

Machine-Level Programming IV: Data

Today

■ Arrays

- One-dimensional
- Multi-dimensional
 - Flat
 - Multi-level

■ Why?

- Allows you to write correct C code (array access calculation)
- Performance implications
- Opens the door to advanced applications
 - Efficient storage of arrays/data
 - Useful for data exchange (python <-> C)

Today

■ Arrays

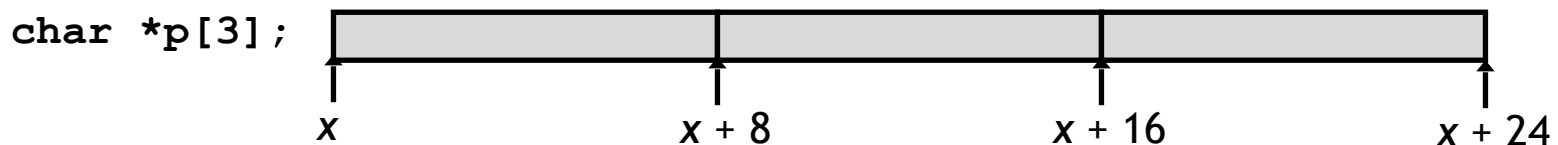
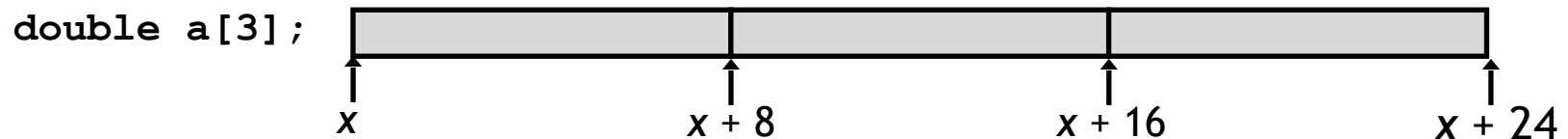
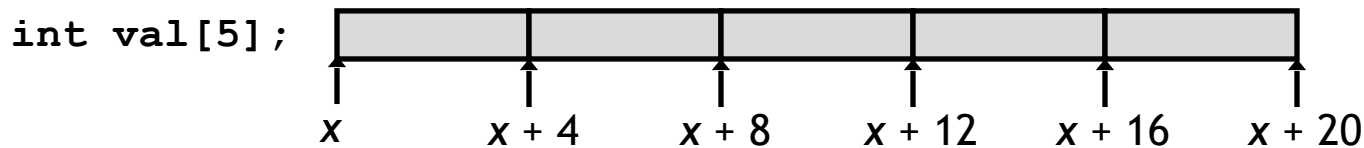
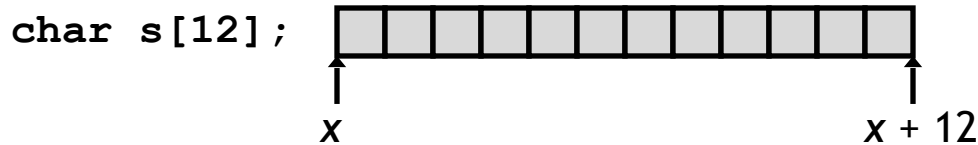
- One-dimensional
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Array Allocation

■ Basic Principle

$T \ A[L];$

- Array of data type T and length L
- Contiguously allocated region of $L * \text{sizeof}(T)$ bytes in memory

