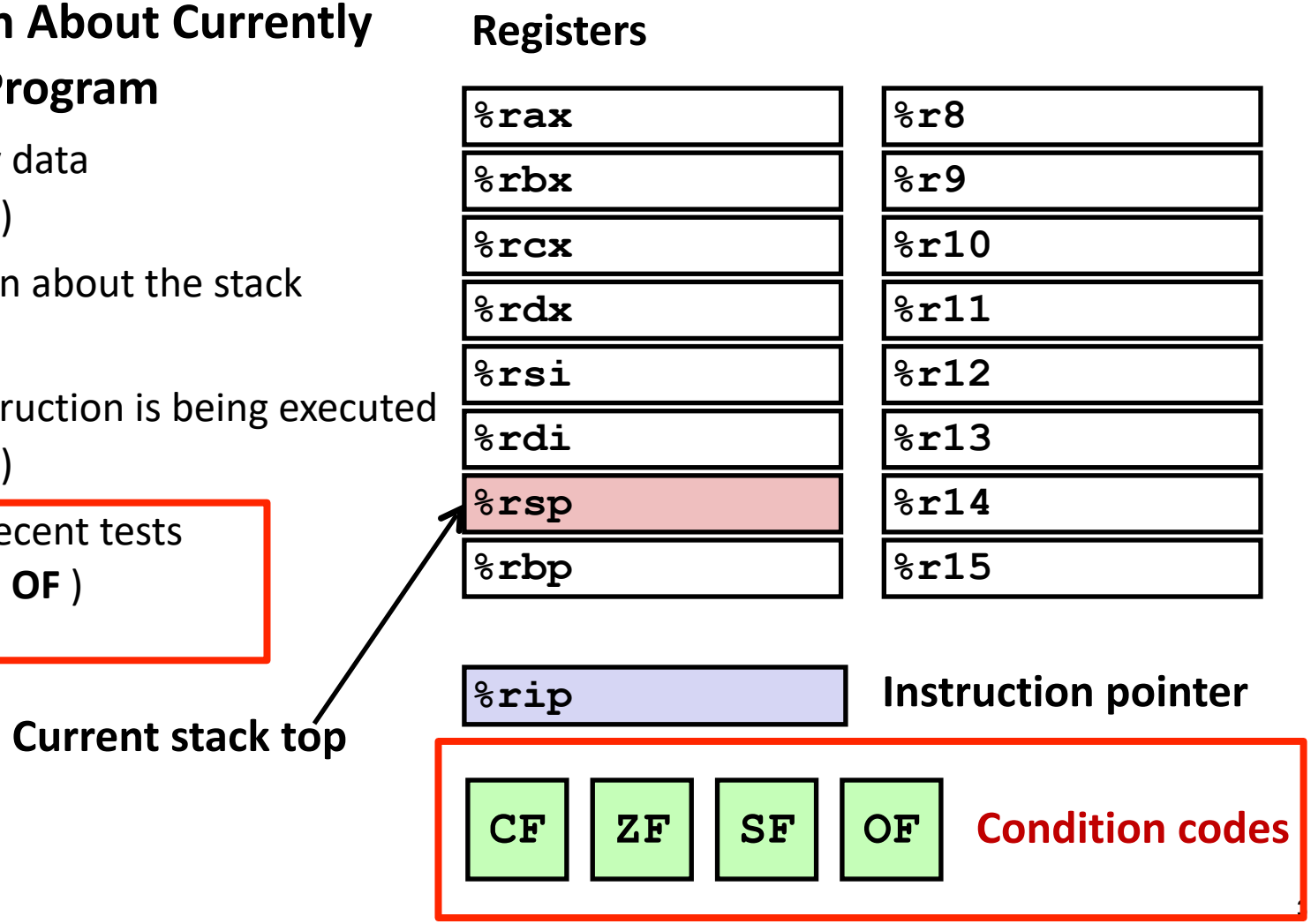
Machine-Level Programming II: Control

Today

- **Control: Condition codes**
- Conditional branches
- Loops
- Switch Statements

Processor State (x86-64, Partial)

- **CPU Needs to Store Information About Currently Executing Program**
  - Temporary data (`%rax`, …)
  - Information about the stack (`%rsp`)
  - Which instruction is being executed (`%rip`, …)
  - Status of recent tests (**CF, ZF, SF, OF**)



