SQL Server

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## Cheat Sheets

### Select

#### Phases

➎ SELECT empid, YEAR(orderdate) AS orderyear, COUNT(\*) AS numorders

➊ FROM Sales.Orders

➋ WHERE custid = 71

➌ GROUP BY empid, YEAR(orderdate)

➍ HAVING COUNT(\*) > 1

➏ ORDER BY empid, orderyear;

|  |  |
| --- | --- |
| ➊ FROM | Specify the table we want to query |
| ➋ WHERE | Uses a predicate to filter the rows returned by the FROM phrase. Where clauses enable the use of indices to improve performance and reduce the network traffic that would occur if we performed a table scan and filtered on the client. |
| ➌ GROUPBY | Produce a group for each unique combination of values specified in this clause. |
| ➍ HAVING | Uses a predicate to filter the groups returned. Can utilise aggregate functions in the predicate. |
| ➎ SELECT | Specify the columns we want to see in the result |
| ➏ ORDER BY | Sort the rows for presentation purposes |

##### GROUP BY

If a query contains a group by phase any subsequent HAVING, SELECT, and ORDERBY clauses work on groups. As such they can only operate on expressions that return a single scalar value per group. Any fields specified in the GROUP BY phase implicitly have this process. Any elements that do not meet this restriction can only be used as inputs to aggregation functions such as COUNT, SUM, AVG, MIN, MAX.

##### ORDER BY

In SQL, a table has no order. By using order by the result is ordered and hence cannot be considered a table. In SQL the ordered rows are referred to as a cursor. Unlike all other phases the order by phase can utilise column alisases defined in the select phase as the order by is the only phase that follows the select phase. Returns a cursor rather than table.

#### Phase Illustrations

##### SELECT

➋ SELECT \*

➊ FROM TelNumber

A picture containing chart

Description automatically generated

##### WHERE

The Where phrase adds a predicate to filter the results. ➌ Select \*

➊ FROM TelNumber

➋ WHERE PersonId = 1

A screenshot of a computer

Description automatically generated with medium confidence

##### GROUP BY

Allows grouping. The select can only work on columns appropriate to the grouping.

➍ Select PersonId, COUNT(Number) AS 'Number Count'

➊ FROM TelNumber

➋ WHERE PersonId IN (1,2)

➌ GROUP BY PersonId

A picture containing diagram

Description automatically generated

##### HAVING

The having Phase filters the results of the Group By phase

➎ Select PersonId, COUNT(Number) AS 'Number Count'

➊ FROM TelNumber

➋ WHERE PersonId IN (1,2)

➌ GROUP BY PersonId

➍ HAVING count(Number) > 1

A picture containing diagram

Description automatically generated

##### ORDER BY

Order By returns a cursor rather than a table.

➎SELECT PersonId, COUNT(Number) AS 'Number Count'

➊FROM dbo.TelNumber

➋WHERE PersonId in (1,2)

➌GROUP BY PersonId

➍Having Count(Number) >= 1

➏ORDER BY [Number Count] DESC

Usage Notes

* It is the only Phase that can access aliases from the SELECT Phase.
* We can order by attributes not in the select phase if we do not use DISTINCT.
* If we use DISTINCT, the ORDER BY elements must be in the SELECT phase.

The reason being is that with DISTINCT there are multiple rows for each distinct value and it is not clear which one to use for the ordering.

### Data Modification

|  |  |  |
| --- | --- | --- |
| Method | Code | Description |
| **INSERT INTO** | INSERT INTO Person  (firstname, secondname)  VALUES  ('Indiana', 'Jones'),  ('Carl', 'Jones') | Non-standard T-SQL. All rows treated as transaction |
| **SELECT INTO** | SELECT \*  INTO dbo.Position2  FROM Position | Non-standard statement that creates a new table if it does not exist. |
| **INSERT SELECT** | INSERT INTO Position2  SELECT \*  FROM Position | Uses a select to define the rows to insert into a table. |
| **INSERT EXEC** | CREATE PROC GetPositions  AS  SELECT \* From Position  GO  INSERT INTO Position2  EXEC GetPositions | Insert from a stored procedure. |
| **BULK INSERT** |  | Insert from a file |
| **Create table on fly.** | SELECT \*  FROM ( VALUES  (1, 1),  (2, 4)  ) AS SQUARES (X, [X^2]) | Non-standard T-SQL |
|  |  |  |

### Procedures

|  |  |  |
| --- | --- | --- |
| Method | Code | Description |
| **Out Params** | DECLARE @id as int  EXEC [dbo].[AddProductType]  @id = @id OUTPUT,  @description = 'OrdinaryShare'  SELECT @id as N'@id' |  |

### Joins

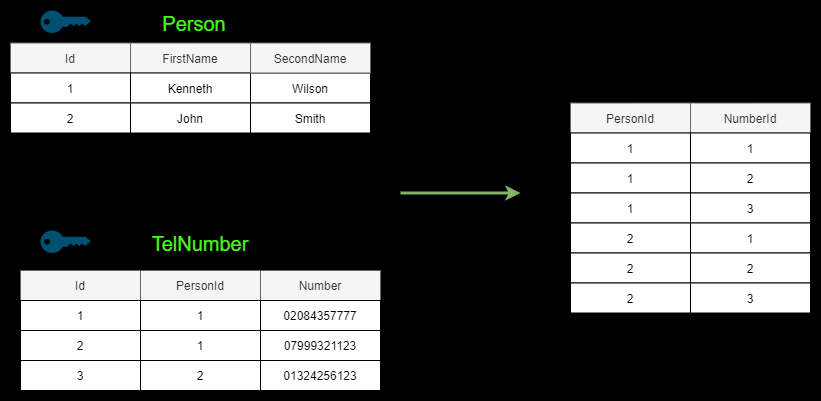
#### Cross Join

A cross join is a simple cartesian product.

SELECT P.id As 'PersonId', T.id AS 'NumberId'

FROM Person P

CROSS JOIN TelNumber T

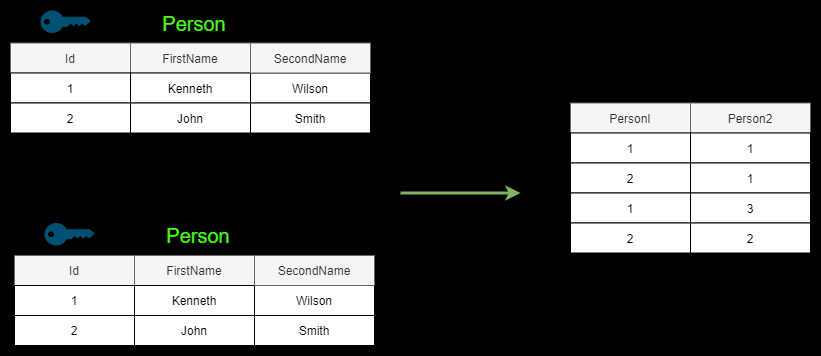


#### Self-Cross Join

SELECT P1.id AS 'Person1', P2.id AS 'Person2'

FROM Person P1

CROSS JOIN Person P2



#### Inner Join

An inner join is implemented in two logical phases; a cartesian product followed by a filtration based on a predicate specified in the ON clause.

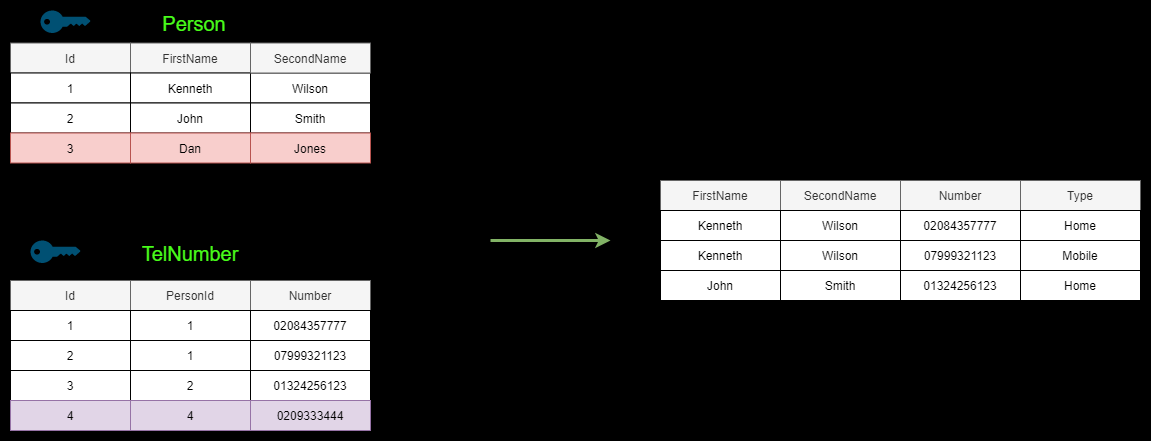
SELECT P.FirstName, P.SecondName, T.Number

FROM Person P

INNER JOIN TelNumber T

ON P.Id = T.PersonId

Note that the third row in the Person table is excluded from the result because it does not have any matching rows in the TelNumber table. Similarly the fourth row in the TelNumber table is excluded from the result as it has no match in the Person table. This the defining feature of an inner join.



#### Composite Join

A composite join is just a join where the AS clause has multiple attributes e.g., something like

SELECT

P.firstName AS 'First Name',

P.secondName AS 'Second Name',

T.telNumber AS 'Num'

FROM Person AS P

INNER JOIN TelNumber AS T

ON P.secondName = T.secondName

ANF P.firstName = T.firstName

#### Non Equi-Join (Join Condition has any operator other than equality)

SELECT

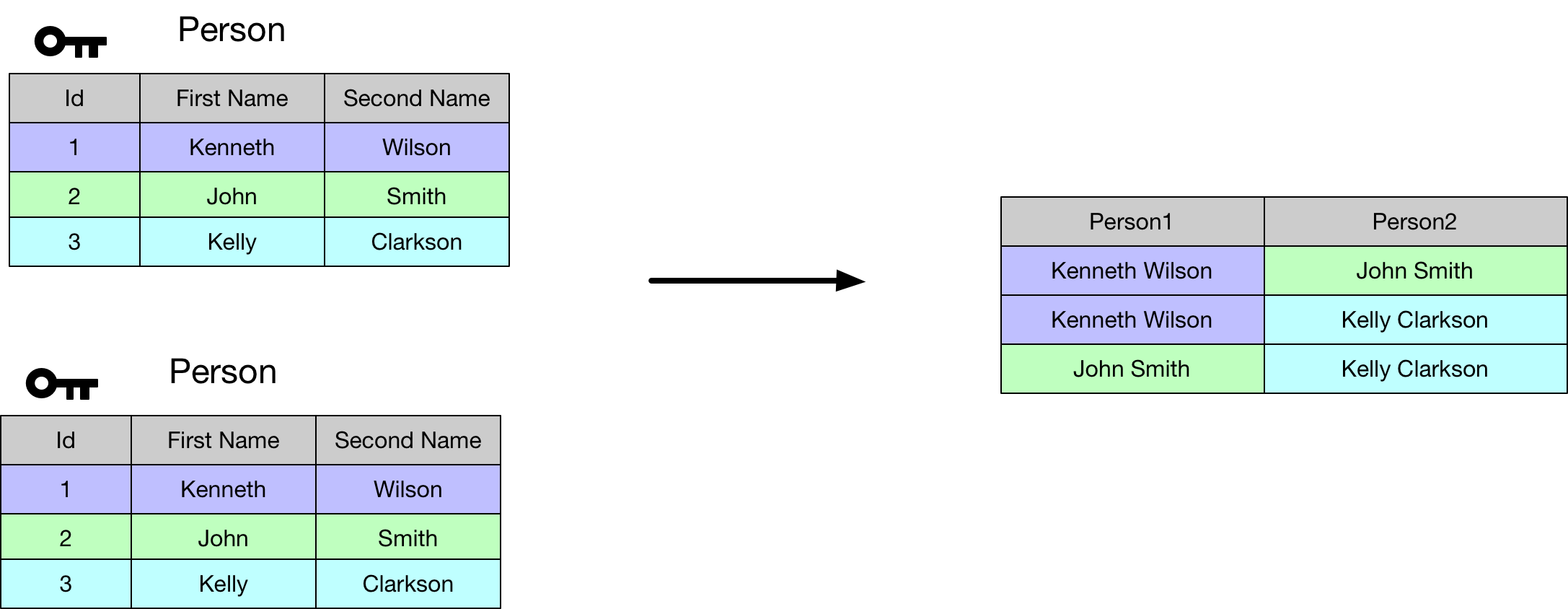
P1.firstName + ' ' + P1.secondName AS 'Person1',

P2.firstName + ' ' + P2.secondName AS 'Person2'

FROM Person AS P1

INNER JOIN Person AS P2

ON P1.id < P2.id



#### Multi-Join

Joining more than one table. Logically multi joins proceed from left to right with the result of the first table operator becoming the left input to the second table operator and so on.

