Tutorial-10

TA-Sinong

Quiz-9

QUIZ: GENERATING TYPES

Generate the polymorphic types for the following function:

```
fun fold (f, a, l) =
case l of
nil => a
| h::t => fold (f, f (h, a), t)
```

Please show the intermedia steps and the equations that you are solving.

```
fun fold (f:A, a:B, l:C):D =
  case I of
  nil => a:B
  | h::t:E list => fold (f:A, f (h:E, a:B):((E*B)->B), t:E list)
```

D=B

C=E list

A=E*B->B

fold: ((E*B)->B)*B*(E list) -> B

Homework-9

Problem 1. Given the following variant of untyped lambda calculus:

```
e::=
x (variables)
| c (constants)
| \x.e
l e1 e2
| e1 bop e2 (binary op)
| uop e (unary op)
| let x = e1 in e2
I if e1 then e2 else e3
| letfun f(x) = e1 in e2 (defining a recursive function f(x) for use in e2)
| {e1, e2}
l e.1
e.2
l inl e
l inr e
| case e1 of inl x \Rightarrow e2 | inr x \Rightarrow e3
| nil
l e1 :: e2
| case e1 of nil => e2 | x1 :: x2 => e3
| (e)
```

(a) Inductively define the constraint generation judgement:

$$G \mid -u ==> e:t, q$$

Solution.

$$\frac{G(x) = t}{G \vdash x \Rightarrow x : t, \{\}}$$

$$\frac{G(x) = t}{G \vdash x \Rightarrow x : t, \{\}}$$

$$\frac{G \vdash c \Rightarrow c : int, \{\}}{G \vdash c \Rightarrow c : bool, \{\}}$$

$$\frac{G \vdash c \Rightarrow c : bool, \{\}}{G \vdash c \Rightarrow c : bool, \{\}}$$

$$\frac{G, x : t_1 \vdash t \Rightarrow t : t_2, q}{G, \lambda x : t_1.t : t_1 \rightarrow a, \ q \cup \{a = t_2\}}$$

$$(CT - Var)$$

$$(CT - Int)$$

$$(CT - Bool)$$

$$(CT - Abs)$$

$$\frac{G \vdash u_{1} \Rightarrow e_{1} : t_{1}, q_{1} \quad G \vdash u_{2} \Rightarrow e_{2} : t_{2}, q_{2}}{G \vdash u_{1} u_{2} \Rightarrow e_{1} e_{2} : a, \ q_{1} \cup q_{2} \cup \{t_{1} = t_{2} \to a\}}$$

$$\frac{G \vdash u_{1} \Rightarrow e_{1} : t_{1}, q_{1} \quad G \vdash u_{2} \Rightarrow e_{2} : t_{2}, q_{2}}{G \vdash u_{1} bop \ u_{2} \Rightarrow e_{1} bop \ e_{2} : a, \ q_{1} \cup q_{2} \cup \{t_{1} = t_{2} = a\}}$$

$$\frac{G \vdash u_{1} \Rightarrow e_{1} : t_{1}, q_{1} \quad G \vdash u_{2} \Rightarrow e_{2} : t_{2}, q_{2}}{G \vdash uop \ u \Rightarrow uop \ e : a, \ q \cup \{t = a\}}$$

$$\frac{G \vdash u_{1} \Rightarrow e_{1} : t_{1}, q_{1} \quad G \vdash u_{2} \Rightarrow e_{2} : t_{2}, q_{2}}{G \vdash let \ x = u_{1} \ in \ u_{2} \Rightarrow let \ x = e_{1} \ in \ e_{2} : a, \ q_{1} \cup q_{2} \cup \{t_{2} = a\}}$$

$$\frac{G \vdash u_{1} \Rightarrow e_{1} : t_{1}, q_{1} \quad G \vdash u_{2} \Rightarrow e_{2} : t_{2}, q_{2}}{G \vdash if \ u_{1} \ then \ u_{2} \ else \ u_{3} \Rightarrow if \ e_{1} \ then \ e_{2} \ else \ e_{3} : a,}$$

$$q_{1} \cup q_{2} \cup q_{3} \cup \{t_{1} = bool, t_{2} = t_{3} = a\}$$

$$\frac{G}{G \vdash let fun \ f(x) = u_{1} \ in \ u_{2} \Rightarrow let fun \ f(x : a) : b = e_{1} \ in \ e_{2} : c,}$$

$$q_{1} \cup q_{2} \cup q_{3} \cup \{t_{1} = b, t_{2} = c\}$$

$$\frac{G \vdash u_{1} \Rightarrow e_{1} : t_{1}, q_{1} \quad G \vdash u_{2} \Rightarrow e_{2} : t_{2}, q_{2}}{G \vdash let fun \ f(x) = u_{1} \ in \ u_{2} \Rightarrow let fun \ f(x : a) : b = e_{1} \ in \ e_{2} : c,}$$

