

Passive SSH Key Compromise via Lattices

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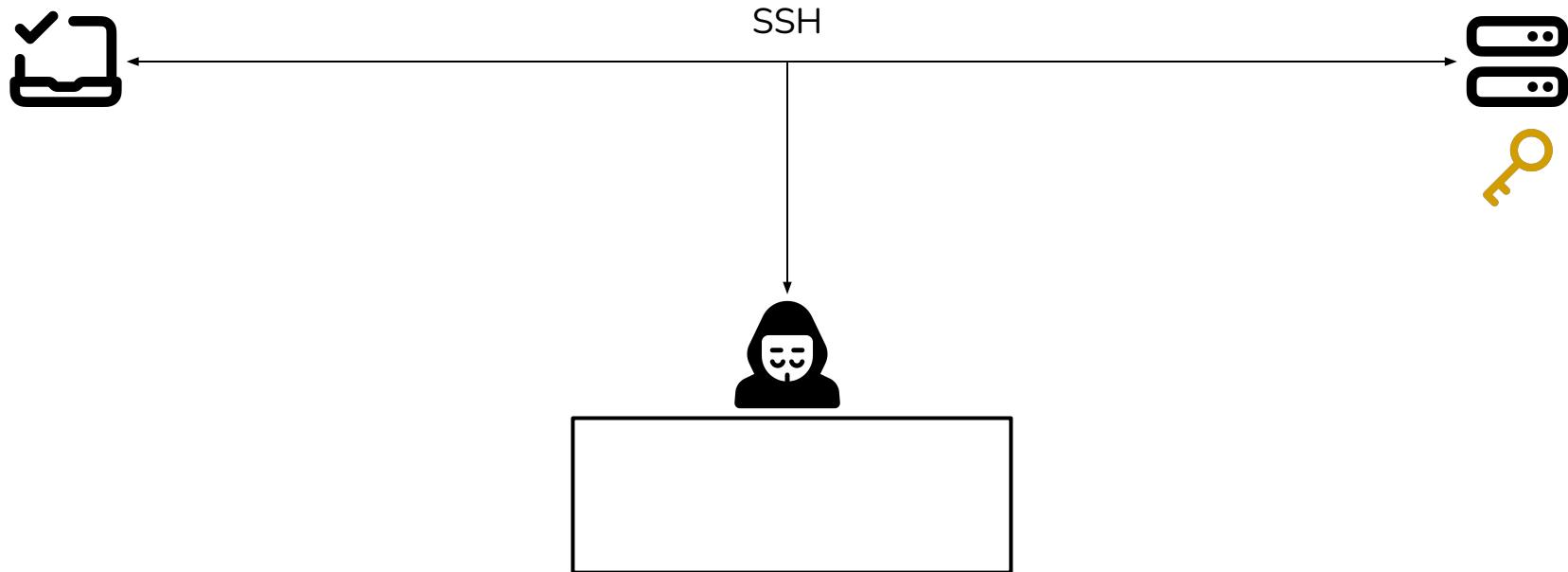
ACM CCS, November 2023

