Embedded OS Implementation, Fall 2020

Project #3 (due January 6, 2021 (Wednesday) 13:00)

[PART I] NPCS Implementation

Objective:

To implement the non-preemptible critical section (NPCS) based on RM scheduler in uC/OS-II.

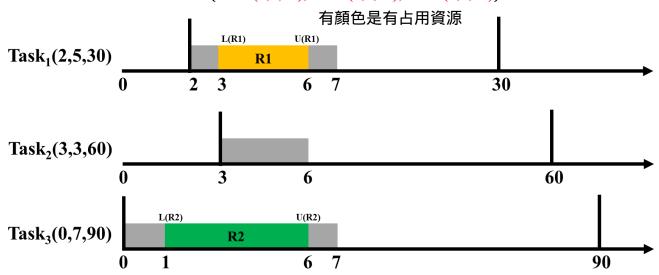
Problem Definition:

uC/OS-II uses a variation of the priority-inherent protocol to deal with priority inversions. In this assignment, you are going to implement the NPCS based on RM scheduler in uC/OS-II.

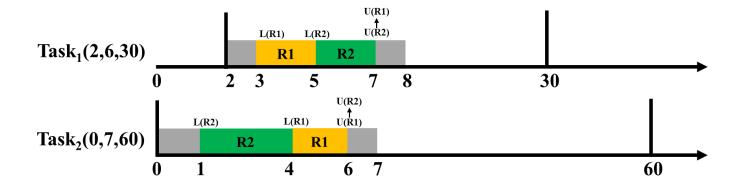
Consider the two periodic task sets and observe how the task suffers the schedule delay.

Periodic Task Set = $\{task_{ID} (arrival time, execution time, period)\}$

Task Set $1 = \{ task_1(2,5,30), task_2(3,3,60), task_2(0,7,90) \}$



Task Set $2 = \{ task_1(2,6,30), task_2(0,7,60) \}$



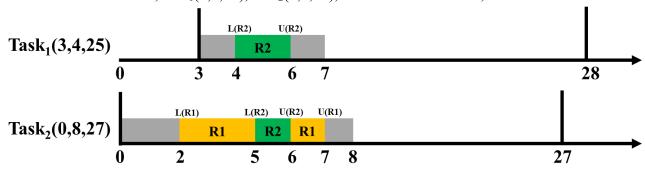
Evaluation:

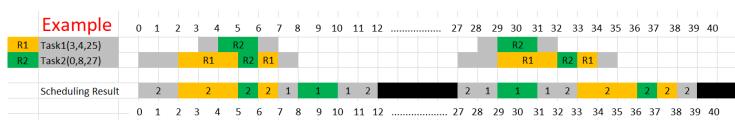
The output format:

Tick	Event
#	Task ID
#	Task ID get R
#	Task ID release R

The NPCS Example of Output Result:

Consider two tasks, task₁(3,4,25), task₂(0,8,27), and two resources R1, R2.





```
ick
               Event
                            get R1
                            get R2
                            release R2
                            release R1
                            get R2
8
10
11
12
27
28
31
32
33
36
37
38
               Task
                            release R2
               Task
               Task
               Task 63
               Task
                            get R2
                            release R2
               Task
              Task 1 Telease R2
Task 2
Task 2 get R1
Task 2 get R2
Task 2 release R2
Task 2 release R1
```

[PART II] CPP Implementation

Objective:

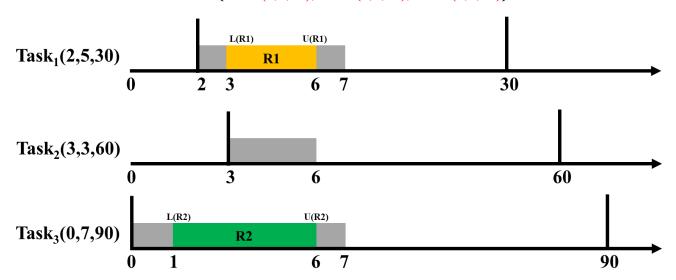
To implement the ceiling-priority protocol (CPP) based on RM scheduler in uC/OS-II.

Problem Definition:

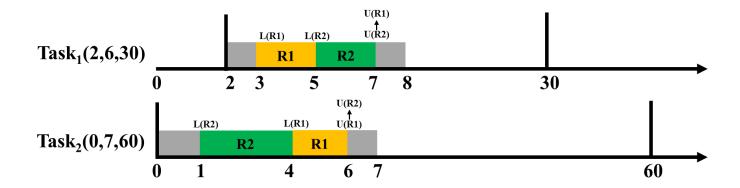
uC/OS-II uses a variation of the priority-inherent protocol to deal with priority inversions. In this assignment, you are going to implement the CPP based on RM scheduler in uC/OS-II.

Consider the two periodic task sets and observe how the task suffers the schedule delay.

Periodic Task Set = $\{task_{ID} (arrival time, execution time, period)\}$ Task Set 1 = $\{task_1 (2,5,30), task_2 (3,3,60), task_2 (0,7,90)\}$



Task Set $2 = \{ task_1(2,6,30), task_2(0,7,60) \}$



Evaluation:

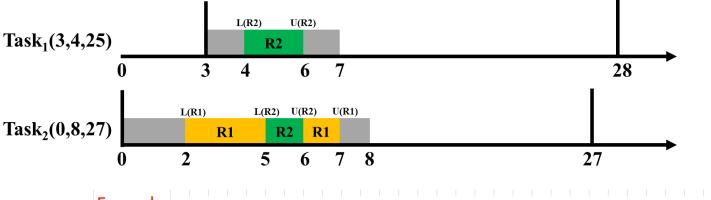
The output format:

Tick	Event	SystemCeiling
#	Task ID	
#	Task ID get R	Ceiling_prio→Ceiling_prio
#	Task ID release R	Ceiling_prio→Ceiling_prio

X If the system ceiling is the lowest, use the symbol " Ω ".