SELECT P.FirstName, P.SecondName, T.Number, NT.[Type]

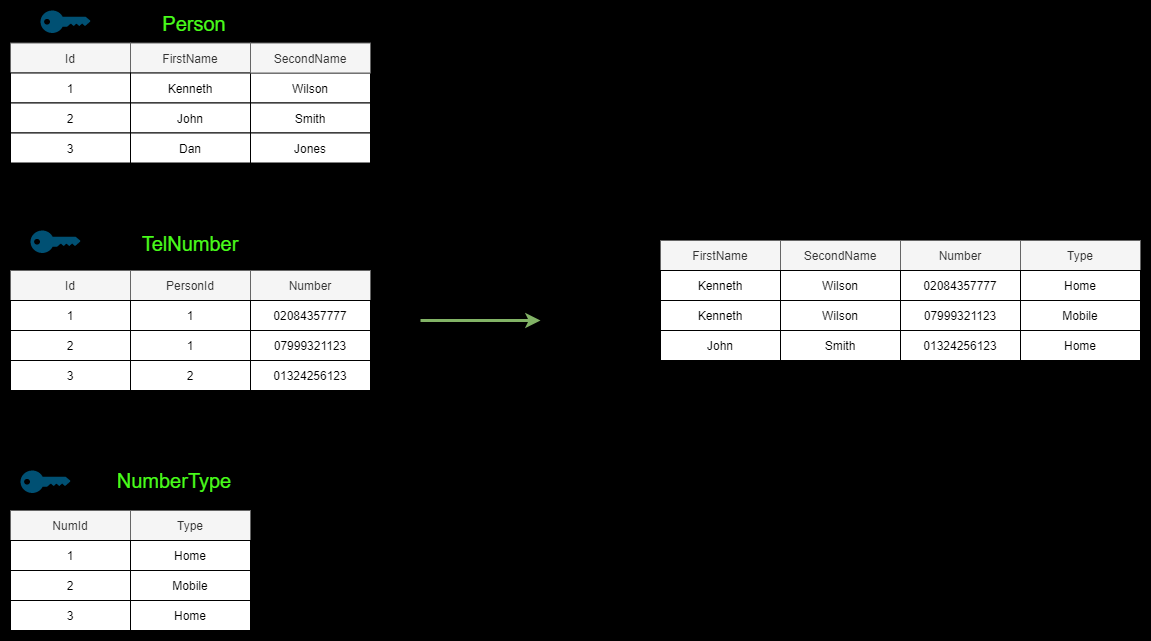
FROM Person P

INNER JOIN TelNumber T

ON P.id = T.personId

INNER JOIN NumberType NT

ON T.id = NT.NumId



#### Outer Join

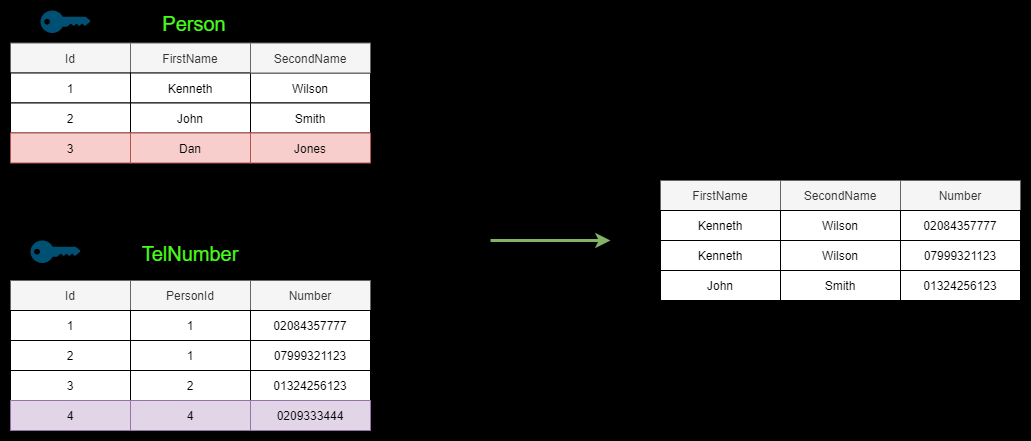
An inner join leaves out any rows that do not match.

SELECT P.FirstName, P.SecondName, T.Number

FROM Person P

INNER JOIN TelNumber T

ON P.Id = T.PersonId



Outer joins enable us to fix this problem.

##### LEFT OUTER JOIN

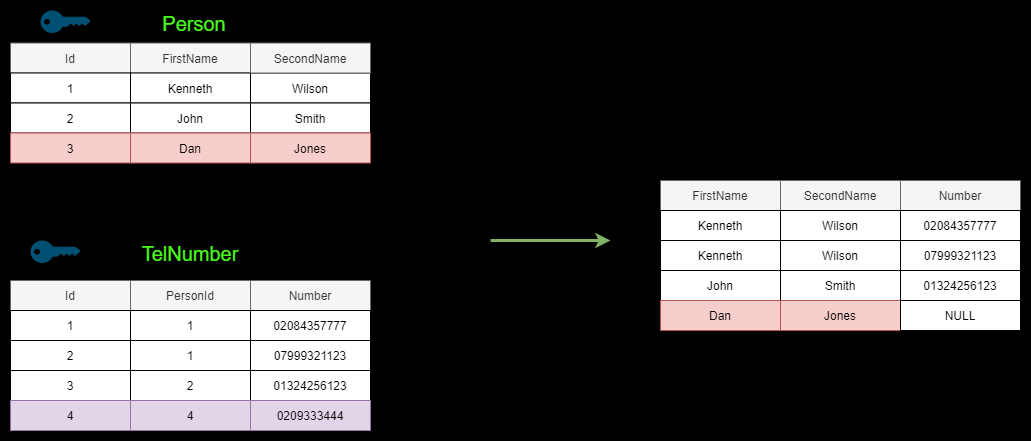
If we want to add back in row 3 in the Person table, we can use a LEFT OUTER JOIN

SELECT P.FirstName, P.SecondName, T.Number

FROM Person P

LEFT OUTER JOIN TelNumber T

ON P.Id = T.PersonId



##### RIGH OUTER JOIN

If we want to add back in row 4 in the TelNumber table, we can use a RIGHT OUTER JOIN.

SELECT

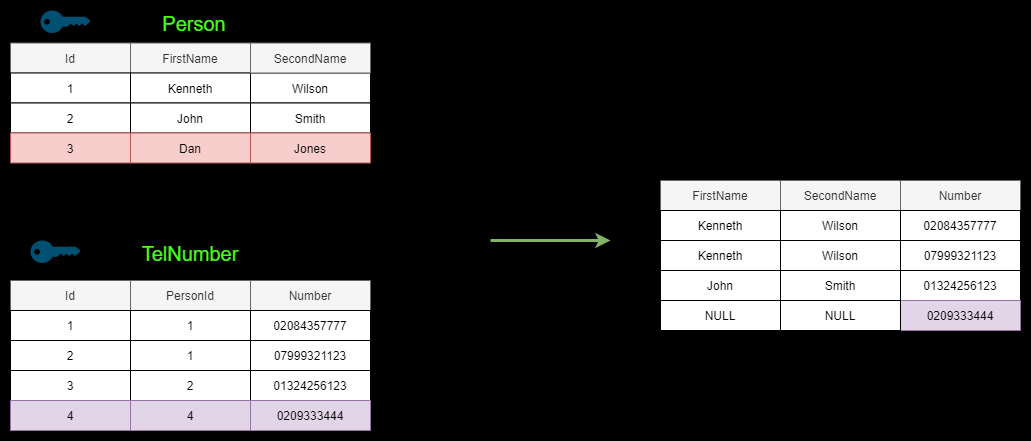
P.firstName + ' ' + P.secondName AS 'Person',

T.telNumber AS 'Number'

FROM Person AS P

LEFT OUTER JOIN TelNumber AS T

On P.id = T.personId



##### FULL OUTER JOIN

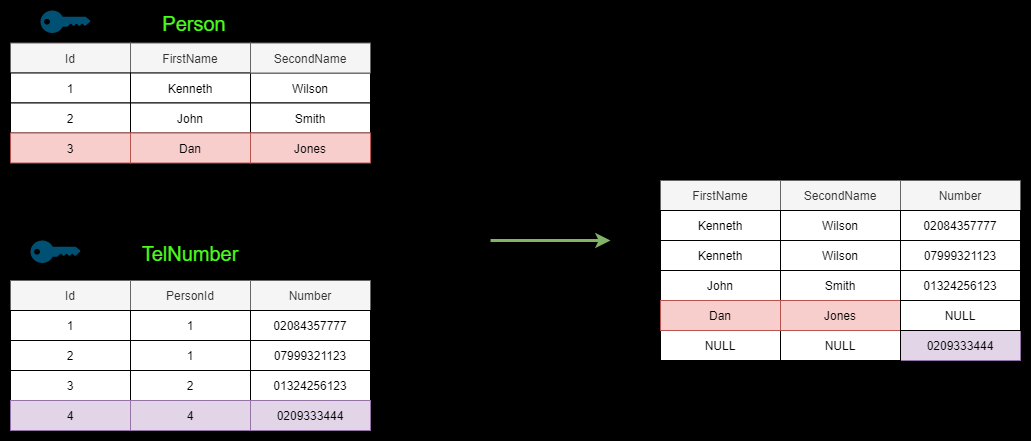
If we want to add back in row 4 in the TelNumber table row and row 3 in the Person table, we can use a FULL OUTER JOIN

SELECT P.FirstName, P.SecondName, T.Number

FROM Person P

FULL OUTER JOIN TelNumber T

ON P.Id = T.PersonId

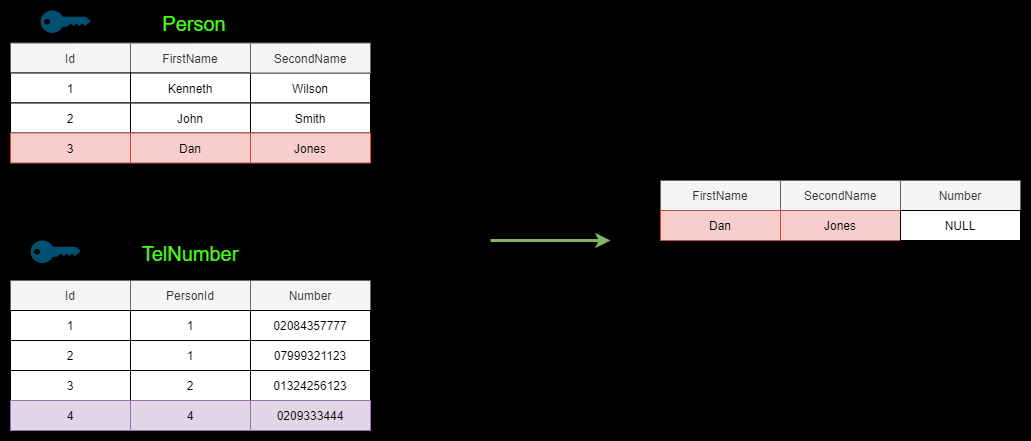


##### OUTER JOIN ON VERSUS WHERE

The ON clause determines which rows match between the two tables. The WHERE clause filters the matched rows.