$$\frac{q_{1} \cup q_{2} \cup q_{3} \cup \{t_{1} = b, t_{2} = c\}}{G \vdash u_{1} \Rightarrow e_{1} : t_{1}, q_{1} \quad G \vdash u_{2} \Rightarrow e_{2} : t_{2}, q_{2}}$$

$$\frac{G \vdash \{u_{1}, u_{2}\} \Rightarrow \{e_{1}, e_{2}\} : a * b, \ q_{1} \cup q_{2} \cup \{t_{1} = a, t_{2} = b\}}$$

$$(CT - App)$$

$$\begin{array}{ll} G \vdash u \Rightarrow \{e_1,e_2\} : t_1 * t_2, q \\ \hline G \vdash u.1 \Rightarrow e_1 : a, \ q \cup \{t_1 = a\} \\ \hline G \vdash u.2 \Rightarrow e_2 : a, \ q \cup \{t_2 = a\} \\ \hline G \vdash u.2 \Rightarrow e_2 : a, \ q \cup \{t_2 = a\} \\ \hline G \vdash u \Rightarrow e : t, q \\ \hline G \vdash int[a + b] \ u \Rightarrow int[a + b] \ e : a + b, \ q \cup \{t = a\} \\ \hline G \vdash u \Rightarrow e : t, q \\ \hline G \vdash inr[a + b] \ u \Rightarrow inr[a + b] \ e : a + b, \ q \cup \{t = b\} \\ \hline G \vdash u \Rightarrow e : t_1 + t_2, q_1 \ G, x_1 : t_1 \vdash u_1 \Rightarrow e_1 : t, q_2 \ G, x_2 : t_2 \vdash u_2 \Rightarrow e_2 : t, q_3 \\ \hline G \vdash (case \ u \ of \ inl \ x_1 \Rightarrow u_1 | inr \ x_2 \Rightarrow u_2) \Rightarrow (case \ e : a + b \ of \ inl \ x_1 \Rightarrow e_1| \\ \hline G \vdash u \Rightarrow nit[t] : t \ list, \\ \hline G \vdash u \Rightarrow nit[t] : t \ list, \\ \hline G \vdash u_1 \Rightarrow e_1 : t, q_1 \ G \vdash u_2 \Rightarrow e_2 : t \ list, q_2 \\ \hline G \vdash u_1 \Rightarrow e_1 : e_2 : a, \ q_1 \cup q_2 \cup \{a = t \ list\} \\ \hline G \vdash u \Rightarrow e : t_1 \ list, q_1 \ G \vdash u_1 \Rightarrow e_1 : t, q_2 \ G, x_1 : t, x_2 : t_1 \ list \vdash u_2 \Rightarrow e_2 : t, q_3 \\ \hline G \vdash (case \ u \ of \ nit[a] \Rightarrow u_1 | x_1 : x_2 \Rightarrow u_2) \Rightarrow (case \ e : a \ list \ of \ nit[a] \Rightarrow e_1| \\ \hline x_1 : x_2 \Rightarrow e_2) : b, q_1 \cup q_2 \cup q_3 \cup \{t_1 = a, t = b\} \\ \hline G \vdash u \Rightarrow e : t, q \\ \hline G \vdash u \Rightarrow e : t, q \\ \hline G \vdash u \Rightarrow e : unit, q \cup \{t = unit\} \\ \hline \end{array} \qquad (CT - Cnit)$$

(b) Give the detailed derivation of the following expressions and obtain the set of equations, then solve these equations to get the principle solution and give the universal polymorphic types:

```
(1) letfun sum(1) = case 1 of nil => 0 | x1 :: x2 => x1 + sum(x2) in sum(12::10::0::nil)
```

```
Solution.
derivation:
  (letfun\ sum(l:a) = case\ l\ of\ nil:b\ list => 0|x1:c::x2:d => x1 + sum(x2):d \rightarrow e
  in \ sum(12 :: 10 :: 0 :: nil), \{\})
  (by CT-Bop)
\rightarrow (let fun sum(l: a) = case l of nil: b list => 0|x1: c:: x2: d => (x1 + sum(x2)): f
  in \ sum(12 :: 10 :: 0 :: nil), \{d = a, c = e, f = c\}
  (by CT-Casel)
\rightarrow (let fun sum(l: a) = (case l of nil list => 0|x1:: x2 => x1 + sum(x2)): g
  in \ sum(12 :: 10 :: 0 :: nil) : e,
  \{d = a, c = e, f = c, a = b \ list, c = int, f = int, int = q\}
  (by CT-Letfun)
\rightarrow ((letfun\ sum(l:a) = (case\ l\ of\ nil\ list => 0|x1::x2=>x1+sum(x2)):g
  in \ sum(12 :: 10 :: 0 :: nil) : e)h,
  \{d = a, c = e, f = c, a = b \ list, c = int, f = int, int = g, g = e, h = e\}
  (by CT-Cons)
\rightarrow (let fun sum(l: a) = (case l of nil list => 0|x1:: x2 => x1 + sum(x2)): g
  in \ sum(12 :: 10 :: 0 :: nil) : e,
  \{d = a, c = e, f = c, a = b \ list, c = int, f = int, int = g, g = e, h = e, d = int \ list\}
```

