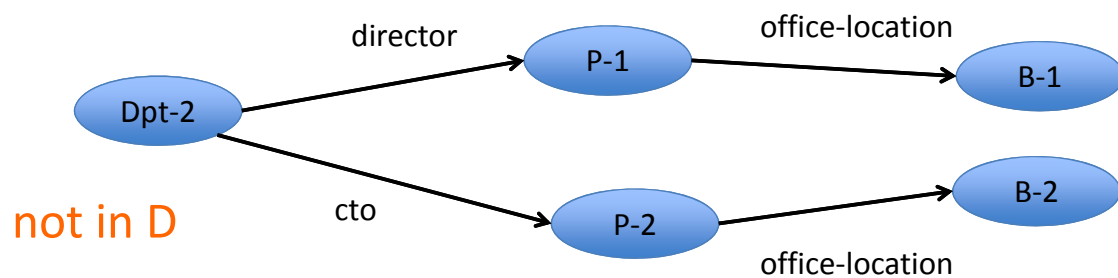
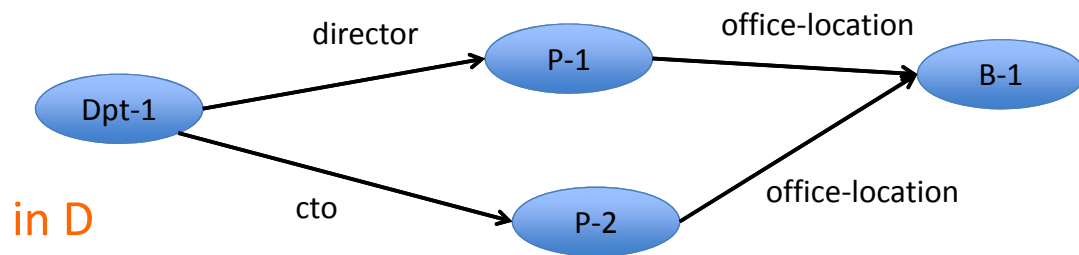


Inference rules in DL and SWRL

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C. Métral

Expressivity of DL

- DLs have the tree model property
 - +/- each set of axioms has a model that is a tree
 - Impossible to specify cyclic models
- Example
 - A department is in the class D if and only if its director and chief technology officer (cto) are located in the same building



In DL (OWL 2)

Impossible to define D

Many other examples cannot be defined in OWL-2

Comes from the fact that most DLs enjoy the **Tree Model Property**.

if a Tbox has a model

then it has a model that doesn't contain cycles

A fact is a consequence of a Tbox if it is true in **every** model of the Tbox, so no "cyclic fact" is a consequence of a TBox.

Inference rules

Rules to produce

- New **type** assertions
 - x is a member of class C
- New **property** assertions
 - x is connected to y through property p

SWRL Rules - syntax

rule ::= antecedant -> consequent

antecedant ::= atom, atom, ...

consequent ::= atom, atom, ...

atom ::= description '(' i-object ')'

| dataRange '(' d-object ')'

| individualvaluedPropertyID '(' i-object i-object ')'

| datavaluedPropertyID '(' i-object d-object ')'

| sameAs '(' i-object i-object ')'

| differentFrom '(' i-object i-object ')'

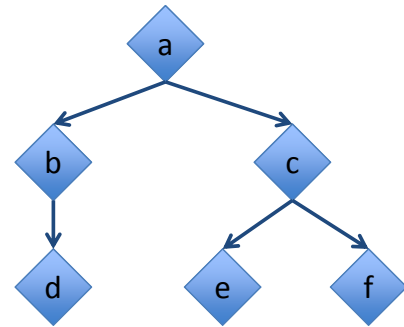
| builtIn '(' builtinID { d-object } ')'

Person(?x), Person(?y), Person(?z), hasChild(?x, ?y), hasChild(?y, ?z) ->
hasGrandChild(?x, ?z)

Interpretation

- Find all the variable bindings that satisfy the antecedent
- For each such binding the consequent must be satisfied

