

# Math 571 - Homework 3

Richard Ketchersid

**Problem 3.1.** Define a metric on  $\mathbb{Z}$  for each integer  $n > 1$  as follows. Let  $s \in \mathbb{Z}$  and define  $e_n(s) = \max\{a \in \mathbb{N} \mid n^a \mid s\}$ . Set  $d_n(s, t) = n^{-e_n(s-t)}$  if  $s \neq t$  and  $d_n(s, s) = 0$ .

- a) Show that  $d_n : \mathbb{Z} \times \mathbb{Z} \rightarrow [0, 1]$  is a metric. (Look at (c) before proving the triangle inequality.)

Symmetry and reflexivity are trivial. The triangle inequality is dealt with below.

- b) Interpret the metric, for example what does it mean to say  $d_n(s, t) < \delta$  ( $s$  and  $t$  are within  $\delta$  of each other.)

$d_n(s, t) < d_n(s, t')$  iff there are more copies of  $n$  in the factorization of  $s - t$  than of  $s - t'$ . So the closer  $t$  is to  $s$  the more factors of  $n$  there are in  $s - t$ .

Suppose  $n^a \mid s, t$  and  $s = n^a s'$  and  $t = n^a t'$ , then  $e_n(s, t) = a + e_n(s', t')$ .

- c)  $d_n(s, t) \leq \max\{d_n(s, r), d_n(r, t)\}$ . ( $d_n$  is an **ultrametric**.)

Suppose  $\max\{d_n(s, r), d_n(s, t)\} = d_n(s, r)$ . Let  $a = e_n(s, r)$  so  $n^a \mid s - r$  and  $n^a \mid s - t$ . Now  $s - t = (s - r) + (r - t)$  so  $n^a \mid s - t$  and hence  $d_n(s, t) \leq \max\{d_n(s, r), d_n(r, t)\} = n^{-a}$ .

**Problem 3.2** (R:2:17). Consider all reals in  $[0, 1]$  whose decimal expansion requires only the digits 3 and 5, no 0's, so there are infinitely many 3's and 5's. Call this set  $Y$ . Prove or disprove each of the following:

It is important to think a bit about the representation of reals in  $[0, 1]$  in decimal form before jumping into this. clearly there is a map  $\phi : 10^{\mathbb{N}} \rightarrow [0, 1]$  given by  $\phi(x) = \sum_{i \in \mathbb{N}} x(i)10^{-(i+1)}$ . So for example

$$(3, 1, 4, 1, 5, 9, \dots) \mapsto_{\phi} \frac{3}{10} + \frac{1}{10^2} + \frac{4}{10^3} + \frac{1}{10^5} + \dots = 0.314159$$

There is a very natural topology and metric on  $10^{\mathbb{N}}$  the basic open sets are  $[s] = \{x \mid x \supset s\}$  for  $s \in 10^{<\mathbb{N}}$ . The metric can be expressed as  $d(x, y) = \frac{1}{10^{i+1}}$  where  $i$  is least so that  $x(i) \neq y(i)$ . (This is an ultrametric again). So  $[s] = N_{10^{(-\text{len}(s))}}(x)$  for any  $x \in [s]$ . Note that this is related to but different from the metric on  $[0, 1]$  given by  $\rho(x, y) = |\sum_{i \in \mathbb{N}} (y(i) - x(i))10^{-(i+1)}|$  (the usual metric). In particular,  $\rho(0.100, 0.099) = 0$  so  $\rho$  is not a metric on  $10^{\mathbb{N}}$ . We could use  $d'(x, y) = \sum_{i \in \mathbb{N}} |y(i) - x(i)|10^{-(i+1)}$  which looks more like  $\rho$ , but is still different since  $d'(0.100, 0.099) = \frac{1}{10} + \sum_{i=2}^{\infty} \frac{9}{10^i} = 0.1 + 0.1 = .2$ . Anyway,  $d'$  records more info than we need in  $10^{\mathbb{N}}$  where all we really care about is the first digit on which  $x$  and  $y$  disagree.

