

# Random Access Problem in Machine-to-Machine

- Source Code: <https://goo.gl/9zUwYw>
- **Before starting to read this article, please install chrome extension: Github with MathJax , to ensure the correctness of formula format.**
- Other requirements: `gnuplot` , `astyle`
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## Talking About Network Modeling

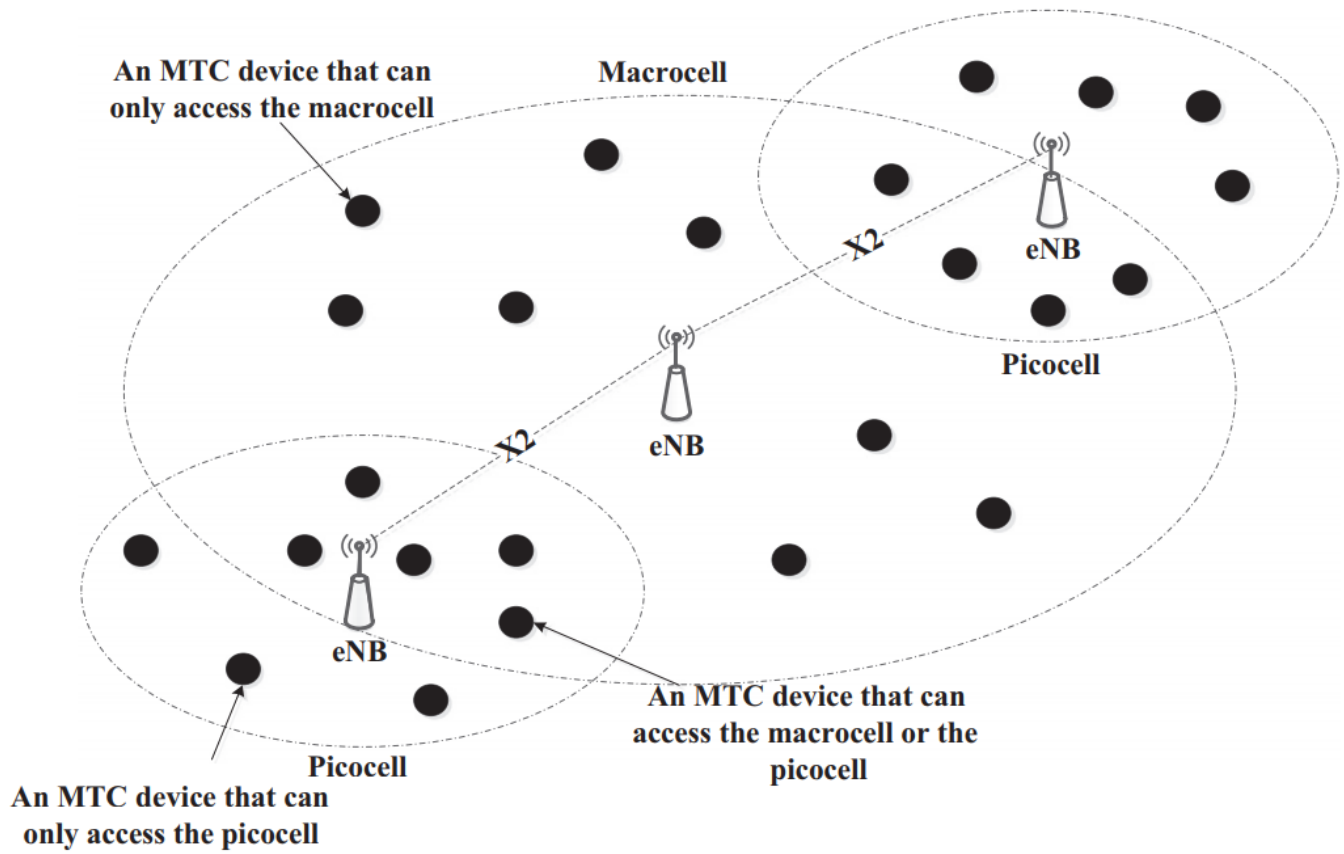
A basic way to build network model:

1. Technique background (From paper ,technical essay, magazine ... etc)
2. System model (Mathematic & Question-Defined)
3. Simulation (Base on Step 1,2 to construct program )
4. Analysis (Measure the result from program)

Base on this four method, we can easily define our question and find an efficient way to locate and find the answer.