`hasChild(?x, ?y), hasChild(?y, ?z)`
`-> hasGrandChild(?x, ?z)`

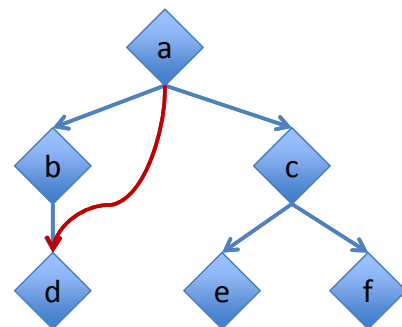


→ hasChild

Interpretation

`hasChild(?x, ?y), hasChild(?y, ?z)`
`-> hasGrandChild(?x, ?z)`

`?x=a, ?y=b, ?z=d -> hasGrandChild(a,d)`



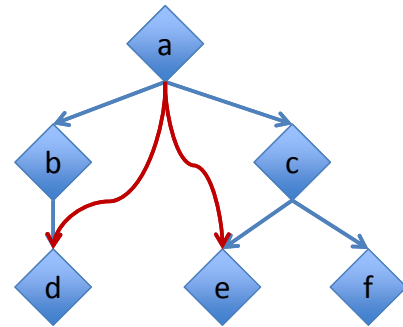
→ hasGrandChild

Interpretation

`hasChild(?x, ?y), hasChild(?y, ?z)`
→ `hasGrandChild(?x, ?z)`

`?x=a, ?y=b, ?z=d` → `hasGrandChild(a,d)`

`?x=a, ?y=c, ?z=e` → `hasGrandChild(a,e)`



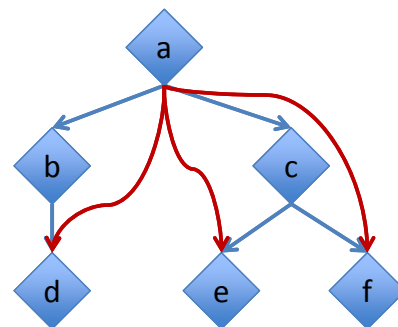
Interpretation

`hasChild(?x, ?y), hasChild(?y, ?z)`
→ `hasGrandChild(?x, ?z)`

`?x=a, ?y=b, ?z=d` → `hasGrandChild(a,d)`

`?x=a, ?y=c, ?z=e` → `hasGrandChild(a,e)`

`?x=a, ?y=c, ?z=f` → `hasGrandChild(a,f)`



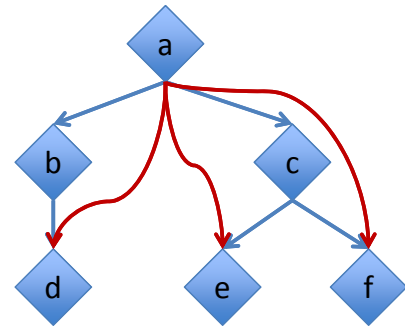
Interpretation

$\text{hasChild}(\text{?x}, \text{?y}), \text{hasChild}(\text{?y}, \text{?z})$
 $\rightarrow \text{hasGrandChild}(\text{?x}, \text{?z})$

$\text{?x}=\text{a}, \text{?y}=\text{b}, \text{?z}=\text{d} \rightarrow \text{hasGrandChild}(\text{a}, \text{d})$

$\text{?x}=\text{a}, \text{?y}=\text{c}, \text{?z}=\text{e} \rightarrow \text{hasGrandChild}(\text{a}, \text{e})$

$\text{?x}=\text{a}, \text{?y}=\text{c}, \text{?z}=\text{f} \rightarrow \text{hasGrandChild}(\text{a}, \text{f})$

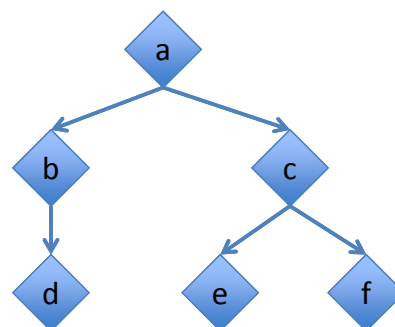


an interpretation that
satisfies the rule

DifferentFrom

Variables with different names may represent the same individual !

$\text{hasChild}(\text{?x}, \text{?y}), \text{hasChild}(\text{?x}, \text{?z})$
 $\rightarrow \text{hasSibling}(\text{?y}, \text{?z})$

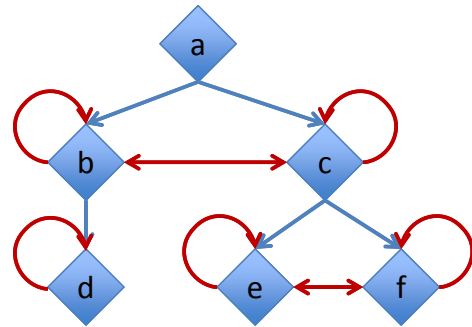


DifferentFrom

Variables with different names may represent the same individual !

hasChild(?x, ?y), hasChild(?x, ?z)

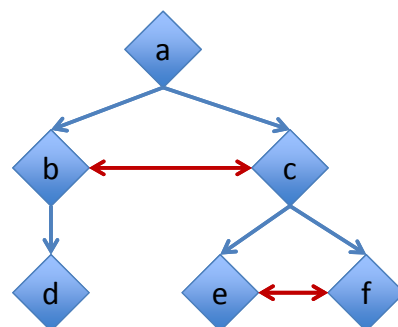
-> **hasSibling**(?y, ?z)