##### SELECT ONLY OUTER ROWS

The following shows how to only return outer rows.



SELECT P.FirstName, P.SecondName, T.Number

FROM Person P

LEFT OUTER JOIN TelNumber T

ON P.Id = T.PersonId

WHERE T.number IS NULL

##### OUTER JOIN BUGS

In general, we should never refer to columns in the non-preserved side of an outer join from the WHERE clause. This will remove any rows without a match in the non-preserved side effectively changing the join to an INNER JOIN.

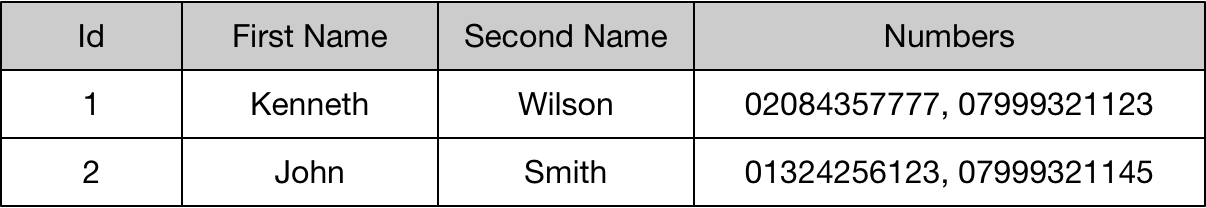
Similarly if we have a multiple table join where we first perform an outer join between two tables and then perform an inner join between a third table and a field on the non-preserved side of the outer join then any outer rows will be discarded.

### Normalization

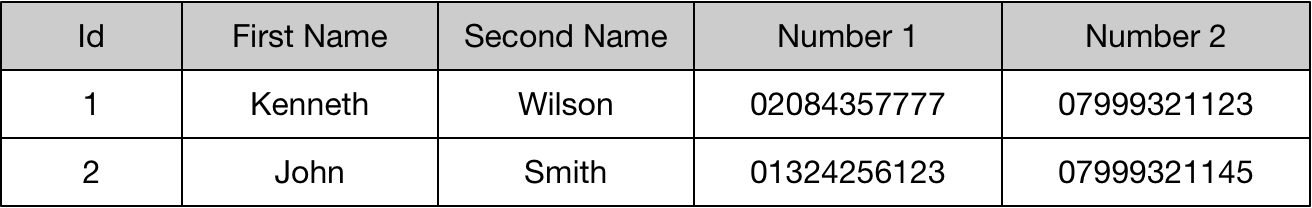
#### Forms

1st Normal Form

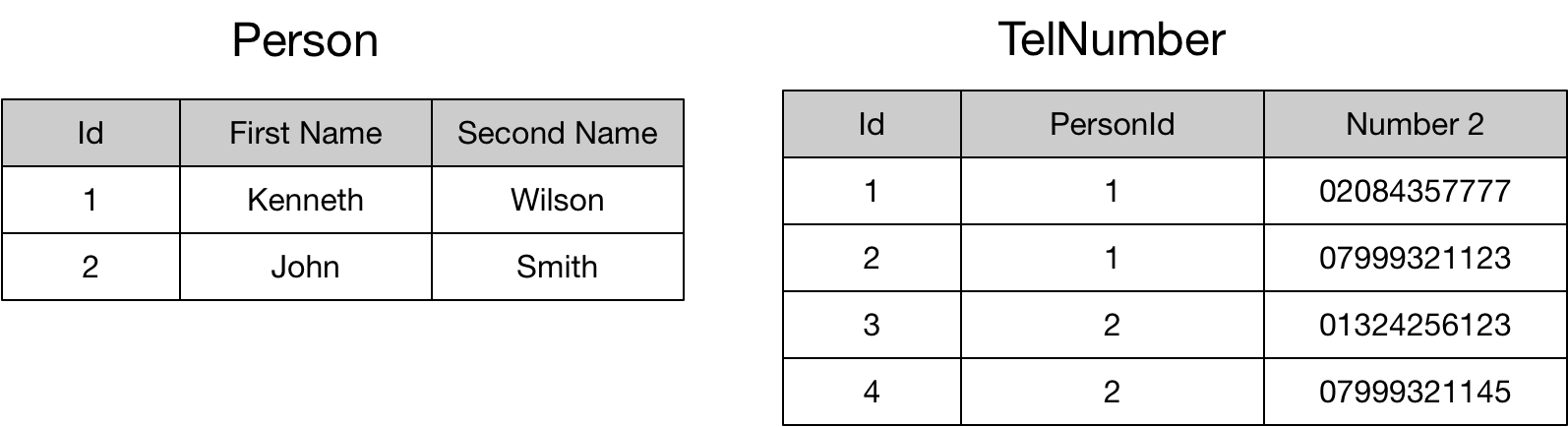
For a table to be in first normal form each cell must contain only one value from the domain. The following table violates this restriction



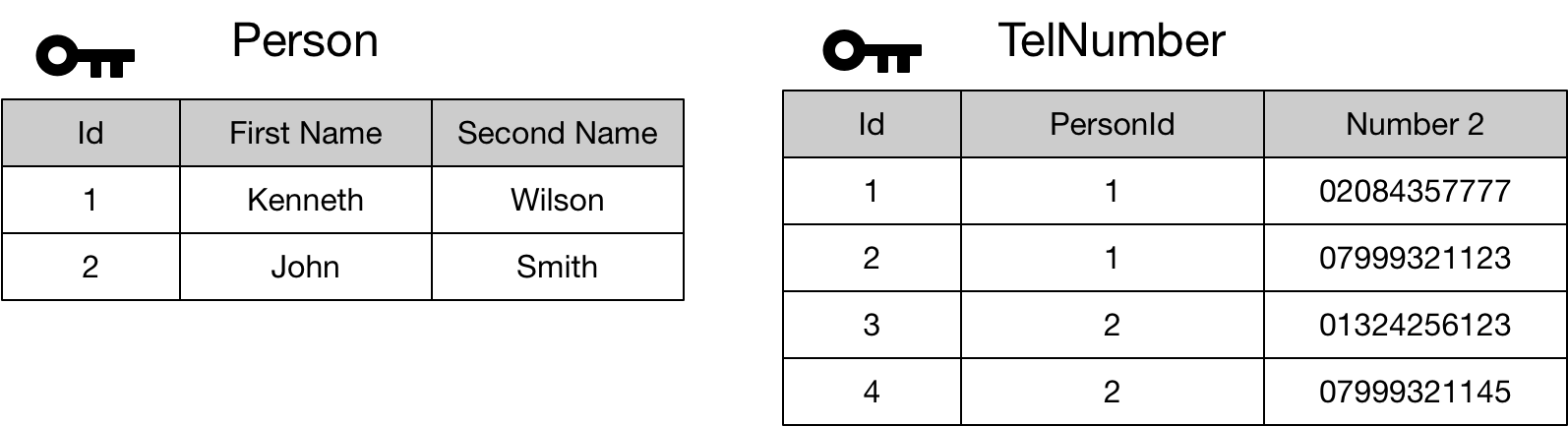
We should note the following table is still in violation of 1st normal form as repeating columns groups are also disallowed.



We can fix this using a schema something like the following. The PersonId field acts as the foreign key that indexes into the person table

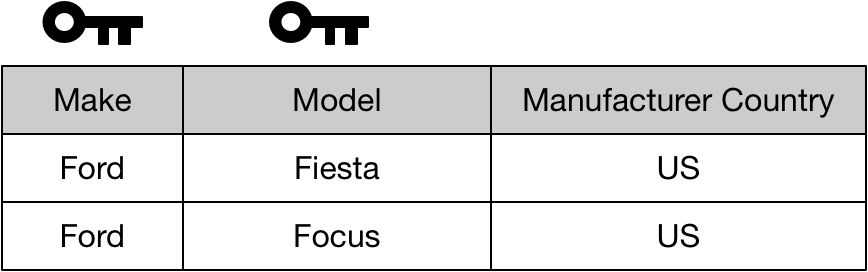


The second restriction on 1st normal form is that every row is unique. We can ensure uniqueness of rows by applying a candidate key to the row.

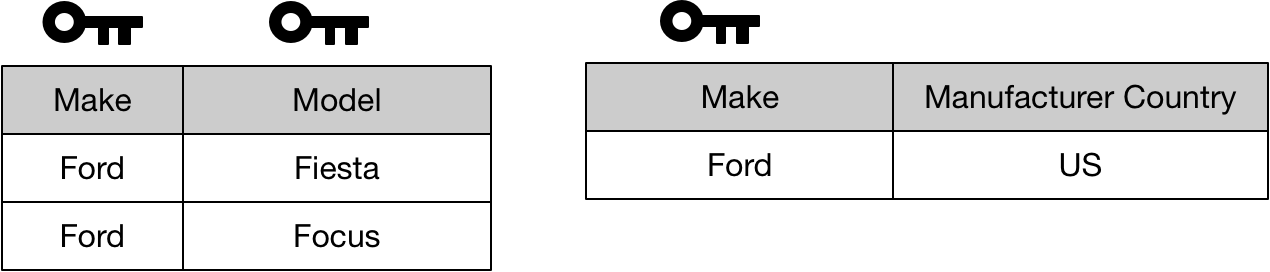


2nd Normal Form

Second normal for applies to relations with composite keys. Where there is a composite key, we should not be able to locate the value of any non-key attribute using only part of the composite key. The following table is in violation of 2nd normal form.



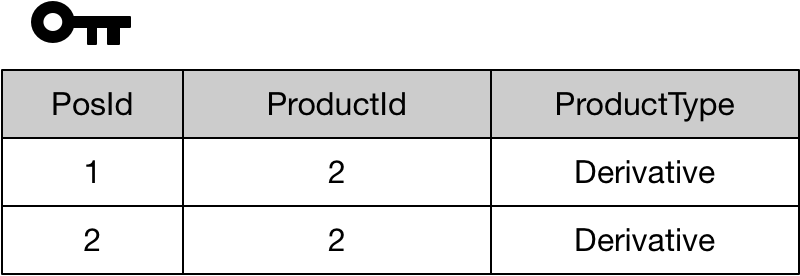
We can fix this as follows



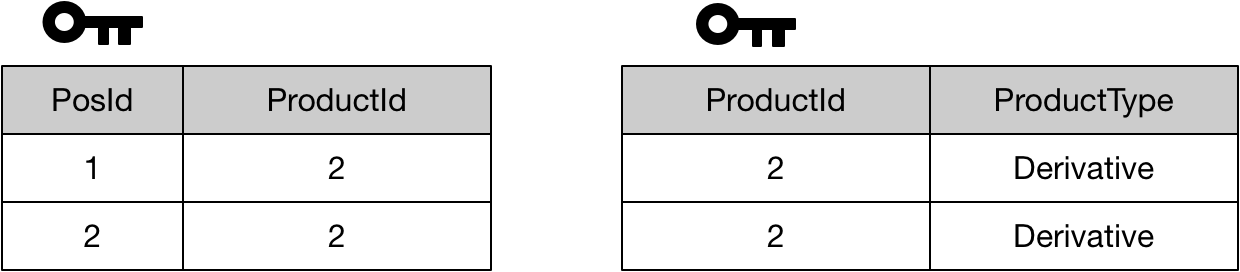
A relation with a single attribute primary key in 1st normal form is automatically in 2nd normal form.

