

Pre-fix Sum Equals K :

case 1

input: $\text{nums} = [1, 1, 1], k = 2$

output: 2

if you ever get a pre-fix sum, where two values are the same, that means the value in between those two values sum to zero.

$$\text{sum}_i - \text{sum}_j = 0$$

case 2

input: $\text{nums} = [1, 2, 3], k = 3$

output: 2

if: $\text{sum}_i - k = \text{sum}_j$ then, we know the difference between sum_i and sum_j is actually k .

deriving the idea

$$\rightarrow \text{sum}_i - \text{sum}_j = k$$

$$\rightarrow \text{sum}_i = k + \text{sum}_j$$

$$\rightarrow \text{sum}_i - k = \text{sum}_j$$

$\begin{matrix} 1 & 2 & 3 \\ \hline [1] & [1] & [1] \end{matrix}, \begin{matrix} 1 & 3 & 6 & 3 \\ \hline [1] & [2, 3] & [-3] \end{matrix}$

Using a hash map, we will store $\text{prefix_sum} - k$ as the key and the # of times we've seen it as the value.

handling case 1: Assume we have an empty hash map: $\{\}$

prefix_sum: $[1, 1, 1]$

1+1

↑ this is an answer, but we only add it if $\text{prefix_sum} - k$ is already in our hash map. So, we're not counting it towards our number of solutions.

intuition: the first time we see a contiguous subarray that equals to k , we won't count that solution.

To handle this case, we need to add $\{0:1\}$ to account for the first solution.

pseudocode :

input: array of integers, nums and target value k

output: # of contiguous subarrays equal to k .

subarray(nums, k):

prefix_sum, result = 0;

prefix_hash = $\{0:1\}$

for num in nums :

prefix_sum += num;

if prefix_sum - k in prefix_hash:

result += prefix_hash[prefix_sum - k]

if prefix_sum not in prefix_hash:

prefix_hash[prefix_sum] += 1;

else:

prefix_hash[prefix_sum] = 1;

return result;