DifferentFrom

hasChild(?x, ?y), hasChild(?x, ?z), **DifferentFrom** (?y, ?z)

-> **hasSibling**(?y, ?z)



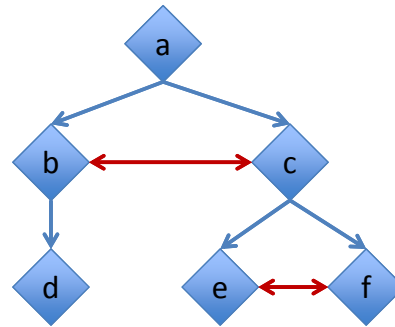
DifferentFrom

$\text{hasChild}(?x, ?y), \text{hasChild}(?x, ?z), \text{DifferentFrom}(?y, ?z)$
→ $\text{hasSibling}(?y, ?z)$

But works only if

$\text{DifferentIndividual}(b, c)$

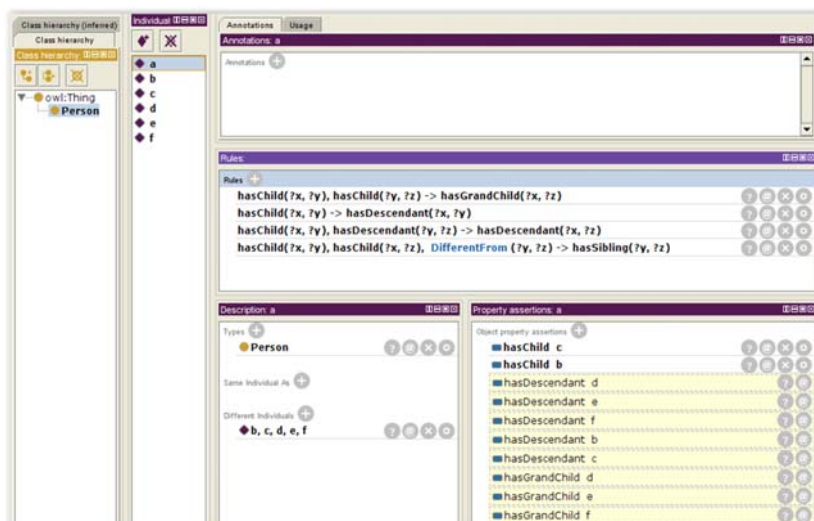
$\text{DifferentIndividual}(e, f)$



Example

$\text{hasChild}(?x, ?y) \rightarrow \text{hasDescendant}(?x, ?y)$

$\text{hasChild}(?x, ?y), \text{hasDescendant}(?y, ?z) \rightarrow \text{hasDescendant}(?x, ?z)$



DL-safe rules

Query answering for DL-axioms + rules is **undecidable**

It is **decidable** if rules are DL-safe

A rule r is called DL-safe if each variable in r occurs in a non-DL-atom in the rule body.

Practically: the variables in rules can only be bound to known individuals

DL-safe rules

Axioms:

TBox: $\text{Parent} \equiv \text{hasChild some Person}$

ABox: $\text{Parent}(a), \text{Parent}(b), \text{Parent}(c), \text{Person}(d), \text{hasChild}(a,d)$

Rule:

$\text{hasChild}(?x, ?y) \rightarrow \text{PersonWithChild}(?x)$

consequence:

?

DL-safe rules

Axioms:

TBox: $\text{Parent} \equiv \text{hasChild some Person}$

ABox: $\text{Parent}(a), \text{Parent}(b), \text{Parent}(c), \text{Person}(d), \text{hasChild}(a,d)$

Rule:

$\text{hasChild}(?x, ?y) \rightarrow \text{PersonWithChild}(?x)$

consequence:

$\text{PersonWithChild}(a)$

without the DL-safe restriction:

$\text{PersonWithChild}(a), \text{PersonWithChild}(b), \text{PersonWithChild}(c)$

Builtin predicates

To deal with numbers, strings, etc.

$\text{Rectangle}(?x), \text{hasWidthInMetres}(?x, ?w), \text{greaterThan}(?w, 10)$

$\rightarrow \text{WideRectangle}(?x)$

$\text{Rectangle}(?x), \text{hasHeightInMetres}(?x, ?h), \text{hasWidthInMetres}(?x, ?w),$
 $\text{multiply}(?a, ?w, ?h), \text{greaterThan}(?a, 100)$

$\rightarrow \text{LargeRectangle}(?x)$

Builtin predicates

swrlb:equal
swrlb:notEqual
swrlb:lessThan
swrlb:lessThanOrEqual
swrlb:greaterThan
swrlb:greaterThanOrEqual

Builtin predicates

swrlb:add
swrlb:subtract
swrlb:multiply
swrlb:divide
swrlb:integerDivide
swrlb:mod
swrlb:pow
swrlb:unaryPlus
swrlb:unaryMinus
swrlb:abs
swrlb:ceiling
swrlb:floor
swrlb:round
swrlb:roundHalfToEven
swrlb:sin
swrlb:cos
swrlb:tan

