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Team Note of ATRI

Ji Chan Kim, Seong Min Jeon, Myeong Geun Hwang

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9.4 문제 풀이 체크리스트
<pre>struct dsu { public: dsu() : _n(0) {} explicit dsu(int n) : _n(n), psz(n, -1) {} int merge(int a, int b) {</pre>

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```
psz[y] = x;
      return x;
    bool same(int a, int b) { return leader(a) == leader(b); }
    int leader(int a) { return (psz[a] < 0 ? a : psz[a] = leader(psz[a])); }</pre>
    int size(int a) { return -psz[leader(a)]; }
    vector<vector<int>> groups() {
      vector<int> buf(_n), gsz(_n);
      for (int i = 0: i < n: i++) {
        buf[i] = leader(i);
        gsz[buf[i]]++;
      }
      vector<vector<int>> result(_n);
      for (int i = 0; i < _n; i++) result[i].reserve(gsz[i]);</pre>
      for (int i = 0; i < _n; i++) result[buf[i]].push_back(i);</pre>
      result.erase(remove_if(result.begin(), result.end(), [&](const vector<int>& v) {
      return v.empty(); }), result.end());
      return result:
    }
  private:
    int _n;
    vector<int> psz;
};
1.2 Fenwick
11 fenwick[N+1]:
void update(int i, ll diff) {
  while (i \leq N) {
    a[i] += diff:
    i += (i \& -i):
 }
// ex. sum
11 query(int i) {
 11 r = 0;
  while (i) {
   r += a[i]:
    i -= (i & -i);
 }
 return r:
1.3 Line Container(cht)
// cht, MinLineContainer<ll> 처럼 사용, ax+b 꼴 점화식
enum Objective {
 MAXIMIZE = +1,
 MINIMIZE = -1.
template <typename T>
struct Line {
  mutable T k, m, p;
  bool operator<(const Line& o) const { return k < o.k; }</pre>
  bool operator<(T x) const { return p < x; }</pre>
```

```
template <typename T>
T lc_inf() { return numeric_limits<T>::max(); }
template <> long double lc_inf<long double>() { return 1 / .0; }
template <typename T> T lc_div(T a, T b) \{ return a / b - ((a ^ b) < 0 and a % b); \}
template <> long double lc_div(long double a, long double b) { return a / b; };
template <typename T, Objective objective>
struct LineContainer : multiset<Line<T>, less<>>> {
 using super = multiset<Line<T>, less<>>;
 using super::begin, super::end, super::insert, super::erase;
 using super::empty, super::lower_bound;
 const T inf = lc inf<T>():
 bool insect(typename super::iterator x, typename super::iterator y) {
   if (y == end()) return x->p = inf, false;
   if (x->k == y->k) x->p = (x->m > y->m ? inf : -inf);
   else x->p = lc_div(y->m - x->m, x->k - y->k);
   return x->p >= y->p;
 void add(T k, T m) {
   auto z = insert(\{k * objective, m * objective, 0\}), y = z++, x = y;
   while (insect(y, z)) z = erase(z);
   if (x != begin() and insect(--x, y)) insect(x, y = erase(y));
   while ((y = x) != begin() and (--x)->p >= y->p) insect(x, erase(y));
 T query(T x) {
   assert(!emptv()):
   auto 1 = *lower_bound(x);
   return (1.k * x + 1.m) * objective;
template <typename T>
using MinLineContainer = LineContainer<T, Objective::MINIMIZE>;
template <typename T>
using MaxLineContainer = LineContainer<T, Objective::MAXIMIZE>;
1.4 Matrix
ll n. s. e:
using Matrix = vector<vector<11>>;
Matrix operator* (const Matrix &op1, const Matrix &op2) {
 Matrix result(op1.size(), vector<ll>(op1[0].size()));
 for (int i = 0; i < 5*n; i++) {
   for (int j = 0; j < 5*n; j++) {
     for (int k = 0; k < 5*n; k++) {
       result[i][j] += op1[i][k] * op2[k][j];
       result[i][j] %= DIV;
   }
 }
 return result;
Matrix matrix_pow(ll k) {
 if (k == 1) return matrix;
 else if (k % 2) return matrix_pow(k-1) * matrix;
```

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```
else {
    Matrix _matrix = matrix_pow(k/2);
    return matrix * matrix:
}
int main() {
  Matrix result;
 11 k;
  scanf("%11d %11d %11d %11d", &n, &s, &e, &k);
  for (int i = 0; i < n; i++) {
    vector<ll> v;
    int x:
    for (int j = 0; j < n; j++) {
      scanf("%1d", &x);
      for (int d = 5; d > 0; d--) {
       if (d == x) v.push_back(1);
        else v.push_back(0);
     }
    }
    for (int d = 0: d < 4: d++) {
      vector<ll> temp;
      for (int ti = 1; ti <= 5*n; ti++) {
        if (5*i + d + 1 == ti - 1) temp.push_back(1); else temp.push_back(0);
      matrix.push_back(temp);
    matrix.push_back(v);
  result = matrix_pow(k);
1.5 PBDS(ordered set)
// 자료형만 수정, less_equal <-> less 변경시 주의
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace "gnu pbds;
#define ordered_set tree<int, null_type, less_equal<int>,
rb_tree_tag, tree_order_statistics_node_update>
// less_equal, greater_equal 사용시 필요
void m_erase(ordered_set &os, int val){
  int idx = osOS.order_of_key(val);
 ordered_set::iterator it = os.find_by_order(idx);
  if(*it == val) os.erase(it);
}
1.6 Rope and Crope
__gnu_cxx::crope s; // string
__gnu_cxx::rope<int> r; // template
insert(pos, string)
```

substr(pos, len)
erase(pos, len)

1.7 Lazy Reversible BBST(Splay)

```
// ref: nachia library, like ACL seg, lazy
struct SplayTreeByIdx{
// ACL segと ほぼ同じ
 struct S { 11 1, r, ans, len; }; // affine monoid
  static S op(S 1, S r) {
    11 nl = (1.1 == 1.len ? 1.len+r.1 : 1.1);
    11 nr = (r.r == r.len ? r.len+l.r : r.r);
    11 len = 1.len+r.len:
    return { nl, nr, max({nl, nr, l.ans, r.ans, l.r+r.l}), len };
  }
  static S e() { return { 0, 0, 0, 0 }; }
  static void reverse_prod(S& x) { swap(x.1, x.r); } // 反時に prod を更新する
  struct F { 11 x; };
  static S mapping(F f, S x) { return x; }
  static F composition(F f, F x) { return x; }
  static F id() { return {0}; }
  struct Node{
    Node *1 = 0, *r = 0, *p = 0;
    Sa = e(): // 頂点が持つ値
    S prod = e(); // 集約(rev==1 のときでも reverse_prod 作用み)
    F f = id(); // 延(a,prod には作用み)
    int i = -1; // 配列の index
    int z = 0; // 頂点の重み ( NIL なら 0 、普通の頂点は 1 )
    int sumz = 0; // 部分木の重み
    int rev = 0: // 反の延
  }:
  using pNode = unique_ptr<Node>;
  pNode pNIL;
  Node *NIL = nullptr;
  vector<pNode> A;
  Node *R;
  SplayTreeByIdx() {
    if(!pNIL){
      pNIL = make_unique<Node>();
      NIL = pNIL.get();
      NIL->1 = NIL->r = NIL->p = NIL;
      R = NIL:
  }
  // 播
  void prepareDown(Node* c){
    if(c->1 != NIL) {
      // a,prod への作用を忘れずに
      c\rightarrow l\rightarrow a = mapping(c\rightarrow f, c\rightarrow l\rightarrow a);
      c \rightarrow l \rightarrow prod = mapping(c \rightarrow f, c \rightarrow l \rightarrow prod);
      c\rightarrow l\rightarrow f = composition(c\rightarrow f, c\rightarrow l\rightarrow f);
    if(c->r != NIL) {
      // a,prod への作用を忘れずに
      c\rightarrow r\rightarrow a = mapping(c\rightarrow f, c\rightarrow r\rightarrow a);
      c->r->prod = mapping(c->f, c->r->prod);
      c \rightarrow r \rightarrow f = composition(c \rightarrow f, c \rightarrow r \rightarrow f);
```

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```
if(c->rev) {
    swap(c->1, c->r);
    if(c->1 != NIL) {
      c->l->rev ^= 1; // 播
     reverse_prod(c->l->prod);
    if(c->r != NIL) {
     c->r->rev ^= 1; // 播
     reverse_prod(c->r->prod);
   }
    c \rightarrow rev = 0;
  c->f = id(); // 播み
// 集約
void prepareUp(Node* c) {
  c->sumz = c->l->sumz + c->r->sumz + 1; // 部分木の重み
  c->prod = op(op(c->1->prod,c->a),c->r->prod); // 集約 c は播み
// (便利)
// p の親が、子として p を照するので、それを書き換えられるようにする
// 根の場合に張して R の照を返す。
Node*& parentchild(Node* p) {
 if(p->p == NIL) return R;
 if(p\rightarrow p\rightarrow l == p) return p\rightarrow p\rightarrow l;
  else return p->p->r;
// 左回
void rotL(Node* c) {
 Node* p = c->p;
  parentchild(p) = c;
  c->p = p->p;
  p->p = c;
  if(c->1 != NIL) c->1->p = p; // 子が NIL かもしれない
 p->r = c->1;
 c->1 = p;
// 右回
void rotR(Node* c) {
 Node* p = c->p;
  parentchild(p) = c;
  c\rightarrow p = p\rightarrow p;
  p->p = c;
  if(c->r != NIL) c->r->p = p; // 子が NIL かもしれない
 p->1 = c->r;
  c->r = p;
// splav 後 c は播み
void splay(Node* c) {
  prepareDown(c); // ループが回らない時のために
  while(c->p != NIL) {
   Node* p = c->p;
   Node* pp = p->p;
    // 播は親から
    if (pp != NIL) prepareDown(pp);
```

