

1.) In the domain of all students, we define predicates

$M(x)$: x is a math major

$C(x)$: x is a computer science major

$A(x)$: x is required to take CS 225 .

Express each of the following English sentences in terms of $M(x)$, $C(x)$, $A(x)$, quantifiers, and logical connectives.

(a) Some computer science majors are required to take CS 225.

$$\exists x | C(x) \wedge A(x)$$

(b) Not all math majors are computer science majors.

$$\sim \forall x (M(x) \wedge C(x))$$

(c) All students who are both computer science and math majors are not required to take CS 225.

$$\forall x ((C(x) \wedge M(x)) \rightarrow \sim A(x))$$

(d) There is a student who is neither a math major and is not required to take CS 225.

$$\exists x | \sim M(x) \wedge \sim A(x)$$

2.) Let $B(x)$, $S(x)$, and $A(x)$ be the predicates

$B(x)$: x is a good basketball player

$S(x)$: x is a good soccer player

$A(x)$: x is a good athlete

Translate each of the following quantified logic expressions (provided in the file) into English considering the domain to consist of all people .

1.) For all people, if someone is a good athlete, then he/she is a good basketball and soccer player.

2.) Not everyone is a good soccer player.

3.) There is a person who is a good soccer and basketball player or not a good athlete.

4.) There is a person who is not both a good soccer player and not a good basketball player.

3.) Negate each of the following statements:

1) Every real number is positive or negative or zero .

There is no number that is positive or negative or zero

2) Some people are not honest .

All people are honest.

3) Not all young people are good athletes.

All young people are good athletes.

4.) Prove or disprove that $\forall x (P(x) \rightarrow Q(x))$ and $\sim \exists x \sim (\sim Q(x) \rightarrow \sim P(x))$ are logically equivalent. (Hint: Start with the double negation rule $\sim(\sim(\forall x (P(x) \rightarrow Q(x))))$)

$\forall x (P(x) \rightarrow Q(x)) \equiv \sim(\sim(\forall x (P(x) \rightarrow Q(x))))$ **by Double Negation**

$\equiv \sim \exists x \sim (P(x) \rightarrow Q(x))$ **by DeMorgan's Law for quantifiers**

$\equiv \sim \exists x \sim (\sim Q(x) \rightarrow \sim P(x))$ **by contrapositive**

Therefore, $\forall x (P(x) \rightarrow Q(x))$ and $\sim \exists x \sim (\sim Q(x) \rightarrow \sim P(x))$ are logically equivalent

5.) True or false: For the set of all negative integers, $\exists x (x + 1 < -x)$

TRUE. Take the case of $x = -6$. $-6 + 1 < -(-6)$

$-5 < 6$

6.) By using direct proof method, show that if k is any odd integer and m is any even integer, then, $k^2 + m^2$ is odd.

Proof:

Suppose k is any odd integer and m is any even integer. [We must show that $k^2 + m^2$ is odd.] By definition of odd, $k = 2a + 1$ where a is some integer. By definition of even, $m = 2b$ where b is some integer. By substitution and algebra,

$$\begin{aligned} k^2 + m^2 &= (2a + 1)^2 + (2b)^2 \\ &= 4a^2 + 4a + 1 + 4b^2 \\ &= 2(2a^2 + 2a + 2b^2) + 1 \end{aligned}$$

Assume $t = (2a^2 + 2a + 2b^2)$, which makes t an integer because it is the sum of the product of integers. Thus, $k^2 + m^2 = 2t + 1$, and by definition of odd, $k^2 + m^2$ is odd [what we needed to show.]

7.) Prove by contraposition that for all integers m and n , if $m + n$ is even then m and n are both even or m and n are both odd.

If $m + n$ is even, then $((m \text{ and } n \text{ are both even}) \text{ OR } (m \text{ and } n \text{ are both odd}))$

Contraposition: If $\sim(m \text{ and } n \text{ are both odd}) \text{ AND } \sim(m \text{ and } n \text{ are both even})$, then $m + n$ is odd.

Proof (by contraposition):

Suppose m and n are not both odd nor both even. [We must show that $m + n$ is odd.] There are two cases when m and n are not both odd nor both even. One case is when m is even and n is odd, and the other case is when m is odd and n is even.

Case 1 (m is odd, n is even):

Suppose m is odd and n is even. [We must show that $m + n$ is odd.] By definition of odd, $m = 2j + 1$ where j is some integer. By definition of even, $n = 2k$ where k is some integer. By substitution and algebra,

$$\begin{aligned} m + n &= 2j + 1 + 2k \\ &= 2(j + k) + 1 \end{aligned}$$

Assume $t = j + k$, which makes t an integer because it is the sum of integers. Thus, $m + n = 2t + 1$, and by definition of odd, $m + n$ is odd [as we needed to show].

Case 2 (m is even, n is odd):

Suppose m is even, and n is odd. [We must show that $m + n$ is odd.] By definition of even, $m = 2a$ where a is some integer. By definition of odd, $n = 2b + 1$ where b is some integer. By substitution and algebra,

$$\begin{aligned} m + n &= 2a + 2b + 1 \\ &= 2(a + b) + 1 \end{aligned}$$

Assume $q = a + b$, which makes q an integer because it is the sum of integers.

Thus, $m + n = 2q + 1$, and by definition of odd, $m + n$ is odd [as we needed to show].

In both cases of m and n not being both odd nor both even, we found that $m + n$ is odd [as we needed to show].