1 2 3 4

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acta	gaac	gaac	aagc
gaac	ctag	aagc	acta
aagc	tcat	tcat	ctag
ctag	ttct	gcgt	gaac
			-
tcat	aagc	acta	gcgt
gcgt	gcgt	ctag	tcat
ttct	acta	ttct	ttct

a.

- b. the size k-mer substring can be calculated to be n-k+1. The time it takes to check if the k-mer substring is in the corpus takes time $\theta(m-k+1)=\theta(m)+\theta(n\log m)=\theta(m+n\log m)$ because the k and 1 are irrelevant as the function grows towards infinity, and the nlogm comes from a tree constructed of k-mer substrings with the worst case scenario being log m, which is then multiplied by the n elements resulting in n*log m.
- c. The problem in b can be solved more efficiently by using radix sort on the k-mer substrings and the corpus. This results in time n-k+1 for the k-mer substring and m-k+1 to check the corpus string. Therefore the worst case time would be $\theta(m-k+1) + \theta(n-k+1) = \theta(m+n)$ due to the fact the k and 1s become irrelevant as the function grows towards infinity.