```
if (p != NIL) prepareDown(p);
    prepareDown(c);
    if (p->1 == c) {
      if(pp == NIL) { rotR(c); }
      else if(pp->l == p) { rotR(p); rotR(c); }
      else if(pp->r == p) { rotR(c); rotL(c); }
    else {
      if(pp == NIL) { rotL(c); }
      else if(pp->r == p) { rotL(p); rotL(c); }
      else if(pp->l == p) { rotL(c); rotR(c); }
   // 集約は子から
   if(pp != NIL) prepareUp(pp);
   if(p != NIL) prepareUp(p);
   prepareUp(c);
 prepareUp(c); // ループが回らない時のために
Node* kth_element(int k) {
 Node* c = R;
 while(true) {
    prepareDown(c);
   if(c->l->sumz == k) break;
   if(c\rightarrow l\rightarrow sumz > k) \{ c = c\rightarrow l; continue; \}
   k = c > 1 - sumz + 1:
    c = c->r;
 prepareDown(c);
 splay(c);
 return c;
void insert_at(int k, S x) {
 pNode pnx = make_unique<Node>(*NIL);
 Node* nx = pnx.get();
 nx->z = nx->sumz = 1;
 nx->i = A.size();
 nx->a = nx->prod = x;
 A.emplace_back(move(pnx));
 if(k == 0) { // 左端
   nx->r = R;
   if(R != NIL) R->p = nx; // 元 O 頂点かもしれない
   prepareUp(nx); // 入したら集約
   return;
 if(k == R->sumz) { // 右端(左端と同)
   nx \rightarrow 1 = R;
   if(R != NIL) R->p = nx;
   R = nx;
   prepareUp(nx);
   return;
 auto p = kth_element(k);
  nx -> 1 = p -> 1;
```

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```
nx->r = p;
  R = nx;
  p->1->p = nx;
 p->p = nx;
  p->1 = NIL;
  prepareUp(p); // split/merge の影響
 prepareUp(nx); //
void erase at(int k) {
  auto p = kth_element(k);
  if(k == 0) { // 左端
   R = p->r:
   if(R!= NIL) R->p = NIL; // O 頂点になるかもしれない
  else if(k == R->sumz-1) { // 右端
   R = p \rightarrow 1;
   if(R != NIL) R->p = NIL;
  else {
    auto 1 = p \rightarrow 1;
    auto r = p->r;
   r\rightarrow p = NIL; // split
    R = r:
    kth_element(0);
                 // merge
   r = R;
    r->1 = 1:
                  //
   1->p = r;
    prepareUp(r); // split/merge の影響
  swap(p->i,A.back()->i); // index が更新されるよ
  swap(A[p->i],A[A.back()->i]); // 後ろに移動して
  A.pop_back(); // 削除
Node* between(int 1, int r) {
  if(1 == 0 && r == R->sumz) return R; // 全域
  if(l == 0) return kth_element(r)->1; // 左端
  if(r == R->sumz) return kth_element(l-1)->r; // 右端
  auto rp = kth_element(r);
  auto lp = rp->1;
  R = lp:
             // split
  lp \rightarrow p = NIL; //
  lp = kth_element(1-1);
                // merge
  R = rp;
  rp->l = lp; //
  lp->p = rp; //
 prepareUp(rp); // split/merge の影響
 return lp->r;
void reverse(int 1, int r) {
  auto c = between(1,r);
  c->rev ^= 1;
 reverse_prod(c->prod);
  splay(c);
void apply(int 1, int r, F f) {
```

```
auto c = between(1,r);
   c->a = mapping(f,c->a);
    c->prod = mapping(f,c->prod);
    c\rightarrow f = composition(f,c\rightarrow f);
    splay(c);
 S prod(int 1, int r) {
    return between(1,r)->prod;
}; // end of struct SplayTreeByIdx
SplayTreeByIdx splay; // insert_at := i-thへデータ投入(splay node)
1.8 Segment Tree
template <class S, auto op, auto e> struct segtree {
    segtree(int n = 0) : segtree(vector<S>(n, e())) {}
    segtree(const vector<S>& v) : _n(v.size()) {
      sz = 1;
      while (sz < _n) sz <<= 1;
      d.assign(2 * sz, e());
     for (int i = 0; i < _n; i++) d[sz + i] = v[i];
      for (int i = sz - 1; i >= 1; i--) update(i);
    void set(int p, S x) {
     p += sz;
     d[p] = x;
     for (int i = 1; i <= __builtin_ctz(sz); i++) update(p >> i);
   S get(int p) { return d[p + sz]; }
   S prod(int 1, int r) {
      S sml = e(), smr = e();
     1 += sz;
     r += sz:
      while (1 < r) {
        if (1 \& 1) sml = op(sml, d[1++]);
        if (r \& 1) smr = op(d[--r], smr);
       1 >>= 1:
        r >>= 1;
      return op(sml, smr);
   S all_prod() { return d[1]; }
 private:
   int _n, sz;
   vector<S> d;
    void update(int k) { d[k] = op(d[2 * k], d[2 * k + 1]); }
};
1.9 Lazy Segment Tree
template <class S, auto op, auto e, class F, auto mapping, auto composition, auto id>
struct lazy_segtree {
   lazy_segtree(int n = 0) : lazy_segtree(vector<S>(n, e())) {}
```

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```
lazv_segtree(const vector<S>& v) : _n(v.size()) {
 sz = 1;
 while (sz < _n) sz <<= 1;
 lg = __builtin_ctz(sz);
 d.assign(2 * sz, e());
 lz.assign(sz, id());
 for (int i = 0; i < n; i++) d[sz + i] = v[i];
 for (int i = sz - 1; i >= 1; i--) update(i);
void set(int p, S x) {
 p += sz;
 for (int i = lg; i \ge 1; i--) push(p >> i);
 d[p] = x;
 for (int i = 1; i <= lg; i++) update(p >> i);
S get(int p) {
 p += sz;
 for (int i = lg; i >= 1; i--) push(p >> i);
 return d[p];
S prod(int 1, int r) {
 if (1 == r) return e();
 1 += sz, r += sz;
 for (int i = lg; i >= 1; i--) {
   if (((1 >> i) << i) != 1) push(1 >> i);
   if (((r >> i) << i) != r) push((r - 1) >> i);
 }
 S sml = e(), smr = e();
 while (1 < r) {
   if (1 \& 1) sml = op(sml, d[1++]);
   if (r \& 1) smr = op(d[--r], smr);
   1 >>= 1, r >>= 1;
 return op(sml, smr);
S all_prod() { return d[1]; }
void apply(int p, F f) {
 p += sz;
 for (int i = lg; i >= 1; i--) push(p >> i);
 d[p] = mapping(f, d[p]);
 for (int i = 1; i <= lg; i++) update(p >> i);
void apply(int 1, int r, F f) {
 if (l == r) return;
 1 += sz, r += sz;
 for (int i = lg; i >= 1; i--) {
   if (((1 >> i) << i) != 1) push(1 >> i);
   if (((r >> i) << i) != r) push((r - 1) >> i);
 int 12 = 1, r2 = r;
 while (1 < r) {
   if (1 & 1) all_apply(1++, f);
   if (r & 1) all_apply(--r, f);
   1 >>= 1, r >>= 1;
 }
```

```
1 = 12, r = r2;
     for (int i = 1; i <= lg; i++) {
        if (((1 >> i) << i) != 1) update(1 >> i);
        if (((r >> i) << i) != r) update((r - 1) >> i);
   }
 private:
   int _n, sz, lg;
   vector<S> d;
   vector<F> lz;
   void update(int k) { d[k] = op(d[2 * k], d[2 * k + 1]); }
   void all_apply(int k, F f) {
     d[k] = mapping(f, d[k]);
     if (k < sz) lz[k] = composition(f, lz[k]);</pre>
   void push(int k) {
      all_apply(2 * k, lz[k]);
      all_apply(2 * k + 1, lz[k]);
     lz[k] = id();
   }
};
```

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2 Monoid

```
2.1 RMQ and RNQ
S op(S a, S b){ return min(a, b); } // or max
S e() { return INF: }
S mapping(F f, S x){ return f+x; }
F composition(F f, F g){ return f+g; }
F id() { return 0: }
2.2 RMQ and RUQ
S op(S a, S b){ return min(a, b); } // or max
S e() { return INF: }
S mapping(F f, S x) \{ return (f == ID ? x : f); \}
F composition(F f, F g){ return (f == ID ? g : f); }
F id(){ return ID; }
2.3 RSQ and RNQ
struct S { 11 x; int sz; };
S \circ p(S a, S b) \{ return \{a.x+b.x, a.sz+b.sz\}; \}
S e(){ return {0, 0}; }
S mapping(F f, S x) { return \{x.x + f*x.x, x.sz\}; }
F composition(F f, F g){ return f+g; }
F id(){ return 0; }
2.4 RSQ and RUQ
struct S { 11 x; int sz; };
S op(S a, S b) { return {a.x+b.x, a.sz+b.sz}; }
S e() { return {0, 0}; }
S mapping(F f, S x) { if(f != ID) x.x = f*x.sz; return x; }
F composition(F f, F g) { return (f == ID ? g : f); }
F id() { return ID; }
      Maximum Array(Struct Segment)
struct S {
  bool check;
 int 1, r, sum, len;
 S(int x = 0, int c = false): sum(0) { 1 = r = x; len = 1; check = c; }
 S(int 1, int r, int sum, int len, int c): 1(1), r(r), sum(sum), len(len) { check = c; }
};
using F = int;
const F ID = -1;
S op(S a, S b) {
  if (!a.check || !b.check) {
    if (!a.check && !b.check) return S():
    if (!a.check) return b;
    else if (!b.check) return a;
  return S(a.1, b.r, a.sum+b.sum+(a.r==b.1), a.len+b.len, true);
S e() { return S(); }
S mapping(F f, S x) {
```

```
if (f == ID) return x;
return S(f, f, x.len-1, x.len, true);
}
F composition(F f, F g) { return (f == ID ? g : f); }
F id() { return ID; }
```

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3 DP

3.1 Doubling(Sparse Table)

```
// Tree 형태에서, 어떤 값의 조상...
int n, f[200000], dp[32][200000];
void doubling() {
  for (int j = 0; j < n; j++) dp[0][j] = f[j];
  for (int i = 1; i < 31; i++) {
    for (int j = 0; j < n; j++) {
      dp[i][j] = dp[i-1][dp[i-1][j]];
    }
  }
}
```

3.2 Streetlight(prefix sum + dp)