3rd Normal Form

To be in third normal form the relation must first be in second normal form. The second rule is that no non-key attribute can be identified by another non-key attribute. This table is in violation of 3rd normal form



We can fix this as follows.



We can summarise the second and third forms as meaning than in order to identify the value of any non-key attribute we need to use the full primary key. Furthermore, we cannot identify the value of any non-key field using another non-key field.

### Variables

|  |  |  |
| --- | --- | --- |
| Method | Code | Description |
| **Declare** | DECLARE @x AS INT; |  |
| **Set Value** | SET @x = 10 |  |
| **Declare and Initialize** | DECLARE @x AS INT = 10; |  |
| **Set value from scalar sub query** | DECLARE @x AS INT;  SET @x = (SELECT VAL  FROM Position  WHERE PositionId = 1  ); |  |
| **Set multiple variables from select** | DECLARE @x AS INT;  DECLARE @y AS INT;  SELECT  @x = [Value],  @y = ClientId  FROM Position  WHERE PositionId = 1 | Non-standard T-SQL extension |

### Control Flow

#### IF...ELSE

DECLARE @x AS INT = 4;

IF @x = 5

BEGIN

PRINT 'Value is 5';

PRINT 'A Second Statement';

END;

ELSE

IF @x = 4

BEGIN

PRINT 'Value is 4';

RINT 'A Second Statement';

END;

ELSE

BEGIN

PRINT 'Value is something else';

PRINT 'A Second Statement';

END;

#### WHILE

DECLARE @x AS INT = 3;

WHILE @x > 0

BEGIN

PRINT @x

SET @x = @x -1

END

#### CURSOR

DECLARE @posId As INT;

DECLARE MyCursor CURSOR FAST\_FORWARD FOR

SELECT PositionId

FROM Position;

OPEN MyCursor;

FETCH NEXT FROM MyCursor INTO @posId;

WHILE @@FETCH\_STATUS = 0

BEGIN

PRINT @posId

FETCH NEXT FROM MyCursor INTO @posId

END;

CLOSE MyCursor;

yg

DEALLOCATE MyCursor;

### Predicates

Predicates can be combined using the AND and OR operators. A predicate is an expression that is either true or false. Predicates can be used to

1. Enforce data integrity.
2. Filter data into subsets.
3. Specify sets by their properties rather than explicit enumeration of elements

|  |  |  |
| --- | --- | --- |
| Predicate | Example | Description |
| IN | SELECT \* from Products  WHERE prodId IN (1, 2, 3) | Check if an expression is equal to at least one element in the set. |
| BETWEEN | SELECT \* from Products  WHERE prodId BETWEEN 1 AND 3 | Check if a value is in a specified range |
| LIKE | SELECT \* from Products  WHERE description LIKE 'E%' | Check if a character string conforms to a patter |

### Date and Time

|  |  |
| --- | --- |
| Type | Description |
| DATETIME | Legacy |
| SMALLDATETIME | Legacy |
| DATE | 3 bytes representing range from 00010101 to 99991231 |
| TIME | 3 to 5 bytes giving accuracy to 100 nanoseconds |
| DATETIME2 | 6 to 8 bytes giving date and time accuracy to same level as DATE and TIME combined |
| DATETIMEOFFSET | 8 to 10 bytes. Like DATETIME2 but includes offset from UTC |

Legacy Datetime type

When dealing with the legacy type DATETIME the convention is that we use a time of midnight if we want to only use the date part and a date of January 1st 1900 if we are only interested in the time part.

Functions

|  |  |  |
| --- | --- | --- |
| Function | Returns | Description |
| GETDATE | DATETIME | Gets current date and time |
| GETUTCDATE | DATETIME | Get current date and time in UTC |
| SYSDATETIMEOFFSET | DATETIMEOFFSET | Get current date and time with UTC offset |
| SWITCHOFFSET |  | Switches to a different UTC Offset (Timezone) |
| DATEADD |  | Add years, months or days to a |
| DATEDIFF |  | Give different between dates in some date part (year, month, day etc) |
| DATEPART |  | Get year, month, day etc |
| YEAR,MONTH,DAY |  | Abbreviations for DATEPART |

As none of the above functions return only the date or only the time, we need to do a little extra to get these.

SELECT

CAST(SYSDATETIME() AS DATE) AS [Date],

CAST(SYSDATETIME() AS TIME) AS [Time]

|  |  |
| --- | --- |
| **Date** | **Time** |
| 27/04/2020 00:00:00 | 20:29:54.8753964 |

#### Adding Dates

insert Trades values(2,'20200103')

select tradeId AS 'Trade Id', tradeDate AS 'Trade Date' from Trades

|  |  |
| --- | --- |
| **Trade Id** | **Trade Date** |
| 2 | 03/01/2020 00:00:00 |

#### Inefficient of manipulation columns in filters

We need to careful when applying filters such as in WHERE phases. If we manipulate the filtered column this can prevent the database server using the index in an efficient manner. The following is an inefficient query.

SELECT tradeId AS 'Trade Id', tradeDate AS 'Trade Date'

FROM Trades

WHERE YEAR(tradeDate) = 2020

We can improve our query using a range filter as follows which enables the efficient use on indices.

SELECT tradeId AS 'Trade Id', tradeDate AS 'Trade Date'

FROM Trades

WHERE tradeDate >= '20200101' AND tradeDate < '20210101'

### Strings

Characters can be either Unicode or regular.

|  |  |
| --- | --- |
| Type | Description |
| CHAR | Fixed length, 1 byte per character |
| VARCHAR |  |
| NCHAR | Fixed length Unicode, 2 bytes per character |
| NVARCHAR |  |

String Function

|  |  |  |
| --- | --- | --- |
| Operation | Example | Description |
| + | SELECT 'Hello' + 'World' | String concatenation operator |
| COALESCE | DECLARE @str AS VARCHAR(10) = NULL;  SELECT COALESCE(@STR, 'Empty') | Return first non-null element in list |
| SUBSTRING | SELECT  SUBSTRING('Hello', 1,3) | Result is Hel |
| LEFT | SELECT LEFT('Hello', 2) | Shorthand form of SUBSTRING. Result is He |
| RIGHT | SELECT Right('Hello', 2) | Shorthand form of SUBSTRING. Result is lo |
| LEN | SELECT LEN('Hello') | Return number of characters. In this case 5 |
| CHARINDEX | SELECT CHARINDEX('l','Hello World',5) | Return the index of first occurrence of substring starting it specified index. In this case result is 10 |
| PATINDEX | SELECT PATINDEX('%[0-9]%','AA5AA') | Return first index of pattern. In this case index 3. |
| REPLACE | SELECT REPLACE('Boy Meets Boy','Boy', 'Girl') | ‘Girl Meets Girl’ |
| REPLICATE | SELECT REPLICATE('Boy', 3) | ‘BoyBoyBoy; |
| STUFF | SELECT STUFF('Hello', 1,1,'Boy') | Remove substring and insert substring |
| RTRIM/LTRIM |  | Remove trailing/leading spaces. |
| FORMAT | SELECT FORMAT(1234, '0000000'); | Format value based on .NET format string. Result is 0001234 |
| STRING\_SPLIT | SELECT \* FROM STRING\_SPLIT('1,2,3,4',',') AS S | 1  2  3  4 |

### Operators

#### Equality and Ordinality

|  |  |
| --- | --- |
| Operator | Description |
| = | Equality |
| > |  |
| < |  |
| >= |  |
| <= |  |
| <> |  |

#### Logical

|  |  |
| --- | --- |
| Operator | Description |
| AND | and |
| OR |  |
| NOT |  |

#### Arithmetic

|  |  |
| --- | --- |
| Operator | Description |
| + |  |
| - |  |
| \* |  |
| / |  |
| % | Modulo |

If two arguments are of the same type the result is of the same type. So, 7 / 2 = 3. We might want to perform a cast in this instance as follows.

SELECT CAST(7 as Numeric(12,2))/ CAST(3 as Numeric(12,2)) as num

### Examples

#### Basic Select

The following queries work on the below Person table.

|  |  |  |
| --- | --- | --- |
| **id** | **firstName** | **secondName** |
| 1 | Kenneth | Wilson |
| 2 | John | Smith |
| 3 | Kelly | Clarkson |

|  |  |  |  |
| --- | --- | --- | --- |
| Name |  | Code |  |
| Select |  | SELECT \*  FROM Person | |  |  |  |  | | --- | --- | --- | --- | | **id** |  | **firstName** | **secondName** | | 1 |  | Kenneth | Wilson | | 2 |  | John | Smith | | 3 |  | Kelly | Clarkson | |
| Select Named Columns |  | SELECT P.firstName  FROM Person P | |  | | --- | | **firstName** | | Kenneth | | John | | Kelly | |
| Select AS |  | SELECT P.firstName AS "First"  FROM Person P | |  | | --- | | **First** | | Kenneth | | John | | Kelly | |

#### Where (Filtration)

The following queries work on the below Person table.

|  |  |  |
| --- | --- | --- |
| **id** | **firstName** | **secondName** |
| 1 | Kenneth | Wilson |
| 2 | John | Smith |
| 3 | Kelly | Clarkson |
| 3 | John | Wilson |

|  |  |  |  |
| --- | --- | --- | --- |
| Name |  | Code |  |
| Equality |  | SELECT firstName  FROM Person P  WHERE firstName = 'John' | |  | | --- | | **firstName** | | John | | John | |
| Inequality |  | SELECT \*  FROM Person  WHERE firstName <> 'John' | |  |  |  | | --- | --- | --- | | **id** | **firstName** | **secondName** | | 1 | Kenneth | Wilson | | 3 | Kelly | Clarkson | |
| Like |  | SELECT \*  FROM Person  WHERE firstName LIKE '%enn%' | |  |  |  | | --- | --- | --- | | **id** | **firstName** | **secondName** | | 1 | Kenneth | Wilson | |
| IN |  | SELECT \*  FROM Person  WHERE firstName  IN ('Kenneth', 'John') | |  |  |  | | --- | --- | --- | | **id** | **firstName** | **secondName** | | 1 | Kenneth | Wilson | | 2 | John | Smith | | 3 | John | Wilson | |
| NOT |  | SELECT \*  FROM Person  WHERE firstName  NOT IN ('Kenneth', 'John') | |  |  |  | | --- | --- | --- | | **id** | **firstName** | **secondName** | | 3 | Kelly | Clarkson | |
| OR |  | SELECT \*  FROM Person  WHERE firstName = 'John'  OR firstName = 'Kelly' | |  |  |  | | --- | --- | --- | | **id** | **firstName** | **secondName** | | 2 | John | Smith | | 3 | Kelly | Clarkson | | 3 | John | Wilson | |
| AND |  | SELECT \*  FROM Person  WHERE firstName = 'John'  AND secondName = 'Smith' | |  |  |  | | --- | --- | --- | | **id** | **firstName** | **secondName** | | 2 | John | Smith | |
|  |  |  |  |

The following queries work on the below Products table.

|  |  |
| --- | --- |
| **productId** | **description** |
| 1 | European Call |
| 2 | Variance Swap |

|  |  |  |  |
| --- | --- | --- | --- |
| Name |  | Code |  |
| Less than |  | SELECT \*  FROM Product  WHERE productId < 2 | |  |  | | --- | --- | | **productId** | **description** | | 1 | European Call | |
| Less than or equal to |  | SELECT \*  FROM Product  WHERE productId <= 2 | |  |  | | --- | --- | | **productId** | **description** | | 1 | European Call | | 2 | Variance Swap | |
| Greater than |  | SELECT \*  FROM Product  WHERE productId > 1 | |  |  | | --- | --- | | **productId** | **description** | | 2 | Variance Swap | |
|  |  |  |  |

## Detailed Information

### Overview

An RDBMS is supposed to implement the relational model and provide the means to store, manage, enforce the integrity of, and query data.