The CPP Example of Output Result:

Consider two tasks, $task_1(3,4,25)$, $task_2(0,8,27)$, and two resources R1, R2.



Ceiling	g	Example	(0 1	1	2 3	3 4	1 !	5	6	7	8	9	10	11	12			. 2	7 2	8 2	9 3	0 3	1 3	2 3	33 3	4 3	5 3	6 37	7 38	8 39	40
2	R1	Task1(3,4,25)						F	12													R	2									
1	R2	Task2(0,8,27)					R1		R2	R1													R1		R2	R1						
		System Ceiling		Ω	Ω	2	2	1	1	2	2	2	1	2	Ω	Ω	Ω	Ω	Ω	Ω	Ω	1	1	Ω	Ω	2	2	2	1	2	Ω	
		Scheduling Result		2		2	1		1	1		2	2	2	2					2	1		L	1	2		2		2	2	2	
				0 :	1	2 3	3 4	1	5	6	7	8	9	10	11	12			2	7 2	8 2	9 3	0 3	31 3	32 3	3 3	34 3	35 3	6 3	7 3	8 39	40

```
        Tick
        Event
        SystemCeiling

        0
        Task 2
        2

        2
        Task 2 get R1
        Ω→2

        3
        Task 1
        2

        4
        Task 1 get R2
        2→1

        6
        Task 1 release R2
        1→2

        7
        Task 2
        2

        9
        Task 2 get R2
        2→1

        10
        Task 2 release R2
        1→2

        11
        Task 2 release R1
        2→Ω

        12
        Task 63
        27

        27
        Task 1
        2

        28
        Task 1
        2

        29
        Task 1 get R2
        Ω→1

        31
        Task 2 get R2
        1→Ω

        32
        Task 2 get R1
        Ω→2

        33
        Task 2 get R2
        2→1

        36
        Task 2 get R2
        2→1

        37
        Task 2 release R2
        1→2

        38
        Task 2 release R1
        2→Ω

        39
        Task 63
```

Crediting:

[PART I] NPCS Implementation [50%]

- The screenshot result (with the given format) of the two task sets. (Time tick 0-100) (10%)
- A report that describes your implementation, including scheduling results of two task sets, modified functions, data structure, etc. (please ATTACH the screenshot of the code and MARK the modified part). (40%)

[PART II] CPP Implementation [40%]

- The screenshot result (with the given format) of the two task sets. (Time tick 0-100) (10%)
- A report that describes your implementation, including scheduling results of two task sets, modified functions, data structure, etc. (please **ATTACH** the screenshot of the code and **MARK** the modified part). (30%)

[PART III] Performance Analysis [10%] 比較分析 優缺點 怎麼避免deadlock

- Compare the scheduling behaviors between NPCS and CPP with the results of PART I and PART II. (5%)
- Explain how NPCS and CPP avoid the deadlock problem. (5%)
- * You must modify the source code.

Project submit:

Submit to Moodle

Submit deadline: January 6, 2021 (Wednesday) 13:00

File name format: RTOS_your student ID_PA3.zip

RTOS_your student ID_PA3.zip includes:

- The report (RTOS_your student ID_PA3.pdf).
- The files you modify(main.c, os core.c, etc.)

Hints:

1. In the application region, we define priorities of tasks and shared resources.

```
#define R1_PRIO 1
#define R2_PRIO 2
#define TASK1_PRIORITY 11
#define TASK2_PRIORITY 12
#define TASK3_PRIORITY 13
#define TASK4_PRIORITY 14
```

2. We also declare shared resource, as follows:

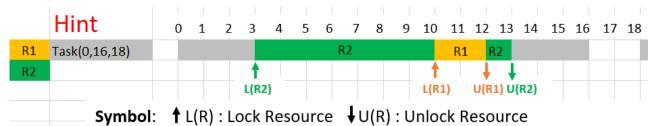
```
OS_EVENT* R1;
OS_EVENT* R2;
```

3. In the main function, we not only create tasks but also create shared resources.

```
INT8U err;
R1 = OSMutexCreate(R1_PRIO, &err);
R2 = OSMutexCreate(R2_PRIO, &err);
```

4. To simulate the duration that a resource is held, we can program a function to implement it:

5. To modeling a task's behavior, we can program the task function as following:



```
void task1(void* pdata)
    INT8U err;
    while (1)
        printf("%d\tTask 1\n", OSTimeGet());
        mywait(3);
        printf("%d\tTask 1 get R2\n", OSTimeGet());
        OSMutexPend(R2, 0, &err);
        mywait(7);
        printf("%d\tTask 1 get R1\n", OSTimeGet());
        OSMutexPend(R1, 0, &err);
        mywait(2);
        printf("%d\tTask 1 release R1\n", OSTimeGet());
        OSMutexPost(R1);
        mywait(1);
        printf("%d\tTask 1 release R2\n", OSTimeGet());
        OSMutexPost(R2);
        mywait(3);
        OSTimeDly( T1_Deadline - OSTimeGet());
```