Condition Codes (Implicit Setting)

Single bit registers

- **SF** Sign Flag (for signed)
- **ZF** Zero Flag

CF Carry Flag (for unsigned)

OF Overflow Flag (for signed)

Implicitly set (think of it as side effect) by arithmetic operations

- Example: **`addq`** *Src*, *Dest* ↔ **`t = a+b`**
- **ZF set** if **`t == 0`**
- **SF set** if **`t < 0`** (as signed)
- **CF set** if carry out from most significant bit (unsigned overflow)
- **OF set** if two's-complement (signed) overflow

```
(a>0 && b>0 && t<0) || (a<0 && b<0 && t>=0)
```

Not set by `leaq` instruction

Condition Codes (Explicit Setting: Compare)

- **Arithmetic Operations Set Condition Codes Implicitly**
- **Explicit Setting by Compare Instruction**
  - `cmpq` *Src2*, *Src1*
  - `cmpq b,a` like computing `a-b` without setting destination

  - **ZF set** if `(a-b) == 0`
  - **SF set** if `(a-b) < 0` (as signed)
  - **CF set** if carry out from most significant bit (used for unsigned comparisons)
  - **OF set** if two's-complement (signed) overflow

```
(a>0 && b<0 && (a-b)<0) || (a<0 && b>0 && (a-b)>0)
```

Condition Codes (Explicit Setting: Test)

- **Explicit Setting by Test instruction**
  - `testq` *Src2, Src1*
    - `testq b,a` like computing `a&b` without setting destination
  - Sets condition codes based on value of ***Src1* & *Src2***
    - **ZF set** when `a&b == 0`
    - **SF set** when `a&b < 0`
  - Why the OF and CF flags are not set?
    - Because there is no overflow in bitwise &
  - Why two operands?
    - Useful to have one of the operands be a mask

Summary

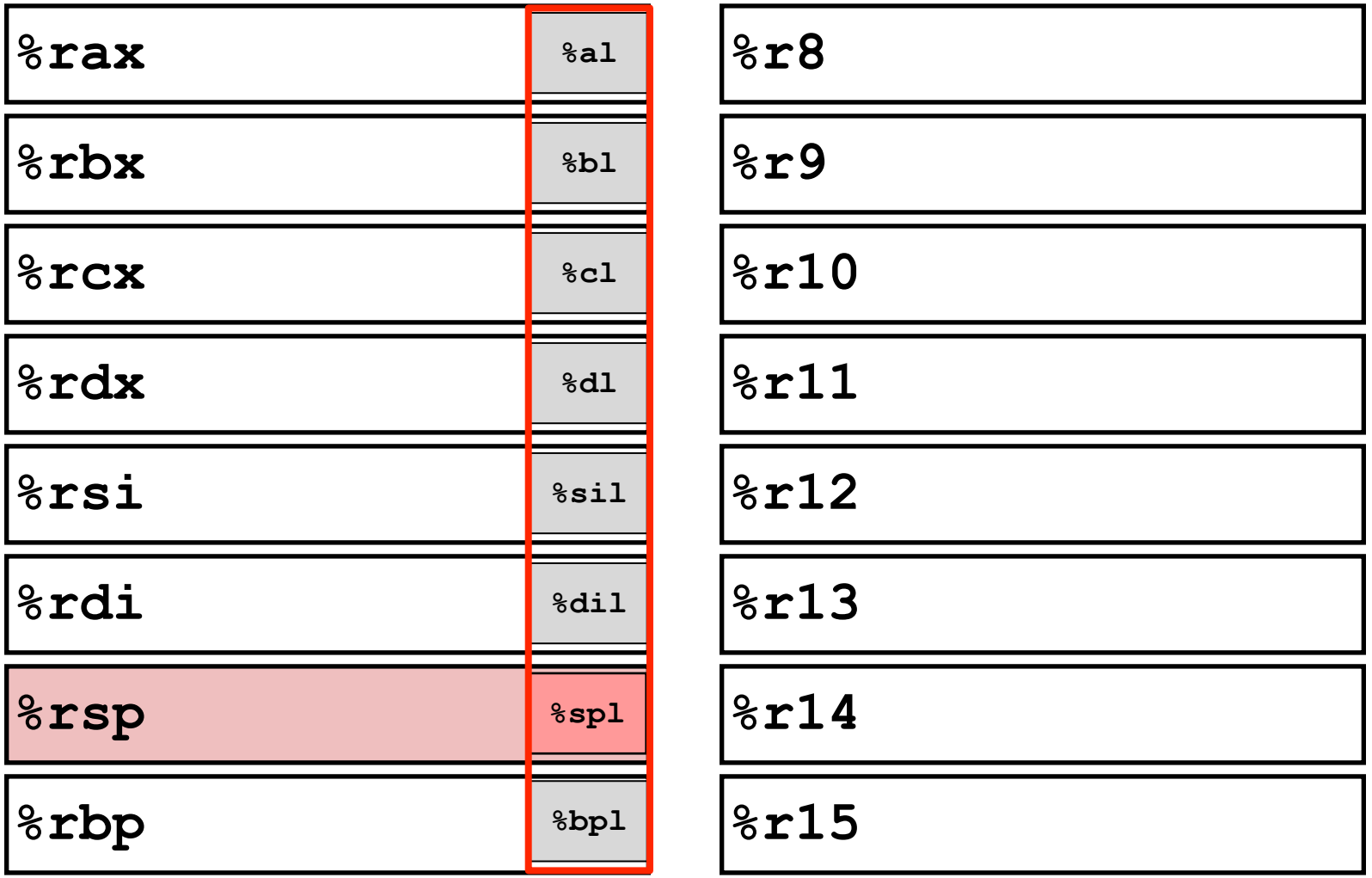
- **How to set condition codes?**
  - **Implicitly**
    - Using arithmetic operations
  - **Explicitly**
    - cmp b,a <=> a - b
    - test b,a <=> b & a

