

# Formalizing double sided auctions in Coq

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## Abstract

In this paper we introduce a formal framework for analyzing double sided auction mechanisms in a theorem prover. In double sided auctions multiple buyers and sellers participate for trade. Any mechanism for double sided auctions to match buyers and sellers should satisfy certain properties of matching. For example, fairness, perceived-fairness, individual rationality are some of the important properties. These are critical properties and to verify them we need a formal setting. We formally define all these notions in a theorem prover. This provides us a formal setting in which we prove some useful results on matching in a double sided auction. Finally, we use this framework to analyse properties of two important class of double sided auction mechanism. All the properties that we discuss in this paper are completely formalized in the Coq proof assistant.

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## 1 Introduction

Trading is a principal component of all modern economy. Over the century more and more complex instruments (for example, index, future, options etc.) are being introduced to trade in the financial markets. With the arrival of computer assisted trading, the volume and liquidity in the markets has improved significantly. Today all big stock exchanges use computer algorithms (matching algorithms) to match buy requests (demands) with sell requests (supplies) of traders. Computer algorithms are also used by many traders to place orders in the markets. This is known as algorithmic trading. As a result of all this the markets has become complex and large. Hence, the analysis of markets is no more feasible without the help of computers.

A potential trader (buyer or seller) places orders in the markets through a broker. These orders are matched by the stock exchange to execute trades. Most stock exchanges divide the trading activity into three main sessions known as pre-markets, continuous markets and post markets. While in the pre-markets session an opening price of a product is discovered through double sided auction. In the continuous markets session the incoming buyers and sellers are continuously matched against each other on a priority basis. In the post-markets session clearing of the remaining orders is done and a closing price is discovered.

A double sided auction mechanism allows multiple buyers and sellers to trade simultaneously [1]. In double sided auctions, auctioneer (e.g. stock exchange) collects buy and sell requests over a period. Each potential trader places the orders with a *limit price*: below which a seller will not sell and above which a buyer will not buy. The exchange at the end of this period matches these orders based on their limit prices. This entire process is completed



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46 using a matching algorithm for double sided auctions.

47 Designing algorithms for double sided auctions is a well studied topic [2, 4, 3, 5]. A major  
 48 emphasis of many of these algorithms is to maximize the number of matches or maximize  
 49 the profit of the auctioneer. Note that an increase in the number of matches increases the  
 50 liquidity in the markets. A matching algorithm can produce a matching with a uniform price  
 51 or a matching with dynamic prices. While an algorithm which clears each matched bid-ask  
 52 pair at a single price is referred as uniform price algorithm. An algorithm which may clear  
 53 each matched bid-ask pair at different prices is referred as dynamic price algorithm. There  
 54 are other important properties besides the number of matches which are considered while  
 55 evaluating the effectiveness of a matching algorithm. For example, fairness, uniform pricing,  
 56 individual rationality are some of the relevant features used to compare these matching  
 57 algorithms. However, it is known that no single algorithm can possess all of these properties  
 58 [4, 2].

59 In this paper, we describe a formal framework to analyze double sided auctions using a  
 60 theorem prover. For this work, we assume that each trader wishes to trade a single unit of  
 61 the product and all the products are indistinguishable as well as indivisible. We have used  
 62 the Coq proof assistant to formally define the theory of double sided auctions. Furthermore,  
 63 we use this theory to validate various properties of matching algorithms. We formally prove  
 64 some important properties of two algorithms; a uniform price algorithm and a dynamic price  
 65 algorithm.

## 66 **2 Modeling double sided auctions**

67 In this section we formally define various concepts involved in a double sided auction  
 68 mechanism. List is a useful data structure for describing various processes and their properties  
 69 in a double sided auction mechanism. In this section, we describe some essential properties  
 70 on lists which are used for stating important results on matching in double sided auctions.