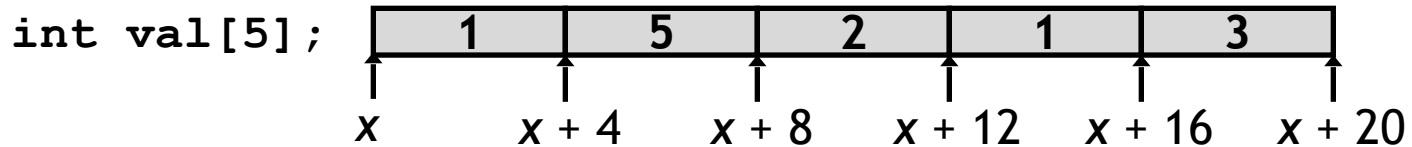


Array Access

Basic Principle

$T \ A[L];$

- Identifier **A** can be used as a pointer to array element 0



Reference Type Value

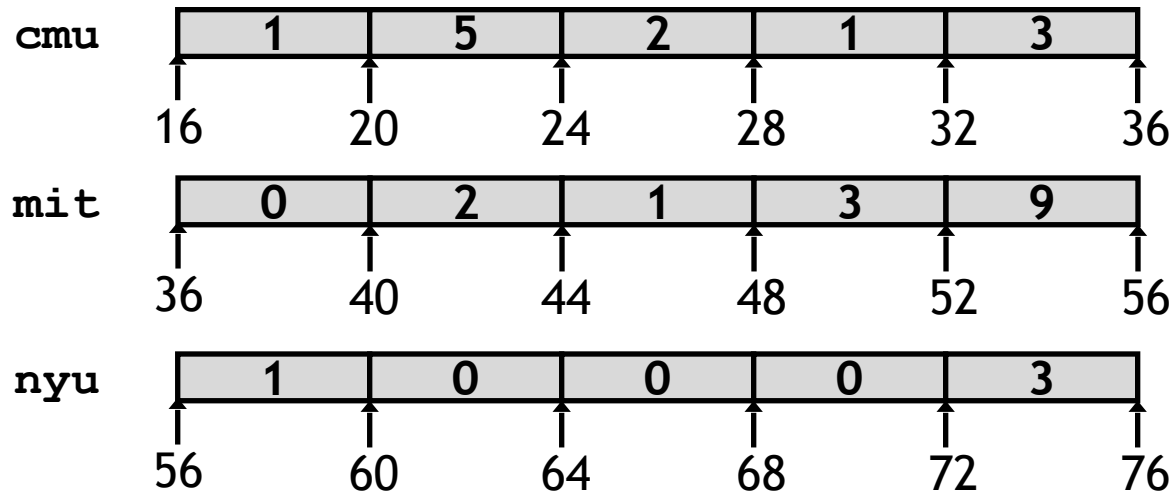
<code>val[4]</code>	<code>int</code>	3
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Array Example

```
#define ZLEN 5

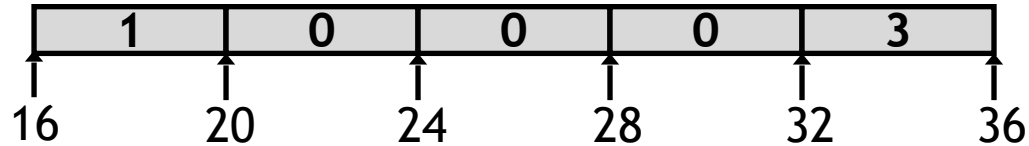
int cmu[ZLEN] = { 1, 5, 2, 1, 3 };
int mit[ZLEN] = { 0, 2, 1, 3, 9 };
int nyu[ZLEN] = { 1, 0, 0, 0, 3 };
```



- Example arrays were allocated in successive 20 byte blocks
 - Not guaranteed to happen in general

Array Accessing Example

```
int nyu[ZLEN];
```



```
int get_digit(int z[ZLEN], int digit)
{
    return z[digit];
}
```

IA32

```
# %rdi = z
# %rsi = digit

movl (%rdi,%rsi,4), %eax # z[digit]
```

- Register `%rdi` contains starting address of array
- Register `%rsi` contains array index
- Desired digit at `%rdi + 4*%rsi`
- Use memory reference `(%rdi,%rsi,4)`

Array Loop Example

```
void zincr(int z[ZLEN]) {  
    size_t i;  
    for (i = 0; i < ZLEN; i++)  
        z[i]++;  
}
```

```
# %rdi = z  
movl    $0, %eax          # i = 0  
jmp     .L3               # goto middle  
.L4:                     # loop:  
addl    $1, (%rdi,%rax,4) # z[i]++  
addq    $1, %rax          # i++  
.L3:                     # middle  
cmpq    $4, %rax          # i:4  
jbe     .L4               # if <=, goto loop  
rep; ret
```


