If not, we swap it back to its original position. Swapping the cells is an $O(1)$ operation since it is accessed by directly indexing into an vector.

The board is stored as a 2D array `sBoard[x][y]` with its values either -1 or the number of the cell it contains, thus selecting two random cells is an $O(1)$ operation.

The cells are saved in a vector of cell structs: `sCells`, which for each cell contains a list of all nets it belongs to. This way calculating the cost of each cell swap is quick since we already have a list of all nets that it belongs to and we do not have to search through the netlist to calculate this cost.

4. Results

Some sample results. It is evident that increasing the annealing time only improves the quality of the placement by a small amount. The fast Anneal has around 8X less iterations, but the final cost is only between 1-2% worse. View `ResultsFastAnneal.nfo` and `ResultsSlowAnneal.nfo` for full results (the output of the `testAllFiles.sh`).

		Slow Anneal		Fast Anneal	
	Initial Cost	Final Cost	T Iterations	Final Cost	T Iterations
Alu2	6637	2417	1482	2445	229
Apex1	41264	16016	1520	16516	236
Apex4	83552	32559	1566	33086	242
C880	7200	2876	1468	2931	228
Cm138a	142	72	1293	74	198
Cm150a	328	162	1331	165	200
Cm151a	227	75	2691	76	505

