

# Chapter 1 (Part 1): Mathematical Logic

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# Outlines of this talk

- Statements
- Predicate, Quantifiers
- Methods of Proof

# Statements

## ■ Statements

A Statement (or proposition) is a sentence that is definitely true or definitely false

### ■ Examples:

- 1- Every number is divisible by 2.
- 2- Every even number is divisible by 2.
- 3- Oran is the capital of Algeria.
- 4- "A" comes after "B" in the English alphabetic.
- 5- Mahrez scored 572 goals during his career.

■ Non-Examples:

- 1-  $x$  is even.
- 2- She is smart.
- 3- He is a good student.
- 4-  $x + 3 \geq 2$

■ More examples of statements:

P: For every integer  $n > 1$ , the number  $2^n - 1$  is prime.

Q: Every polynomial of degree  $n$  has at most  $n$  roots.

# Statements

If  $P$  is a statement and  $Q$  is another statement, we will define new statements constructed from  $P$  and  $Q$

## A logical conjunction

The conjunction of  $P$  and  $Q$  (read:  $P$  and  $Q$ ) is the statement  $P \wedge Q$  which asserts that  $P$  and  $Q$  are both true. The " $P$  and  $Q$ " statement is false otherwise.

We summarize this in a truth table:

$P$	$Q$	$P \wedge Q$	$Q \wedge P$
T	T	T	T
T	F	F	F
F	T	F	F
F	F	F	F

Figure: Truth table of  $P \wedge Q$

Example:

1-  $(2 + 1 = 5) \wedge (3 + 6 = 1)$

2- P: The number 1 is even and the number 3 is odd.

3- Q : The number 3 is even and the number 2 is odd.

4- R: The number 2 is even and the number 4 is odd.

# Statements

## A logical disjunction

The disjunction of  $P$  and  $Q$  (read:  $P$  or  $Q$ ) is the statement  $\vee$  which asserts that either  $P$  is true, or  $Q$  is true, or both are true, and is false otherwise.

We summarize this in a truth table:

$p$	$q$	$p \vee q$	$q \vee p$
T	T	T	T
T	F	T	T
F	T	T	T
F	F	F	F

Figure: Truth table of  $P \vee Q$

Examples:  $P : \sqrt{2}$  is rational and  $Q : \frac{1+\sqrt{5}}{2} < 2$ .

The statement  $P \vee Q$  is true because  $Q$  is true.

P: The number 2 is even or the number 3 is odd.

Q: The number 1 is even or the number 3 is odd.

R: The number 2 is even or the number 4 is odd.

S: The number 3 is even or the number 2 is odd.



# Statements

## The negation "Not"

A negation is an operator on the logical value of a proposition that sends true to false and false to true. The negation (or logical Not) of  $P$ , denoted  $\neg P$ .

$p$	$\neg p$
T	F
F	T

Figure: Truth table of  $\neg P$

# Statements

## The implication (conditional statement)

The statement  $\neg P \vee Q$  is denoted by  $P \Rightarrow Q$ . We say :

- 1)  $P$  implies  $Q$
- 2) if  $P$ , then  $Q$

- The statement  $P$  is called the hypothesis of the implication, and the statement  $Q$  is called the conclusion of the implication.
- The truth table is:

P	Q	$\neg P$	$P \Rightarrow Q (\neg P \vee Q)$
T	T	F	T
T	F	F	F
F	T	T	T
F	F	T	T

# Statements

## Remark

- 1)  $Q \rightarrow P$ , this is the converse of  $P \rightarrow Q$
- 2)  $\neg P \rightarrow \neg Q$ , this is the inverse of  $P \rightarrow Q$ ,
- 3)  $\neg Q \rightarrow \neg P$ , this is the contrapositive of  $P \rightarrow Q$

Examples:

- 1)  $P: 5^2 < 0$  and  $Q: 1 < 2$ . ( $P \rightarrow Q$  is true)
- 2) " $2 + 2 = 5 \rightarrow \sqrt{2} = 2$ " (true)(if  $P$  is false, the statement  $P \rightarrow Q$  is always true".

# Statements

## The biconditional or double implication

(read:  $P$  if and only if  $Q$ ) is defined by

$$P \iff Q \text{ " is the statement " } (P \rightarrow Q) \text{ and } (Q \rightarrow P) \text{"}$$

This statement is true when  $P$  and  $Q$  are true or when  $P$  and  $Q$  are false.