Overview

1. Can faulty RSA signatures reveal the private SSH host key to eavesdroppers?

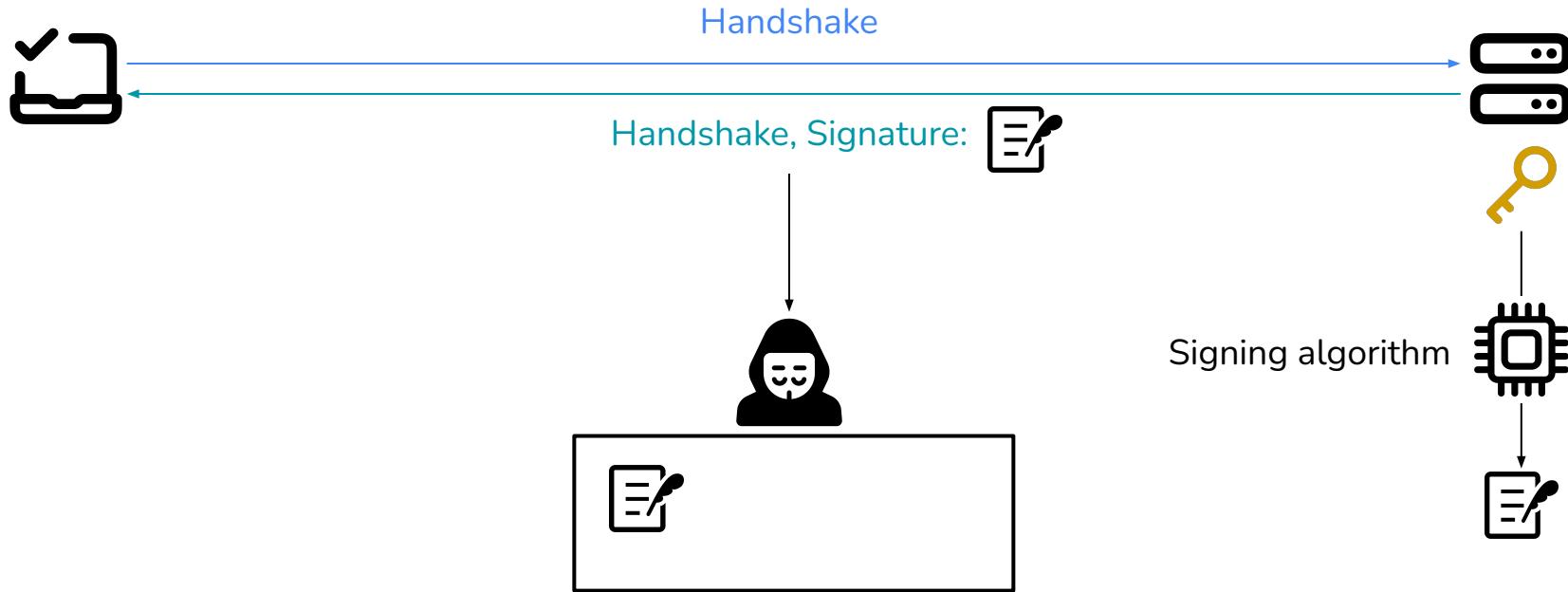
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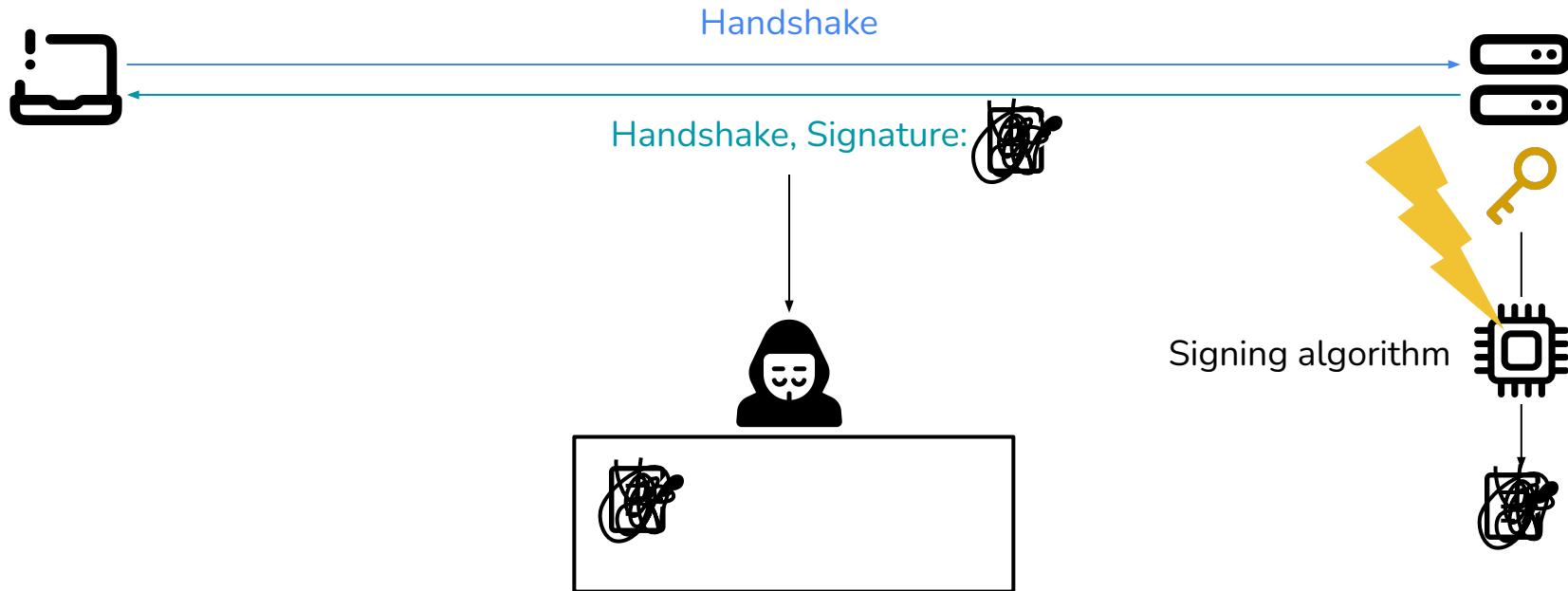
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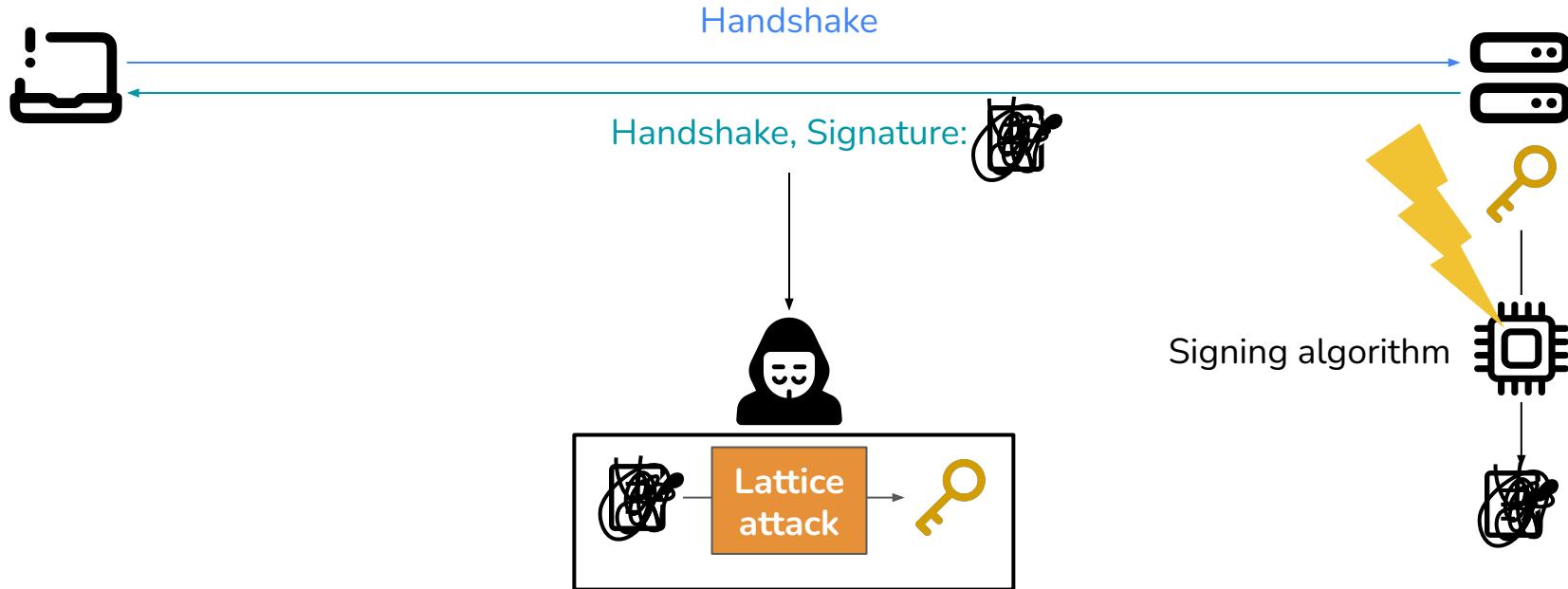
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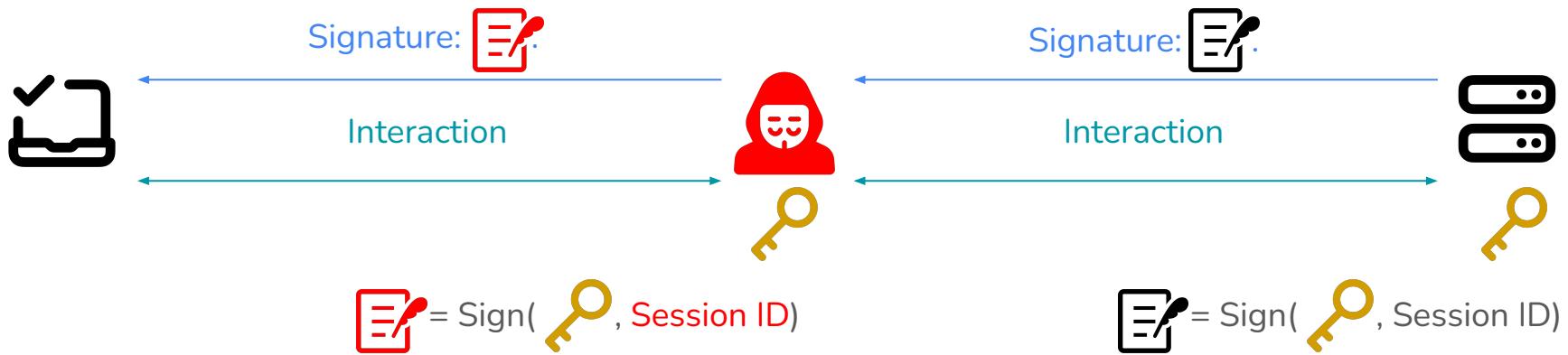
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Implications of host key compromise