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Multi-dimensional Arrays

■ Declaration

$T \ A[R][C];$

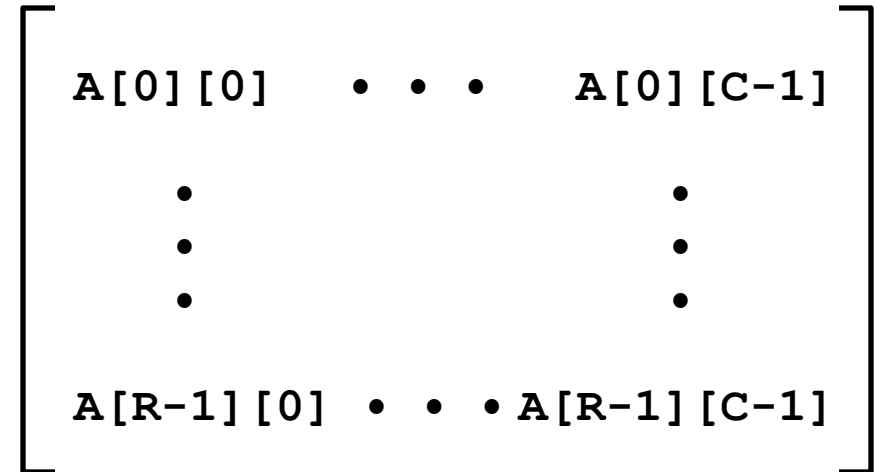
- 2D array of data type T
- R rows, C columns
- Type T element requires K bytes

■ Array Size

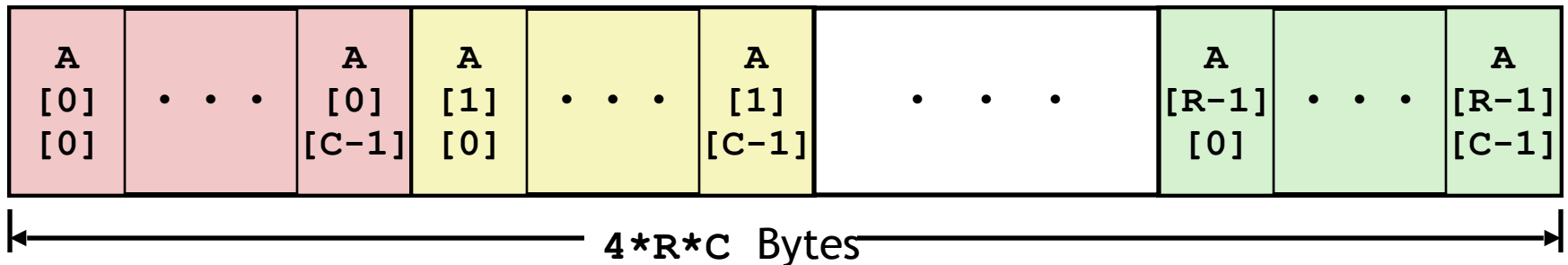
- $R * C * K$ bytes

■ Arrangement

- Row-Major Ordering

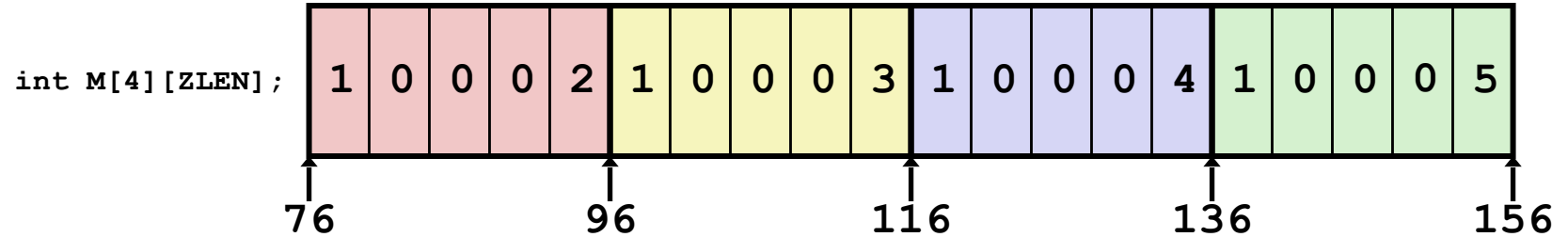


`int A[R][C];`



Multi-dimensional Array Example

```
int M[4][ZLEN] =  
    {{1, 0, 0, 0, 2},  
     {1, 0, 0, 0, 3},  
     {1, 0, 0, 0, 4},  
     {1, 0, 0, 0, 5}};
```