```
ll x[1002], w[1002], wsum[1002], dp[1002][1002][2];
int main() {
  cin >> n >> m:
  for (int i = 1; i <= n; i++) {
    cin >> x[i] >> w[i];
    wsum[i] = wsum[i-1] + w[i];
  for (int i = 0; i \le 1001; i++) for (int j = 0; j \le 1001; j++) dp[i][j][0] = dp[i][j][1]
  = 1'000'000'000'000'000LL;
  dp[m][m][0] = dp[m][m][1] = 0;
  for (int len = 1; len <= n; len++) {</pre>
    for (int 1 = max(1, m-len+1); 1 \le m; 1++) {
      int r = 1 + len - 1;
      if (r > n) continue;
      if (l-1 >= 1) dp[l-1][r][0] = min({dp[l-1][r][0]},
      dp[1][r][0]+(wsum[n]-wsum[r]+wsum[1-1])*(x[1]-x[1-1]),
      dp[1][r][1]+(wsum[n]-wsum[r]+wsum[1-1])*(x[r]-x[1-1]));
      if (r+1 \le n) dp[1][r+1][1] = min(\{dp[1][r+1][1],
      dp[1][r][0]+(wsum[n]-wsum[r]+wsum[1-1])*(x[r+1]-x[1]),
      dp[1][r][1]+(wsum[n]-wsum[r]+wsum[1-1])*(x[r+1]-x[r]));
 }
  cout << min(dp[1][n][0], dp[1][n][1]);</pre>
```

4 Geometry

4.1 vector2

```
double pi = acos(-1);
template<class T>
struct vector2 {
 T x, y; // double or ll, dot, cross, fix req
 explicit vector2(T _x = 0, T _y = 0): x(_x), y(_y)
 vector2 operator + (const vector2 &rhs) const {
   return vector2(x + rhs.x, y + rhs.y);
 vector2 operator - (const vector2 &rhs) const {
   return vector2(x - rhs.x, y - rhs.y);
 vector2 operator * (const double k) const {
   return vector2(k*x, k*y);
 bool operator == (const vector2 &rhs) const {
   return x == rhs.x && y == rhs.y;
 bool operator < (const vector2 &rhs) const {</pre>
   return x != rhs.x ? x < rhs.x : y < rhs.y;
 T dot(const vector2 &rhs) const {
   return x*rhs.x + y*rhs.y;
 T cross(const vector2 &rhs) const {
   return x * rhs.y - y * rhs.x;
 double norm() const {
   return hypot(x, y);
 T norm2() const {
   return x*x+y*y;
 }
 vector2 basis() const {
   double d = norm(); // req T=double
   return vector2(x/d, y/d);
 vector2 reverse() const {
   return vector2(-1*x, -1*y);
 }
int doubleCompare(double x, double y, double eps = 1e-8) {
 if (fabs(x-y) <= eps) return 0; // 0 is true
 return (x-y > 0) ? 1 : -1;
ll ccw(vector2 v1, vector2 v2, vector2 v3) {
 if (doubleCompare((v2-v1).cross(v3-v2), 0) == 0) return 0;
 else return (v2-v1).cross(v3-v2); // 1(>0), -1(<0)
bool pointYCompare(const vector2 &a, const vector2 &b) {
 return a.x != b.x ? a.x < b.x : a.y < b.y;
```

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```
bool pointThetaCompare(const vector2 &a, const vector2 &b) {
  if (ccw(vector2(), a, b) == 0) return a.norm() < b.norm();</pre>
  else return ccw(vector2(), a, b) > 0;
double safe_acos(double a) {
  if (a \le -1.0) return pi; // reg acos(-1) = pi
  else if (a \ge 1.0) return 0;
  else return acos(a):
}
// like cross
double f(vector2 p1, vector2 q1, vector2 p2, vector2 q2) {
 return ((p2-p1).cross(q2-p2) / (q1-p1).cross(q2-p2));
double k = f(p1, q1, p2, q2); // k = 
double cx = p1.x + k*(p2.x-p1.x), cy = p1.y + k*(p2.y-p1.y); // p1->p1, p2->q2 교차점
4.2 vector3
// fix vector2 -> 3. this template only v3
vector3 cross(const vector3 &rhs) const {
  return vector3(y*rhs.z-z*rhs.y, z*rhs.x-x*rhs.z, x*rhs.y-y*rhs.x);
double norm() const {
  return sqrt(x*x + y*y + z*z);
     Line Intersection
// reg vector2
struct line2 {
 vector2 p1, p2;
  void setLeftRight() {
    if (p2 < p1) {
      vector2 temp = p2;
      p2 = p1;
      p1 = temp;
 }
};
int line_intersection(vector2 s1, vector2 e1, vector2 s2, vector2 e2) {
  if (e1 < s1) swap(s1, e1);
  if (e2 < s2) swap(s2, e2);
  int c1s = iccw(s1, e1, s2), c1e = iccw(s1, e1, e2),
    c2s = iccw(s2, e2, s1), c2e = iccw(s2, e2, e1);
  if (c1s * c1e == -1 && c2s * c2e == -1) return 2;
  else if (c1s == 0 && c1e == 0 && c2s == 0 && c2e == 0) {
    if (s1.x == e1.x) {
     // 세로 두 직선
      if (s1.y > s2.y) swap(s1, s2), swap(e1, e2);
      if (e1.y < s2.y) return 0;
      if (doubleCompare(e1.y, s2.y) == 0) return 1;
    } else {
```

```
// 가로 두 직선
     if (s1.x > s2.x) swap(s1, s2), swap(e1, e2);
     if (e1.x < s2.x) return 0;
     if (doubleCompare(e1.x, s2.x) == 0) return 1;
   }
   return 3:
   if (s1.x > s2.x) swap(s1, s2), swap(e1, e2);
   if (e1.x < s2.x) return 0;
   if (s1.y > s2.y) swap(s1, s2), swap(e1, e2);
   if (e1.y < s2.y) return 0;
   if (e1.y == s2.y) return 1;
   if (s1.x > s2.x) swap(s1, s2), swap(e1, e2);
   if (e1.x == s2.x) return 1;
   return 3;
 } else if (c1s * c1e * c2s * c2e == 0) {
   if ((c1s * c1e == 0 \&\& c2s * c2e == 1)
      || (c1s * c1e == 1 && c2s * c2e == 0)
   ) return 0; else return 1;
 } else return 0:
double f(vector2 p1, vector2 q1, vector2 p2, vector2 q2) {
 return ((p2-p1).cross(q2-p2) / (q1-p1).cross(q2-p2));
// 교차점 계산
double w = f(v[i], v[i], v[k], v[nk]);
double cx = v[i].x + w*(v[j].x-v[i].x), cy = v[i].y + w*(v[j].y-v[i].y);
// 반사
vector2 p, np, fp;
p.input();
double k = (e-s).dot(p-s)/(e-s).get_size();
np = s + (e-s).basis()*k;
fp = p+(np-p)*2; // x1의 경우 사영
cout << fp.x << " " << fp.y << "\n";
4.4 Convex Hull
// cps: 컨헐집합, [0]이 제일 외곽, O(nlogn)
vector<vector2> p;
vector<vector2>> cps;
int n, check[1050], ci = 1, cnt = 0;
double oa = 0, ia = 0;
int main () {
 cin >> n:
 p = vector<vector2>(n);
 for (int i = 0; i < n; i++) {
   check[i] = 0;
   cin >> p[i].x >> p[i].y;
   p[i].i = i;
 while (p.size() >= 3) {
   sort(p.begin(), p.end(), pointYCompare);
   vector2 save = p.front();
   for (int i = 1; i < p.size(); i++) {</pre>
     p[i].x = p[i].x - p[0].x;
     p[i].y = p[i].y - p[0].y;
```

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```
sort(next(p.begin()), p.end(), pointThetaCompare);
    vector<vector2> np, cp;
    p[0] = vector2();
    p[0].i = save.i;
    cp.push_back(p[0]);
    cp.push_back(p[1]);
    for (int i = 2; i < p.size(); i++) {</pre>
      while (cp.size() > 1) {
        if (ccw(cp[cp.size()-2], cp[cp.size()-1], p[i]) \le 0) {
          np.push_back(cp.back());
          cp.pop_back();
       } else break;
      cp.push_back(p[i]);
    if (cp.size() < 3) break;</pre>
    for (auto &e: cp) check[e.i] = ci, e.x += save.x, e.y += save.y;
    for (auto &e: np) e.x += save.x, e.y += save.y;
    swap(p, np);
    cps.push_back(cp);
    ci++;
}
      Rotating Callipers
// 가장 먼 두 점, D(n)
int i2 = 1, p1, p2;
11 maxdist = 0;
for (int i1 = 0; i1 < cp.size();) {
  auto dist = (cp[i2]-cp[i1]).normll();
  if (maxdist <= dist) {</pre>
    maxdist = dist;
    p1 = i1;
    p2 = i2;
  auto ccwv = ccw(cp[i1], cp[(i1+1)%cp.size()], cp[i2], cp[(i2+1)%cp.size()]);
  if (ccwv > 0) i2 = (i2+1)\%cp.size();
  else i1++;
     Angle Sort
sort(a, a+n, [&](const vector2 &a, const vector2 &b){
  if((tie(a.x, a.y) > tie(0, 0)) ^ (tie(b.x, b.y) > tie(0, 0))) return tie(a.x, a.y) >
  tie(b.x, b.y);
  if(ccw(a, b) != 0) return ccw(a, b) > 0; // ccw 잘 수정하셈
  return hypot(a) < hypot(b);</pre>
});
```

5 Graph

5.1 Floyd warshall