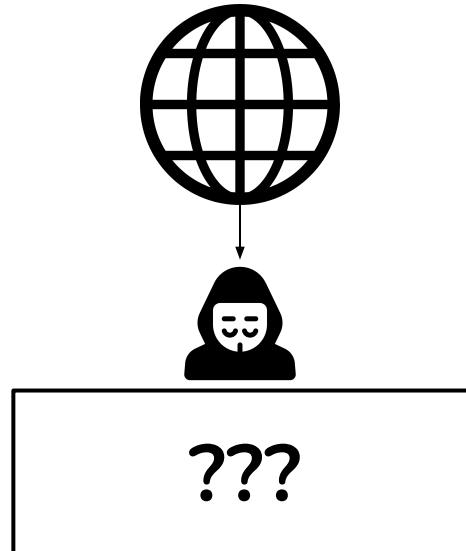
Active attacker with the key can impersonate the server

- Can then intercept and arbitrarily tamper with connection



Overview

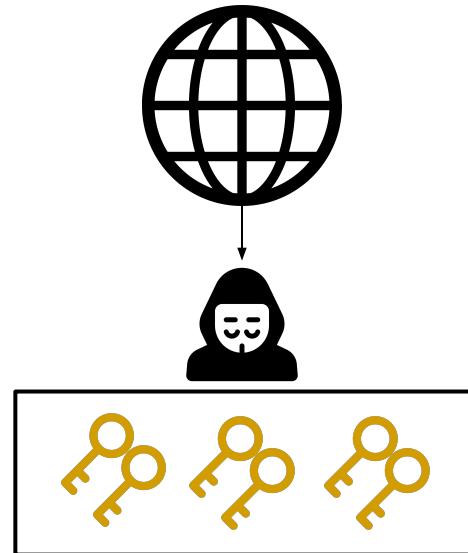
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2. Are such keys leaked frequently over the internet?



Overview

1. Can faulty RSA signatures reveal the private SSH host key to eavesdroppers?
2. Are such keys leaked frequently over the internet?

We analyzed **3.2 billion** SSH
signatures on the internet
(1.2 billion *RSA signatures*)



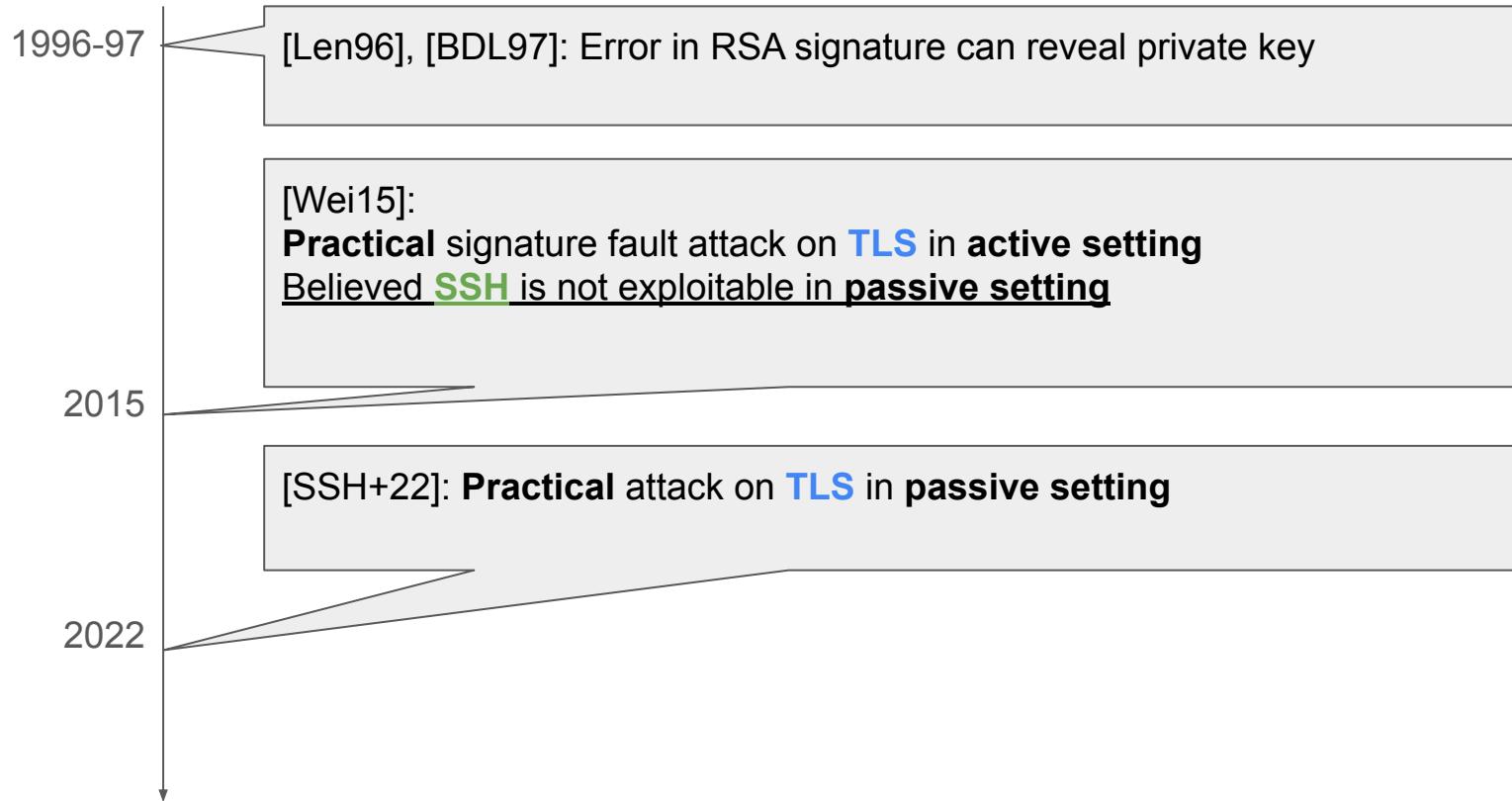
4,962 signatures:
Revealed **private keys**