- T-SQL Fundamentals

The relational model is based on set theory and predicate logic. SQL is a declarative language in that we describe what we want done and leave the details to the RDBMS. A relation in SQL consists of a heading that specifies the set of attributes (columns) and a body that consists of a set of tuples (rows).

### Architecture

A single SQL server can hold multiple user databases in addition to a set of system databases (tempdn, model etc)

### SQL Commands

All SQL commands can be split into three buckets.

|  |  |
| --- | --- |
|  | Column Header |
| Data Definition Language | CREATE, ALTER, DROP |
| Data Manipulation Language | SELECT, INSERT, UPDATE, DELETE, TRUNCATE, MERGE |
| Data Control Language | GRANT and REVOKE |

### Tables, heaps, and indexes

#### Heap

A table which has no clustered index is known as a heap. The rows for a heap are stored with no order. This means if we want to specify the order of the rows, we need to use an ORDER BY clause.

DROP TABLE IF EXISTS dbo.Products;

CREATE TABLE dbo.Products

(

prodId INT NOT NULL,

description VARCHAR(30) NOT NULL

);

We should avoid using a heap in the following situations.

* If data is frequently accessed in sorted order.
* If data is frequently grouped.
* If range queries are often used.

### Data Definition

The following creates a table.

DROP TABLE IF EXISTS dbo.Products;

CREATE TABLE dbo.Products

(

prodId INT NOT NULL,

description VARCHAR(30) NOT NULL

);

#### Constraints

A RDBMS allows one to model data integrity by specifying constraints. A candidate key is specified on an attribute (column) to enforce uniqueness of tuples (rows). One of the candidate keys is chosen as the primary key and forms the preferred way of uniquely identifying rows. A Foreign key is specified on a referencing relation. A foreign key constraint can be used to enforce referential integrity by ensuring only values that exist in the referenced relations are allowed in the referencing relations foreign key.

If we apply key constraints to a table, each element is unique, and it can be considered a set (Otherwise it is a bag or a multiset). Order is unimportant in a set. For this reason, the result of a query has no order unless we explicitly give some sorting criteria.

#### Primary Key Constraint

A primary key constraint enforces uniqueness of rows and disallows nulls in constraint attributes. A table can have at most one primary key. Behind the scenes SQLServer will use a unique index to efficiently enforce uniqueness.

* Each table can have only one primary key.
* The fields making up the primary key cannot be null.
* The server creates an index to efficiently enforce uniqueness and retrieval.

ALTER TABLE dbo.Products

ADD CONSTRAINT PK\_Products

PRIMARY KEY(prodId);

#### Unique Constraints

Although a table can have at most one primary key constraint, it can have multiple unique constraints. Like primary key constraints, unique constraints also enforce unique rows. By using unique constraints SQL Server supports having alternative keys on a table. Unlike primary key constraints, unique constraints allow nulls. Internally SQLServer implements unique constraints using unique indices.

### Three Valued Logic

In SQL predicates can evaluate to TRUE, FALSE or UNKNOWN. If one of the arguments in a logical expression is NULL, then the result is UNKNOWN. If a logical expression is used in a query filter, then any result of UNKNOWN leads to a rejection (accept true). If a logical expression used in a check constraint returns UNKNOWN, the value is accepted (reject false)

### Subqueries

Subqueries can be single value or multi-valued. They can also be self-contained or correlated.

#### Self-Contained

select \* from TelNumber T

where T.PersonId IN

( Select MIN(P.Id)

From Person P

);

#### Correlated

The following uses correlated subqueries to calculate the percentage of portfolio for each position.

Select

P1.PortfolioId,

P1.PosId,

P1.Val,

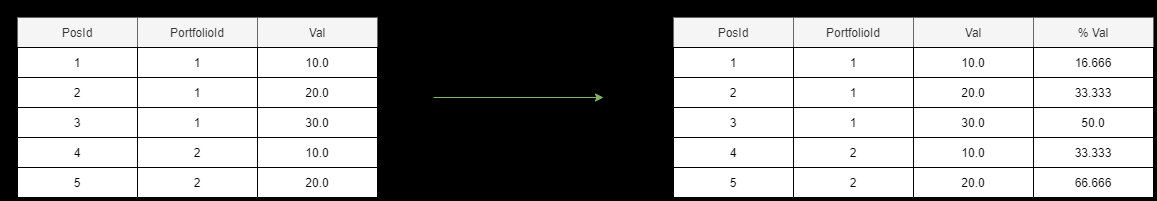
100 \* Val / ( SELECT SUM(Val)

FROM Positions P2

WHERE P1.PortfolioId = P2.PortfolioId

) AS '%Val'

FROM Positions P1



#### Self-Contained

Return everyone who does not have a telephone number.

select P.FirstName, P.SecondName from Person P

where P.Id not in

(

select distinct T.PersonId

from TelNumber T

)

I used distinct here but in general we can expect the database engine to make these kinds of optimisations for us.

#### Exists

We could re-write the previous query as follows.

select P.FirstName, P.SecondName from Person P

where NOT EXISTS

(

select \*

from TelNumber T

where T.PersonId = P.Id

)

Note that exists has particularly good performance as the database engine can optimise.

#### Previous value

We can use correlated subqueries to calculate previous values.

SELECT PosId, PortfolioId, Val,

(

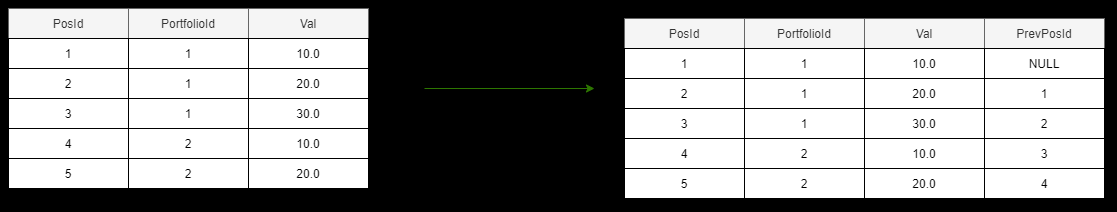
SELECT MAX(P2.PosId)

FROM Positions AS P2

WHERE P2.PosId < P1.PosId

) As PrevPosId

FROM Positions AS P1



#### Next value

We can use correlated subqueries to calculate previous values.

SELECT PosId, PortfolioId, Val,

(

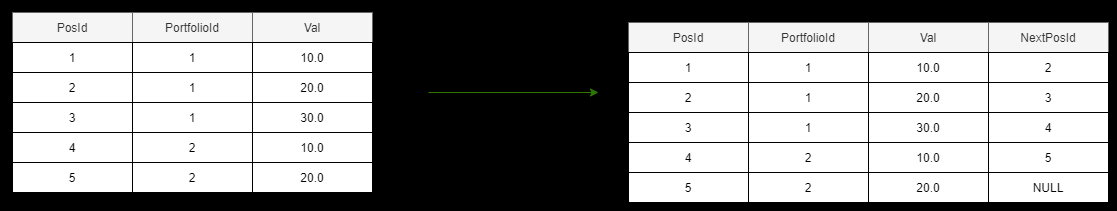
SELECT MIN(P2.PosId)

FROM Positions AS P2

WHERE P2.PosId > P1.PosId

) As NextPosId

FROM Positions AS P1



#### Running Total

SELECT PosId, PortfolioId, Val,

(

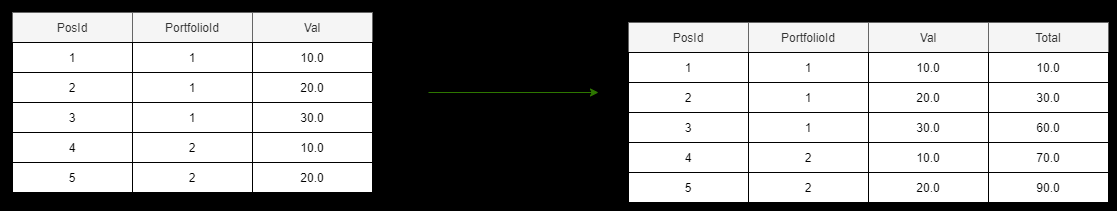
SELECT SUM(P2.VAL)

FROM Positions AS P2

WHERE P2.PosId <= P1.PosId

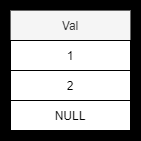
) As RunningTotal

FROM Positions AS P1



#### NOT IN BUgs

We need to be careful when using NOT IN with a subquery that might return NULL. Consider the following table.



We might expect the following query to return “True” but in fact it returns nothing. This is because the result of the NOT IN is the empty set.

SELECT 'True'

WHERE 10 NOT IN

(

SELECT VAL FROM NUMS

WHERE VAL IS NOT NULL

)

Effectively the NOT IN evaluated to

10 NOT IN (1,2,NULL)

Which is

NOT (10=1 OR 10 = 2 OR 10 =NULL)

Which is

NOT (FALSE OR FALSE OR UNKNOWN)

Which is

NOT (UNKNWON)

Which is

UNKNOWN

To fix the bug we need to exclude NULL from the subquery.

### Table Expressions

Table expressions are expressions whose values are relational tables. Any manipulations that expect a table can work with table expressions. There are four types

* Derived Tables
* Common Table Expressions (CTEs)
* Views
* Inline Table Value Functions

It is worth noted that table expressions are virtual and are hence used to improve readability of code. They are unnested by the database engine.

#### Derived Tables

Derived tables or table subqueries are defined in the FROM clause of an outer query.

select \*

from

( SELECT \*

FROM TelNumber

WHERE PersonId = 1

) As PersonOneNumbers

#### Common Table Expressions

WITH Person1Numbers AS

(

SELECT \*

FROM TelNumber T

WHERE T.PersonId = 1

)

SELECT \* FROM Person1Numbers;

Aliases can be specified inline.

WITH Person1Numbers AS

(

SELECT T.Id As NumberId,T.Number As Num

FROM TelNumber T

WHERE T.PersonId = 1

)

SELECT \* FROM Person1Numbers;

Or external

WITH Person1Numbers(NumberId, Num) AS

(

SELECT T.Id,T.Number

FROM TelNumber T

WHERE T.PersonId = 1

)

SELECT \* FROM Person1Numbers;

We can refer to one CTE from another. In such case we use a comma

WITH Person1Numbers(NumberId, Num) AS

(

SELECT T.Id,T.Number

FROM TelNumber T

WHERE T.PersonId = 1

),

Person1FirstNumber(NumberId, Num) AS

(

SELECT P.NumberId, P.Num

FROM Person1Numbers P

WHERE P.NumberId = 1

)

SELECT \* FROM Person1FirstNumber;

Because a CTE is named and then used we can refer to the same CTE from multiple places in the outer query

WITH Person1Numbers(NumberId, Num) AS

(

SELECT T.Id,T.Number

FROM TelNumber T

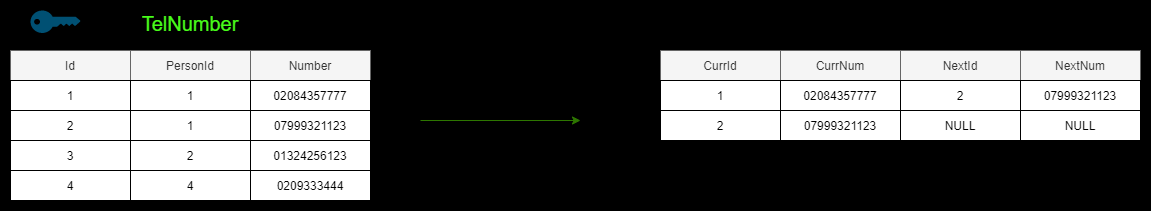
WHERE T.PersonId = 1

)

SELECT P1.NumberId AS CurrId, P1.Num AS CurrNum,P2.NumberId AS NexId, P2.Num AS NextNum FROM Person1Numbers AS P1

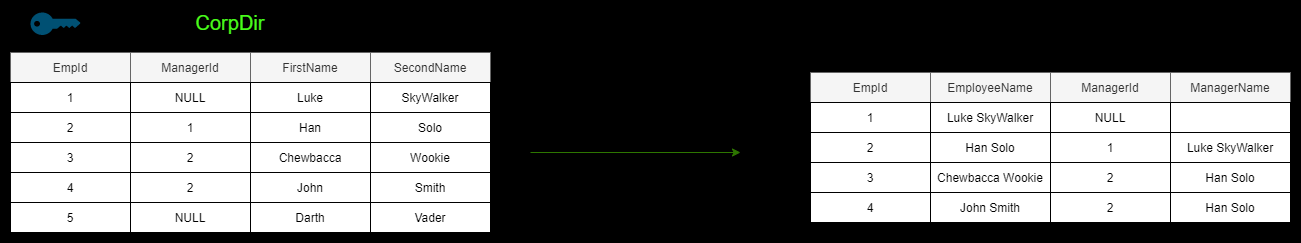
LEFT OUTER JOIN Person1Numbers P2

ON P1.NumberId = P2.NumberId – 1



##### Recursion

CTEs also support recursion. This is an advanced technique and one which is quite hard to understand. Consider an example.