```
Time Complexity: O(n<sup>3</sup>)
int citys[101][101];
int main() {
 int n, m;
 cin >> n >> m;
 for (int i = 1; i <= 100; i++)
   for (int j = 1; j \le 100; j++)
     citys[i][j] = 1e6 * 100;
 int a, b, c;
 for (int i = 0; i < m; i++) {
   cin >> a >> b >> c:
   citys[a][b] = min(citys[a][b], c);
 for (int k = 1; k \le n; k++) {
   for (int i = 1; i <= n; i++) {
     for (int j = 1; j \le n; j++) {
       if (i != j)
          citys[i][j] = min(citys[i][j], citys[i][k]+citys[k][j]);
   }
 }
 for (int i = 1; i <= n; i++)
   for (int j = 1; j \le n; j++)
     cout << (citys[i][j] != 1e6 * 100 ? citys[i][j] : 0) << (j != n ? ' ' : '\n');</pre>
5.2 scc
// 1-index, initialize: scc_graph(int size), necessary make_graph
struct scc_graph {
 int n, id = 0;
 vector<int> ids, finished;
 vector<vector<int>> path, scc;
 stack<int> stk:
 scc_graph(int _n) {
   n = n:
   ids = finished = vector<int>(n+1);
   path = vector<vector<int>>(n+1);
 }
 void make_graph() { for (int i = 1; i <= n; i++) if (!ids[i]) dfs(i); }</pre>
 void add_edge(int u, int v) { path[u].push_back(v); }
 int dfs(int s) {
   ids[s] = ++id;
   stk.push(s);
   int parent = ids[s];
   for (auto &i: path[s]) {
     if (!ids[i]) parent = min(parent, dfs(i));
```

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```
else if (!finished[i]) parent = min(parent, ids[i]);
    }
    if (parent == ids[s]) {
      vector<int> findset;
      while (true) {
        auto x = stk.top();
        stk.pop();
        findset.push_back(x);
        finished[x] = 1;
        if (s == x) break:
      sort(findset.begin(), findset.end());
      scc.push_back(findset);
    return parent;
 }
};
      2-sat
5.3
struct csr {
  vector<int> s, el;
  csr(int n, vector<pii>& edges): s(n + 1), el(edges.size()) {
    for (auto e: edges) s[e.first+1]++;
    for (int i = 1; i \le n; i++) s[i] += s[i-1];
    auto cnt = s;
    for (auto e : edges) el[cnt[e.first]++] = e.second;
 }
};
struct scc_graph {
  scc_graph(int n) : _n(n) {}
  void add(int from, int to) { edges.push_back({from, to}); }
  pair<int, vector<int>> scc_ids() {
    auto g = csr(_n, edges);
    int cord = 0, gnum = 0;
    vector\langle int \rangle vis, low(n), ord(n, -1), ids(n);
    vis.reserve(_n);
    auto dfs = [&](auto self, int v) -> void {
      low[v] = ord[v] = cord++;
      vis.push_back(v);
      for (int i = g.s[v]; i < g.s[v + 1]; i++) {
        auto to = g.el[i];
        if (ord[to] == -1) {
          self(self, to);
          low[v] = min(low[v], low[to]);
        } else low[v] = min(low[v], ord[to]);
      if (low[v] == ord[v]) {
        while (true) {
          int u = vis.back();
          vis.pop_back();
```

```
ord[u] = _n, ids[u] = gnum;
         if (u == v) break:
        gnum++;
   }:
   for (int i = 0; i < n; i++) if (ord[i] == -1) dfs(dfs, i);
   for (auto& x : ids) x = gnum - 1 - x;
   return {gnum, ids};
 }
 vector<vector<int>> scc() {
   auto ids = scc_ids();
   int gnum = ids.first;
   vector<int> cnt(gnum);
   for (auto x : ids.second) cnt[x]++;
   vector<vector<int>> g(ids.first);
   for (int i = 0; i < gnum; i++) g[i].reserve(cnt[i]);</pre>
   for (int i = 0; i < _n; i++) g[ids.second[i]].push_back(i);</pre>
   return g;
 }
 int n:
 vector<pii> edges;
};
struct two_sat {
 int n;
 vector<bool> ans:
 scc_graph scc;
  explicit two_sat(int _n): n(_n), ans(_n), scc(2*_n) {}
 void add(int i, bool f, int j, bool g) {
   scc.add(2 * i + (f ? 0 : 1), 2 * j + (g ? 1 : 0));
   scc.add(2 * j + (g ? 0 : 1), 2 * i + (f ? 1 : 0));
 bool satisfiable() {
   auto id = scc.scc_ids().second;
   for (int i = 0: i < n: i++) {
     if (id[2 * i] == id[2 * i + 1]) return false;
     ans[i] = id[2 * i] < id[2 * i + 1];
   }
   return true;
 vector<bool> answer() { return ans; }
};
5.4 hld
// require ACL seg
template <typename T = int>
struct Edge {
 int from, to;
 T cost;
```

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```
Edge(int _to, T _cost = 1): from(-1), to(_to), cost(_cost) {} // from is array index
  Edge(int _from, int _to, T _cost): from(_from), to(_to), cost(_cost) {} // only edge
  save, ex) use in prim
  Edge & operator = (const int &x) {
   to = x;
   return *this;
  operator int() const { return to; }
};
// weighted graph
template <typename T>
using WeightedEdges = vector<Edge<T>>;
template <typename T>
using WeightedFIDXEdges = vector<vector<pair<int, T>>>;
using UnweightedEdges = vector<vector<int>>;
template <typename T>
pair<WeightedEdges<T>, WeightedFIDXEdges<T>> w_fidx_graph(int n, int m = -1, bool
is_directed = false, bool is_1idx = true) {
 WeightedEdges<T> wg;
 WeightedFIDXEdges<T> g(n), c(n);
  if (m == -1) m = n-1:
  int ui = 1;
  while (m--) {
   int u. v:
   Tw:
    cin >> u >> v;
    // cin >> v:
    // u = ++ui;
    w=0;
    if (is_1idx) u--, v--;
    wg.emplace_back(u, v, w);
    g[u].push_back({v, w});
    if (!is_directed) g[v].push_back({u, w});
  queue<pair<int,int>> q;
 q.push(\{0, -1\});
  while (q.size()) {
    auto [v, pv] = q.front();
    q.pop();
    for (auto &[e, w]: g[v]) {
     if (e != pv) {
        c[v].emplace_back(e, w);
        q.push({e, v});
     }
   }
  return {wg, c};
template <typename T, class S, auto Sop, auto Se>
struct HeavyLightDecomposition {
 // private:
 public:
    int root, eid, fid;
    vector<int> sz, d, p, top, in, out, vis, id;
```

```
vector<T> x;
  WeightedFIDXEdges<T> c;
  atcoder::segtree<S, Sop, Se> seg;
  void dfs1(int v, int pv = -1, int cid = -1) {
    vis[v] = 1:
    sz[v] = 1:
    id[v] = cid;
    for (auto &e: c[v]) {
      if (e.first == pv) continue;
      d[e.first] = d[v]+1;
      p[e.first] = v;
      dfs1(e.first, v, cid);
      sz[v] += sz[e.first];
      if (sz[e.first] > sz[c[v][0].first]) swap(e, c[v][0]);
  }
  void dfs2(int v, int pv = -1) {
    vis[v] = 2;
    in[v] = eid++:
    for (auto &[e, w]: c[v]) {
      if (e == pv) continue;
      top[e] = e == c[v][0].first ? top[v] : e;
      dfs2(e, v);
      x[e] = w;
    out[v] = eid:
// public:
  HeavyLightDecomposition(WeightedFIDXEdges<T> g, int _root = 0)
    : c(g),
      root(_root),
      eid(0),
      fid(0),
      sz(g.size(), 0),
      d(g.size(), 0),
      x(g.size(), 0),
      in(g.size(), -1),
      out(g.size(), -1),
      vis(g.size(), 0),
      id(g.size(), 0),
      seg(atcoder::segtree<S, Sop, Se>(g.size()+1)),
      top(g.size(), root),
      p(g.size(), root) {
    for (int i = 0; i < g.size(); i++) if (vis[i] != 1) dfs1(i, -1, fid++);
    for (int i = 0; i < g.size(); i++) if (vis[i] != 2) dfs2(i);
    // init_seg();
  }
  void init_seg(vector<T> v) {
    assert(x.size() == v.size());
    for (int i = 0; i < x.size(); i++) seg.set(in[i], x[i] = v[i]);</pre>
  int lca(int u, int v) {
    while (top[u] != top[v]) {
```

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```
if (d[top[u]] < d[top[v]]) swap(u, v);
        int st = top[u];
        u = p[st];
      }
      if (d[u] > d[v]) swap(u, v);
      return u:
    }
    S query(int u, int v, int inc_start = 0) {
      Sr = Se();
      while (top[u] != top[v]) {
       if (d[top[u]] < d[top[v]]) swap(u, v);</pre>
        int st = top[u];
        r = Sop(r, seg.prod(in[st], in[u]+1));
        u = p[st];
      }
      if (d[u] > d[v]) swap(u, v);
      r = Sop(r, seg.prod(in[u]+inc_start, in[v]+1));
      return r:
    }
};
template <typename T, class S, auto Sop, auto Se, class F, auto mapping, auto composition,
struct LazyHeavyLightDecomposition {
  // private:
  public:
    int root, eid;
    vector<int> sz, d, p, top, in, out, vis;
    vector<T> x;
    WeightedFIDXEdges<T> c;
    atcoder::lazy_segtree<S, Sop, Se, F, mapping, composition, id> seg;
    void dfs1(int v, int pv = -1) {
      vis[v] = 1;
      sz[v] = 1:
      for (auto &e: c[v]) {
        if (e.first == pv) continue;
        d[e.first] = d[v]+1;
        p[e.first] = v;
        dfs1(e.first, v);
        sz[v] += sz[e.first];
        if (sz[e.first] > sz[c[v][0].first]) swap(e, c[v][0]);
      }
    }
    void dfs2(int v, int pv = -1) {
      vis[v] = 2:
      in[v] = eid++;
      for (auto &[e, w]: c[v]) {
       if (e == pv) continue:
        top[e] = e == c[v][0].first ? top[v] : e;
        dfs2(e, v);
        x[e] = w:
     }
      out[v] = eid;
```

```
// public:
 LazyHeavyLightDecomposition(WeightedFIDXEdges<T> g, int _root = 0)
   : c(g),
      root(_root),
      eid(0),
      sz(g.size(), 0),
      d(g.size(), 0),
      x(g.size(), 0),
      in(g.size(), -1),
      out(g.size(), -1),
      vis(g.size(), 0),
      seg(atcoder::lazy_segtree<S, Sop, Se, F, mapping, composition, id>(g.size()+1)),
      top(g.size(), root),
     p(g.size(), root) {
    for (int i = 0; i < g.size(); i++) if (vis[i] != 1) dfs1(i);</pre>
   for (int i = 0; i < g.size(); i++) if (vis[i] != 2) dfs2(i);</pre>
   // init seg():
 void init_seg(vector<S> v) {
   assert(x.size() == v.size());
   for (int i = 0; i < x.size(); i++) seg.set(in[i], v[i]);</pre>
  int lca(int u, int v) {
   while (top[u] != top[v]) {
     if (d[top[u]] < d[top[v]]) swap(u, v);</pre>
      int st = top[u];
      u = p[st];
   if (d[u] > d[v]) swap(u, v);
   return u;
 S query(int u, int v, int inc_start = 0) {
   S 1 = Se():
   Sr = Se():
    while (top[u] != top[v]) {
     // if (d[top[u]] < d[top[v]]) swap(u, v);</pre>
      if (d[top[u]] > d[top[v]]) {
        int st = top[u];
       1 = Sop(seg.prod(in[st], in[u]+1), 1);
       u = p[st]:
     } else {
        int st = top[v];
       r = Sop(seg.prod(in[st], in[v]+1), r);
        v = p[st];
   if (d[u] < d[v]) r = Sop(seg.prod(in[u]+inc_start, in[v]+1), r);</pre>
   else 1 = Sop(seg.prod(in[v]+inc_start, in[u]+1), 1);
    swap(1.1, 1.r);
   return Sop(1, r);
```