History of signature fault attacks

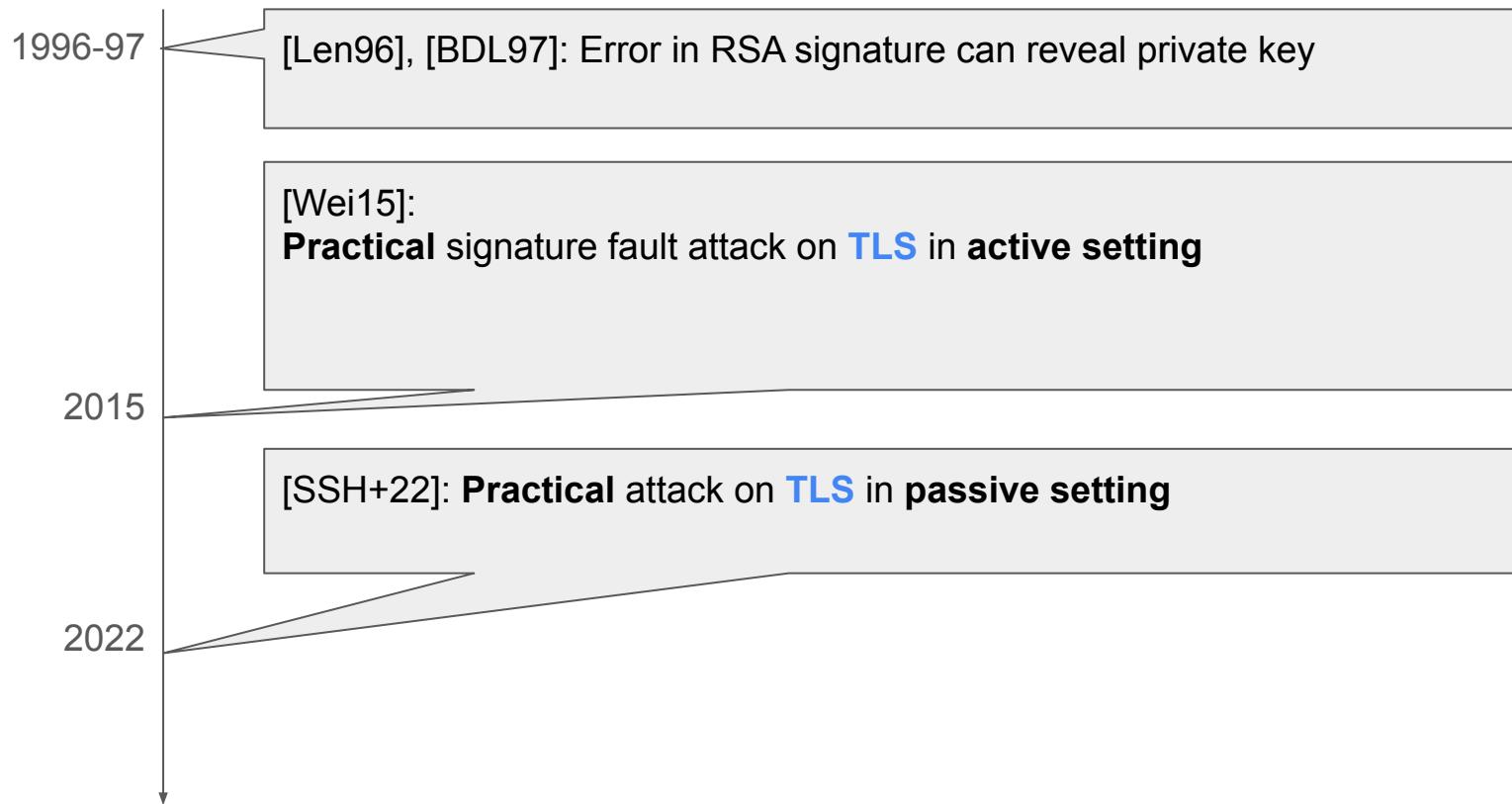
1996-97

[Len96], [BDL97]: Error in RSA signature can reveal private key

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(only if signed message is fully known)

[Wei15]:

Practical signature fault attack on **TLS** in **active setting**

Believed **SSH** is not exploitable in **passive setting**

Message hash includes **DH shared secret** → unknown to eavesdropper

2015

[SSH+22]: **Practical** attack on **TLS** in **passive setting**

Message **fully computable** by eavesdropper

2022

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Practical signature fault attack on **TLS** in **active setting**

Believed SSH is not exploitable in passive setting

Message hash includes **DH shared secret** → unknown to eavesdropper

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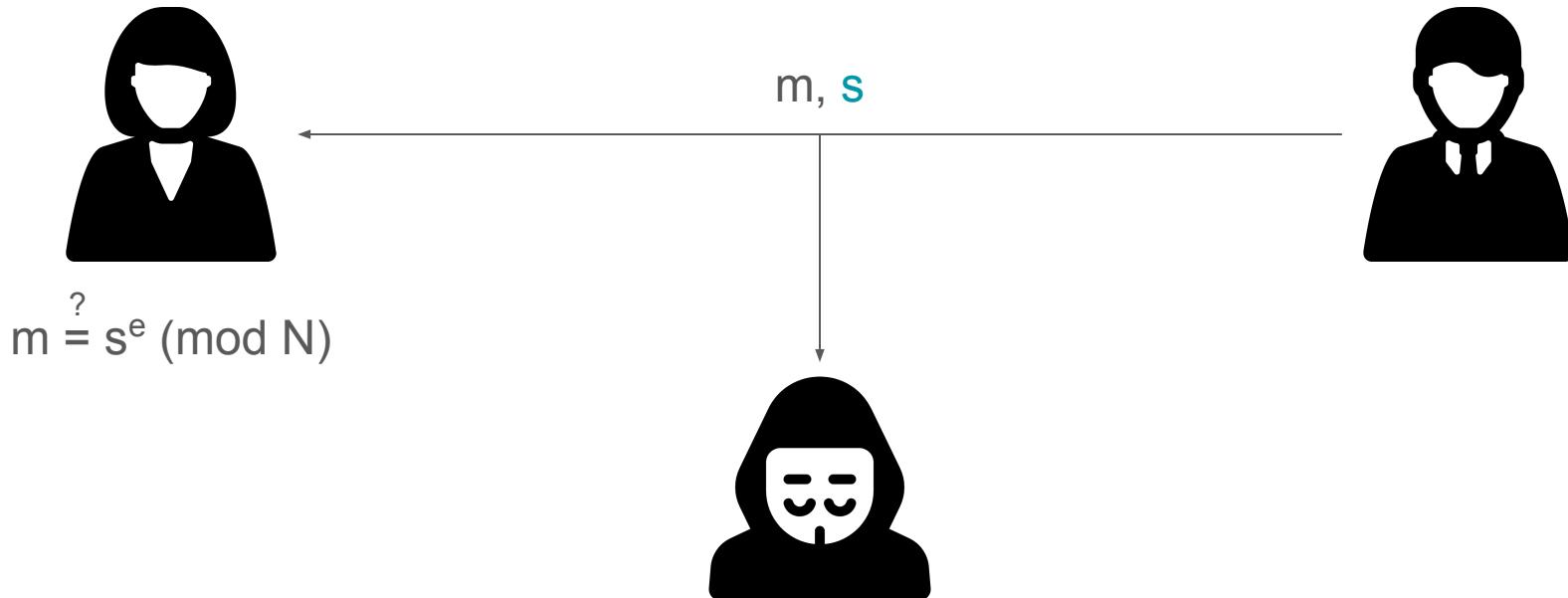
Textbook RSA signing

Public key = (N, e)

Private key = (p, q, d)