■ `int M[4][ZLEN]`

- Variable **M**: array of 4 elements, allocated contiguously
- Each element is an array of 5 `int`'s, allocated contiguously

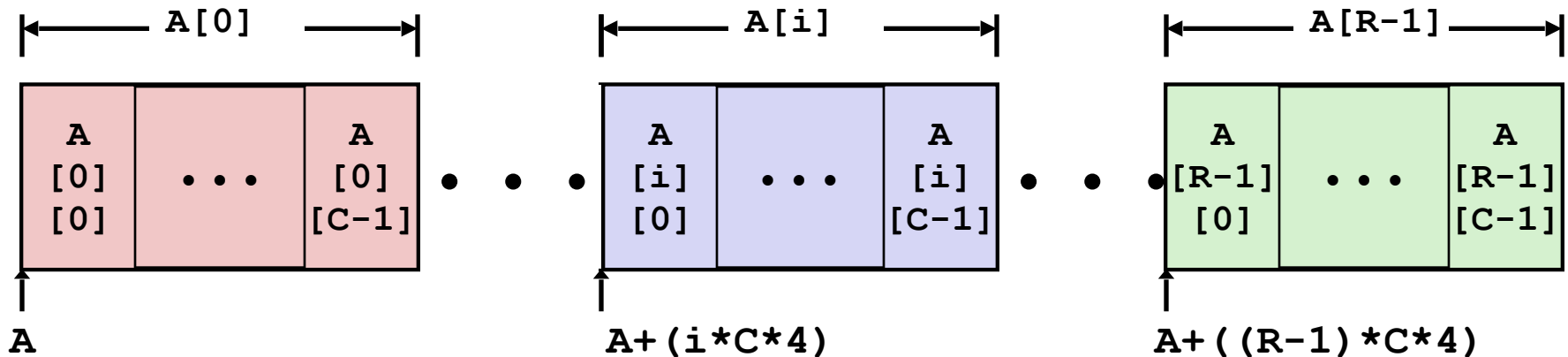
■ “Row-Major” ordering of all elements in memory

Multi-dimensional Array - Row Access

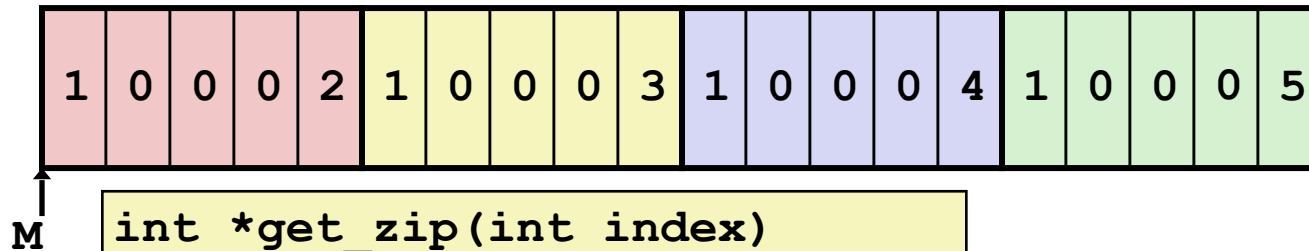
■ Row Vectors

- $A[i]$ is array of C elements
- Each element of type T requires K bytes
- Starting address $A + i * (C * K)$

```
int A[R][C];
```



Row Access Code



```
int *get_zip(int index)
{
    return M[index];
}
```

Assembly code?