solve constraint set:

```
(I, \{d = a, c = e, f = c, a = b \text{ list}, c = int, f = int, int = g, g = e, h = e, d = int \text{ list}\}
\rightarrow([d = a] \circ I, \{c = e, f = c, a = b \text{ list}, c = int, f = int, int = g, g = e, h = e, a = int \text{ list}\})
\rightarrow([c = e] \circ [d = a] \circ I, \{f = e, a = b \text{ list}, e = int, f = int, int = g, g = e, h = e, a = int \text{ list}\})
\rightarrow([f = e] \circ [c = e] \circ [d = a] \circ I, \{a = b \text{ list}, e = int, int = g, g = e, h = e, a = int \text{ list}\})
\rightarrow([a = b \text{ list}] \circ [f = e] \circ [c = e] \circ [d = a] \circ I, \{e = int, int = g, g = e, h = e, b \text{ list} = int \text{ list}\})
\rightarrow([e = int] \circ [a = b \text{ list}] \circ [f = e] \circ [c = e] \circ [d = a] \circ I, \{int = g, h = int, b \text{ list} = int \text{ list}\})
\rightarrow([g = int] \circ [e = int] \circ [a = b \text{ list}] \circ [f = e] \circ [c = e] \circ [d = a] \circ I, \{h = int, b \text{ list} = int \text{ list}\})
\rightarrow([h = int] \circ [g = int] \circ [e = int] \circ [a = b \text{ list}] \circ [f = e] \circ [c = e] \circ [d = a] \circ I, \{b \text{ list} = int \text{ list}\})
\rightarrow([b = int] \circ [h = int] \circ [g = int] \circ [e = int] \circ [a = b \text{ list}] \circ [f = e] \circ [c = e] \circ [d = a] \circ I, \{b \text{ list} = int \text{ list}\})
```

pincipal solution: S(b)=S(c)=S(e)=S(f)=S(g)=S(h) int, S(d)=S(a) int list

universal polymorphic types:

```
\begin{array}{l} letfun\ sum(l:int\ list):int=case\ l\ of\ nil:int\ list=>0 \\ |x1:int::x2:int=>x1+sum(x2)\\ in\ sum(12::10::0::nil)\\ :int \end{array}
```

```
(2) let x = inr (5::4::3) in
  case x of inl y => y.1 + y.2 |
  inr y => (case y of nil => 0 | h::1 => h)
```

```
(let x : b = inr(5 :: 4 :: 3) in

case x : b of inl \ y : c \Rightarrow y.1 + y.2|

inr \ y' : d \Rightarrow (case \ y' \ of \ nil : e \ list \Rightarrow 0 | h : f :: l : g \Rightarrow h), \{\})

(by CT-Proj1 CT-Proj2 and CT-Bop)

→(let x : b = inr(5 :: 4 :: 3) in

case x : b of inl \ y : c \Rightarrow (y.1 + y.2) : j|

inr \ y' : d \Rightarrow (case \ y' \ of \ nil : e \ list \Rightarrow 0 | h : f :: l : g \Rightarrow h), \{c = c_1 * c_2, c_1 = c_2, j = c_1\})

(by CT-Casel)

→(let x : b = inr(5 :: 4 :: 3) in

case x : b of inl \ y : c \Rightarrow (y.1 + y.2) : j|

inr \ y' : d \Rightarrow (case \ y' \ of \ nil : e \ list \Rightarrow 0 | h : f :: l : g \Rightarrow h) : k,

\{c = c_1 * c_2, c_1 = c_2, j = c_1, d = e \ list, g = e \ list, f = int, k = int\})
```

CT-Case \rightarrow (let x:b=inr(5::4::3) in case x:b of inl $y:c \Rightarrow (y.1+y.2):j$ $inr \ y': d \Rightarrow (case \ y' \ of \ nil: e \ list \Rightarrow 0 | h: f:: l: q \Rightarrow h): k,$ $\{c = c_1 * c_2, c_1 = c_2, j = c_1, d = e \text{ list}, q = e \text{ list}, f = int, k = int, b = n + m,$ c = n, d = m, j = kCT-inr $\rightarrow (let \ x : b = inr(5 :: 4 :: 3) \ in$ case x:b of inl $y:c \Rightarrow (y.1+y.2):j$ $inr\ y': d \Rightarrow (case\ y'\ of\ nil: e\ list \Rightarrow 0|h: f:: l: q \Rightarrow h): k,$ $\{c = c_1 * c_2, c_1 = c_2, j = c_1, d = e \ list, g = e \ list, f = int, k = int, b = n + m, \}$ $c = n, d = m, j = k, m = int list\}$ CT-Let \rightarrow ((let x : b = inr(5 :: 4 :: 3)) incase x:b of inl $y:c \Rightarrow (y.1+y.2):i$ $inr \ y': d \Rightarrow (case \ y' \ of \ nil: e \ list \Rightarrow 0 | h: f:: l: g \Rightarrow h): k): o,$ $\{c = c_1 * c_2, c_1 = c_2, j = c_1, d = e \text{ list, } q = e \text{ list, } f = int, k = int, b = n + m,$ $c = n, d = m, i = k, m = int \ list, o = k$