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```
void apply(int u, int v, F x, int inc_start = 0) {
      while (top[u] != top[v]) {
       if (d[top[u]] < d[top[v]]) swap(u, v);</pre>
       int st = top[u];
       seg.apply(in[st], in[u]+1, x);
       u = p[st];
     }
      if (d[u] > d[v]) swap(u, v);
      seg.apply(in[u]+inc_start, in[v]+1, x);
};
     Bellman Ford
// require HLD-WeightedEdges,
template <tvpename T = int>
vector<T> bellman_ford(int n, kENN3N::WeightedEdges<T> edges, int sv = 0, int ev = -1) {
 T MAX = numeric_limits<T>::max() / 2;
  vector<T> d(n, MAX):
 d[sv] = 0:
  for (int i = 0; i < n; i++) {
    bool upd = false;
    for (auto &e: edges) {
     if (d[e.from] == MAX) continue;
      if (d[e.to] > d[e.from] + e.cost) upd = true, d[e.to] = d[e.from] + e.cost;
    if (!upd) return d;
  if (ev == -1) return vector<T>();
  vector<bool> nega(n, false);
  for (int i = 0; i < n; i++) {
    for (auto &e: edges) {
      if (d[e.from] == MAX) continue;
      if (d[e.to] > d[e.from] + e.cost) nega[e.to] = true, d[e.to] = d[e.from] + e.cost;
      if (nega[e.from]) nega[e.to] = true;
 }
  if (nega[ev]) return vector<T>();
  else return d;
벸만 포드 부등식
방법은, xj - xi <= T일때, i -> j로 가는 가중치 T의 간선을 만들어주는 것이다.
이러한 일련의 작업이 끝났다면, 임의의 정점 i에서 벨만 포드를 돌렸을 때, xj = xi + ShortestPath(i,
j) 라는 값을 대입했을 때 저것을 만족함을 보장할 수 있다.
*/
WeightedEdges<11> g;
for (int i = 0; i < q; i++) {
  int a, b, x;
  string op;
  cin >> a >> b >> op >> x;
  if (op == "<=") g.push_back(Edge<11>(b+n, a, x));
  else g.push_back(Edge<11>(a, b+n, -x));
```

```
for (int i = 1; i <= n+m; i++) g.push_back(Edge<11>(0, i, 0));
auto y = bellman_ford(n+m+1, g);
cout << (y.size() ? "Possible" : "Impossible") << "\n";</pre>
5.6 Maxflow
 Time Complexity: O(n+m)\sqrt{n}, O(n^2m)
template <class Cap> struct mf_graph {
 mf_graph(int n = 0) : _n(n), g(n) {}
 void add_edge(int from, int to, Cap cap) {
   int from_id = g[from].size();
   int to_id = g[to].size();
   g[from].push_back({to, to_id, cap});
   g[to].push_back({from, from_id, 0});
 Cap flow(int s, int t, Cap limit = numeric_limits<Cap>::max()) {
   Cap flow = 0;
   vector<int> level( n). iter( n):
   auto bfs = [&]() {
     fill(level.begin(), level.end(), -1);
     level[s] = 0;
     queue<int> q;
     q.push(s);
     while (!q.empty()) {
       int v = q.front();
       q.pop();
       for (auto& e : g[v]) {
         if (e.cap > 0 && level[e.to] < 0) {
           level[e.to] = level[v] + 1;
           q.push(e.to);
       }
     return level[t] != -1;
   };
   auto dfs = [&](auto self, int v, Cap up) -> Cap {
     if (v == t) return up;
     Cap res = 0;
     for (int& i = iter[v]; i < g[v].size(); i++) {</pre>
        _edge &e = g[v][i], &re = g[e.to][e.rev];
       if (e.cap > 0 && level[v] < level[e.to]) {</pre>
          Cap d = self(self, e.to, min(up - res, e.cap));
          if (d > 0) {
           e.cap -= d;
           re.cap += d;
           res += d;
            if (res == up) break;
         }
     return res;
   };
   while (flow < limit && bfs()) {
     fill(iter.begin(), iter.end(), 0);
```

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```
Cap f = dfs(dfs, s, limit - flow);
      if (!f) break;
      flow += f:
   }
    return flow;
private:
 struct _edge { int to, rev; Cap cap; };
 int _n;
 vector<vector<_edge>> g;
};
mf_graph<int> mf;
5.7 Min cost Max Flow
 Time Complexity: O(F(n+m)\log(n+m))
template <class Cap, class Cost> struct mcf_graph {
 public:
    mcf_graph(int n = 0) : _n(n), g(n) {}
    void add_edge(int from, int to, Cap cap, Cost cost) {
      g[from].push_back({to, (int)g[to].size(), cap, cost});
     g[to].push_back({from, (int)g[from].size() - 1, 0, -cost});
    pair<Cap, Cost> flow(int s, int t, Cap flow_limit = numeric_limits<Cap>::max()) {
      Cap flow = 0:
      Cost cost = 0;
      vector<Cost> h(_n, 0);
      vector<int> prev_v(_n), prev_e(_n);
      while (flow < flow_limit) {</pre>
       vector<Cost> dist(_n, numeric_limits<Cost>::max());
       dist[s] = 0;
       priority_queue<pair<Cost, int>, vector<pair<Cost, int>>, greater<pair<Cost, int>>>
        que.push(\{0, s\});
        while (!que.empty()) {
          auto [d, v] = que.top();
          que.pop();
          if (dist[v] < d) continue;</pre>
          for (int i = 0; i < g[v].size(); i++) {</pre>
            _edge& e = g[v][i];
            if (e.cap > 0 \&\& dist[e.to] > dist[v] + e.cost + h[v] - h[e.to]) {
              dist[e.to] = dist[v] + e.cost + h[v] - h[e.to];
              prev_v[e.to] = v;
              prev_e[e.to] = i;
              que.push({dist[e.to], e.to});
         }
       }
        if (dist[t] == numeric_limits<Cost>::max()) break;
        for (int i = 0; i < _n; i++) h[i] += dist[i];</pre>
```

```
Cap d = flow_limit - flow;
    for (int v = t; v != s; v = prev_v[v]) d = min(d, g[prev_v[v]][prev_e[v]].cap);
    flow += d;
    cost += d * h[t];
    for (int v = t; v != s; v = prev_v[v]) {
        _edge& e = g[prev_v[v]][prev_e[v]];
        e.cap -= d;
        g[v][e.rev].cap += d;
    }
    }
    return {flow, cost};
}
private:
    struct _edge { int to, rev; Cap cap; Cost cost; };
    int _n;
    vector<vector<_edge>> g;
};
```

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6 Math

6.1 Combination

```
// ncr
struct Combination {
  const 11 MAX, MOD;
  vector<ll> inv, fac, fi;
  Combination(ll max, ll mod): MAX(max), MOD(mod) {
    inv = vector<ll>(max+1);
    fac = vector<ll>(max+1):
    fi = vector<ll>(max+1):
    fac[1] = inv[1] = fi[1] = 1;
    for (ll i = 2: i <= max: i++) {
      fac[i] = (fac[i-1]*i) \% mod;
      inv[i] = (mod - mod/i) * inv[mod%i] % mod;
      fi[i] = fi[i-1]*inv[i] % mod:
    }
  }
  ll ncr(ll n, ll r) {
    if (r == 0 || n == r) return 1;
    return (fac[n]*fi[r]%MOD)*fi[n-r]%MOD;
 }
};
6.2 FFT - int
using pll = pair<ll, ll>;
using ull = unsigned long long;
// X.v 와 같이 접근해야함. int 범위 내 NTT 전용
namespace NTT {
  constexpr int MOD = 998244353;
  constexpr int G = 3; // Primitive root for MOD
  struct mint {
    int v:
    mint(11 x = 0) : v(x \% MOD) { if (v < 0) v += MOD; }
    mint& operator+=(const mint& o) { v += o.v; if (v >= MOD) v -= MOD; return *this; }
    mint& operator-=(const mint& o) { v -= o.v; if (v < 0) v += MOD; return *this; }
    mint& operator*=(const mint& o) { v = (int)((11)v * o.v % MOD); return *this; }
    mint operator+(const mint& o) const { return mint(*this) += o; }
    mint operator-(const mint& o) const { return mint(*this) -= o; }
    mint operator*(const mint& o) const { return mint(*this) *= o; }
    mint pow(ll n) const {
      mint res = 1, x = *this;
      while (n > 0) {
       if (n & 1) res *= x;
        x *= x;
       n >>= 1;
      }
      return res;
    mint inv() const { return pow(MOD - 2); }
  };
  int ceil_pow2(int n) {
    int x = 0;
```

```
while ((1U \ll x) \ll (unsigned int)(n)) x++;
 return x:
int bsf(unsigned int n) { return __builtin_ctz(n); }
void butterfly(vector<mint>& a) {
 int n = a.size():
 int h = ceil_pow2(n);
 static bool first = true;
 static vector<mint> sum e:
 if (first) {
   first = false;
   sum e.resize(30):
   mint es[30], ies[30];
   int cnt2 = bsf(MOD - 1);
   mint e = mint(G).pow((MOD - 1) >> cnt2), ie = e.inv();
   for (int i = cnt2; i >= 2; i--) {
     es[i - 2] = e;
     ies[i-2] = ie;
     e *= e:
      ie *= ie:
   mint now = 1;
   for (int i = 0; i < cnt2 - 2; i++) {
      sum_e[i] = es[i] * now;
      now *= ies[i];
 for (int ph = 1; ph <= h; ph++) {
   int w = 1 << (ph - 1), p = 1 << (h - ph);
   mint now = 1:
   for (int s = 0; s < w; s++) {
     int offset = s \ll (h - ph + 1);
     for (int i = 0; i < p; i++) {
       auto 1 = a[i + offset]:
       auto r = a[i + offset + p] * now;
       a[i + offset] = 1 + r:
        a[i + offset + p] = 1 - r:
      now *= sum_e[bsf(~(unsigned int)(s))];
 }
void butterflv inv(vector<mint>& a) {
 int n = a.size();
 int h = ceil_pow2(n);
 static bool first = true;
 static vector<mint> sum ie:
 if (first) {
   first = false;
   sum ie.resize(30):
   mint es[30], ies[30]:
   int cnt2 = bsf(MOD - 1);
   mint e = mint(G).pow((MOD - 1) >> cnt2), ie = e.inv();
   for (int i = cnt2: i \ge 2: i--) {
      es[i - 2] = e; ies[i - 2] = ie;
```