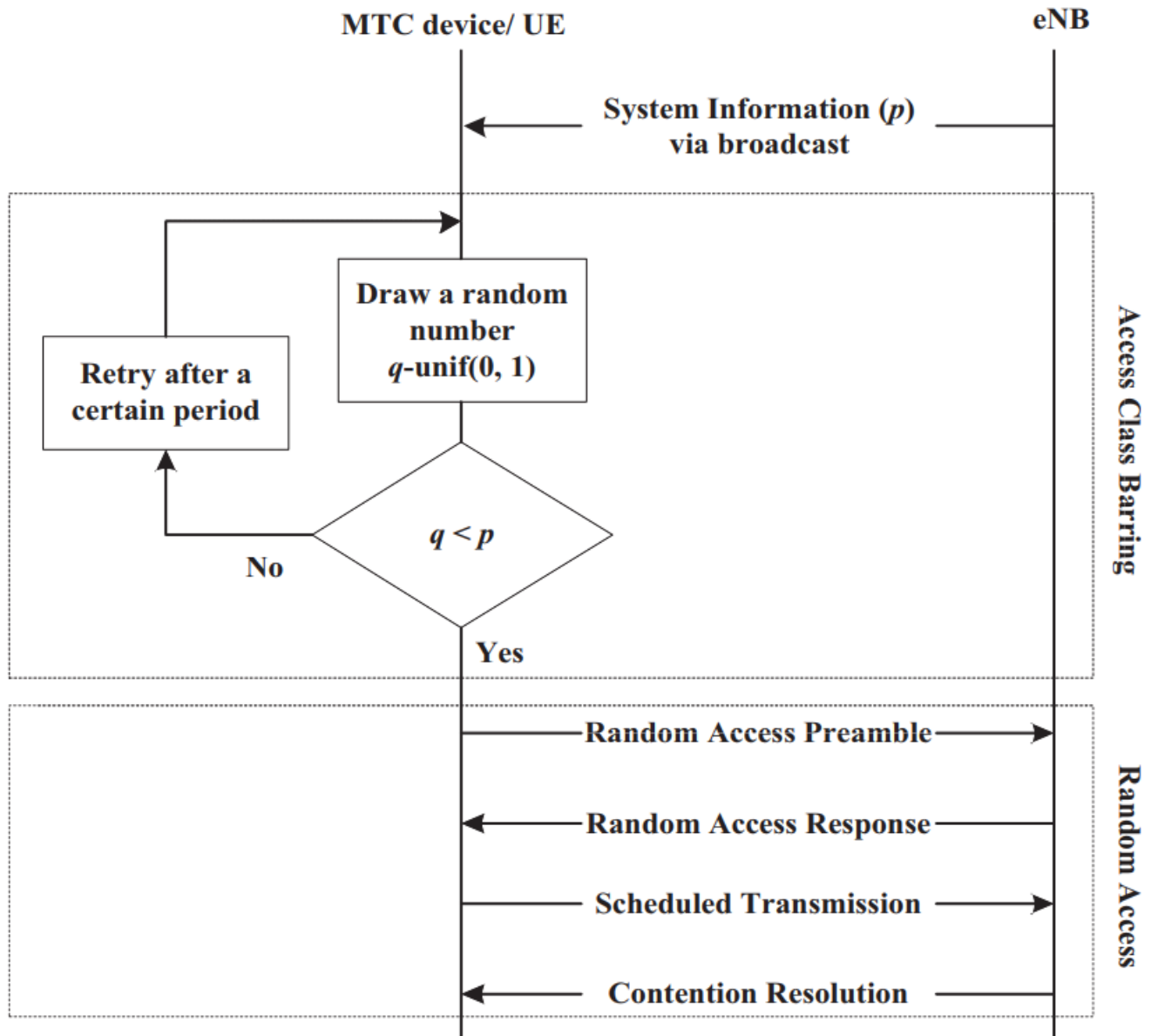
## Our Environment: LTE-A cellular network

- LTE-A cellular network [\[1\]](#)



**Fig. 1.** Architecture of an LTE-A cellular network [8].

- Access class barring procedure for an MTC device or UE [1]



**Fig. 2.** Access class barring procedure for an MTC device or UE [10].

## Part A: Access Class Barring

### Technique Background

- Here we can refer the **Fig 2.** to see the Access Class Barring part.
- We need to create a simulation and mathematic model to represent this part.
- Assumption was made:
  - Draw a random number  $q \sim \text{uniform}(0, 1)$  before performing the random access procedure.
  - Assume that devices are blocked for 1 ms.
  - Calculate the access delay by simulation and mathematical analysis.

### System Model

Here we construct our mathematic model from technique background:

- Before calculate access delay, consider:
  - How many times do we need to complete this job? (All devices pass the access class barring)
  - Find the worst case as upperbound.
  - Then we can summation the result from best case to worse case.
- Our upperbound is generated after running the simulation model( will mention later ), we use the worse case in simulation process as our worse case in mathematic model, so that the result can plot together in well organization.