### 71 **2.1 Bid, Ask and limit price**

72 An auction is a competitive event, where goods and services are sold to the highest bidders.  
 73 In a double sided auction multiple buyers and sellers place their orders to buy or sell an item  
 74 to an agent. The agent, known as auctioneer, matches these buy-sell requests based on their  
 75 *limit prices*. While the limit price for a buy order (bid), is the price above which the buyer  
 76 doesn't want to buy one quantity of the item. The limit price of a sell order (ask), is the  
 77 price below which the seller doesn't want to sell one quantity of the item. The notions of bid  
 78 as well as ask can be represented as a record with two fields.

```
79 Record Bid: Type := Mk_bid { bp:> nat;    idb: nat }.
80 Record Ask: Type := Mk_ask { sp:> nat;    ida: nat }.
```

81 For a bid **b**, (**bp b**) is the limit price and (**idb b**) is the unique identity of bid **b**. Similarly  
 82 for an ask **a**, (**sp a**) is the limit price and (**ida a**) is the unique identity of ask **a**. In this  
 83 work we assume that each bid is a buy request for one unit of the item. Similarly each ask is  
 84 a sell request for one unit of the item. If a trader wishes to buy or sell multiple units, he can  
 85 create multiple bids or asks with different ids.

86 Note: Mention about the coercion and Bid and asks are attached to eqType. Mention  
 87 the lemma and canonical structures.

## 2.2 Matching in DSA

In a double sided auction (DSA), the auctioneer collects all the buy and sell requests for a fixed duration. All the buy requests (bids) can be assumed to be present in a list  $B$ . Similarly, all the sell requests (asks) can be present in a list  $A$ . At the end of the duration, the auctioneer matches bids in  $B$  against asks in  $A$ . Furthermore, the auctioneer also assigns a trade price to each matched bid-ask pair. The result of this process is a matching  $M$ , which consists of all the matched bid-ask pairs together with their trade prices. is also represented using list.

In any matching  $M$ , a bid (ask) appears at most once in  $M$ . A bid  $b$  can be matched against an ask  $a$  if  $bp(b) \geq sp(a)$ . We say, a bid-ask pair  $(b,a)$  is matchable if  $bp(b) \geq sp(a)$ . Note that, there can be bids (asks) that are not matched in  $M$ . The collection of bids present in  $M$  is denoted as  $B_M$  and collection of asks present in  $M$  is denoted as  $A_M$ . More precisely, for a given list of bids  $B$  and list of asks  $A$ ,  $M$  is a matching iff, (1) All the bid-ask pairs in  $M$  are matchable, (2)  $B_M$  is duplicate-free, (3)  $A_M$  is duplicate-free, (4)  $B_M \subseteq B$ , and (5)  $A_M \subseteq A$ . Formally, Matching is,

```

Definition matching (M: list fill_type):=
  (All_matchable M) /\ (NoDup (bids_of M)) /\ (NoDup (asks_of M)).

Definition matching_in (B:list Bid) (A:list Ask) (M:list fill_type):=
  (matching M) /\ ((bids_of M) [<=] B) /\ ((asks_of M) [<=] A).

```

108 ► **Lemma 1** (Lorem ipsum). *Vestibulum sodales dolor et dui cursus iaculis. Nullam ullam-*  
 109 *corper purus vel turpis lobortis eu tempus lorem semper. Proin facilisis gravida rutrum.*  
 110 *Etiam sed sollicitudin lorem. Proin pellentesque risus at elit hendrerit pharetra. Integer at*  
 111 *turpis varius libero rhoncus fermentum vitae vitae metus.*

112 **Proof.** Cras purus lorem, pulvinar et fermentum sagittis, suscipit quis magna.

113 ▷ **Claim 2.** content...

114 Proof. content... ◁

115 ◀

116 ► **Corollary 3** (Curabitur pulvinar,). *Nam liber tempor cum soluta nobis eleifend option congue*  
 117 *nihil imperdiet doming id quod mazim placerat facer possim assum. Lorem ipsum dolor sit*  
 118 *amet, consectetur adipiscing elit, sed diam nonummy nibh euismod tincidunt ut laoreet*  
 119 *dolore magna aliquam erat volutpat.*