$N = pq, d = e^{-1} \pmod{\phi(N)}$

$s = m^d \pmod{N}$

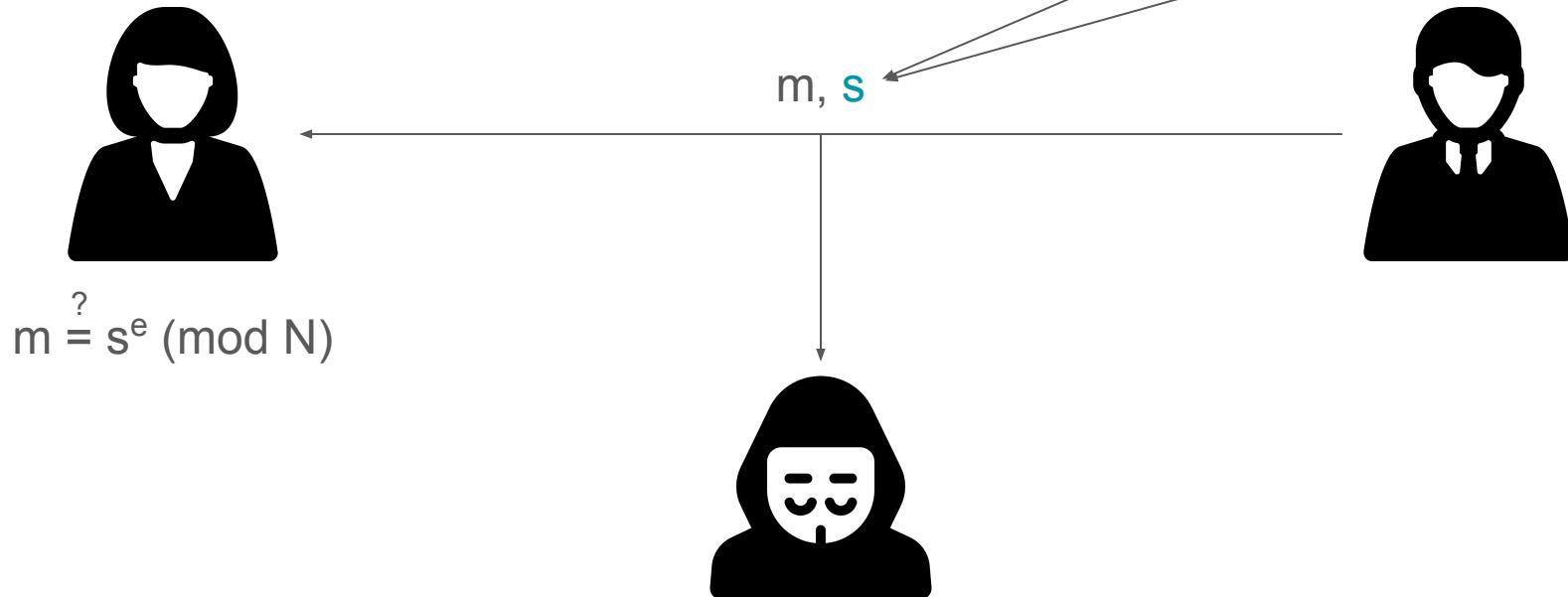


Textbook RSA signing with CRT optimization

Private key = (p, q, d_p, d_q)
 $N = pq, d = e^{-1} \pmod{\phi(N)}$

Public key = (N, e)

$$s_p = m^{d_p} \pmod{p}$$
$$s_q = m^{d_q} \pmod{q}$$



RSA-CRT signing, with a fault

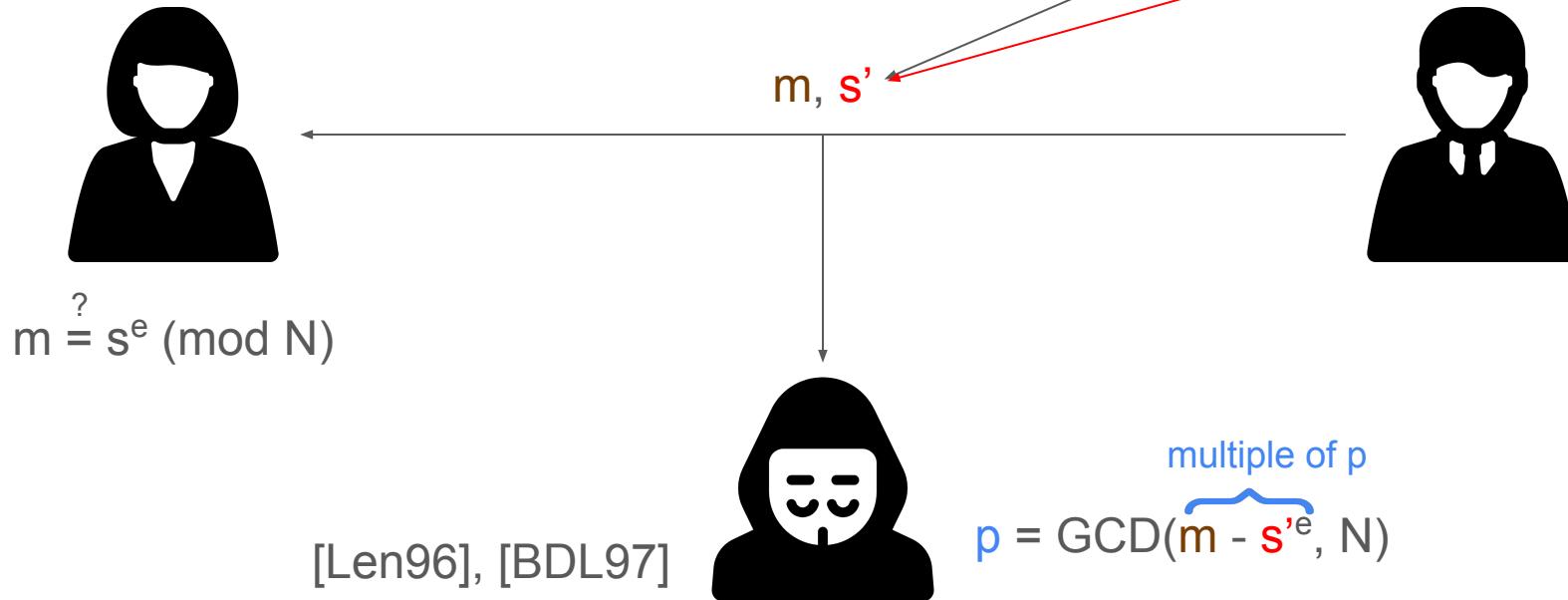
Public key = (N, e)

Private key = (p, q, d_p, d_q)