## Simulation

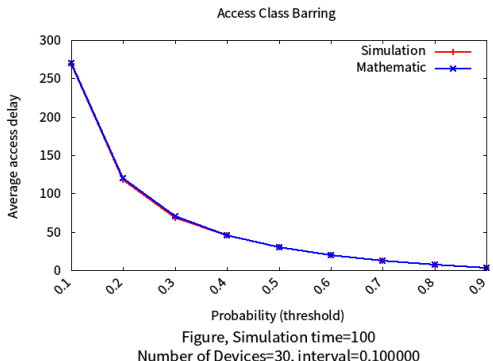
In simulation part, I choose probability "p"(threshold) as observed value; Adjust this value to see the variation among different p value.

In each p case, we run s times of simulation routine. Each routine will finish when all devices have passed the access class barring; and each routine will record its delay times at the end of routine.

After running s simulation routines, we can get an average access delay! We take this value as **access delay** of simulation.

## Analysis

Run with command `make test_a` to generate the test data of Part A.( `make test` will run all the testcase, include A and B.)

case	simulation times	Number of device	P	result
1	100	30	0.1~0.9	

case	simulation times			result
2	100	40	0.1~0.9	<p>Figure, Simulation time=100 Number of Devices=40, interval=0.100000</p>
3	10000	20	0.1~0.9	<p>Figure, Simulation time=10000 Number of Devices=20, interval=0.100000</p>
4	100000	20	0.1~0.9	<p>Figure, Simulation time=100000 Number of Devices=20, interval=0.100000</p>

As the result shown above, we can see curves of **simulation** and **mathematic** are matching with each other.

## Part B: Random Access

### Technique Background

- Here we can refer the **Fig 2.** to see the Random Access part.
- We need to create a simulation and mathematic model to represent this part.
- Assumption was made:
  - Simulate a contention-based random access procedure.

- Collision occurs when devices are trying to transmit the same preamble.
- There is no retransmission in your simulation.
- Compute the success probability after the random access procedure by simulation and mathematical analysis.

## System Model

Assume there have `M` devices (e.g. `UE / MTC device` ), and `N` preamble.

We can find the successful nodes from this paper. [\[2\]](#)

$$u(n) = n(1 - \frac{1}{p})^{n-1}$$

*n = number of devices, p = number of preamble*

And we can easily calculate the successful probability:

$$success\ probability = (1 - \frac{1}{p})^{n-1}$$

We take this as our mathematic model.

## Simulation

In simulation part, I choose `number of preamble` as observed value; Adjust this value to see the variation among different `preamble` value with specified `interval` to iterate, and in each `preamble` case, we run `s` times of simulation routine.

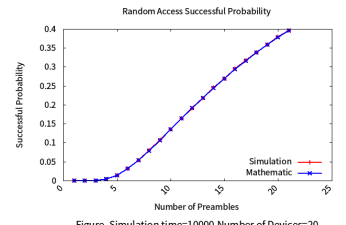
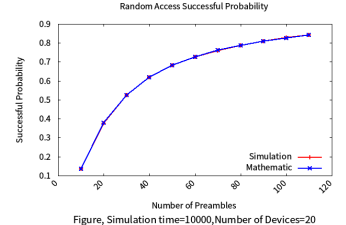
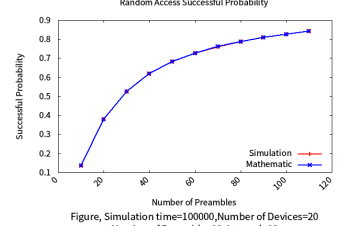
Consider that we don't need to care about "**retransimisson**", so in each simulation routine, we only need to calculate the success devices(which using unique preamble).

And at the end of routine, using success devices to calculate **successful probability**.

## Analysis

Run with command `make test_b` to generate the test data of `Part B`.(`make test` will run all the testcase, include A and B.)

case	simulation times	<i>Number of device</i>	<i>Preamble range</i>	result
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case	simulation times			result
1	10000	20	1 ~ 20	 <p>Figure, Simulation time=10000, Number of Devices=20 Number of Preamble=1, interval=1</p>
2	10000	20	10 ~ 100	 <p>Figure, Simulation time=10000, Number of Devices=20 Number of Preamble=10, interval=10</p>
3	100000	20	10 ~ 100	 <p>Figure, Simulation time=100000, Number of Devices=20 Number of Preamble=10, interval=10</p>

As the result shown above, we can see curves of **simulation** and **mathematic** are matching with each other.

## Reference

[1] Efficient cooperative access class barring with load balancing and traffic adaptive radio resource management for M2M communications over LTE-A. **Yi-Huai Hsu, Kuochen Wang, Yu-Chee Tseng.**

*Department of Computer Science, National Chiao Tung University, Hsinchu 300, Taiwan*

[2] Lower Bounds on the LTE-A Average Random Access Delay Under Massive M2M Arrivals.

**Mehmet Koseoglu.** *IEEE, member*