solve constraint set:

```
(I, \{c = c_1 * c_2, c_1 = c_2, j = c_1, d = e \ list, g = e \ list, f = int, k = int, b = n + m,
                    c = n, d = m, i = k, m = int \ list, o = k
 \rightarrow ([c = c_1 * c_2] \circ I, \{c_1 = c_2, j = c_1, d = e \text{ list}, g = e \text{ list}, f = int, k = int, b = n + m,
                    c_1 * c_2 = n, d = m, i = k, m = int \ list, o = k
\rightarrow ([c_1 = c_2] \circ [c = c_1 * c_2] \circ I, \{j = c_2, d = e \text{ list}, g = e \text{ list}, f = int, k = int, b = n + m, d = int, k = int, 
                    c_2 * c_2 = n, d = m, i = k, m = int \ list, o = k
\to ([j=c_2] \circ [c_1=c_2] \circ [c=c_1*c_2] \circ I, \{d=e\ list, g=e\ list, f=int, k=int, b=n+m, s=1, list, f=int, list, f=int, list, list, f=int, list, f=in
                    c_2 * c_2 = n, d = m, c_2 = k, m = int \ list, o = k
\rightarrow ([d = e \ list] \circ [j = c_2] \circ [c_1 = c_2] \circ [c = c_1 * c_2] \circ I, \{g = e \ list, f = int, k = int, b = n + m, j = 1, k = int, 
                    c_2 * c_2 = n, e \ list = m, c_2 = k, m = int \ list, o = k
 \rightarrow ([g = e \ list] \circ [d = e \ list] \circ [j = c_2] \circ [c_1 = c_2] \circ [c = c_1 * c_2] \circ I, \{f = int, k = int, b = n + m, b = n
                    c_2 * c_2 = n, e \ list = m, c_2 = k, m = int \ list, o = k
 \rightarrow ([f = int] \circ [g = e \ list] \circ [d = e \ list] \circ [j = c_2] \circ [c_1 = c_2] \circ [c = c_1 * c_2] \circ I, \{k = int, c_1 = c_2 = c_1 * c_2 = c_1 * c_2 = c_2 = c_1 * c_2 = c_2 = c_1 * c_2 = 
                    b = n + m, c_2 * c_2 = n, e \ list = m, c_2 = k, m = int \ list, o = k
 \rightarrow ([k = int] \circ [f = int] \circ [g = e \ list] \circ [d = e \ list] \circ [j = c_2] \circ [c_1 = c_2] \circ [c = c_1 * c_2] \circ I,
                    \{b = n + m, c_2 * c_2 = n, e \ list = m, c_2 = int, m = int \ list, o = int\}\}
 \rightarrow ([b=n+m]\circ [k=int]\circ [f=int]\circ [g=e\ list]\circ [d=e\ list]\circ [j=c_2]\circ [c_1=c_2]\circ
                     [c = c_1 * c_2] \circ I, \{c_2 * c_2 = n, e \ list = m, c_2 = int, m = int \ list, o = int\})
```

pincipal solution: $S(o)=S(e)=S(c_2)=S(k)=S(f)=S(j)=S(c_1)=int$, S(m)=S(g)=S(d)=int list, S(n)=S(c)=int*int, S(b)=int*int+int list

universal polymorphic types:

```
let \ x: int*int + int \ list = inr(5::4::3) \ in case \ x \ of \ inl \ y: int*int \Rightarrow y.1 + y.2| inr \ y': int \ list \Rightarrow (case \ y' \ of \ nil: int \ list \Rightarrow 0|h: int :: l: int \ list \Rightarrow h) : int
```

```
(3) let x = inl {3::2::1, nil} in
  case x of inl y => (if y.2 == nil
  then case y.1 of nil => 0 | h::1 => h
  else 0)
  inr y => (case y of nil => 0 | h::1 => h)
```