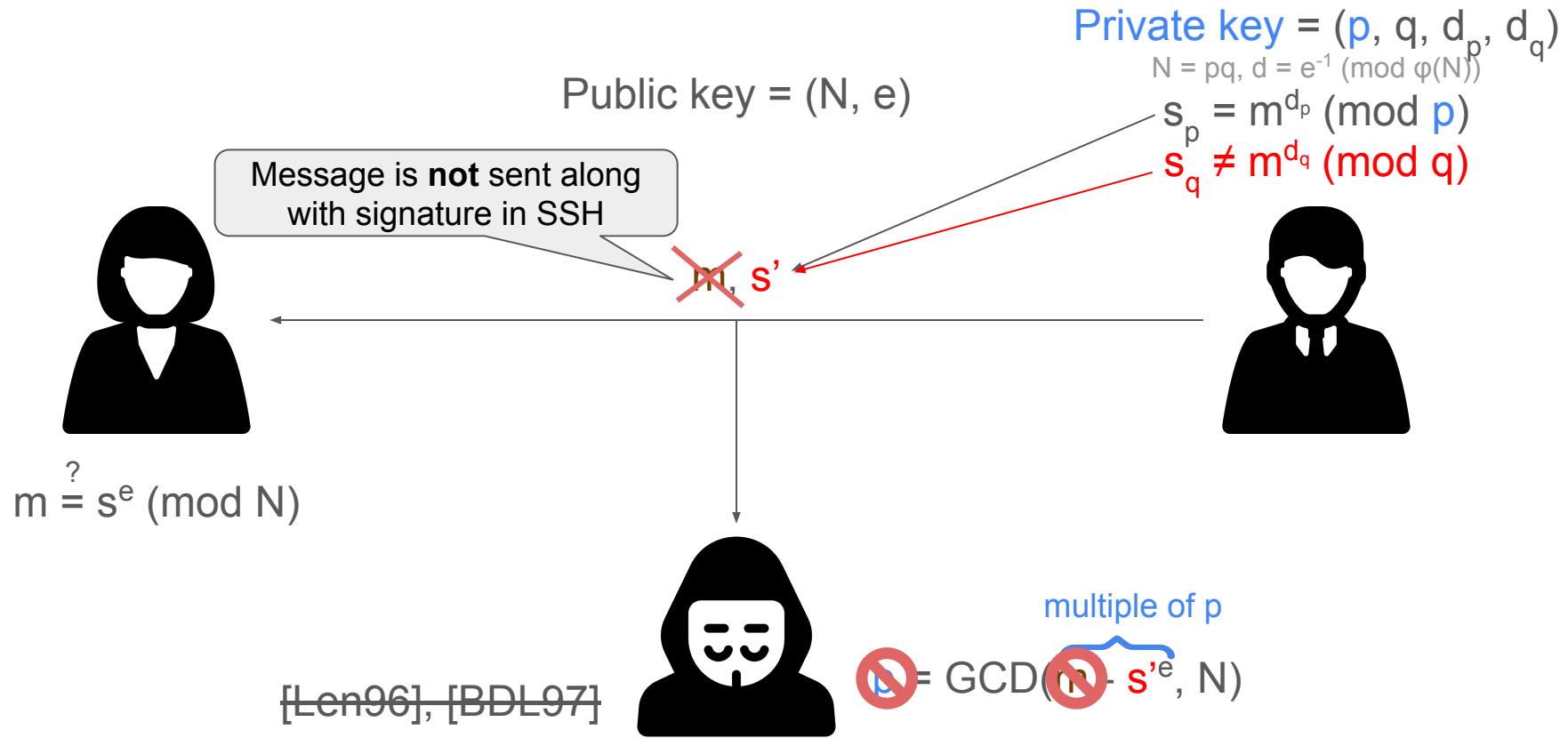
$$N = pq, d = e^{-1} \pmod{\phi(N)}$$

$$s_p = m^{d_p} \pmod{p}$$

$$s_q \neq m^{d_q} \pmod{q}$$



RSA-CRT fault in SSH



PKCS#1 v1.5 RSA signature padding

Message format for SSH:

$m = 0x0001FF\dots FF00 \{hash\} \{hash\}$

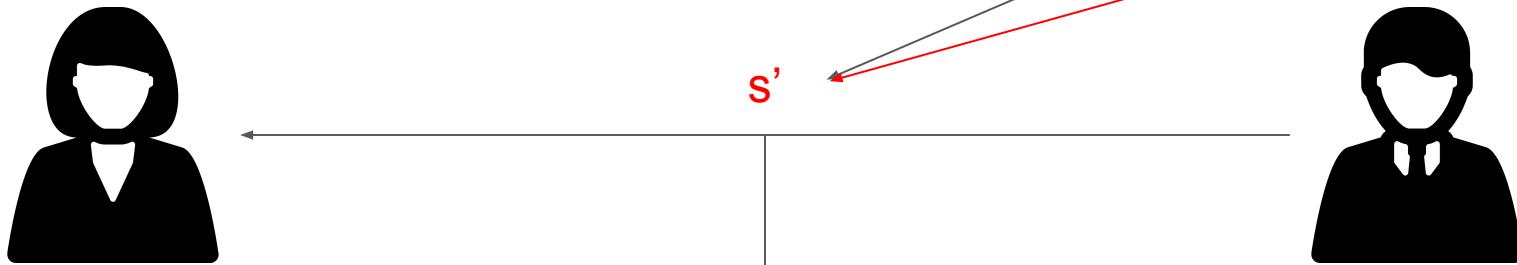
- **Padding** is deterministic and known to the attacker
- Attacker knows almost all of message **m** , except for the **hash**

RSA-CRT fault in SSH, passive attack

Public key = (N, e)

Private key = (p, q, d_p, d_q)
 $N = pq, d = e^{-1} \pmod{\phi(N)}$

$$s_p = m^{d_p} \pmod{p}$$
$$s_q \neq m^{d_q} \pmod{q}$$



Message format for SSH: (PKCS#1 v1.5)
 $m = 0x0001FF\dotsFF00 \{algorithm\} \{hash\}$

[CJK+09]

known small unknown

Knows that $m = \overbrace{a}^{\text{known}} + \overbrace{r}^{\text{small unknown}}$

$p = \text{Approx-GCD}(a - s'^e, N)$

Lattice attack

Message format for SSH: (PKCS#1 v1.5)

$m = 0x0001FF\dots FF00 \{hash\} \{hash\} < N$

Experimented with attacking faulty SSH signatures:

- Generated instances that correspond to SSH parameter choices
 - e.g., RSA-2048, SHA-256

For all common parameter choices (RSA modulus length > 4 × hash length):

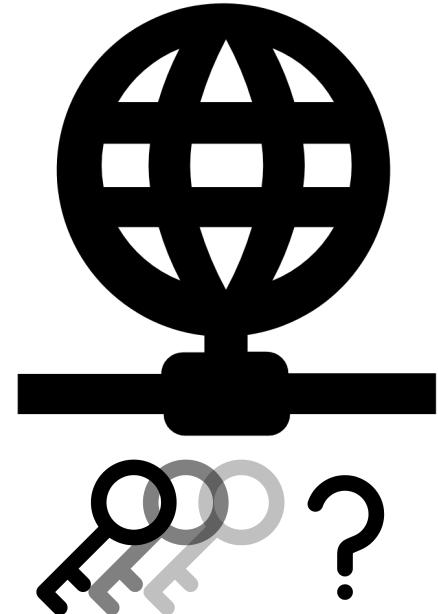
- Attack averages under 0.2 seconds per signature
- Recovers correct key for every generated signature