How to Read the Condition Codes?

Reading Condition Codes

SetX Instructions

- Set low-order byte of destination to 0 or 1 based on condition codes

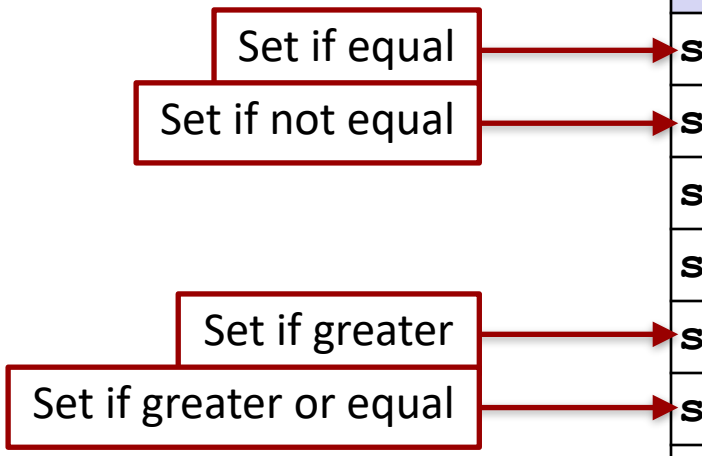


Reading Condition Codes

SetX Instructions

- Does not alter remaining 7 bytes

```
sete    %al     # If ZF == 1, set %al to 1
                # otherwise set it to 0
```



SetX	Description
sete	Equal / Zero
setne	Not Equal / Not Zero
sets	Negative
setns	Nonnegative
setg	Greater (Signed)
setge	Greater or Equal (Signed)
setl	Less (Signed)
setle	Less or Equal (Signed)
seta	Above (unsigned)
setb	Below (unsigned)

Reading Condition Codes

Example

```
int gt(long x, long y)
{
  return x > y;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

```
cmpq    %rsi, %rdi    # (%rdi - %rsi) <=> (x - y)
setg    %al           # Set when > 0
```

SetX Instructions

- Once the lower byte is set, use **movzbq** to set the remaining bits to 0

Practice

```
int eq(long x, long y)
{
  return (x == y);
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

```
    cmpq    %rsi, %rdi    # (%rdi - %rsi) <=> (x - y)
    sete    %al           # Set when = 0
    movzbq  %al, %rax     # Zero rest of %rax
    ret
```

Today

- Control: Condition codes
- **Conditional branches**
- Loops
- Switch Statements

Expressing with Goto Code

- C allows `goto` statement
- Jump to position designated by label

```
long absdiff
  (long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