WITH Hierarchy AS

(

--The terminating case

SELECT

EmpId,

ManagerId,

FirstName,

SecondName,

CAST('' AS VARCHAR(30)) ManagerName

FROM

CorpDir

WHERE

EmpId = 1

UNION ALL

SELECT

Subordinates.EmpId,

Subordinates.ManagerId,

Subordinates.FirstName,

Subordinates.SecondName,

CAST( Superiors.FirstName + ' ' + Superiors.SecondName AS

VARCHAR(30)) AS ManagerName

FROM

Hierarchy AS Superiors

INNER JOIN CorpDir AS Subordinates

ON Subordinates.ManagerId = Superiors.EmpId

)

select EmpId, FirstName + ' ' + SecondName AS EmployeeName,ManagerId, ManagerName

FROM Hierarchy

#### Views

Whereas derived tables are scoped to a single expression, views are stored as objects in the database meaning they can be reused. As an object we can manage access to the view in the same was as other database objects.

One point of note is that a View is compiled. So if we use SELECT \* in a view definition it will only include columns that exist in the table at the point it is compiled. Any subsequent table alterations will not be included. The sp\_refreshview procedure can be used to refresh a view in such cases.

##### Schema Binding

We can bind a view’s schema to the schemas of referenced objects. This mean those object schemas cannot be altered or deleted.

##### Check Option

It is possible to make edits through a view. These edits might not conform to the filter in the view. If we want to prevent this we need to use the WITH CHECK OPTION when we create the view.

#### Inline Table Valued Functions

Inline TVF are like views that accept parameters.

CREATE Function dbo.GetDirectReports

(@mgrId AS INT) RETURNS TABLE

AS

RETURN

SELECT \*

FROM CorpDir C

WHERE C.ManagerId = @mgrId

GO

We call the Inline TVF as follows.

SELECT \* FROM GetDirectReports(1)

### Apply

Apply is a non-standard operator.

#### CROSS APPLY

The CROSS APPLY operator works on two input tables. Typically, the RHS is a derived table or TVF. A cross apply statement looks very much like a CROSS JOIN, but it is actually different. With CROSS APPLY the left side is evaluated first. Then the rhs is evaluated per row from the left. In this way the rhs can refer to elements of the lhs.

### Transactions

A Transaction is implicitly started using BEGIN TRAN(SACTION)

A Transaction is committed using COMMIT TRAN(SACTION)

A Transaction is rolled back using ROLLBACK TRAN(SACTION

Concurrency is handled differently in an In Memory OLTP Database

Disk based box SQL Server uses locks as default concurrency control

Locks are either exclusive or shared.

A transaction that modifies data obtains exclusive locks on any resources it updates.

No other transaction can obtain an exclusive lock on a resource while another transaction has an exclusive lock on the same resource.

In SQL Server box product the default ISOLATION level is READ COMMITED.

The disk-based box SQL-Server instance uses locking as the default concurrency control. Locking support two modes: exclusive and shared. When modifying data inside a transaction the transaction obtains an exclusive lock on any resources and holds this lock until the transaction completes or is rolled back. While one transaction holds an exclusive lock on a resource no other transaction can obtain an exclusive lock on the same resource until the first transaction completes. Whether another transaction read from the same object while the first transaction holds an exclusive lock depends on isolation level.

In A SQ

## Set Operations

Set operators work on two query result sets. The queries cannot have ORDER BY although this can be added to the set operator result. The LFS and RHS queries must have the same number of columns and data types.

UNION (DISTINCT)

UNION ALL

INTERSECT

INTERSECT (DISTINCT)

EXCEPT

EXCEPT (DISTINCT)

## Window Functions

To highlight how window functions work, we will work on the following simple table.

|  |  |  |
| --- | --- | --- |
| **PositionId** | **ClientId** | **VALUE** |
| 1 | 1 | 10.0000 |
| 2 | 1 | 10.0000 |
| 3 | 1 | 10.0000 |
| 4 | 2 | 10.0000 |
| 5 | 2 | 10.0000 |

The following query shows a quite simple Window Function against this table.

SELECT P.ClientId, P.PositionId, P.VALUE,

SUM(VALUE) OVER( PARTITION BY P.ClientId

ORDER BY P.PositionId

ROWS BETWEEN UNBOUNDED PRECEDING

AND CURRENT ROW) As Total

FROM Position P

The window is specified by the OVER clause which consists of three parts. First the PARTITION BY clause restricts the window to a subset of positions whose values in the partitioning columns match the current row. In our case we restrict the window to the rows that have the same ClientId as the current row (PARTITION BY P.ClientId). The ORDER BY clause gives meaning to rank (ORDER BY P.PositionId). Finally the window frame clause filters a subset of rows from the preceding expression.

|  |  |  |  |
| --- | --- | --- | --- |
| **ClientId** | **PositionId** | **VALUE** | **Total** |
| 1 | 1 | 10.0000 | 10.0000 |
| 1 | 2 | 10.0000 | 20.0000 |
| 1 | 3 | 10.0000 | 30.0000 |
| 2 | 4 | 10.0000 | 10.0000 |
| 2 | 5 | 10.0000 | 20.0000 |

## Appendices

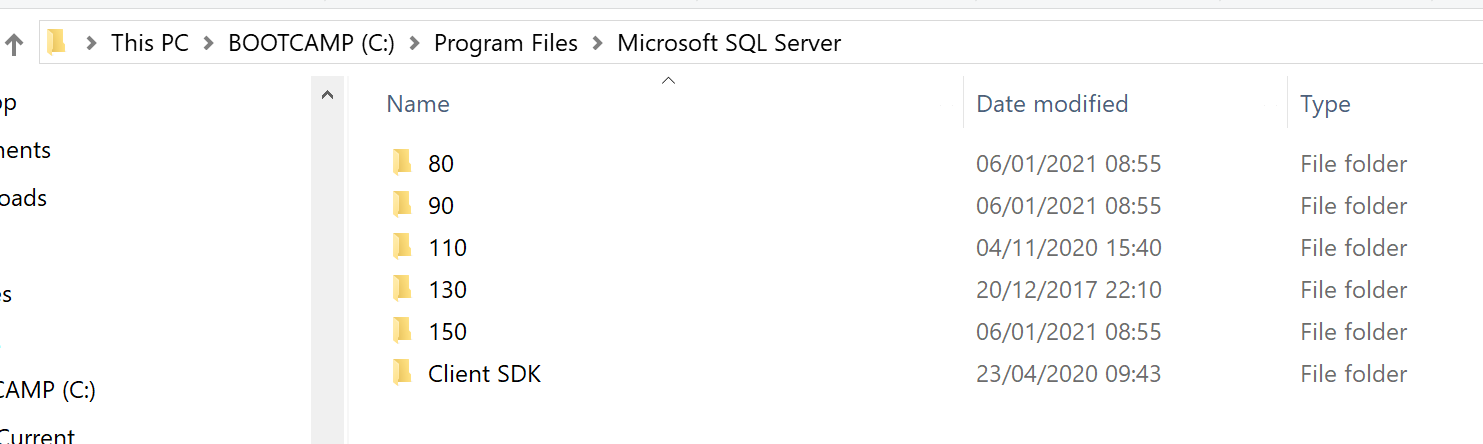
### SQLServer LocalDB Install

#### Install

Use the following link and choose the third one from the top

<https://www.hanselman.com/blog/download-sql-server-express>

The default locations of the installs are C:\Program Files\Microsoft SQL Server On my machine I see multiple versions.

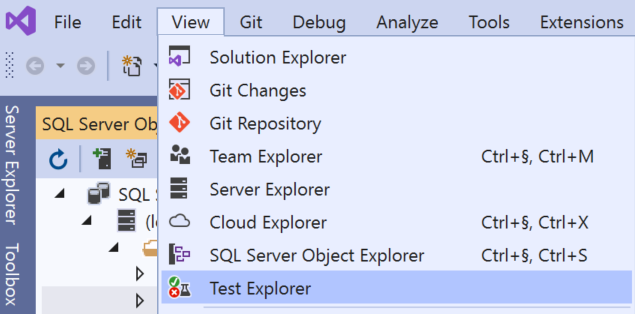


#### Start/Manage

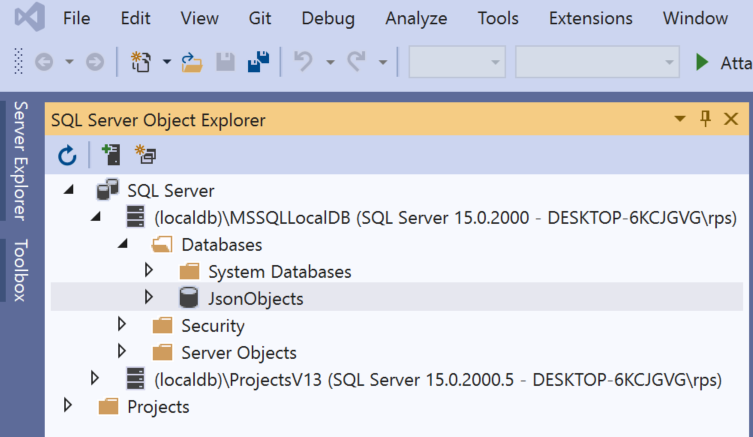
LocalDb supports two kinds on instance: Automatic instances and named instances. The automatic instance for SQLServer is called MSSQLLocalDb.

##### Visual Studio

The easiest way to manage the local DB instances is to use visual studio.



And when open it looks like this.



##### Command Prompt

Open a command prompt. We can list all instances of SQLServer using the command.

C:\Users\rps>"C:\Program Files\Microsoft SQL Server\150\Tools\Binn\SqlLocalDB.exe" info

>> MSSQLLocalDB

We create a named instance as.

C:\Users\rps>"C:\Program Files\Microsoft SQL Server\150\Tools\Binn\SqlLocalDB.exe" create KennysLocalDb

We can view a named install.

"C:\Program Files\Microsoft SQL Server\150\Tools\Binn\SqlLocalDB.exe" info KennysLocalDb

>> Name: KennysLocalDb

>> Version: 15.0.2000.5

>> Shared name:

>> Owner: DESKTOP-6KCJGVG\rps

>> Auto-create: No

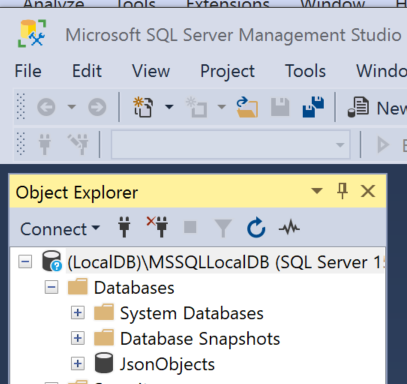
>> State: Stopped

>> Last start time: 19/01/2021 17:47:37

>> Instance pipe name:

#### Microsoft SQL Server Management Studio

If we open management studio it should just automatically connect to the local DB.