Trivially, each  $[s]$  is clopen, since  $[s]$  is open and

$$[s]^c = \bigcup \{t \mid t|n-2 = s|n-2 \wedge t(n-1) \neq s(n-1) \wedge n = \text{len}(s)\}$$

so  $[s]^c$  is open. This means that  $10^{\mathbb{N}}$  is **totally disconnected**, namely,  $\{[s] \mid s \in 10^{<\mathbb{N}}\}$  is a base of clopen sets, so the only connected sets are singletons and the empty set.

Notice that  $\phi|Y' : Y' \rightarrow Y$  where  $Y' = \{3, 5\}^{\mathbb{N}} \subseteq 10^{\mathbb{N}}$  is bijective and continuous. In fact,  $Y$  and  $Y'$  are homeomorphic, so any topological property that  $Y$  has is shared by  $Y'$  and vice-versa. Note by the way that  $Y$  has no finite length decimals, for example,  $0.33 = 0.33\bar{0} = 0.32\bar{9} \notin Y$ .

a )  $Y$  is dense in  $[0, 1]$ ?

It is clear that  $Y'$  is not dense in  $10^{\mathbb{N}}$  since for example  $[1] \cap Y' = \emptyset$ . This can easily be turned into an argument that  $Y$  is not dense on  $[0, 1]$ , namely,  $N_{10^{-1}}(0) \cap Y = \emptyset$ , since the closest element in  $Y$  to 0 is  $0.33333\cdots$ .

b )  $Y$  nowhere dense in  $[0, 1]$ ?

Again, it is trivial to see that  $Y' = \{3, 5\}^{\mathbb{N}}$  is nowhere-dense in  $10^{\mathbb{N}}$ . Take any open set  $O$  and  $[s] \subset O$ . Then  $[s0] \subset [s]$  and  $[s0] \cap Y' = \emptyset$ .

Again a variant of this works in  $[0, 1]$ . Let  $O$  be open, take  $N_{\delta}(x) \subset O$ . Say  $x = 0.d_0d_1\cdots d_i\cdots$  get  $y = 0.d_0\cdots d_i0\cdots$  where  $10^{-(i+2)} < \delta/2$ . Now consider  $N_{10^{-(i+3)}}(y) \subset O$  and disjoint from  $Y$ .

c )  $Y$  is countable?

The standard diagonalization argument shows that  $Y$  ( $Y'$ ) is uncountable.

d )  $Y$  is closed?

Again it is trivial to see that  $Y'$  is closed in  $10^{\mathbb{N}}$ . We can see this in  $[0, 1]$  with essentially the same argument. Suppose  $x_i \in Y$  and  $x_i \rightarrow x$ . Then it is clear that for all  $n$ ,  $x \mid i \subset x \mid n$  for all sufficiently large  $i$ . Thus  $x \in Y$ .

e )  $Y$  is compact?

This is easy to show directly, but  $Y$  is a closed subset of a compact space, hence compact.

f )  $Y$  is perfect?

Again this is trivial looking at  $Y'$  inside of  $10^{\mathbb{N}}$  and the same argument works in  $[0, 1]$ .

**Problem 3.3** (R:2:20\*). If  $E$  is connected is  $\text{Cl}(E)$  and/or  $\text{Int}(E)$  necessarily connected?

Of course, give a proof or a counterexample.

The easiest here is  $\text{Int}(E)$ . Take  $E \subseteq \mathbb{R}^2$  to be  $N_1((-1, 0)) \cup N_1((0, 1))$ . So  $E$  consists of unit discs centered at  $(-1, 0)$  and  $(0, 1)$  that just “kiss” at  $(0, 0)$ .  $\text{Int}(E)$  is the union of the interiors of the two discs, and these form their own separating sets. So if  $E$  is connected,  $\text{Int}(E)$  **need not be** connected.