```
long absdiff_j
  (long x, long y)
{
    long result;
    int ntest = x <= y;
    if (ntest) goto Else;
    result = x-y;
    goto Done;
 Else:
    result = y-x;
 Done:
    return result;
}
```

Jumping

jX Instructions

- Jump to different part of code depending on condition codes

jX	Description
`jmp`	Unconditional
`je`	Equal / Zero
`jne`	Not Equal / Not Zero
`js`	Negative
`jns`	Nonnegative
`jg`	Greater (Signed)
`jge`	Greater or Equal (Signed)
`jl`	Less (Signed)
`jle`	Less or Equal (Signed)
`ja`	Above (unsigned)
`jb`	Below (unsigned)

Conditional Branch Example (Old Style)

- Generation

```
gcc -Og -S -fno-if-conversion control.c
```

```
long absdiff
   (long x, long y)
{
   long result;
   if (x > y)
      result = x-y;
   else
      result = y-x;
   return result;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

```
absdiff:
    cmpq    %rsi, %rdi   # x - y
    jle     .L4
    movq    %rdi, %rax
    subq    %rsi, %rax
    ret
.L4:                     # x <= y
    movq    %rsi, %rax
    subq    %rdi, %rax
    ret
```

Practice

```
long f2
   (long x, long y)
{
   long result;
   if (x < y)
      result = x-y;
   else
      result = y-x;
   return result;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

```
f2:
    cmpq    %rsi, %rdi   # x - y
    jge     .L4
    movq    %rdi, %rax
    subq    %rsi, %rax
    ret
.L4:                     # x >= y
    movq    %rsi, %rax
    subq    %rdi, %rax
    ret
```

General Conditional Expression Translation (Using Branches)

C Code

```
val = Test ? Then_Expr : Else_Expr;
```

```
val = x>y ? x-y : y-x;
```

Goto Version

```
  ntest = !Test;
  if (ntest) goto Else;
  val = Then_Expr;
  goto Done;
Else:
  val = Else_Expr;
Done:
  . . .
```

- Create separate code regions for then & else expressions
- Execute appropriate one

Using Conditional Moves

Conditional Move Instructions

- Instruction supports:

  if (Test) Dest ←Src

- GCC tries to use them
  - But, only when known to be safe

Why use conditional moves?

- Branches are very disruptive to instruction flow through pipelines
- Conditional moves do not require control transfer

C Code

```
val = Test
    ? Then_Expr
    : Else_Expr;
```

Goto Version

```
val = Then_Expr;
eval = Else_Expr;
nt = !Test;
if (nt) val = eval;
```

Practice

C Code

```
val = (x>0) ? 1 : -1;
```

Conditional Move Version

```
val = 1;
eval = -1;
if (x<=0) val = eval;
```

Conditional Move Example

```
long absdiff
   (long x, long y)
{
    long result;
    if (x > y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

```
absdiff:
    movq    %rdi, %rax   # x
    subq    %rsi, %rax   # result = x-y
    movq    %rsi, %rdx
    subq    %rdi, %rdx   # eval = y-x
    cmpq    %rsi, %rdi   # x - y
    cmovle  %rdx, %rax   # if <=, result = eval
    ret
```

Bad Cases for Conditional Move

Expensive Computations

```
val = Test(x) ? Hard1(x) : Hard2(x);
```

- Both values get computed
- Only makes sense when computations are very simple

Risky Computations

```
val = p ? *p : 0;
```

- Both values get computed
- May have undesirable effects

Computations with side effects

```
val = x > 0 ? x*=7 : x+=3;
```

- Both values get computed
- Must be side-effect free

Practice

```
long absdiff
   (long x, long y)
{
    long result;
    if (x == y)
        result = x-y;
    else
        result = y-x;
    return result;
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	Return value

```
absdiff:
    movq    %rdi, %rax   # x
    subq    %rsi, %rax   # result = x-y
    movq    %rsi, %rdx
    subq    %rdi, %rdx   # eval = y-x
    cmpq    %rsi, %rdi   # x - y
    cmovne   %rdx, %rax   # if !=, result = eval
    ret
```

Condition Codes (Implicit Setting)

- **Single bit registers**
  - **SF** Sign Flag (for signed)
  - **ZF** Zero Flag

CF Carry Flag (for unsigned)

OF Overflow Flag (for signed)

- **Set**
  - **Implicitly set**
    - Using arithmetic operations
  - **Explicitly set**
    - cmp
    - test
- **Read**
  - **Copy**
    - set
  - **Take action**
    - jmp
    - cmov