Solution.

```
derivation:
    (let \ x : a = inl \ \{3 :: 2 :: 1, nil : b \ list\} \ in
    case x of inly: c => (if y.2 == nil)
    then case y.1 of nil => 0 | h : d :: l : e => h
    else 0)
    inr \ y': f => (case \ y \ of \ nil => 0|h': d':: l': e'=> h), \{\})
    CT-Case and CT-Proj1 and CT-cons
 \rightarrow (let \ x : a = inl \ \{3 :: 2 :: 1, nil : b \ list\} \ in
    case x of inl y: c \Rightarrow (if y.2 == nil)
    then (case y.1 of nil = > 0 | h : d :: l : e = > h) : t_1
    else 0)
    inr \ y': f => (case \ y \ of \ nil => 0 | h': d':: l': e'=> h), \{c=c_1*c_2, c_1=b \ list \}
    e = d \ list, c_1 = d \ list, d = int, t_1 = int
    CT-If and CT-Bop
 \rightarrow (let \ x : a = inl \ \{3 :: 2 :: 1, nil : b \ list\} \ in
    case x of inly: c => (if y.2 == nil)
    then (case y.1 of nil = > 0 | h : d :: l : e = > h) : t_1
    else\ 0):t_2
    inr \ y': f => (case \ y \ of \ nil => 0 | h': d':: l': e'=> h), \{c=c_1*c_2, c_1=b \ list \}
    e = d \ list, c_1 = d \ list, d = int, t_1 = int, c_2 = b \ list, t_2 = int\}
```

CT-Case and CT-cons $\rightarrow (let \ x : a = inl \ \{3 :: 2 :: 1, nil : b \ list\} \ in$ case x of inl $y: c \Longrightarrow (if \ y.2 == nil)$ then (case y.1 of nil => 0 | h : d :: l : e => h): t_1 $else\ 0):t_2$ $inr \ y': f => (case \ y \ of \ nil => 0 | h' :: l' => h): t_3, \{c = c_1 * c_2, c_1 = b \ list, e = d \ list \}$ $c_1 = d \ list, d = int, t_1 = int, c_2 = b \ list, t_2 = int, f = b \ list, e' = d' \ list, d' = int, t_1 = int, d' = int,$ $d' list = b list, t_3 = int \}$ CT-Case $\rightarrow (let \ x : a = inl \ \{3 :: 2 :: 1, nil : b \ list\} \ in$ $(case\ x\ of\ inl\ y:c=>(if\ y.2==nil)$ then (case y.1 of nil = > 0 | h : d :: l : e = > h): t_1 $else\ 0):t_2$ $inr \ y': f => (case \ y \ of \ nil => 0 | h' :: l' => h): t_3): t_4, \{c = c_1 * c_2, c_1 = b \ list, e = d \ list\}$ $, c_1 = d \ list, d = int, t_1 = int, c_2 = b \ list, t_2 = int, f = b \ list, e' = d' \ list, d' = int,$ d' list = b list, $t_3 = int$, $a = a_1 + a_2$, $a_1 = c$, $a_2 = f$, $t_2 = t_3$, $t_4 = t_2$

CT-inl and CT-Cons

```
\rightarrow (let \ x : a = inl \ \{3 :: 2 :: 1, nil : b \ list\} \ in
        (case x of inl y: c \Rightarrow (if y.2 == nil)
       then (case y.1 of nil => 0 | h : d :: l : e => h): t_1
       else\ 0):t_{2}
       inr\ y': f => (case\ y\ of\ nil => 0|h':: l'=> h): t_3): t_4, \{c=c_1*c_2, c_1=b\ list, e=d\ list\}
       c_1 = d \ list, d = int, t_1 = int, c_2 = b \ list, t_2 = int, f = b \ list, e' = d' \ list, d' = int, t_1 = int, d' = int, 
       d' list = b list, t_3 = int, a = a_1 + a_2, a_1 = c, a_2 = f, t_2 = t_3, t_4 = t_2, a_1 = int list * b list \})
       CT-Let
\rightarrow ((let x : a = inl \{3 :: 2 :: 1, nil : b \ list) in
        (case x of inl y: c \Rightarrow (if y.2 == nil)
       then (case y.1 of nil => 0 | h : d :: l : e => h): t_1
       else\ 0):t_2
       e = d list, c_1 = d list, d = int, t_1 = int, c_2 = b list, t_2 = int, f = b list, e' = d' list, d' = int,
       d' list = b list, t_3 = int, a = a_1 + a_2, a_1 = c, a_2 = f, t_2 = t_3, t_4 = t_2, a_1 = int list * b list,
       t_5 = t_4
```