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```
e *= e; ie *= ie;
     }
      mint now = 1;
      for (int i = 0; i < cnt2 - 2; i++) { sum_ie[i] = ies[i] * now; now *= es[i]; }
    for (int ph = h; ph >= 1; ph--) {
      int w = 1 \ll (ph - 1), p = 1 \ll (h - ph);
      mint inow = 1;
      for (int s = 0: s < w: s++) {
       int offset = s << (h - ph + 1);</pre>
       for (int i = 0; i < p; i++) {
          auto l = a[i + offset]:
          auto r = a[i + offset + p];
         a[i + offset] = 1 + r;
          a[i + offset + p] = (l - r) * inow;
       }
       inow *= sum_ie[bsf(~(unsigned int)(s))];
     }
   }
  vector<mint> convolution(vector<mint> a, vector<mint> b) {
    int n = a.size(), m = b.size();
    if (!n || !m) return {};
    if (min(n, m) \le 60) {
     if (n < m) { swap(n, m); swap(a, b); }
      vector<mint> ans(n + m - 1):
     for (int i = 0; i < n; i++) for (int j = 0; j < m; j++) ans [i + j] += a[i] * b[j];
      return ans;
    int z = 1 << ceil_pow2(n + m - 1);
    a.resize(z); b.resize(z);
    butterfly(a); butterfly(b);
    for (int i = 0; i < z; i++) a[i] *= b[i];
    butterfly_inv(a);
    a.resize(n + m - 1);
    mint iz = mint(z).inv();
   for (int i = 0; i < n + m - 1; i++) a[i] *= iz;
    return a;
} // namespace NTT
vector<NTT::mint> a;
6.3 FFT - ll (Dnc + CRT)
using ull = unsigned long long;
// fft ll = DnC + CRT
// ====== AtCoder Convolution Library =========
namespace atcoder {
namespace internal {
  int ceil_pow2(int n) { int x = 0; while ((1U \ll x) \ll (ui)(n)) x++; return x; }
 int bsf(ui n) { return __builtin_ctz(n); }
  constexpr ll safe_mod(ll x, ll m) { x \% = m; if (x < 0) x += m; return x; }
  constexpr pair<11, 11> inv_gcd(11 a, 11 b) {
    a = safe_mod(a, b);
    if (a == 0) return \{b, 0\};
```

```
11 s = b, t = a;
   11 m0 = 0, m1 = 1;
   while (t) {
     s -= t * u:
     auto tmp = s; s = t; t = tmp;
     tmp = m0; m0 = m1; m1 = tmp;
   if (m0 < 0) m0 += b / s;
   return {s, m0};
 constexpr ll pow_mod(ll x, ll n, int m) {
   if (m == 1) return 0;
   ui _m = (ui)(m);
   ull r = 1, y = safe_mod(x, m);
   while (n) { if (n & 1) r = (r * y) \% _m; y = (y * y) \% _m; n >>= 1; }
   return r;
 }
 constexpr bool is_prime(int n) {
   if (n <= 1) return false;
   if (n == 2 || n == 7 || n == 61) return true;
   if (n % 2 == 0) return false:
   11 d = n - 1;
   while (d \% 2 == 0) d /= 2;
   constexpr ll bases[3] = \{2, 7, 61\};
   for (ll a : bases) {
     11 t = d, y = pow_mod(a, t, n);
     while (t != n - 1 \&\& y != 1 \&\& y != n - 1) { y = y * y \% n; t <<= 1; }
     if (y != n - 1 && t % 2 == 0) return false;
   }
   return true;
 constexpr int primitive_root(int m) {
   if (m == 998244353) return 3;
   if (m == 2) return 1:
   int divs[20] = \{2\}, cnt = 1;
   int x = (m - 1) / 2;
   while (x \% 2 == 0) x /= 2;
   for (int i = 3; (11)(i) * i <= x; i += 2) {
     if (x \% i == 0) \{ divs[cnt++] = i; while <math>(x \% i == 0) \times /= i; \}
   if (x > 1) divs[cnt++] = x:
   for (int g = 2; g++) {
     bool ok = true;
     for (int i = 0; i < cnt; i++) {
       if (pow_mod(g, (m - 1) / divs[i], m) == 1) { ok = false; break; }
     if (ok) return g;
   }
 template <class T> using is_integral = typename std::is_integral<T>;
} // namespace internal
template <int m>
struct static_modint {
```

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```
using mint = static_modint;
  static constexpr int mod() { return m; }
  static_modint(ll x = 0) : v(x \% m) \{ if (v < 0) v += m; \}
  mint& operator+=(const mint& o) { v += o.v; if (v >= m) v -= m; return *this; }
  mint& operator-=(const mint& o) { v -= o.v; if (v < 0) v += m; return *this; }
 mint& operator*=(const mint& o) { v = (int)((11)v * o.v % m); return *this; }
  mint& operator/=(const mint& o) { return *this = *this * o.inv(); }
  mint operator+(const mint& o) const { return mint(*this) += o; }
  mint operator-(const mint& o) const { return mint(*this) -= o; }
  mint operator*(const mint& o) const { return mint(*this) *= o; }
  mint operator/(const mint& o) const { return mint(*this) /= o: }
 mint pow(ll n) const {
    mint res = 1, x = *this;
    while (n > 0) { if (n \& 1) res *= x; x *= x; n >>= 1; }
    return res:
  mint inv() const {
    if (prime) { assert(v); return pow(m - 2); }
    else { auto eg = internal::inv_gcd(v, m); assert(eg.first == 1); return eg.second; }
  static constexpr bool prime = internal::is_prime(m);
}:
namespace internal {
  template <class mint> void butterfly(vector<mint>& a) {
    constexpr int g = primitive root(mint::mod());
    int n = a.size(), h = ceil_pow2(n);
    static vector<mint> sum_e;
    if (sum e.emptv()) {
      sum e.resize(30):
     mint es[30], ies[30];
      int cnt2 = bsf(mint::mod() - 1);
      mint e = mint(g).pow((mint::mod() - 1) >> cnt2), ie = e.inv();
      for (int i = cnt2: i >= 2: i--) {
        es[i - 2] = e; ies[i - 2] = ie; e *= e; ie *= ie;
      mint now = 1:
      for (int i = 0; i < cnt2 - 2; i++) {
        sum_e[i] = es[i] * now; now *= ies[i];
    }
    for (int ph = 1; ph <= h; ph++) {
      int w = 1 \ll (ph - 1), p = 1 \ll (h - ph);
      mint now = 1:
      for (int s = 0; s < w; s++) {
       int offset = s \ll (h - ph + 1);
       for (int i = 0; i < p; i++) {
          auto l = a[i + offset], r = a[i + offset + p] * now;
          a[i + offset] = 1 + r; a[i + offset + p] = 1 - r;
        now *= sum e[bsf(~(ui)(s))]:
   }
  template <class mint> void butterfly_inv(vector<mint>& a) {
```

```
constexpr int g = primitive_root(mint::mod());
   int n = a.size(), h = ceil_pow2(n);
   static vector<mint> sum_ie;
   if (sum_ie.empty()) {
      sum_ie.resize(30);
     mint es[30], ies[30]:
     int cnt2 = bsf(mint::mod() - 1);
     mint e = mint(g).pow((mint::mod() - 1) >> cnt2), ie = e.inv();
     for (int i = cnt2: i >= 2: i--) {
        es[i - 2] = e; ies[i - 2] = ie; e *= e; ie *= ie;
     mint now = 1:
     for (int i = 0; i < cnt2 - 2; i++) {
        sum_ie[i] = ies[i] * now; now *= es[i];
   }
   for (int ph = h; ph >= 1; ph--) {
     int w = 1 \ll (ph - 1), p = 1 \ll (h - ph);
     mint inow = 1:
     for (int s = 0; s < w; s++) {
       int offset = s \ll (h - ph + 1);
       for (int i = 0; i < p; i++) {
         auto l = a[i + offset], r = a[i + offset + p];
         a[i + offset] = 1 + r; a[i + offset + p] = (1 - r) * inow;
        inow *= sum ie[bsf(~(ui)(s))]:
   }
} // namespace internal
template <class mint>
vector<mint> convolution(vector<mint> a, vector<mint> b) {
 int n = a.size(), m = b.size();
 if (!n || !m) return {}:
 int z = 1 \ll internal::ceil_pow2(n + m - 1);
 a.resize(z); b.resize(z);
 internal::butterfly(a); internal::butterfly(b);
 for (int i = 0; i < z; i++) a[i] *= b[i];
 internal::butterfly_inv(a);
 a.resize(n + m - 1):
 mint iz = mint(z).inv();
 for (int i = 0; i < n + m - 1; i++) a[i] *= iz;
template <ui mod = 998244353, class T, enable_if_t<internal::is_integral<T>::value>* =
vector<T> convolution_int(const vector<T>& a, const vector<T>& b) {
 int n = a.size(), m = b.size();
 if (!n || !m) return {}:
 using mint = static_modint<mod>;
 vector<mint> a2(n), b2(m);
 for (int i = 0; i < n; i++) a2[i] = mint(a[i]);
 for (int i = 0; i < m; i++) b2[i] = mint(b[i]);
 auto c2 = convolution(move(a2), move(b2));
 vectorT > c(n + m - 1);
```