The truth table is:

$p$	$q$	$p \Rightarrow q$	$q \Rightarrow p$	$p \iff q$
T	T	T	T	T
T	F	F	T	F
F	T	T	F	F
F	F	T	T	T

Figure: Truth table of  $P \iff Q$

Example:  $P$ : " $1 = 2$ ",  $Q$ : "the number of primes is finite"  
 $P \iff Q$  is true because both  $P$  and  $Q$  are false.

# Statements

## Proposition

Let  $P, Q, R$  be three statements. We have the following (true) equivalence:

- $P \iff \neg(\neg P)$
- $P \wedge Q \iff Q \wedge P$  (commutative law for conjunction)
- $P \vee Q \iff Q \vee P$  (commutative law for disjunction)
- $(P \wedge Q) \wedge R \iff P \wedge (Q \wedge R)$  (associative law for conjunction)
- $(P \vee Q) \vee R \iff P \vee (Q \vee R)$  (associative law for disjunction)
- $\neg(P \wedge Q) \iff (\neg P) \vee (\neg Q)$  (De Morgan's law)
- $\neg(P \vee Q) \iff (\neg P) \wedge (\neg Q)$  (De Morgan's law)
- $(P \wedge (Q \vee R)) \iff (P \wedge Q) \vee (P \wedge R)$  (distributive law)
- $(P \vee (Q \wedge R)) \iff (P \vee Q) \wedge (P \vee R)$  (distributive law)
- $P \rightarrow Q \iff \neg Q \rightarrow \neg P$  (law of the contrapositive)

### Tautology, Contradiction, contingency

- 1) A tautology, if it is always true. Example:  $p \vee \neg p$ .
- 2) A contradiction, if it always false. Example:  $p \wedge \neg p$ .
- 3) A contingency, if it is neither a tautology nor a contradiction. Example:  $p$ .

# Predicates and Quantifiers

## Predicates

A predicate is a statement that contains variables. A predicate may be true or false depending on the values of these variables. A predicate is noted by  $P(x)$ ,  $Q(x)$ , .. where  $x$  is in the set  $D$ .

Example:

- 1-  $P(x)$ :  $x$  is human.
- 2-  $Q(x, y)$ :  $x$  is the parent of  $y$ .
- 3-  $R(x)$ :  $x + 2 = x^2$



# Quantifiers

## The universal quantifier

$$\forall x (\text{for all } x), P(x)$$

means that the predicate  $P(x)$  is true for all possible values of  $x$ . An element for which  $P(x)$  is false is called a counterexample of  $\forall x, P(x)$ .

Example:

- "For all integers  $n$ , the integer  $n(n + 1)$  is even".
- We could take a first step towards a symbolic representation of this statement by writing

$$\forall n \in \mathbb{Z}, n(n + 1) \text{ is even}$$

# Quantifiers

## The existential quantifier

$$\exists x (\text{exists } x), P(x)$$

Means that there exists an  $x$  where  $P(x)$  is true.

Sometimes, we will use also

$$\exists! x (\text{exists } x), P(x)$$

It means that there exists a unique  $x$  where  $P(x)$  is true.

- Synonyms for "there exists" include "there is", "there are", "some", and "at least one".

Example:

- "there exists an integer  $n$  such that  $n^2 - n + 1 = 0$ ."
- A symbolic representation of this statement is obtained by writing

$$\exists n \in \mathbb{Z}, n^2 - n + 1 = 0,$$

# Quantifiers

## Negating statements involving quantifiers:

- 1) The negating of " $\forall x \in E \ p(x)$ " is  $\exists x \in E \neg p(x)$ .
- 2) The negating of " $\forall x \in E \ p(x)$ " is  $\exists x \in E \neg p(x)$ .

Example:

- 1) The negation of " $\forall x \in [1, +\infty[ \ (x^2 \geq 1)$ " is the assertion " $\exists x \in [1, +\infty[ \ (x^2 < 1)$ ".
- 2) The negation of  $\exists z \in \mathbb{C} \ (z^2 + z + 1 = 0)$  is  $\forall z \in \mathbb{C} \ (z^2 + z + 1 \neq 0)$ .

It is not more difficult to write the negation of complex sentences. For the statement " $\forall x \in \mathbb{R}, \exists y > 0 \ (x + y > 10)$ " its negation is " $\exists x \in \mathbb{R}, \forall y > 0 \ (x + y \leq 10)$ ".