solve constraint set:

```
(I, \{c = c_1 * c_2, c_1 = b \ list, e = d \ list, c_1 = d \ list, d = int, t_1 = int, c_2 = b \ list, t_2 = int, t_3 = int, t_4 = int, t_5 = int, t_7 = int, t_8 = int, t_8 = int, t_9 = 
                f = b \ list, e' = d' \ list, d' = int, d' \ list = b \ list, t_3 = int, a = a_1 + a_2, a_1 = c_1 * c_2, a_2 = f,
               t_2 = t_3, t_4 = t_2, a_1 = int \ list * b \ list, t_5 = t_4
\rightarrow ([c = c_1 * c_2] \circ I, \{c_1 = b \ list, e = d \ list, c_1 = d \ list, d = int, t_1 = int, c_2 = b \ list, t_2 = int,
                f = b \ list, e' = d' \ list, d' = int, d' \ list = b \ list, t_3 = int, a = a_1 + a_2, a_1 = c_1 * c_2, a_2 = f,
                t_2 = t_3, t_4 = t_2, a_1 = int \ list * b \ list, t_5 = t_4 \}
\rightarrow ([c_1 = b \ list] \circ [c = c_1 * c_2] \circ I, \{e = d \ list, b \ list = d \ list, d = int, t_1 = int, c_2 = b \ list, d = int, t_1 = int, c_2 = b \ list, d = int, t_2 = int, c_2 = b \ list, d = int, t_3 = int, c_4 = int, c_5 = int, c_6 = int, c_7 = int, c_8 = int, c_9 = int, 
               t_2 = int, f = b \ list, e' = d' \ list, d' = int, d' \ list = b" list, t_3 = int, a = a_1 + a_2, a_1 = b \ list * c_2
                a_1, a_2 = f, t_2 = t_3, t_4 = t_2, a_1 = int \ list * b \ list, t_5 = t_4
\rightarrow ([e=d\ list] \circ [c_1=b\ list] \circ [c=c_1*c_2] \circ I, {b\ list=d\ list, d=int, t_1=int, c_2=b\ list, t_1=int, t_2=b\ list, t_2=b\ list, t_1=int, t_2=b\ list, t_2=b\
                t_2 = int, f = b \ list, e' = d' \ list, d' = int, d' \ list = b \ list, t_3 = int, a = a_1 + a_2, a_1 = b' \ list * c_2
                a_1, a_2 = f, t_2 = t_3, t_4 = t_2, a_1 = int \ list * b \ list, t_5 = t_4
\rightarrow ([b=d] \circ [e=d \ list] \circ [c_1=b \ list] \circ [c=c_1*c_2] \circ I, \{d=int, t_1=int, c_2=d \ list, t_1=int, c_2=d \ list, t_1=int, t_2=d \ list, t_2=d \ lis
                t_2 = int, f = d \ list, e' = d' \ list, d' = int, d' \ list = d \ list, t_3 = int, a = a_1 + a_2, a_1 = d \ list * c_2
                a_1, a_2 = f, t_2 = t_3, t_4 = t_2, a_1 = int \ list * d \ list, t_5 = t_4
 \rightarrow ([d=int]\circ [b=d]\circ [e=d\ list]\circ [c_1=b\ list]\circ [c=c_1*c_2]\circ I, \{t_1=int,c_2=int\ list,
                t_2 = int, f = int \ list, e' = d' \ list, d' = int, d' \ list = int \ list, t_3 = int, a = a_1 + a_2,
                a_1 = int \ list * c_2, a_2 = f, t_2 = t_3, t_4 = t_2, a_1 = int \ list * int \ list, t_5 = t_4 \}
```

```
\rightarrow ([c_2 = int \ list] \circ [t_1 = int] \circ [d = int] \circ [b = d] \circ [e = d \ list] \circ [c_1 = b \ list] \circ [c = c_1 * c_2] \circ I, \{
   t_2 = int, f = int \ list, e' = d' \ list, d' = int, d' \ list = int \ list, t_3 = int, a = a_1 + a_2,
   a_1 = int \ list * int \ list, a_2 = f, t_2 = t_3, t_4 = t_2, t_5 = t_4
\rightarrow ([t_2 = int] \circ [c_2 = int \ list] \circ [t_1 = int] \circ [d = int] \circ [b = d] \circ [e = d \ list] \circ [c_1 = b \ list] \circ
   [c = c_1 * c_2] \circ I, \{
   f=int\ list, e'=d'\ list, d'=int, d'\ list=int\ list, t_3=int, a=a_1+a_2, a_1=int\ list*int\ list
   a_{2} = f, int = t_{3}, t_{4} = int, t_{5} = t_{4}
\rightarrow ([f = int \ list] \circ [t_2 = int] \circ [c_2 = int \ list] \circ [t_1 = int] \circ [d = int] \circ [b = d] \circ [e = d \ list] \circ
   [c_1 = b \ list] \circ [c = c_1 * c_2] \circ I, \{
   e'=d' list, d'=int, d' list = int list, t_3=int, a=a_1+a_2, a_1=int list * int list
   a_1, a_2 = int \ list, int = t_3, t_4 = int, t_5 = t_4
\rightarrow ([e'=d'\ list] \circ [f=int\ list] \circ [t_2=int] \circ [c_2=int\ list] \circ [t_1=int] \circ [d=int] \circ [b=d] \circ
   [e = d \ list] \circ [c_1 = b \ list] \circ [c = c_1 * c_2] \circ I, \{
   d' = int, d'  list = int  list, t_3 = int, a = a_1 + a_2, a_1 = int  list * int  list
   ,a_{2}=int\ list,int=t_{3},t_{4}=int,t_{5}\ \overline{\overline{q}}\ t_{4}\})
```

```
\rightarrow ([d'=int]\circ [e'=d'\ list]\circ [f=int\ list]\circ [t_2=int]\circ [c_2=int\ list]\circ [t_1=int]\circ [d=int]\circ 
                       [b = d] \circ [e = d \ list] \circ [c_1 = b \ list] \circ [c = c_1 * c_2] \circ I