Address: $M + index * C * K = M + index * 5 * 4$

```
# %rdi = index
leaq (%rdi,%rdi,4),%rax # 5 * index
leaq M(,%rax,4),%rax    # M + (20 * index)
```

■ Row Vector

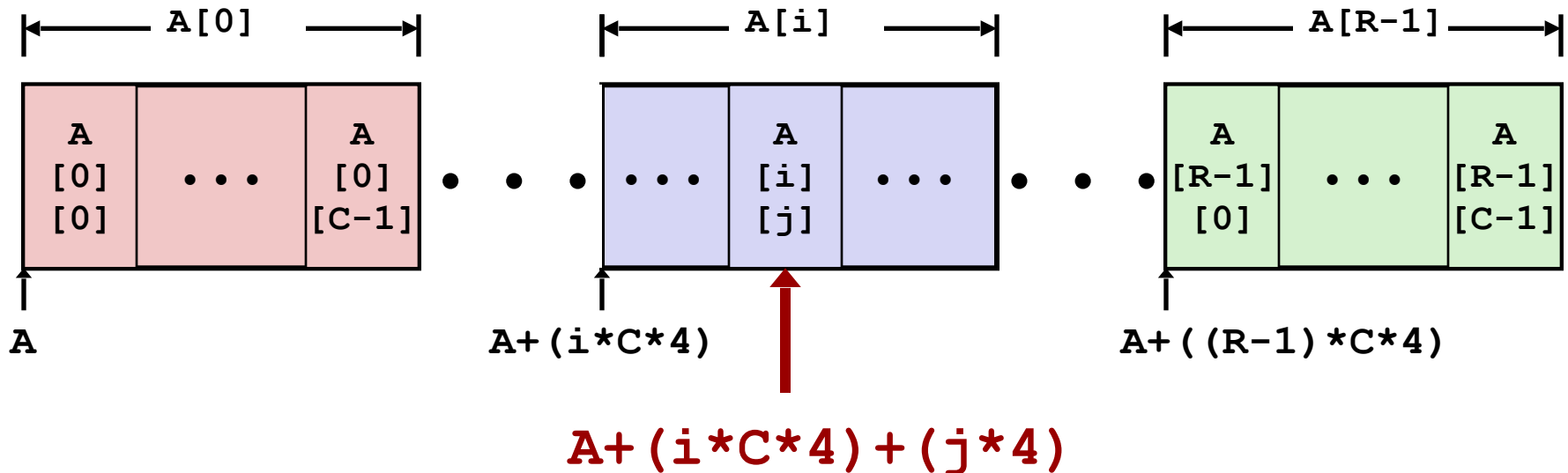
- $M[index]$ is array of 5 int's
- Starting address $M+20*index$

Multi-dimensional Array - Element Access

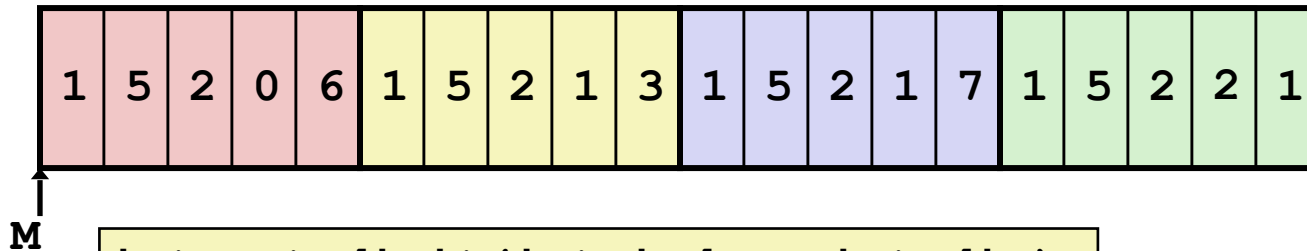
■ Array Elements

- $A[i][j]$ is element of type T , which requires K bytes
- Address $A + i * (C * K) + j * K = A + (i * C + j) * K$

```
int A[R][C];
```



Element Access Code



```
int get_digit(int index, int dig)
{
    return M[index][dig];
}
```