120 ► **Proposition 4.** *This is a proposition*

121 Proposition 4 and Proposition 4 ...

## 122 2.3 Curabitur dictum felis id sapien

123 Curabitur dictum felis id sapien mollis ut venenatis tortor feugiat. Curabitur sed velit diam.  
 124 Integer aliquam, nunc ac egestas lacinia, nibh est vehicula nibh, ac auctor velit tellus non arcu.  
 125 Vestibulum lacinia ipsum vitae nisi ultrices eget gravida turpis laoreet. Duis rutrum dapibus  
 126 ornare. Nulla vehicula vulputate iaculis. Proin a consequat neque. Donec ut rutrum urna.  
 127 Morbi scelerisque turpis sed elit sagittis eu scelerisque quam condimentum. Pellentesque  
 128 habitant morbi tristique senectus et netus et malesuada fames ac turpis egestas. Aenean  
 129 nec faucibus leo. Cras ut nisl odio, non tincidunt lorem. Integer purus ligula, venenatis et  
 130 convallis lacinia, scelerisque at erat. Fusce risus libero, convallis at fermentum in, dignissim  
 131 sed sem. Ut dapibus orci vitae nisl viverra nec adipiscing tortor condimentum. Donec non  
 132 suscipit lorem. Nam sit amet enim vitae nisl accumsan pretium.

## 133 2.4 Proin ac fermentum augue

134 Proin ac fermentum augue. Nullam bibendum enim sollicitudin tellus egestas lacinia euismod  
 135 orci mollis. Nulla facilisi. Vivamus volutpat venenatis sapien, vitae feugiat arcu fringilla ac.  
 136 Mauris sapien tortor, sagittis eget auctor at, vulputate pharetra magna. Sed congue, dui  
 137 nec vulputate convallis, sem nunc adipiscing dui, vel venenatis mauris sem in dui. Praesent  
 138 a pretium quam. Mauris non mauris sit amet eros rutrum aliquam id ut sapien. Nulla  
 139 aliquet fringilla sagittis. Pellentesque eu metus posuere nunc tincidunt dignissim in tempor  
 140 dolor. Nulla cursus aliquet enim. Cras sapien risus, accumsan eu cursus ut, commodo vel  
 141 velit. Praesent aliquet consectetur ligula, vitae iaculis ligula interdum vel. Integer faucibus  
 142 faucibus felis.

143 ■ Ut vitae diam augue.

144 ■ Integer lacus ante, pellentesque sed sollicitudin et, pulvinar adipiscing sem.

145 ■ Maecenas facilisis, leo quis tincidunt egestas, magna ipsum condimentum orci, vitae  
 146 facilisis nibh turpis et elit.

147 ► **Remark 5.** content...

### 3 Pellentesque quis tortor

Nec urna malesuada sollicitudin. Nulla facilisi. Vivamus aliquam tempus ligula eget ornare. Praesent eget magna ut turpis mattis cursus. Aliquam vel condimentum orci. Nunc congue, libero in gravida convallis, orci nibh sodales quam, id egestas felis mi nec nisi. Suspendisse tincidunt, est ac vestibulum posuere, justo odio bibendum urna, rutrum bibendum dolor sem nec tellus.

► **Lemma 6** (Quisque blandit tempus nunc). *Sed interdum nisl pretium non. Mauris sodales consequat risus vel consectetur. Aliquam erat volutpat. Nunc sed sapien ligula. Proin faucibus sapien luctus nisl feugiat convallis faucibus elit cursus. Nunc vestibulum nunc ac massa pretium pharetra. Nulla facilisis turpis id augue venenatis blandit. Cum sociis natoque penatibus et magnis dis parturient montes, nascetur ridiculus mus.*

Fusce eu leo nisi. Cras eget orci neque, eleifend dapibus felis. Duis et leo dui. Nam vulputate, velit et laoreet porttitor, quam arcu facilisis dui, sed malesuada risus massa sit amet neque.

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### A Styles of lists, enumerations, and descriptions

List of different predefined enumeration styles:

■ `\begin{itemize}...\end{itemize}`

■ ...

■ ...

1. `\begin{enumerate}...\end{enumerate}`

2. ...

3. ...

(a) `\begin{alphaenumerate}...\end{alphaenumerate}`

(b) ...

(c) ...

(i) `\begin{romanenumerate}...\end{romanenumerate}`

189 (ii) ...

190 (iii) ...

191 (1) \begin{bracketenumerate}...\end{bracketenumerate}

192 (2) ...

193 (3) ...

194 **Description 1** \begin{description} \item[Description 1] ... \end{description}

195 **Description 2** Fusce eu leo nisi. Cras eget orci neque, eleifend dapibus felis. Duis et leo dui.

196 Nam vulputate, velit et laoreet porttitor, quam arcu facilisis dui, sed malesuada risus

197 massa sit amet neque.

198 **Description 3** ...