Let us add a JSON table.

create table MYDOCS (

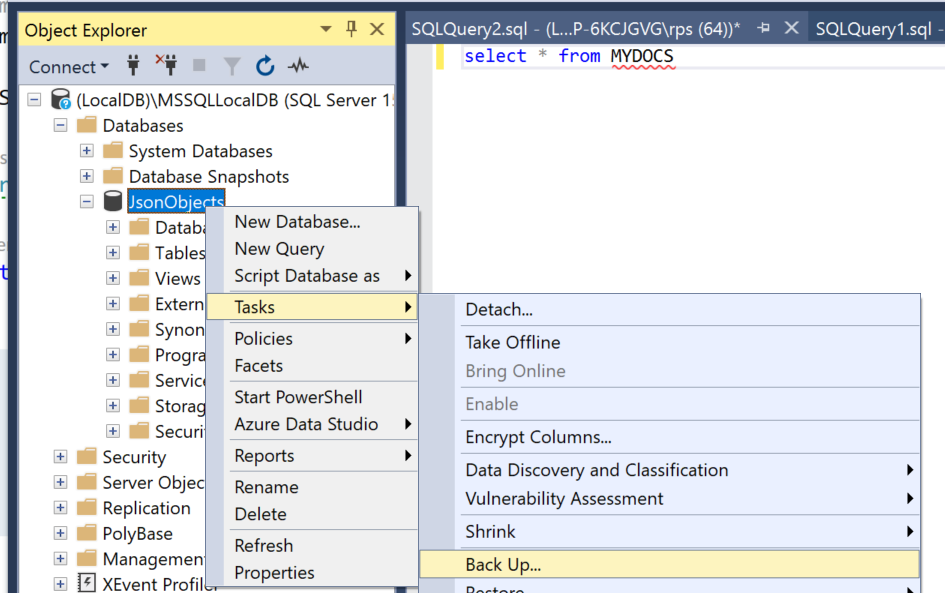
ID bigint primary key identity,

DOC nvarchar(max)

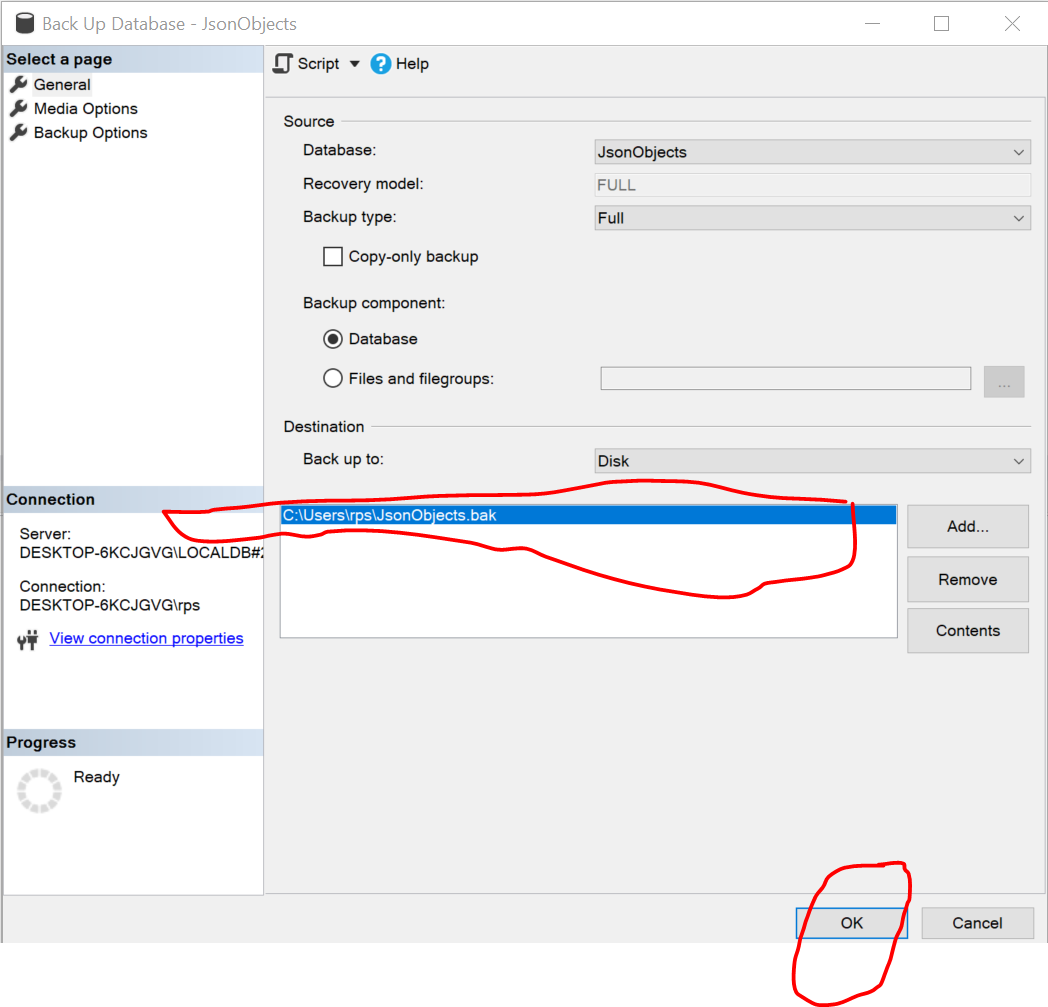
);

### SQLServer Backup/Restore

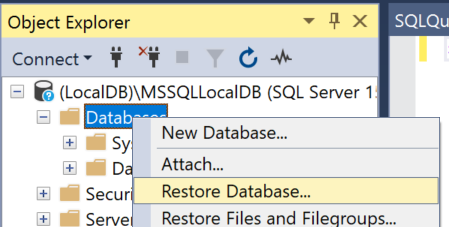
The easiest way to backup and restore our database is to use Microsoft SQL Server Management Studio. Right click the database in the Object Explorer and select as follows.



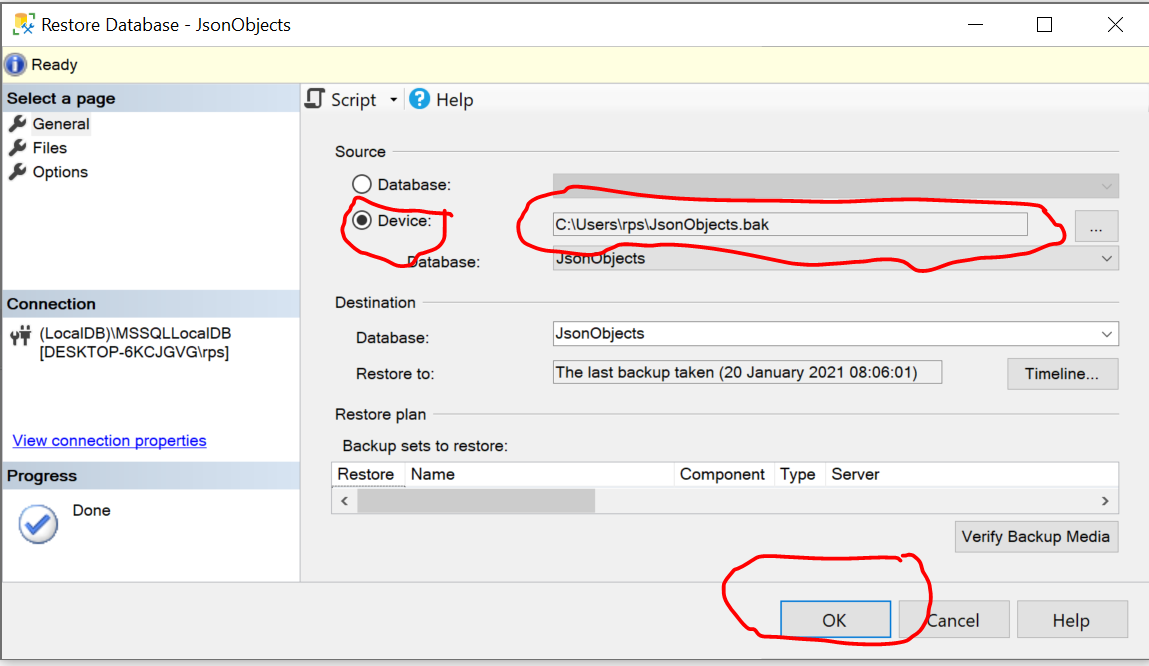
You will then see the following. Just click ok which backs it up to the current user’s home directory.



We can now safely delete our database safe in the knowledge we can back it up. One deleted restore it as follows. In the Object Explorer right-click on databases and select Restore Database.



No in the dialogue select Device and enter the location of the file

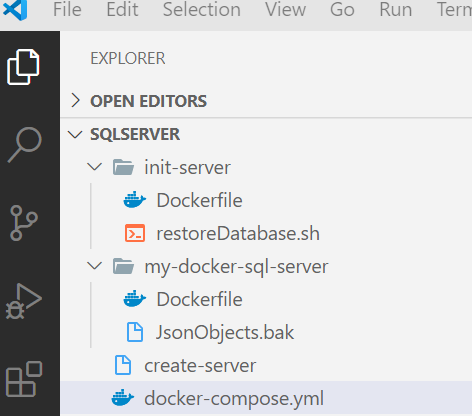


### Running SQLServer In Docker

We can run SQLServer 2019 inside docker. The images are [here](https://hub.docker.com/_/microsoft-mssql-server). For details on how to run the SQL Server Instance see [See](https://bitbucket.org/kennyrnwilson/dockerbasics/src/master/examples/sqlserver/startserver/README.md)

#### Create and Populate Database.

The following docker logic shows how to get an instance running and copy in a backup file to restore a database and its objects. Note that for the following example we already need backup file taking from somewhere that we can use to do the restore.



Listing 1 my-docker-sql-server/Dockerfile

FROM mcr.microsoft.com/mssql/server:2019-latest

USER root

COPY JsonObjects.bak /var/opt/mssql/data/JsonObjects.bak

CMD /opt/mssql/bin/sqlservr

Listing 2 init-server/Dockerfile

FROM mcr.microsoft.com/mssql-tools

USER root

COPY restoreDatabase.sh /restoreDatabase.sh

CMD "./restoreDatabase.sh"

Listing 3 init-server/restoreDatabase.sh

echo "sleeping to allow db to start up"

sleep 10

echo 'resoting database JsonObject'

/opt/mssql-tools/bin/sqlcmd -S my-docker-sql-server -U 'SA' -P $SA\_PASSWORD -Q "RESTORE DATABASE [JsonObjects] FROM DISK = N'/var/opt/mssql/data/JsonObjects.bak' WITH REPLACE"

Note: file ending

For the above shell script make sure the line ending in VS code is set to LF or it will not work.

Listing 4 docker-compose.yml

version: '3'

services:

  init-server:

    build:

      context: init-server

      dockerfile: Dockerfile

    environment:

      SA\_PASSWORD: "Pa!ssWordTwo!"

    links:

      - my-docker-sql-server

  my-docker-sql-server:

    build:

      context: my-docker-sql-server

      dockerfile: Dockerfile

    environment:

      SA\_PASSWORD: "Pa!ssWordTwo!"