Builtin predicates

swrlb:stringEqualIgnoreCase
swrlb:stringConcat
swrlb:substring
swrlb:stringLength
swrlb:normalizeSpace
swrlb:upperCase
swrlb:lowerCase
swrlb:translate
swrlb:contains
swrlb:containsIgnoreCase
swrlb:startsWith
swrlb:endsWith
swrlb:substringBefore
swrlb:substringAfter
swrlb:matches
swrlb:replace
swrlb:tokenize

When you don't need SWRL: DL rules

Some SWRL rules can be encoded in OWL expressions

Example

$\text{Man}(\text{?x}) \wedge \text{hasBrother}(\text{?x}, \text{?y}) \wedge \text{hasChild}(\text{?y}, \text{?z}) \rightarrow \text{Uncle}(\text{?x})$

becomes

$\text{Man} \sqcap \exists \text{hasBrother} . \exists \text{hasChild} . \top \sqsubseteq \text{Uncle}$

$\text{Man and (hasBrother some (hasChild some Thing))} \sqsubseteq \text{Uncle}$

it's sometimes tricky ...

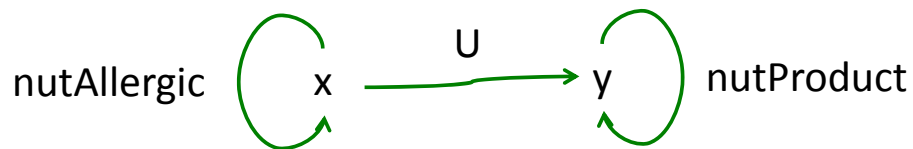
$\text{NutAllergic}(\text{?x}) \wedge \text{NutProduct}(\text{?y}) \rightarrow \text{dislikes}(\text{?x}, \text{?y})$

$\text{NutAllergic} \equiv \exists \text{nutAllergic}.\text{Self}$

$\text{NutProduct} \equiv \exists \text{nutProduct}.\text{Self}$

$\text{nutAllergic} \text{ o } \text{U} \text{ o } \text{nutProduct} \sqsubseteq \text{dislikes}$

U = universal property ($x \text{ U } y$ is always true)

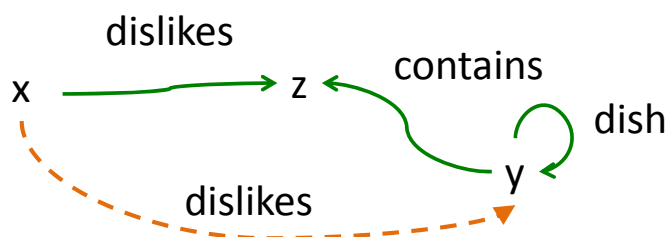


... more

$\text{dislikes}(\text{?x}, \text{?z}) \wedge \text{Dish}(\text{?y}) \wedge \text{contains}(\text{?y}, \text{?z}) \rightarrow \text{dislikes}(\text{?x}, \text{?y})$

becomes

- **$\text{Dish} \equiv \exists \text{dish}.\text{Self}$**
- **$\text{dislikes} \text{ o } \text{contains}^- \text{ o } \text{dish} \sqsubseteq \text{dislikes}$**



Rules versus SPARQL queries

- Rules are “executed” globally
 - all rules must be satisfied simultaneously
- Rules may have interactions
 - the outcome of a rule may trigger another one
- SPARQL queries are executed independently

Simulating rules with queries

define a 'construct' query for each rule

repeat

- execute each query
- add the results to the RDF graph

until nothing new is created

$\text{parent}(?x, ?y) \wedge \text{ancestor}(?y, ?z) \rightarrow \text{ancestor}(?x, ?z)$

construct { $?x$ ancestor $?z$.}

where { $?x$ parent $?y$. $?y$ ancestor $?z$.}

SWRL and Protégé

There is a “rule” view in Protégé
to activate it:

- menu Window -> Views -> Ontology Views -> Rules
- (a black dot appears)
- click the Class Annotation | Class Usage pane

The syntax uses ',' for the logical **and** (not '^')

$C(?x), p(?x, ?y) \rightarrow D(?y)$

SameAs and **DifferentFrom** must start with an uppercase letter.

SWRL inference in Protégé

- Pellet and HermiT support SWRL inference
 - simply run the reasoner to perform swrl inference
 - menu Reasoner -> Start Reasoner or Classify or Synchronize
- HermiT does not support the builtin atoms: add, multiply, lessThan, ... (=> hardly usable)
- Reasoners make the rules DL-safe by binding the variable only to explicitly asserted individuals
 - => reasoning is not complete with respect to the axioms