Suppose  $A, B$  witness that  $\text{Cl}(E)$  is disconnected. Then

- i)  $A' = A \cap E \neq \emptyset \neq B \cap E = B'$ ,
- ii)  $A' \cup B' = E$ , and

iii)  $A' \cap B' = \emptyset$ .

So  $E$  must also be disconnected. (i) uses the following:

For any open  $O$ :

$$O \cap E \neq \emptyset \iff O \cap \text{Cl}(E) \neq \emptyset$$

( $\implies$ ) is trivial. For the other direction argue the contrapositive

$$O \cap E = \emptyset \implies O \cap \text{Cl}(E) = \emptyset$$

This is true since

$$O \cap E = \emptyset \implies O \subseteq \text{Int}(E^c) \implies O \cap (\text{Int}(E^c))^c = O \cap \text{Cl}(E) = \emptyset$$

So if  $E$  is connected, then  $\text{Cl}(E)$  is connected.

**Problem 3.4.** Show that  $E$  is connected iff for all  $p, q \in E$  there is a connected open relative to  $E$  set  $A \subseteq E$  with  $p, q \in A$ .

The “only if” part is trivial since  $E$  is open in  $E$ ,  $E$  is connected, and  $p, q \in E$ .

For the “if” part show the contrapositive. Suppose  $E$  is not connected, then there are open (relative to  $E$ )  $A, B \subset E$  such that  $A, B \neq \emptyset$  and  $A \cap B = \emptyset$ , and  $A \cup B = E$ . Let  $p \in A$  and  $q \in B$ . Let  $O \subset E$  be open in  $E$  with  $p, q \in O$ . Then  $A' = A \cap O$  and  $B' = B \cap O$  witness that  $O$  is not connected. So there is no open connected subset of  $E$  containing  $p$  and  $q$ .

**Problem 3.5** (R:2:21\*). Prove that every convex subset of  $\mathbb{R}^k$  is connected.

The original problem in Rudin is a four part problem with this being the last part. You might use the original problem as a hint/guide here.

Suppose  $G$  is our convex set and let  $A, B \subset G$  be non-empty and open in  $G$  with  $A \cup B = G$ . Let  $a \in A$  and  $b \in B$ . Then  $\{ta + (1-t)b \mid t \in [0, 1]\} \subseteq G$ . Look at  $\{t \in [0, 1] \mid ta + (1-t)b \in A\} = A'$  and  $B' = \{t \mid ta + (1-t)b \in B\}$ .

**Claim:**  $A', B'$  witness that  $[0, 1]$  is not connected.

The only thing that requires argument is that  $A'$  and  $B'$  are open in  $[0, 1]$ . Let  $c = ta + (1-t)b \in A$ . Then  $B_\epsilon(c) \subset A$  for some  $\epsilon > 0$ . Consider,  $(t+h)a + (1-(t+h))b = ta + (1-t)b + h(a-b) = c + h(a-b)$ . If  $|h| < \epsilon/|a-b| = \delta$ , then  $t+h \in A'$ . This means that  $(t-\delta, t+\delta) \cap [0, 1] \subseteq A'$  and so  $A'$  is open in  $[0, 1]$ .

**Fact:**  $[0, 1]$  is connected. (Rudin 2.47)

**Problem 3.6** (R:2:26). Let  $X$  be a metric space in which every infinite set has a limit. Show that  $X$  is compact.

I prove this in the notes. It is an important and very useful characterization of compactness in a metric space, namely, **sequential compactness**. I do not want you to reproduce the proof I give. Use the hint from Rudin and try it the way he suggests. This builds on some problems you did last week.

Let  $\mathcal{O}$  be an open cover of  $X$ . Our goal is to produce a finite subcover of  $X$ . Problems 6 from homework 2 gives us that  $X$  has a countable base. Thus we can easily get a countable subcover  $\mathcal{O}' \subseteq \mathcal{O}$ , simply assign to each  $x$  a base set  $U_x$  and  $O_{U_x}$  so that  $x \in U_x \subseteq O_{U_x} \in \mathcal{O}$ . Since there are only countably many  $U_x$ , there are also only countable many  $O_{U_x} \in \mathcal{O}$  required. So we may assume  $\mathcal{O} = \{G_i \mid i \in \mathbb{N}\}$ .

