

COS511: Theoretical Machine Learning  
Homework 1

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**Problem. 2a**

*Solution.* Let us repeat steps 1-7 of the proof of learning bounds from lecture 4, to begin to bound the probability that  $h_A$  is consistent and  $\epsilon$ -bad, by showing that  $Pr[h_A \text{ is consistent and } \epsilon\text{-bad}] \leq \sum_{h \in B} (1 - \epsilon)^m$ . These steps have been omitted from the proof for brevity.

For every  $x$ ,  $1 + x \leq e^x$ , so

$$\sum_{h \in B} (1 - \epsilon)^m \leq \sum_{h \in B} e^{-\epsilon m}$$

We know  $\epsilon = \frac{1}{m}(\ln(1/g(h)) - \ln(1/\delta))$  by problem statement.

Thus  $-\epsilon m = \ln(1/g(h)) - \ln(1/\delta)$  by algebra.

And  $-\epsilon m = \ln(g(h)) + \ln(\delta)$  by log rules.

Therefore,

$$\sum_{h \in B} e^{-\epsilon m} = \sum_{h \in B} e^{\ln(g(h)) + \ln(\delta)} = \delta \sum_{h \in B} g(h)$$

And since  $B$  is a subset of  $H$  by definition,

$$\delta \sum_{h \in B} g(h) \leq \delta \sum_{h \in H} g(h)$$

And since  $\sum_{h \in H} g(h) \leq 1$  by problem statement,

$$\delta \sum_{h \in H} g(h) \leq \epsilon$$

Thus we have shown that,

$$\forall h_A \in H, Pr[err_D(h_A) > \epsilon] \leq \delta$$

Which means, with probability  $\geq 1 - \delta$ , and  $\forall h \in H$  if  $h$  is consistent, then

$$err_D(h) \leq \frac{\ln(\frac{1}{g(h)}) + \ln(\frac{1}{\delta})}{m}$$

QED

**Problem. 2b**

*Solution.* Since we know that  $err_D(h) \leq \frac{\ln(\frac{1}{g(h)}) + \ln(\frac{1}{\delta})}{m}$ , to prove that  $err_D(h) \leq O(\frac{|h| + \ln(\frac{1}{\delta})}{m})$ , we need to choose a  $g$  such that  $O(|h|) = \ln(\frac{1}{g(h)})$ .

So suppose we construct  $O(|H|)$  for some scalar  $k$

$$k|h| = O(|H|) = \ln(\frac{1}{g(h)})$$

Then

$$e^{-k|h|} = g(h)$$

And now we need  $\sum_{h \in H} g(h) = \sum_{h \in H} e^{-k|h|} \leq 1$ .

Since  $H$  has  $2^{|h|}$  bitstrings for each  $|h|$ ,

$$\sum_{h \in H} e^{-k|h|} \leq \sum_{|h|=1} 2^{|h|} e^{-k|h|}$$

By geometric series formula,

$$\sum_{|h|=1} 2^{|h|} e^{-k|h|} = \frac{2}{e^k} \frac{1}{1 - \frac{2}{e^k}} = \frac{1}{\frac{e^k}{2} - 1}$$

Now, if we set  $k$  to 4,

$$\frac{1}{\frac{e^k}{2} - 1} = 1 \leq 1$$

Therefore, we have found our  $k$ .

Then,  $g(h) = e^{-\ln(4)|h|} = \frac{1}{4^{|h|}}$ .

Choosing  $g$  to be  $\frac{1}{4^{|h|}}$  satisfies that  $g(h) \in (0, 1]$  for all  $h$ .

QED.

### **Problem. 2c**

*Solution.* In (b), the bound shows that simpler hypotheses probabilistically result in less generalization error than more complex hypotheses, since lowering  $|h|$  lowers error.

In (a), the bound shows that more training data (more priors) probabilistically result in less generalization error, since increasing  $m$  lowers error.