Key leakage in the wild

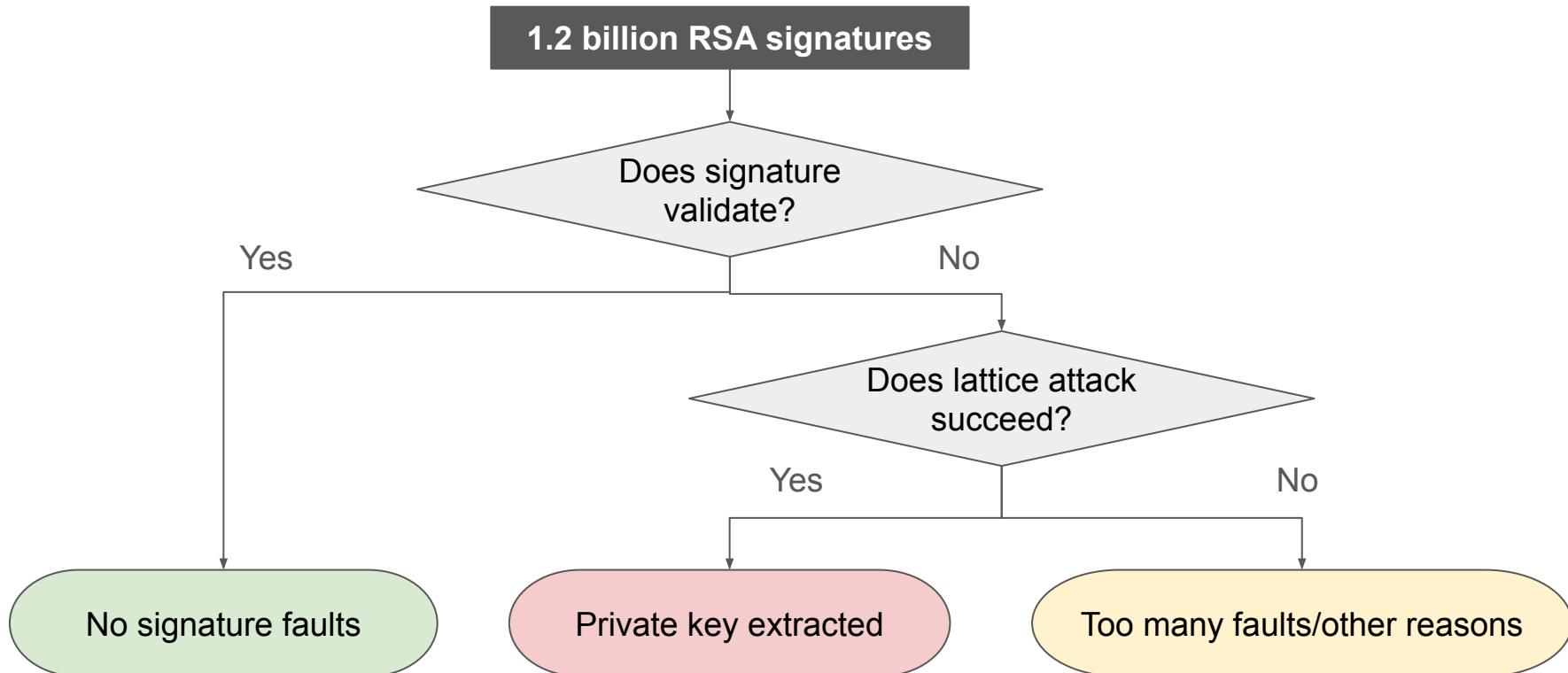
How frequently do keys leak on the internet?

We analyzed data from:

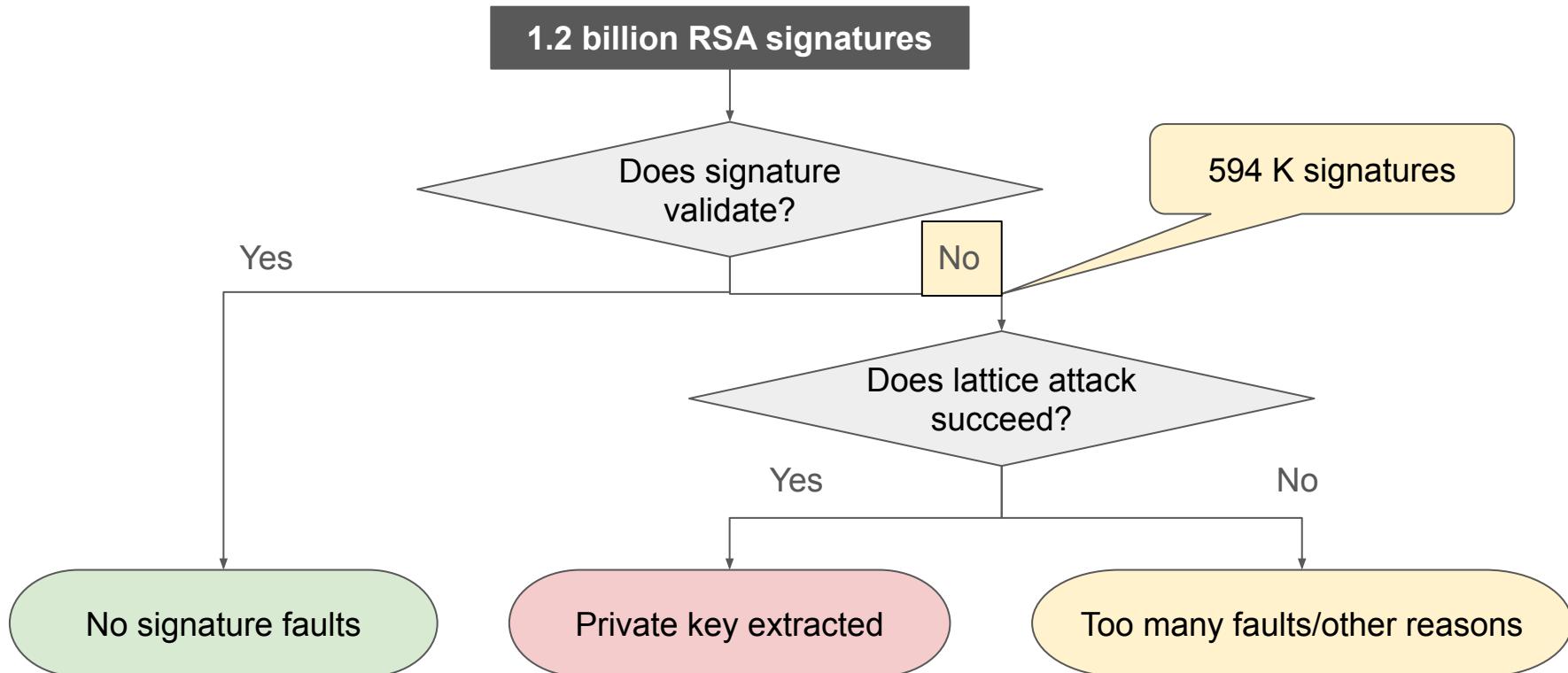
- Zmap scans performed at UCSD
- Passive network traffic through UCSD
- Censys
- University of Michigan