                      int\ list = int\ list, t_3 = int, a = a_1 + a_2, a_1 = int\ list * int\ list
                      , a_2 = int \ list, int = t_3, t_4 = intt_5 = t_4 \})
\rightarrow ([d'=int] \circ [e'=d'\ list] \circ [f=int\ list] \circ [t_2=int] \circ [c_2=int\ list] \circ [t_1=int] \circ
                       [d = int] \circ [b = d] \circ [e = d \ list] \circ [c_1 = b \ list] \circ [c = c_1 * c_2] \circ I, \{
                    t_3 = int, a = a_1 + a_2, a_1 = int \ list * int \ list
                       a_{2} = int \ list, int = t_{3}, t_{4} = int, t_{5} = t_{4}
\rightarrow (\lceil t_3 = int \rceil \circ \lceil d' = int \rceil \circ \lceil e' = d' \ list \rceil \circ \lceil f = int \ list \rceil \circ \lceil t_2 = int \rceil \circ \lceil c_2 = int \ list \rceil \circ \lceil t_1 = int \rceil \circ \lceil c_2 = int \ list \rceil \circ \rceil \circ \lceil c_2 = int \ list \rceil \circ \rceil \circ 
                       [d = int] \circ [b = d] \circ [e = d \ list] \circ [c_1 = b \ list] \circ [c = c_1 * c_2] \circ I, \{
                      a = a_1 + a_2, a_1 = int \ list * int \ list
                      , a_2 = int \ list, int = int, t_4 = int, t_5 = t_4 \})
\rightarrow ([a=a_1+a_2]\circ [t_3=int]\circ [d'=int]\circ [e'=d'\ list]\circ [f=int\ list]\circ [t_2=int]\circ
                        [c_2 = int \ list] \circ [t_1 = int] \circ [d = int] \circ [b = d] \circ [e = d \ list] \circ [c_1 = b \ list] \circ [c = c_1 * c_2] \circ I, \{ c_1 = b \ list = c_1 * c_2 = c_2 * c_2 = c_1 * c_2 = c_1 * c_2 = c_1 * c_2 = c_2 * c_2 = c_2
                      a_1 = int \ list * int \ list
                       a_{2} = int \ list, int = int, t_{4} = int, t_{5} = t_{4}
```

```
\rightarrow ([a_1 = int \ list * int \ list] \circ [a = a_1 + a_2] \circ [t_3 = int] \circ [d' = int] \circ [e' = d' \ list] \circ [f = int \ list] \circ
        [t_2=int]\circ [c_2=int\ list]\circ [t_1=int]\circ [d=int]\circ [b=d]\circ [e=d\ list]\circ [c_1=b\ list]\circ
        [c = c_1 * c_2] \circ I, \{a_2 = int \ list, int = int, t_A = int, t_5 = t_A\})
\rightarrow ([a_2 = int \ list] \circ [a_1 = int \ list * int \ list] \circ [a = a_1 + a_2] \circ [t_3 = int] \circ [d' = int] \circ [e' = d' \ list] \circ
        [f = int \ list] \circ [t_2 = int] \circ [c_2 = int \ list] \circ [t_1 = int] \circ [d = int] \circ [b = d] \circ [e = d \ list] \circ
        [c_1 = b \ list] \circ [c = c_1 * c_2] \circ I, \{int = int, t_4 = int, t_5 = t_4\})
\rightarrow ([a_2 = int \ list] \circ [a_1 = int \ list * int \ list] \circ [a = a_1 + a_2] \circ [t_3 = int] \circ [d' = int] \circ
        [e'=d'\ list]\circ [f=int\ list]\circ [t_2=int]\circ [c_2=int\ list]\circ [t_1=int]\circ [d=int]\circ [b=d]\circ
        [e = d \ list] \circ [c_1 = b \ list] \circ [c = c_1 * c_2] \circ I, \{t_A = int, t_5 = t_A\})
\rightarrow ([t_4 = int][a_2 = int \ list] \circ [a_1 = int \ list * int \ list] \circ [a = a_1 + a_2] \circ [t_3 = int] \circ [d' = int] \circ
        [e'=d'\ list]\circ [f=int\ list]\circ [t_2=int]\circ [c_2=int\ list]\circ [t_1=int]\circ [d=int]\circ [b=d]\circ
        [e = d \ list] \circ [c_1 = b \ list] \circ [c = c_1 * c_2] \circ I, \{t_5 = int\})
\rightarrow ([t_5 = int] \circ [t_4 = int] [a_2 = int \ list] \circ [a_1 = int \ list * int \ list] \circ [a = a_1 + a_2] \circ
        [t_3 = int] \circ [d' = int] \circ [e' = d' \ list] \circ [f = int \ list] \circ [t_2 = int] \circ [c_2 = int \ list] \circ [t_1 = int] \circ
        [d = int] \circ [b = d] \circ [e = d \ list] \circ [c_1 = b \ list] \circ [c = c_1 * c_2] \circ I, \{\})
```

```
pincipal solution:
S(a)=int list*int list+int list
S(b)=S(d)=S(d') = int
S(f)=S(e)=S(e')=int list
S(c) = int list*int list
universal polymorphic types:
        let \ x : int \ list * int \ list + int \ list = inl \ \{3 :: 2 :: 1, nil : int \ list \} \ in
        case x of inl y: int list * int list => (if y.2 == nil
        then case y.1 of nil => 0|h:int::l:int\ list => h
        else 0)
        inr\ y': int\ list => (case\ y\ of\ nil => 0|h': int:: l': int\ list => h)
         : int
```