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```
for (int i = 0; i < n + m - 1; i++) c[i] = c2[i].v;
 return c:
vector<ll> convolution_l1(const vector<ll>& a, const vector<ll>& b) {
  int n = a.size(), m = b.size();
  if (!n || !m) return {}:
  static constexpr ui MOD1 = 754974721; // 2^24 * 45 + 1
  static constexpr ui MOD2 = 167772161; // 2^25 * 5 + 1
  static constexpr ui MOD3 = 469762049: // 2^26 * 7 + 1
  auto c1 = convolution_int<MOD1>(a, b);
  auto c2 = convolution_int<MOD2>(a, b);
  auto c3 = convolution int<MOD3>(a, b):
 11 m1_m2 = (11)MOD1 * MOD2;
 11 m1_inv_m2 = internal::inv_gcd(MOD1, MOD2).second;
 11 m1_m2_inv_m3 = internal::inv_gcd(m1_m2, MOD3).second;
  vector<ll> c(n + m - 1):
  for (int i = 0: i < n + m - 1: i++) {
    11 v1 = c1[i], v2 = c2[i], v3 = c3[i];
    11 x = v1:
    11 v = (v2 - x) * m1_inv_m2 \% MOD2;
    if (y < 0) y += MOD2;
    x += y * MOD1;
    11 z = (v3 - x \% MOD3) * m1 m2 inv m3 \% MOD3;
    if (z < 0) z += MOD3;
    c[i] = x + z * m1_m2;
 return c;
}} // namespace atcoder
6.4 Pollard Rho (with fast prime check)
// \text{ time: } O(n^{(1/4)})
using i64 = __int128_t;
i64 \text{ my\_abs}(i64 \text{ x}) \{ \text{return } (x < 0 ? -1*x : x); \}
i64 primes[] = { 2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 41 };
// i64 primes[] = { 2, 7, 61 }; // int only
i64 f(i64 x, i64 p, i64 mod) {
 i64 r = 1:
 x \%= mod;
  while (p) {
   if (p\%2) r *= x, r %= mod;
   p /= 2, x *= x, x %= mod;
  return r;
// a is test number, a^d = 1 \pmod{N}, a^d = 1 \pmod{N}, a^d = 1 \pmod{N}, a^d = 1 \pmod{N}
bool miller_rabin(i64 n, i64 a) {
 if (n == a) return true;
  if (a % n == 0) return false:
  i64 k = n-1;
  while (true) {
    i64 t = f(a, k, n):
    if (t == n-1) return true;
    if (k\%2) return (t == 1 || t == n-1);
```

k /= 2:

```
bool isp(i64 x) {
 bool p = true;
 for (auto &e: primes) {
   p &= miller_rabin(x, e);
   if (!p) return false;
 return true:
void g(i64 x, vector<i64> &v) {
 if (x == 1) return:
 if (x\%2 == 0) \{ v.push_back(2); g(x/2, v); return; \}
 if (isp(x)) { v.push_back(x); return; }
 i64 a, b, c, w = x;
 auto h = [\&] (i64 y) { return (c + y*y%x) % x; };
   if (w == x) a = b = rand()%(x-2) + 2, c = rand() % 20 + 1;
   a = h(a), b = h(h(b)), w = __gcd(my_abs(a-b), x);
 } while (w == 1):
 g(w, v);
 g(x/w, v);
// source: jinhan, factorize code
vector<pair<i64, int>> pollard_rho(i64 x) {
 vector<i64> v:
 vector<pair<i64, int>> r;
 g(x, v);
 sort(v.begin(), v.end());
 for (auto e: v) {
   if (r.size() && r.back().first == e) r.back().second++;
   else r.push_back({ e, 1 });
 return r;
int main() {
 11 x:
 while (cin >> x) {
   11 r = 1;
   for (auto [e, c]: pollard_rho(x)) r *= (c+1);
   cout << r << "\n";
6.5 Permutation and Combination
중복 순열
서로 다른 n개의 원소 중에서 중복을 허용하면서 r개를 순서있게 뽑는 방법의 수
즉, 한 원소를 여러번 뽑을 수 있고, 뽑은 순서를 고려함
print(n**r)
```

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```
중복 조합
서로 다른 n개의 원소 중에서 중복을 허용하면서 r개를 순서 없이 뽑는 방법의 수
즉, 한 원소를 여러번 뽑을 수 있지만, 순서는 고려하지 않음
math.comb(n+r-1, r)
중복 원소가 있는 순열
동일한 원소가 여러 번 있는 집합(예: AAB)에서 원소들을 모두 사용하여 순열을 만드는 방법의 수
전체 원소의 개수가 n이고, 중복된 원소의 개수가 n1,n2,...,nk일 때
math.factorial(n) / (math.factorial(n1) * math.factorial(n2) * ...)
완전 순열
순열의 일종으로, 일렬로 배열한 대상들의 위치를 재조정했을 때,
모든 대상이 자기 위치에 있지 않도록 하는 배열 방법
D[n] = (n-1)(D[n-1]+D[n-2])
D[0] = 1
D[1] = 0
D[2] = 1
D[3] = 2
D\lceil 4\rceil = 9
D[5] = 44
```

6.6 Pairwise Product Sum

```
# 모든 i<i에 대해 arr[i]*arr[i]의 합
# 1*2 + 1*3 + 1*4 + 2*3 + 2*4 + 3*4
arr = [1,2,3,4]
S = sum(arr)
sum_sq = sum(x*x for x in arr)
result = (S*S - sum_sq) // 2
print(result) # 35
```

6.7 Modular Divide

```
\# (a/b)\%M \Rightarrow (a\%M * pow(b,M-2,M)) \% M
M = 1000000007
a=7
print((a%M * pow(b,M-2,M)) % M)
```

7 String

7.1 KMP

```
vector<int> get_pi(string p) {
 vector<int> pi(p.length(), 0);
 for (int i = 1, j = 0; i < p.length(); i++) {
   while (j \&\& p[i] != p[j]) j = pi[j-1];
   if (p[i] == p[j]) pi[i] = ++j;
 return pi;
int kmp(string s, string p) {
 auto pi = get_pi(p);
 int lj = 0;
 for (int i = 0, j = 0; i < s.length(); i++) {
   1j = j;
   while (j \&\& s[i] != p[j]) j = pi[j-1];
   if (s[i] == p[j]) {
     if (j == p.length() - 1) {
       // result.push_back(i-p.length()+1); if want result string
       j = pi[j];
     } else j++;
   }
 }
 return lj;
7.2 Z
// 접미사와 문자열전체접두사와 가장 긴 공통접두사
template <class T> vector<int> z_algorithm(const vector<T>& s) {
 int n = int(s.size()):
 if (n == 0) return {}:
 vector<int> z(n);
 z[0] = 0;
 for (int i = 1, j = 0; i < n; i++) {
   int & k = z[i];
   k = (j + z[j] \le i) ? 0 : std::min(j + z[j] - i, z[i - j]);
   while (i + k < n \&\& s[k] == s[i + k]) k++;
   if (j + z[j] < i + z[i]) j = i;
 z[0] = n;
 return z:
vector<int> z_algorithm(const string& s) {
 int n = int(s.size()):
 vector<int> s2(n);
 for (int i = 0; i < n; i++) s2[i] = s[i];
 return z_algorithm(s2);
```

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7.3 Trie and Ahocorasick

```
// include Trie
struct Trie {
  map<char, Trie*> ch:
  Trie* fail;
  bool end:
 Trie() { end = false; }
  ~Trie() {
    for (auto trie: ch) delete trie.second;
    ch.clear():
  void insert(string s, int i = 0) {
    if (i == s.length()) {
      this->end = true;
      return:
    if (!ch[s[i]]) ch[s[i]] = new Trie();
    ch[s[i]]->insert(s, i+1):
 }
};
Trie trie;
void set_fail() {
  queue<Trie*> q;
 trie.fail = ≜
  q.push(&trie);
  while (q.size()) {
    Trie* cur = q.front();
    q.pop();
    for (auto [c, nxt]: cur->ch) {
      if (cur == &trie) nxt->fail = ≜
      else {
        auto dest = cur->fail;
        while (dest != &trie && dest->ch.find(c) == dest->ch.cend()) dest = dest->fail;
        if (dest->ch.find(c) != dest->ch.cend()) dest = dest->ch[c]:
        nxt->fail = dest:
      if (nxt->fail->end) nxt->end = true;
      q.push(nxt);
 }
int main() {
  // put all string to trie
  set_fail();
  cin >> n:
  while (n--) {
    cin >> s;
    bool result = false:
    Trie* cur = ≜
    for (auto c: s) {
      while (cur != &trie && cur->ch.find(c) == cur->ch.cend()) cur = cur->fail:
      if (cur->ch.find(c) != cur->ch.cend()) cur = cur->ch[c];
      if (cur->end) result = true;
```

```
cout << (result ? "YES" : "NO") << (n ? "\n" : "");
 }
7.4 Suffix Automaton
struct Suffix Automaton {
 struct Node {
   int len. link. last. cnt = 1:
   int minpos = 0, maxpos = 0;
   // bool isleaf = true;
   map<int, int> nxt:
   // int nxt[26]:
   Node(int len = 0, int link = 0, int last = 0): len(len), link(link), last(last) {}
 };
 vector<Node> v;
 vector<pii> sv; // len / original i
 int total;
 // ll substr_cnt = 0;
 Suffix_Automaton() { total = 0; v.push_back(Node(0, -1)); }
 void add(int c, int last = 0) {
   // int c = _c - 'A';
   v.push_back(Node(v[total].len+1, 0, last));
   int p = total;
   v[p].minpos = v[p].maxpos = v[p].len-1;
   total = v.size()-1;
   // cout << "[debug]" << (char)c << ": " << v[p].len << "." << v[p].link << "." <<
   v[p].minpos << "," << v[p].maxpos << "\n";</pre>
   while (p != -1 \&\& v[p].nxt.find(c) == v[p].nxt.cend()) {
     v[p].nxt.insert({ c, total });
     // v[p].nxt[c] = total;
     // if (v[p].link != -1) substr_cnt += v[p].len-v[v[p].link].len;
     // else substr_cnt++;
     p = v[p].link;
   if (p != -1) {
     int q = v[p].nxt[c];
     int upd = a:
     if (v[p].len+1 < v[q].len) {
        upd = v.size();
       Node clone = v[q];
       clone.len = v[p].len + 1;
        clone.cnt = 0;
       v.push_back(clone);
       v[a].link = upd:
       for (int j = p; j != -1 && v[j].nxt[c] == q; j = v[j].link) v[j].nxt[c] = upd;
      v[total].link = upd;
 }
 // len, i
 pii lcs(string s) {
   int cur = 0;
   Node x = v[0];
   int 1 = 0, rl = 0, ri = 0;
```