Assume there is no finite subcover from  $\mathcal{O}$ , then  $\bigcup_{i=0}^n G_i$  fails to cover  $X$  and so  $F_n = X - \bigcup_{i=0}^n G_i$  is closed and non-empty. Further,  $F_{n+1} \subseteq F_n$  and by assumption  $\bigcap_{i \in \mathbb{N}} F_i = \emptyset$ . Let  $x_i \in F_i$  for each  $i$ . By assumption there is  $x \in \text{Lim}(\{x_i\}_{i \in \mathbb{N}})$ , but then  $x \in F_i$  for all  $i$  since  $\{x_k\}_{k \geq i} \subseteq F_i$  and  $F_i$  is closed. This is a contradiction so the assumption that there is no finite subcover must be false.

**Problem 3.7** (R:2:28). Show that every closed set,  $F$ , in a separable metric space can be written as  $F = P \cup C$  where  $P$  is perfect (perhaps empty) and  $C$  is countable.

A different hint from Rudin's: I gave you a sort of hint in class, define  $F' = F - \text{Iso}(F)$ , recall  $\text{Iso}(F)$  is the set of isolated points of  $F$ .  $F'$  is called the derivative of the set  $F$ . Argue that  $\text{Iso}(F)$  is countable, in some natural sense  $F'$  is *closer to perfection*, since we have removed some isolated points. Notice that  $F'$  is closed. If you haven't reached perfection repeat the process. In this way you build a sequence of closed sets  $F \supset F_1 \supset F_2 \cdots$  and countable sets  $C_i$  so that  $F = \bigcap F_i \cup \bigcup C_i$ . If  $\bigcap F_i = F_\omega$  still has isolated points, continue! A transfinite recursion!

**Proof 1:** Continue as suggested in the hint. There is a strictly descending sequence  $F_\alpha \supset F_\beta$  for  $\alpha < \beta < \gamma$  with the additional property that  $C_\beta = \text{Iso}\left(\bigcap_{\alpha < \beta} F_\alpha\right) \neq \emptyset$  and  $F_\beta \cup C_\beta = \bigcup_{\alpha < \beta} F_\alpha$ . But let  $c_\beta \in C_\beta$ , then there is open  $O_\beta$  with  $c_\beta \in O_\beta$  and  $s_\beta \in O_\beta \cap S$ , where  $S$  is the separable set. Clearly, for  $\alpha < \beta$ ,  $s_\alpha \neq s_\beta$  so this must halt after countably many steps. That is we reach  $\gamma$  so that  $\bigcap_{\alpha < \gamma} F_\alpha = F$  and  $\text{Iso}(F) = \emptyset$ , for  $F$  is perfect and  $C = \bigcup_{\alpha < \gamma} C_\alpha$  is countable.

**Proof 2:** (Follow the text.) Let  $P$  be the set of condensation points of  $F$ . If  $x \notin P$ , then  $x \in O$  for some open  $O$  with  $O \cap F$  countable. We can choose this  $O$  from a countable base,  $\mathcal{B}$  and thus  $C = \bigcup \{O \cap F \in \mathcal{B} \mid |O \cap P| \leq \aleph_0\}$ .  $C$  is countable and  $F = C \cup P$ .

If  $x \in P$ , then  $x$  is a condensation point of  $P$ , not just a condensation point of  $F$ . For suppose there is  $O$  with  $x \in O$  and  $|O \cap P| \leq \aleph_0$ , then  $|O \cap F| = |(O \cap P) \cup (O \cap C)| \leq \aleph_0$ . So clearly,  $\text{Iso}(P) = \emptyset$  and  $P \subset \text{Lim}(P)$ , so  $P$  is closed.