Lemma 1. If a set of constraints q has a solution, then it has a most general one.

Prove this lemma.

Proof.

Lemma 2. If a set of constraint q has a solution, the unification algorithm always return the principal solution

The proof is in the reference book Types and Programming Languages, page 328, 22.4.5

Then lemma 1 is proved.

THEOREM: The algorithm *unify* always terminates, failing when given a non-unifiable constraint set as input and otherwise returning a principal unifier. More formally:

- 1. unify(C) halts, either by failing or by returning a substitution, for all C;
- 2. if $unify(C) = \sigma$, then σ is a unifier for C;
- 3. if δ is a unifier for C, then $unify(C) = \sigma$ with $\sigma \subseteq \delta$.

Proof: For part (1), define the *degree* of a constraint set C to be the pair (m, n), where m is the number of distinct type variables in C and n is the total size of the types in C. It is easy to check that each clause of the *unify* algorithm either terminates immediately (with success in the first case or failure in the last) or else makes a recursive call to *unify* with a constraint set of lexicographically smaller degree.

Part (2) is a straightforward induction on the number of recursive calls in the computation of unify(C). All the cases are trivial except for the two involving variables, which depend on the observation that, if σ unifies $[X \mapsto T]D$, then $\sigma \circ [X \mapsto T]$ unifies $\{X = T\} \cup D$ for any constraint set D.

Part (3) again proceeds by induction on the number of recursive calls in the computation of unify(C). If C is empty, then unify(C) immediately returns the trivial substitution []; since $\delta = \delta \circ$ [], we have [] $\subseteq \delta$ as required. If C is non-empty, then unify(C) chooses some pair (S,T) from C and continues by cases on the shapes of S and T.

Case: S = T

Since δ is a unifier for C, it also unifies C'. By the induction hypothesis, $unify(C) = \sigma$ with $\sigma \subseteq \delta$, as required.

Case: $S = X \text{ and } X \notin FV(T)$

Since δ unifies S and T, we have $\delta(X) = \delta(T)$. So, for any type U, we have $\delta(U) = \delta([X \mapsto T]U)$; in particular, since δ unifies C', it must also unify $[X \mapsto T]C'$. The induction hypothesis then tells us that $unify([X \mapsto T]C') = \sigma'$, with $\delta = \gamma \circ \sigma'$ for some γ . Since $unify(C) = \sigma' \circ [X \mapsto T]$, showing that $\delta = \gamma \circ (\sigma' \circ [X \mapsto T])$ will complete the argument. So consider any type variable Y. If $Y \neq X$, then clearly $(\gamma \circ (\sigma' \circ [X \mapsto T]))Y = (\gamma \circ \sigma')Y = \delta Y$. On the other hand, $(\gamma \circ (\sigma' \circ [X \mapsto T]))X = (\gamma \circ \sigma')T = \delta X$, as we saw above. Combining these observations, we see that $\delta Y = (\gamma \circ (\sigma' \circ [X \mapsto T]))Y$ for all variables Y, that is, $\delta = (\gamma \circ (\sigma' \circ [X \mapsto T]))$.

Case: T = X and $X \notin FV(S)$

Similar.

Case:
$$S = S_1 \rightarrow S_2$$
 and $T = T_1 \rightarrow T_2$

Straightforward. Just note that δ is a unifier of $\{S_1 \rightarrow S_2 = T_1 \rightarrow T_2\} \cup C'$ iff it is a unifier of $C' \cup \{S_1 = T_1, S_2 = T_2\}$.

If none of the above cases apply to S and T, then unify(C) fails. But this can happen in only two ways: either S is Nat and T is an arrow type (or vice versa), or else S = X and $X \in T$ (or vice versa). The first case obviously contradicts the assumption that C is unifiable. To see that the second does too, recall that, by assumption, $\delta S = \delta T$; if X occurred in T, then δT would always be strictly larger than δS . Thus, if unify(C) fails, then C is not unifiable, contradicting our assumption that δ is a unifier for C; so this case cannot occur.