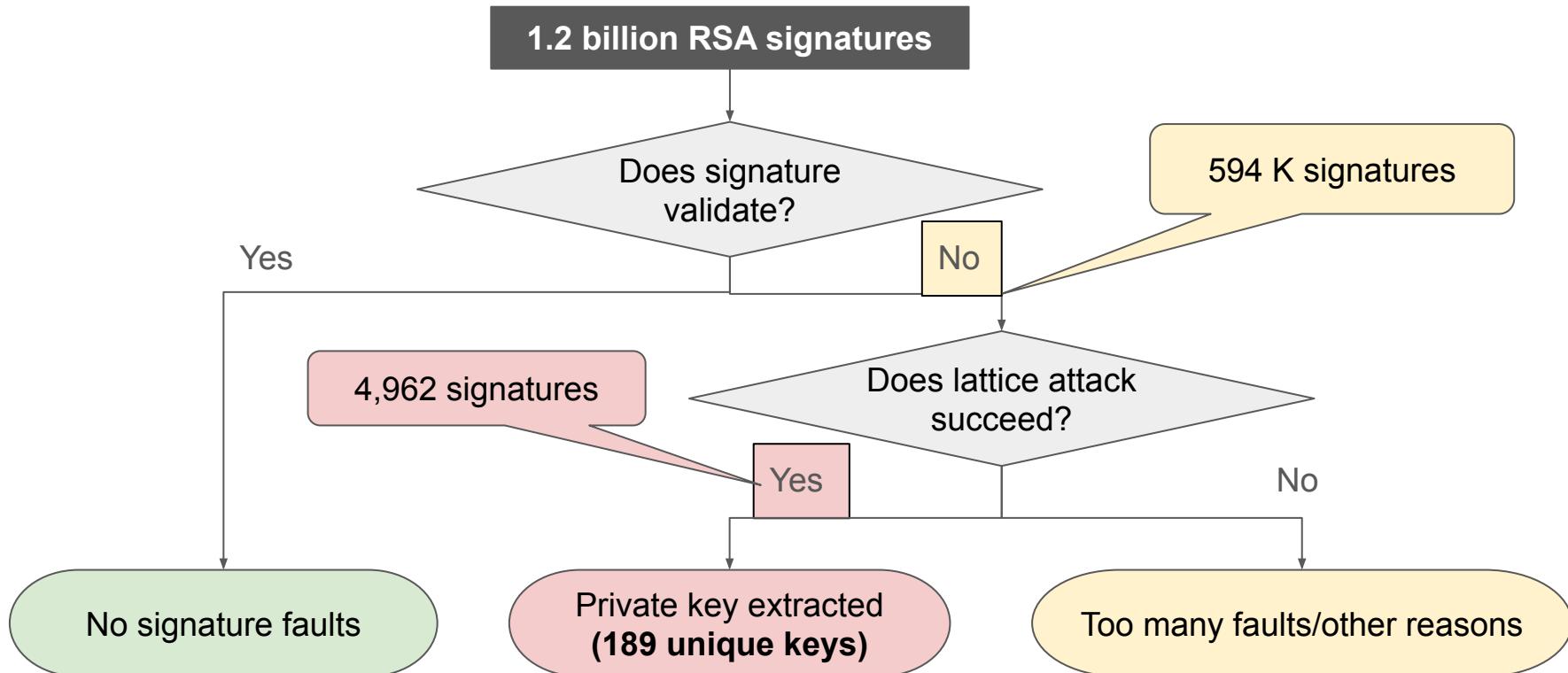
Analysis



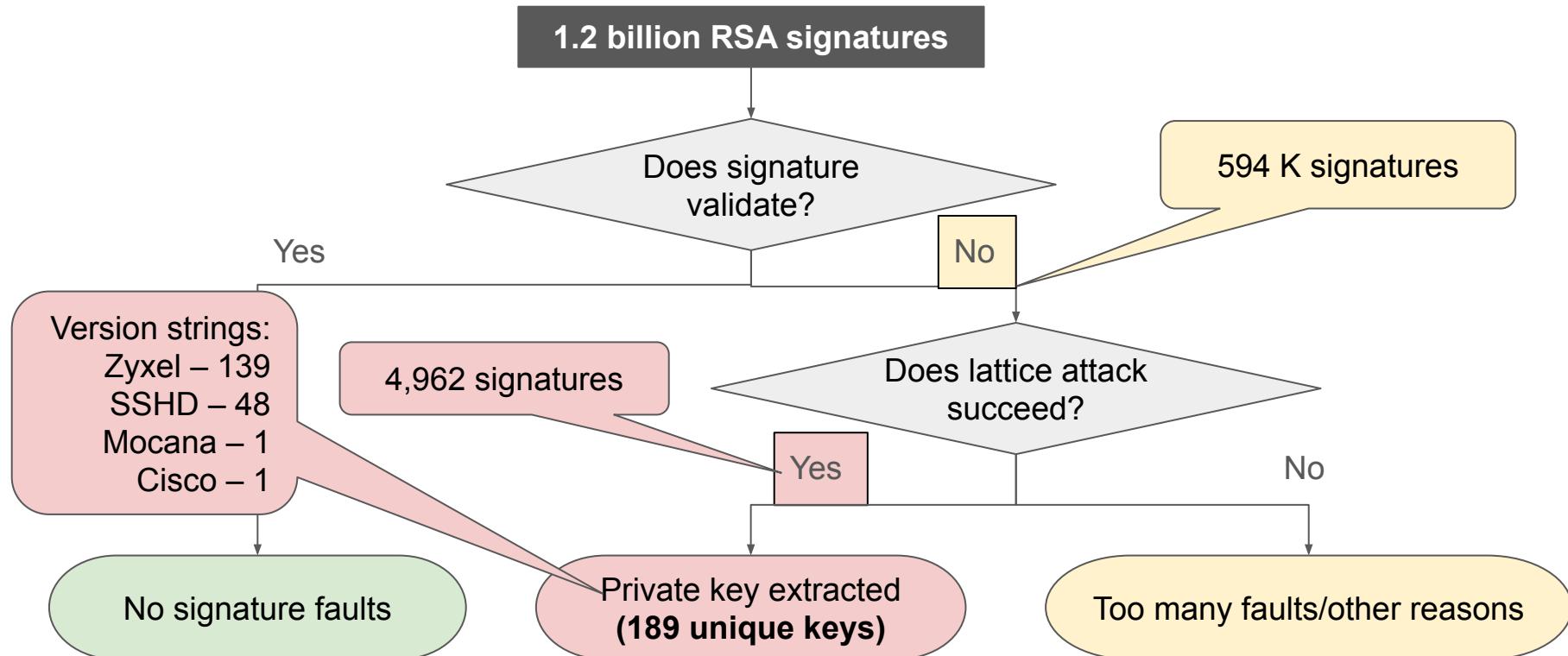
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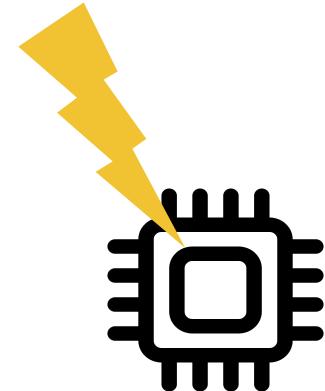
Lessons

Hypothesis: signature faults originate from hardware failures

Include random hardware faults in the threat model

Countermeasure: Validate signatures before sending

- OpenSSL and derivatives include this check
- Vulnerable implementations we observed do not have this countermeasure



Future directions

- SSH: Collect rsa-sha2-* signatures
 - Potentially more vulnerable hosts are out there
- Collect more data passively
- Study similar key leaks on IPsec
 - Our visibility into IPsec hosts is limited

Summary

A single faulty signature reveals the private SSH host key to eavesdroppers

- Private host key allows attacker to later impersonate server
- About 1 out of 1 million analyzed signatures on the internet are vulnerable

We disclosed the vulnerability to device vendors:

- Mitigations confirmed for Cisco, Zyxel, and Hillstone Networks
- Unable to contact Mocana

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<https://ia.cr/2023/1711>

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References

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