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```
for (int i = 0; i < s.length(); i++) {</pre>
   x = v[cur]:
   // cout << i << ": " << x.len << ", " << x.link << "\n";
   // cout << i << "(" << s[i] << "): " << x.len << " " << x.link << " => ";
   int idx = s[i] - 'a':
    // for (auto e: x.nxt) cout << e << " "; cout << " => ";
    while (x.len && x.nxt[idx] == -1) {
     // 1 -= (x.len-v[x.link].len);
     1 = v[x.link].len;
     cur = x.link;
     x = v[x.link]:
    if (x.nxt[idx] != -1) cur = v[cur].nxt[idx], l++;
    if (1 > rl) rl = 1, ri = i;
   // cout << x.len << " " << x.link << "\n";
  return { rl, ri };
void clear() {
 total = 0;
 v.clear();
 v.push_back(Node(0, -1, 0));
// vector<int> search(string t) {
// vector<int> pos:
// int cur = 0, 1 = 0;
// for (int i = 0; i < t.length(); i++) {</pre>
      while (v[cur].len && v[cur].nxt.find(t[i]) == v[cur].nxt.cend()) {
11
        cur = v[cur].link;
//
        1 = v[cur].len;
      if (v[cur].nxt.find(t[i]) != v[cur].nxt.cend()) cur = v[cur].nxt[t[i]], 1++;
//
//
      if (t.length() == 1) pos.push_back(i-1+2);
// }
// return pos;
// bool search(string t) {
// int cur = 0, 1 = 0;
// for (int i = 0; i < t.length(); i++) {</pre>
11
      while (v[cur].len && v[cur].nxt.find(t[i]) == v[cur].nxt.cend()) {
11
        cur = v[cur].link;
        1 = v[cur].len:
//
      if (v[cur].nxt.find(t[i]) != v[cur].nxt.cend()) cur = v[cur].nxt[t[i]], 1++;
      if (t.length() == 1) pos.push_back(i-1+2);
//
// }
// return pos;
// }
void sort() {
  sv = vector<pii>(v.size());
  for (int i = 0; i < v.size(); i++) sv[i] = { v[i].len, i };</pre>
 auto comp = [](pii a, pii b) -> bool {
   return a.first > b.first:
 };
```

```
std::sort(sv.begin(), sv.end(), comp);
 }
};
Suffix_Automaton sa;
```

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8 etc

8.1 Deque Trick

```
int n, 1;
deque<pri>deque<pri>dq;
cin >> n >> 1;
for (int i = 0; i < n; i++) {
   int x;
   cin >> x;
   while (dq.size() && dq.back().first >= x) dq.pop_back();
   dq.push_back({x, i});
   while (dq.size() && dq.front().second <= i-1) dq.pop_front();
   cout << dq.front().first << " ";
}</pre>
```

8.2 imos rectangle and diamond

```
const int FIX = 5. DIV = 1e9:
int ty, px, py, qx, qy, r;
for (int i = 0; i < k; i++) {
  cin >> ty;
  if (tv == 1) {
    cin >> px >> py >> qx >> qy;
    px += FIX;
    py += FIX;
    qx += FIX;
    qy += FIX;
    rect[py][px] += 1;
    rect[qy+1][px] -= 1;
    rect[py][qx+1] -= 1;
    rect[qy+1][qx+1] += 1;
    cin >> px >> py >> r;
    px += FIX;
    py += FIX;
    dia[py-r][px] += 1;
    dia[py-r+1][px] += 1;
    dia[py+1][px-r-1] = 1;
    dia[py+1][px-r] -= 1;
    dia[py+1][px+r+1] -= 1;
    dia[py+1][px+r] = 1;
    dia[py+r+1][px] += 1;
    dia[py+r+2][px] -= 1;
 }
}
for (int i = 0; i < h+FIX; i++) for (int j = 1; j < w+FIX; j++) rect[i][j] += rect[i][j-1];
for (int i = 0; i < w+FIX; i++) for (int j = 1; j < h+FIX; j++) rect[j][i] += rect[<math>j-1][i];
for (int i = 1; i < h+FIX; i++) for (int j = 1; j < w+FIX; j++) dia[i][j] += dia[i-1][j-1];
for (int i = 1; i < h+FIX; i++) for (int j = 0; j < w-1+FIX; j++) dia[i][j] +=
dia[i-1][j+1];
for (int i = 0+FIX; i < h+FIX; i++) {</pre>
  for (int j = 0+FIX; j < w+FIX; j++) {</pre>
    if ((rect[i][j] + dia[i][j]) % 2) cout << "#";</pre>
    else cout << ".";</pre>
```

```
cout << "\n";
8.3 imos triangle
int x, y, s; // (x, y), (x+s, y), (x+s, y+s) = 정점으로 하는 삼각형
for (int i = 0; i < m; i++) {
 cin >> x >> y >> s;
 pos[x][y]++;
 pos[x][y+1]--;
 pos[x+s+1][y]--;
 pos[x+s+2][v+1]++:
 pos[x+s+1][y+s+2]++;
 pos[x+s+2][y+s+2]--;
for (int i = 1; i \le n; i++) for (int j = 1; j \le n; j++) pos[i][j] += pos[i][j-1];
for (int i = 1; i \le n; i++) for (int j = 1; j \le n; j++) pos[i][j] += pos[i-1][j];
for (int i = 1; i \le n; i++) for (int j = 1; j \le n; j++) pos[i][j] += pos[i-1][j-1];
int cnt = 0:
for (int i = 1: i \le n: i++) for (int i = 1: i \le i: i++) if (pos[i][i] > 0) cnt++:
8.4 mo's algorithm
// O(nsqrt(n))
struct Query {
 11 1, r;
 int i;
};
ll n, m, cnt[100001], cntn[100001], ans[100001], sqrtn, a[100001];
Query query[100000];
int main() {
 cin >> n;
 for (int i = 1; i <= n; i++)
   cin >> a[i]:
 cin >> m;
 for (int i = 0; i < m; i++) {
   cin >> query[i].l >> query[i].r;
   query[i].i = i;
 sqrtn = sqrt(n);
 sort(query, query+m, [](Query q1, Query q2){
   if (q1.1/sqrtn == q2.1/sqrtn) return q1.r < q2.r;
   else return q1.1/sqrtn < q2.1/sqrtn;</pre>
 }):
 cntn[0] = n;
 11 mcnt = 0, s = query[0].1, e = query[0].r;
 for (int i = s; i <= e; i++) {
   cntn[cnt[a[i]]]--;
   cnt[a[i]]++:
   cntn[cnt[a[i]]]++;
   mcnt = max(mcnt, cnt[a[i]]);
```

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```
ans[query[0].i] = mcnt;
  int temp;
  for (int i = 1; i < m; i++) {
    while (s < query[i].1) {</pre>
      cntn[cnt[a[s]]]--;
      cnt[a[s]]--;
      cntn[cnt[a[s]]]++;
      s++:
      while (cntn[mcnt] == 0) mcnt--;
    while (query[i].1 < s) {</pre>
      s--;
      cntn[cnt[a[s]]]--;
      if (cnt[a[s]] == mcnt) mcnt++;
      cnt[a[s]]++;
      cntn[cnt[a[s]]]++;
    while (e < query[i].r) {</pre>
      e++:
      cntn[cnt[a[e]]]--;
      if (cnt[a[e]] == mcnt) mcnt++;
      cnt[a[e]]++:
      cntn[cnt[a[e]]]++;
    while (query[i].r < e) {</pre>
      cntn[cnt[a[e]]]--;
      cnt[a[e]]--;
      cntn[cnt[a[e]]]++:
      while (cntn[mcnt] == 0) mcnt--;
    ans[query[i].i] = mcnt;
}
```

8.5 random generator

```
// init
random_device rd;
mt19937 gen(rd());
uniform_int_distribution<int> dis(0, n);
// use
int x = dis(gen);
```

Note and Checklist

9.1 etc formular

```
// Area = I + B/2 - 1, B: 경계격자점, I: 내부격자점
// C^2 = A^2 + B^2 - 2ABcos
// 카탈란수: Cn = { 1/(n+1) } * (2n C n)
```

9.2 pisano

// 주기의 길이를 P라고 하면 N번째 피보나치 수를 M으로 나눈 나머지는 N%P번째 피보나치 수를 M으로 나눈 나머지와 같다.

// 주기는 M = 10^k (k>2)일때, 항상 15 * 10^(k-1)이다.

9.3 Ternary Search

```
/*
while (1+2 < r)로 해두고
i = 1 ~ r까지 순회해서 찾기
*/
```

9.4 문제 풀이 체크리스트

- 비슷한 문제를 풀어본 적이 있던가?
- 단순한 방법에서 시작할 수 있을까? (Brute Force)
- 내가 문제를 푸는 과정을 수식화할 수 있을까? (예제를 직접 해결해보면서)
- 문제를 단순화할 수 없을까?
- 그림으로 그려볼 수 있을까?
- 수식으로 표현할 수 있을까?
- 문제를 분해할 수 있을까?
- 뒤에서부터 생각해서 풀 수 있을까?
- 순서를 강제할 수 있을까?
- 특정 형태의 답만을 고려할 수 있을까? (정규화)
- 구간을 통째로 가져간다 : 플로우 + 적당한 자료구조 (i, i+1, k, 0), (s, e, 1, w), (N, T, k, 0)
- a = b : a만 움직이기, b만 움직이기, 두 개 동시에 움직이기, 반대로 움직이기
- 말도 안 되는 것들을 한 번은 생각해보기 / "당연하다고 생각한 것" 다시 생각해보기
- 확률 : DP, 이분 탐색(NYPC 2019 Finals C)
- 최대/최소 : 이분 탐색, 그리디(Prefix 고정, Exchange Argument), DP(순서 고정)