# Quantifiers

## Remarks :

- 1) The order of quantifiers is very important. For example the two logical sentences are different.

$$"\forall x \in A, \exists y \in B, P(x, y)"$$

$$"\exists y \in B, \forall x \in B, P(x, y)"$$

Example: " In every universty there is a library "

$\forall U$  university,  $\exists I$  library :  $I \in U$  ( Proposition true)

" There is a library which is located in every university "

$\exists I, \forall U : I \in U$

## Remark

$$1) \forall x \in A, \forall y \in B \Leftrightarrow \forall y \in B, \forall x \in A$$

$$2) \exists x \in A, \exists y \in B \Leftrightarrow \exists y \in B, \exists x \in A$$

A proof is an argument that demonstrates why a conclusion is true. A mathematical proof is an argument that demonstrates why a mathematical statement is true, following the rules of mathematics.

# Methods of Proofs

## Direct proof

To prove a statment of the form

$$\text{If } P, \text{ then } Q$$

Assum that  $P$  is true, then show that  $Q$  must be true as well.

**Example:** For each  $x \in \mathbb{Z}$ , if  $x$  is even, then  $5x + 3$  is odd.

**Proof:**(exercise)

## Remark

- 1) Let  $n \in \mathbb{Z}$ . We say  $n$  is an even integer if  $n = 2k$  for some  $k \in \mathbb{Z}$ .
- 2) Let  $n \in \mathbb{Z}$ . We say that  $n$  is an odd integer if  $n = 2k + 1$  for some  $k \in \mathbb{Z}$ .

## Contrapositive proof

Let  $P$  and  $Q$  be statements. Then

$$P \Rightarrow Q \Leftrightarrow \neg Q \Rightarrow \neg P$$

**Example1:** Given  $x \in \mathbb{Z}$ , if  $3x - 7$  is even, then  $x$  is odd.

**Proof:**(exercise)

## Example 2

If the integer  $n^2$  is even, then  $n$  is even.

**Proof:** Suppose that the integer  $n$  is not even, that is, it is odd. We want to show that  $n^2$  is odd. Since  $n$  is odd, there exists an integer  $k$  so that  $n = 2k + 1$ . Then,

$$n^2 = (2k + 1)^2 = 4k^2 + 1 + 4k = 4k^2 + 4k + 1 = 2(2k^2 + 2k) + 1.$$

Since  $2k^2 + 2k$  is an integer,  $n^2$  is odd.

proof by cases looks as follows:

### Proof by cases

proof  $P \Rightarrow Q$  by cases:

- 1 Write  $P \Leftrightarrow r_1 \vee r_2 \vee r_3 \vee \dots \vee r_k$
- 2 Separately prove  $r_i \Rightarrow Q$  for each  $i$ , using any method.
- 3 Conclude that  $P \Rightarrow Q$ , since  $P \Rightarrow Q \Leftrightarrow (r_1 \Rightarrow Q)(r_2 \Rightarrow Q)\dots(r_k \Rightarrow Q)$ .

### Example

If the integer  $n^2$  is a multiple of 3, then  $n$  is a multiple of 3.

Remark In the following example we make use the fact that every integer  $n$  can be uniquely written in the form  $3k + r$ , where  $k$  is an integer and  $r$  equals 0, 1, 2



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### Example

If the integer  $n^2$  is a multiple of 3, then  $n$  is a multiple of 3.

Remark In the following example we make use the fact that every integer  $n$  can be uniquely written in the form  $3k + r$ , where  $k$  is an integer and  $r$  equals 0, 1, 2 The integer  $r$  is the remainder on dividing  $n$  by 3. When the remainder equals 0 we have  $n = 3k$ , so that  $n$  is a multiple of 3.

**Proof** We prove the contrapositive: if  $n$  is not a multiple of 3, then  $n^2$  is not a multiple of 3. Suppose  $n$  is not a multiple of 3. Then the remainder when  $n$  is divided by 3 equals 1 or 2. This leads to two cases:

- Case 1. The remainder on dividing  $n$  by 3 equals 1. Then, there exists an integer  $k$  so that  $n = 3k + 1$ . Hence  $n^2 = (3k + 1)^2 = 9k^2 + 6k + 1 = 3(3k^2 + 2k) + 1$ . Since  $(3k^2 + 2k)$  is an integer, the remainder on dividing  $n^2$  by 3 equals 1. Therefore  $n^2$  is not a multiple of 3.
- Case 2. The remainder on dividing  $n$  by 3 equals 2. Then, there exists an integer  $k$  so that  $n = 3k + 2$ . Hence  $n^2 = (3k + 2)^2 = 9k^2 + 12k + 4 = 3(3k^2 + 4k + 1) + 1$ . Since  $(3k^2 + 4k + 1)$  is an integer, the remainder on dividing  $n^2$  by 3 equals 1. Therefore  $n^2$  is not a multiple of 3.

