


Build parser using Bison parser generator

→ Accepts language generated by mini GLSL

Implement true parser functionality of the compiler

make man in starter

No Semantic Analysis / AST required for now

Complete implementation of scanner.l, dummy parser.y file

CPG for Mini GLSL

Program → Scope

Scope → '{' declaration statements '}'

declarations → declarations declaration
→ ε

Statements → Statements statement
→ ε

declaration → type < identifier > ';'

→ type < identifier > '=' expression ';'

→ 'const' type < identifier > '=' expression ';'

→ ε

Statement → Variable = 'expression' ';'

→ 'if' '(' expression ')' statements else_statement

→ 'while' '(' expression ')' statement

→ Scope
→ ';'

else_statement → 'else' statement
→ ε

type → 'int' | 'ivec2' | 'ivec3' | 'ivec4'

→ 'bool' | 'bvec2' | 'bvec3' | 'bvec4'

→ 'float' | 'vec2' | 'vec3' | 'vec4'

expressions → constructor

→ function

→ < integer literal >

→ floating point literal

→ variable

→ unary_op - expression

→ expression binary_op expression

→ 'true' | 'false'

→ '(' expression ')'

variable → < identifier >

→ < identifier > '(' < integer literal > ')'

unary_op → '!' | '-'

binary_op → '&&' | '||' | '==' | '!=' | '<' | '<=' | '>' | '>=' | '+' | '-' | '*' | '/' | '%'

Constructor → type '(' arguments ')'

function → function_name '(' arguments_opt ')'

function_name → 'dp3' | 'lit' | 'req'

arguments_opt → arguments | ε

arguments → arguments ';' expression | expression

Precedence		
0	[] () { }	L to R
1	- !	
2	^	R to L
3	* /	
4	+ -	
5	= < > <= >=	
6	& & AND	
7	OR	

Integers $2^{21}, 2^{21}$ base 10 $\rightarrow 0 \mid [1-10] [0-10]^*$

Floats $\cdot [0-10]^+$

Identifiers $[a-zA-Z_][a-zA-Z0-9_]^*$

Parser is only responsible for syntax only

Implement trace scanner functionality - In switch

Build lexical analyzer Mini GLSO

make men in starter dir

Compiler 462.c

Makedir

Scanner.h - skeleton flex Scanner

Parser.y - skeleton bison parser

Compiler 467.man - man page / manual

globalvars.c - global vars

common.h - global definitions

Store appropriate info for each token

identified in a tracing file, errors in error file

When -In on, global var trace Scanner is TRUE

Send traces to global trace file FILE*

Use y TRACE(x) to output the trace

Output tokens in same order that they appear in the program

Report error on illegal input

=> y ERROR(x) to report

check for out of bounds integers & identifiers that exceed the allowed length

Define tokens in parser.y file

Bison only takes the tokens and defines them in a header file - parser.tab.h

Associate info with tokens to store important information

yyval union

Bison defines yyval to be of union type from parser.y file

Use yyval in Scanner.h to store the value of each token

Parser - ask Scanner for new tokens

Add new tokens defined to grammar in
parser.y

What is Lex / Flex

Takes .l file \Rightarrow generates a lexical analyzer
implemented in C lex.yy.c

Lex uses reg expressions

header
% %
body
% %
helper functions

Header Section

place header file
% {
#include "helper.h"
% }

place definitions
Syntax <name> <definitions>

DIGIT [0-9]

Body Section

Lexical rules placed here

<rule> <action>

[0-9] +

{ DIGIT } + ^ . { DIGIT } *

Helper Section

C functions here

Lex func & Variables

char * yytext - current matched text
int yyleng - length of matched text

int yyloc? Scanner function

char input C? - get next char in input stream

int yyterminate C? - terminate scanner - return 0

Scanner - breaks input into tokens

Parser - checks grammar & builds structure from those tokens

Scanner.l

comment handler /*

catch all error \rightarrow { y ERROR(...) }

y TRACE

Add tokens def in parser.y

Write regex patterns in scanner.l to recognize each token

Return correct token, call y TRACE

Regex - $[0-9]^+$ for INT-C

Flr matches pattern in order - if error pattern comes first, it catch everything

Parser.y \Rightarrow Token declarations %token INT-C - read char, recognize pattern, return token types

Scanner.l \Rightarrow Pattern + action { DIGIT } + { ... functions } \rightarrow receive tokens from Scanner & check grammar

C style comments work for both

Flr rules

Flr matches longest possible string

when lengths are equal, first rule wins

Keywords come BEFORE identifiers *

Float comes BEFORE int checks *

Keywords

Bool

Type keywords

operators

Float

Integers

Identifiers

Single char tokens

White space

New line

Token # - Bison assigns to each token

When Bison processes parser.y - it generates parser.tab.h

Defines each token as an integer constant

Single char tokens use their ASCII values

input & output FILE defined

yy parse() → yy lex()

get opts → parse command line arguments

file open

Source Dump

Scanner.L

include "parser.tab.h"

yyin = input file

flex input → yyinput

yyTRACE

yy ERROR → yy terminate

yyline = 1 initially

% option noyywrap

% pattern {action} %

"string" { program }

IF defined in parser.tab.h

- matches any single char
⇒ catch all

Define tokens in parser.y

Assign each token an integer ⇒ Scanner returns these values

Grammar rules

parser.out ⇒ LR parser state machine

Bison → parser.y ⇒ Parser.C
parser.tab.h

Scanner.L → Scanner.C

.C → .o

→ compiler program