Assembly code?

```
leaq    (%rdi,%rdi,4), %rax    # 5*index
addl    %rax, %rsi             # 5*index+dig
movl    M(,%rsi,4), %eax       # M + 4*(5*index+dig)
```

■ Array Elements

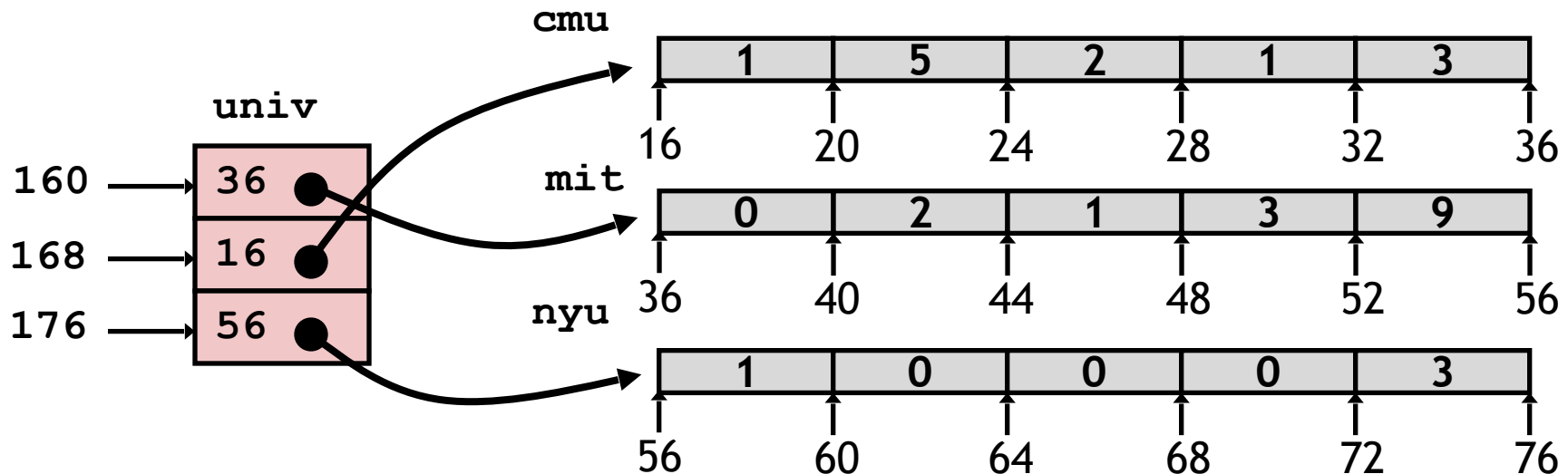
- `M[index][dig]` is `int`
- Address: $M + 20 \cdot \text{index} + 4 \cdot \text{dig}$
 - $= M + 4 \cdot (5 \cdot \text{index} + \text{dig})$
 - $M[\text{index}][\text{dig}] = M + K \cdot (C \cdot \text{index} + \text{dig})$
where K is 4 and C is 5

Multi-Level Array Example

```
int cmu[ZLEN] = { 1, 5, 2, 1, 3 };  
int mit[ZLEN] = { 0, 2, 1, 3, 9 };  
int nyu[ZLEN] = { 1, 0, 0, 0, 3 };
```