## 199 **B** Theorem-like environments

200 List of different predefined enumeration styles:

201 ► **Theorem 7.** *Fusce eu leo nisi. Cras eget orci neque, eleifend dapibus felis. Duis et leo*  
 202 *dui. Nam vulputate, velit et laoreet porttitor, quam arcu facilisis dui, sed malesuada risus*  
 203 *massa sit amet neque.*

204 ► **Lemma 8.** *Fusce eu leo nisi. Cras eget orci neque, eleifend dapibus felis. Duis et leo dui.*  
 205 *Nam vulputate, velit et laoreet porttitor, quam arcu facilisis dui, sed malesuada risus massa*  
 206 *sit amet neque.*

207 ► **Corollary 9.** *Fusce eu leo nisi. Cras eget orci neque, eleifend dapibus felis. Duis et leo dui.*  
 208 *Nam vulputate, velit et laoreet porttitor, quam arcu facilisis dui, sed malesuada risus massa*  
 209 *sit amet neque.*

210 ► **Proposition 10.** *Fusce eu leo nisi. Cras eget orci neque, eleifend dapibus felis. Duis et leo*  
 211 *dui. Nam vulputate, velit et laoreet porttitor, quam arcu facilisis dui, sed malesuada risus*  
 212 *massa sit amet neque.*

213 ► **Exercise 11.** *Fusce eu leo nisi. Cras eget orci neque, eleifend dapibus felis. Duis et leo*  
 214 *dui. Nam vulputate, velit et laoreet porttitor, quam arcu facilisis dui, sed malesuada risus*  
 215 *massa sit amet neque.*

216 ► **Definition 12.** *Fusce eu leo nisi. Cras eget orci neque, eleifend dapibus felis. Duis et leo*  
 217 *dui. Nam vulputate, velit et laoreet porttitor, quam arcu facilisis dui, sed malesuada risus*  
 218 *massa sit amet neque.*

219 ► **Example 13.** Fusce eu leo nisi. Cras eget orci neque, eleifend dapibus felis. Duis et leo  
 220 dui. Nam vulputate, velit et laoreet porttitor, quam arcu facilisis dui, sed malesuada risus  
 221 massa sit amet neque.

222 ► **Note 14.** Fusce eu leo nisi. Cras eget orci neque, eleifend dapibus felis. Duis et leo dui.  
 223 Nam vulputate, velit et laoreet porttitor, quam arcu facilisis dui, sed malesuada risus massa  
 224 sit amet neque.

225 ► **Note.** Fusce eu leo nisi. Cras eget orci neque, eleifend dapibus felis. Duis et leo dui. Nam  
 226 vulputate, velit et laoreet porttitor, quam arcu facilisis dui, sed malesuada risus massa sit  
 227 amet neque.

228 ► Remark 15. Fusce eu leo nisi. Cras eget orci neque, eleifend dapibus felis. Duis et leo dui.  
229 Nam vulputate, velit et laoreet porttitor, quam arcu facilisis dui, sed malesuada risus massa  
230 sit amet neque.

231 ► Remark. Fusce eu leo nisi. Cras eget orci neque, eleifend dapibus felis. Duis et leo dui.  
232 Nam vulputate, velit et laoreet porttitor, quam arcu facilisis dui, sed malesuada risus massa  
233 sit amet neque.

234 ▷ Claim 16. Fusce eu leo nisi. Cras eget orci neque, eleifend dapibus felis. Duis et leo dui.  
235 Nam vulputate, velit et laoreet porttitor, quam arcu facilisis dui, sed malesuada risus massa  
236 sit amet neque.

237 ▷ Claim. Fusce eu leo nisi. Cras eget orci neque, eleifend dapibus felis. Duis et leo dui.  
238 Nam vulputate, velit et laoreet porttitor, quam arcu facilisis dui, sed malesuada risus massa  
239 sit amet neque.

240 **Proof.** Fusce eu leo nisi. Cras eget orci neque, eleifend dapibus felis. Duis et leo dui. Nam  
241 vulputate, velit et laoreet porttitor, quam arcu facilisis dui, sed malesuada risus massa sit  
242 amet neque. ◀

243 Proof. Fusce eu leo nisi. Cras eget orci neque, eleifend dapibus felis. Duis et leo dui. Nam  
244 vulputate, velit et laoreet porttitor, quam arcu facilisis dui, sed malesuada risus massa sit  
245 amet neque. ◀