Both cases have now been considered. In each of them, we have shown that  $n^2$  is not a multiple of 3. It now follows that if  $n$  is not a multiple of 3, then  $n^2$  is not a multiple of 3. This completes the proof.

### Exercise

If the  $x \in \mathbb{Z}$ , then  $x^2 + x$  is even.

## Contradiction:

Another technique is proof by contradiction.

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Let  $R$  and  $S$  be statements.

# Contradiction:

Another technique is proof by contradiction.

## Proof by Contradiction

Let  $R$  and  $S$  be statements. If  $(\neg R) \Rightarrow S$  is true and  $S$  is false, then  $R$  must be true.

## Example

$\sqrt{2}$  is irrational.

**Proof:** Assume  $\sqrt{2}$  is rational, so  $\sqrt{2} = \frac{a}{b}$  with  $a, b$  have no common factors other than 1 and  $b \neq 0$ .

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## Proof by Contradiction

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## Example

$\sqrt{2}$  is irrational.

**Proof:** Assume  $\sqrt{2}$  is rational, so  $\sqrt{2} = \frac{a}{b}$  with  $a, b$  have no common factors other than 1 and  $b \neq 0$ .

$$\sqrt{2} = \frac{a}{b} \Rightarrow 2b^2 = a^2$$

so  $b^2$  is even

so by Example 2,  $a$  is even.

$\exists k \in \mathbb{Z} | a = 2k$ . It now follows that  $2b^2 = a^2 = (2k)^2 = 4k^2$ , so that  $b^2 = 2k^2$ .

Therefore  $b^2$  is even so by Example 2,  $b$  is even. We have now derived the contradiction ( $a$  and  $b$  have no common factors) and ( $a$  and  $b$  are both even).

Therefore,  $\sqrt{2}$  is not rational.

# Counterexample:

## Proof by Counterexample

To show that a proposition  $P : \forall x \in D, P(x)$  is false, we must show that  $\neg P$  is true.

## Example

Let  $f : \mathbb{R} \rightarrow \mathbb{R}$  a function defined by  $f(x) = x^2$

Show that the proposition :

$$\forall (x, y) \in \mathbb{R}^2, f(x) = f(y) \Rightarrow x = y$$

is false.

## Proof by Induction:

**Proof by induction** is a way of proving that a certain statement is true for every positive integer. Proof by induction has four steps:

- 1) Prove the base case: this means proving that the statement is true for the initial value, normally  $n = 1$  or  $n = 0$

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- 2) Assume that the statement is true for the value  $n = k$ . This is called the inductive hypothesis.
- 3) Prove the inductive step: prove that if the assumption that the statement is true for  $n = k$ , it will also be true for  $n = k + 1$ .
- 4) Write a conclusion to explain the proof, saying: "If the statement is true for  $n = k$ , the statement is also true for  $n = k + 1$ . Since the statement is true for  $n = 1$ , it must also be true for  $n = 2$ ,  $n = 3$  and for any other positive integer."

### Example

Prove that for all positive integers  $n$ ,  $3^{2n+2} + 8n - 9$  is divisible by 8.

## Proof by Induction:

**Proof:** First define  $P(n) = 3^{2n+2} + 8n - 9$  is divisible by 8, for  $n \geq 1$

- **Base Case.** For  $n = 1$  we have  $P(1) = 3^{2+2} + 8 - 9 = 3^4 - 1 = 81 - 1 = 80$ . Then, 80 is clearly divisible by 8, hence  $P(1)$  is true.
- **Inductive Step.** Let us set  $n \geq 1$ . Suppose  $P(n)$  is true. We will show that  $P(n+1)$  is true.

$$\begin{aligned}P(n+1) &= 3^{2(n+1)+2} + 8(n+1) - 9 \\&= 3^{2n+4} + 8n + 8 - 9 = 3^{2n+4} + 8n - 9 + 8 \\&= 9 \times 3^{2n+2} + 8n - 9 + 8 \\&= 3^{2n+2} + 8 \times 3^{2n+2} + 8n - 9 + 8 = 3^{2n+2} + 8n - 9 + 8 \times 3^{2n+2} + 8 \\&= P(n) + 8 \times 3^{2n+2} + 8\end{aligned}$$

Since this is the sum of different terms that are all divisible by 8,  $P(n+1)$  must also be divisible by 8 too.

- **Conclusion.** Hence, by the principle of mathematical induction,  $P(n)$  is true for all integers  $n \geq 1$ .