      ACCEPT\_EULA: "Y"

    ports:

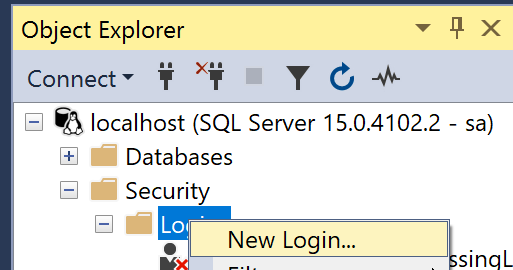
      - 1433:1433

### Logins, Users and Roles

A login is a server level entity and users are database level entities. We can have a login with no user associated. In this case we can log onto the server but not use any of the databases on the server.

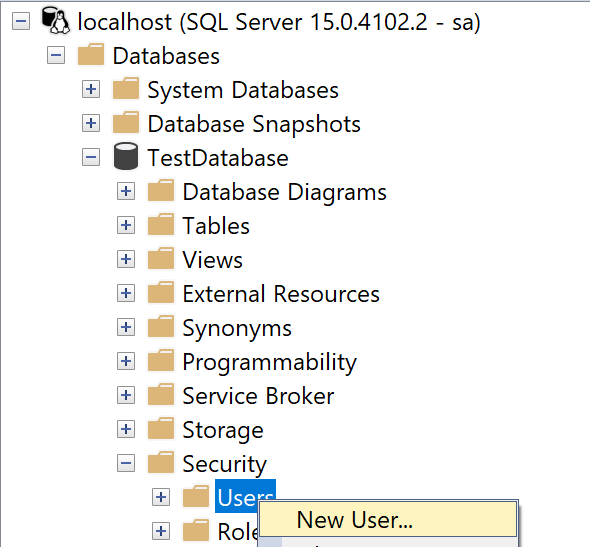
#### Create Login

On the folder <ServerName>/Security/Logins right click and select New Login.



#### Create User

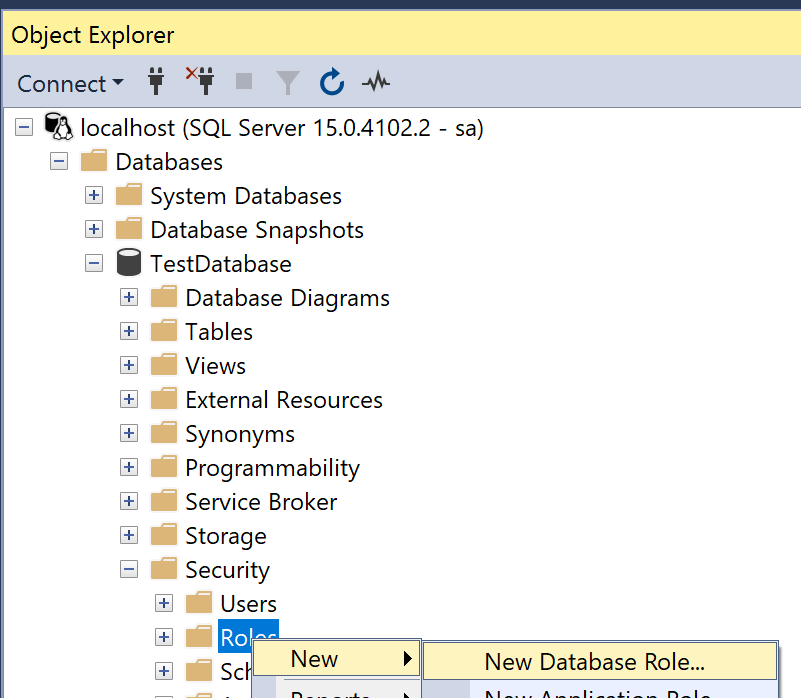
On the folder Databases/<DatabaseName>/Security/Users right click and select New User.



On the setup dialogue enter the name and the login we want to associate the user with. You should now be able to log on to the server with the login and then access the database.

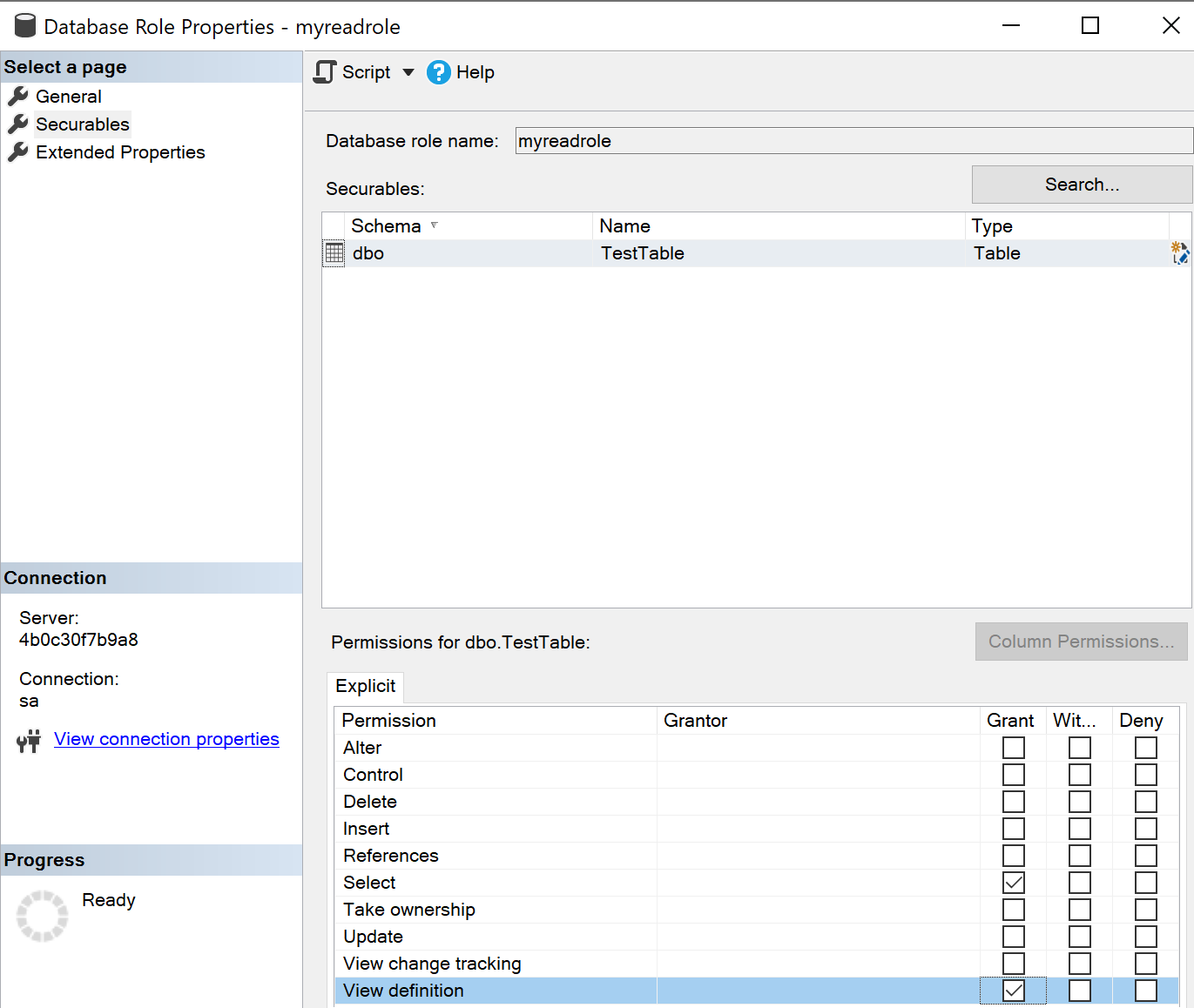
#### Create Role

To be able to do anything we need to create a role that can read our table. On the folder Databases/<Database>/Security/Roles select New Database Role.



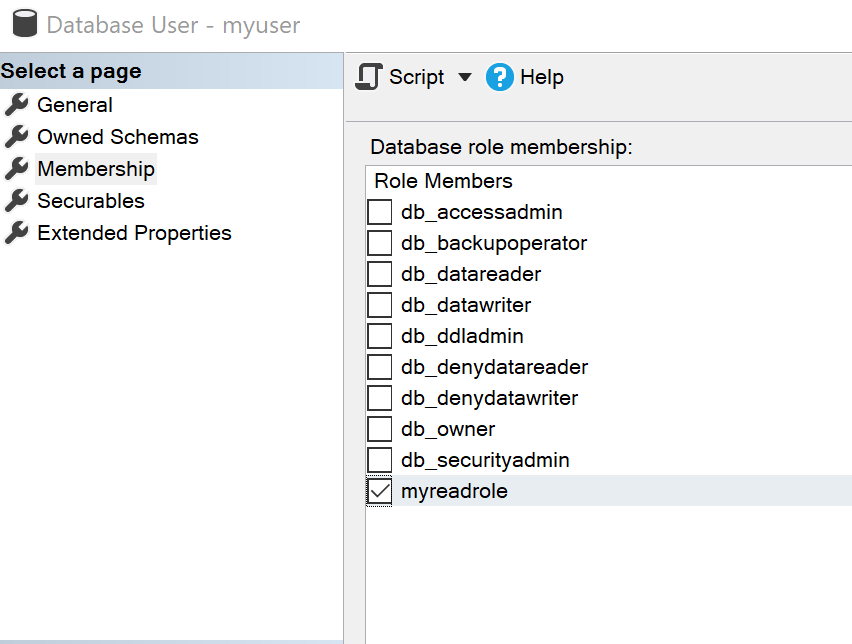
##### Add permissions to role.

Right click on role and select properties and then on the dialogue select Securables.



#### Associate role with user

Right click on user and select properties then on the Membership page add the new role we created.



We can now read the table.

### SQL Server and .NET Core

We can connect to our automatic Local DB from code as follows.

SqlConnection connection = new SqlConnection(@"Data Source=(localdb)\MSSQLLocalDB;Database=JsonObjects;");

connection.Open();

We can insert into our table as follows

var json = @"{""Hello"":""World2""}";

var sql = $"insert into MYDOCS (DOC) Values('{json}')";

var command = new SqlCommand(sql, connection);

command.ExecuteNonQuery();

Questions

Select

Phases

List the phases of a select query.

SELECT

FROM

WHERE

GROUP BY

HAVING

ORDER BY

What is the logical order of the phases and what does each do?

1. FROM select the table we want to query
2. WHERE filter the rows returned.
3. GROUP BY produce group for each combination.
4. HAVING filter the groups.
5. SELECT Specify the columns for result
6. ORDER BY sort results for presentation

What is the key benefit of the where clause?

Enable indices to improve performance and reduce network traffic.

Given the following table write a query that first filters out everyone over 40. Then group by age and countryId and show the count for each group. Only include groups with countryId of 1 or 10

|  |  |  |  |
| --- | --- | --- | --- |
| **firstname** | **secondname** | **age**ΞΞ | **countryId** |
| John | Smith | 25 | 1 |
| Dave | Jones | 25 | 1 |
| Aaro | Litmanen | 30 | 10 |
| Kimi | Raikonen | 30 | 10 |
| Han | Solo | 40 | 2 |
| Luke | Skywalker | 35 | 2 |
| Indiana | Jones | 50 | 3 |
| Carl | Jones | 50 | 3 |

SELECT age, countryId, COUNT(\*)

FROM Person

WHERE age < 50

GROUP BY age, countryId

HAVING countryId IN (1,10)

ORDER BY age, countryId

Basic Select

List the firstname of all rows but rename it “First Name”

SELECT firstname AS 'First Name'

FROM Person

Appendices

Setup Select Database

create table dbo.TelNumber

(

id INT NOT NULL,

PersonId int NULL,

Number int NULL

)

INSERT INTO dbo.TelNumber VALUES

(1, 1, 0208435777),

(2, 1, 07999321123),

(3, 2, 01324256123),

(4, 2, 07999321145),

(5, NULL, 02074257777)