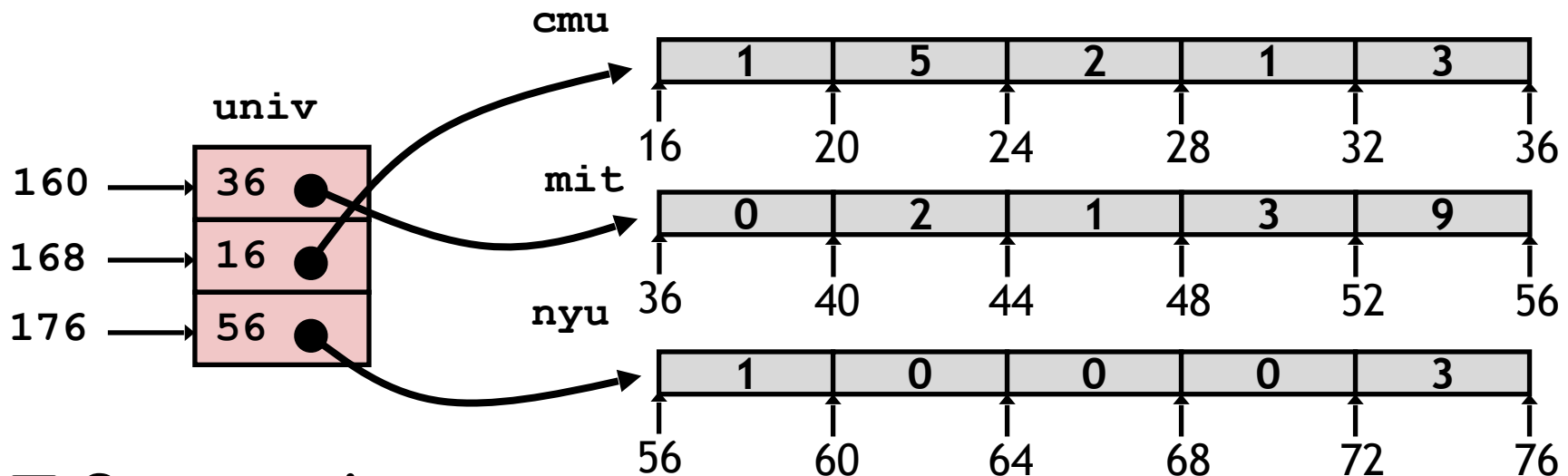
```
#define UCOUNT 3  
int *univ[UCOUNT] = {mit, cmu, nyu};
```

- Variable `univ` denotes array of 3 elements
- Each element is a pointer
 - 8 bytes
- Each pointer points to array of `int`'s



Element Access in Multi-Level Array

```
int get_digit(size_t index, size_t digit)
{
    return univ[index][digit];
}
```



■ Computation

- Must do two memory reads
 - First get pointer to row array
 - Then access element within array
- Element access $\text{Mem}[\text{Mem}[\text{univ} + 8 * \text{index}] + 4 * \text{digit}]$

Element Access in Multi-Level Array

```
int get_digit(size_t index, size_t digit)
{
    return univ[index][digit];
}
```

■ Computation

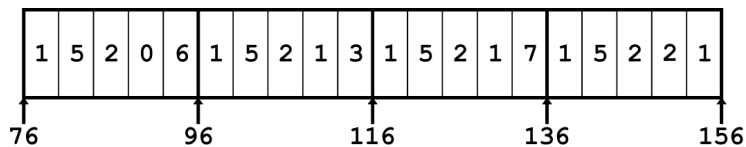
- Element access $\text{Mem}[\text{Mem}[\text{univ} + 8 * \text{index}] + 4 * \text{digit}]$

```
salq    $2, %rsi           # 4*digit
addq    univ(,%rdi,8), %rsi # p = univ[index] + 4*digit
movl    (%rsi), %eax        # return *p
ret
```

Array Element Accesses

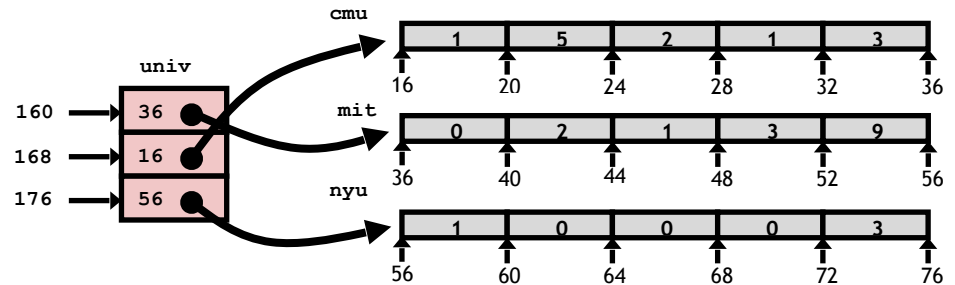
Flat

```
int get_digit
(size_t index, size_t digit)
{
    return M[index][digit];
}
```



Multi-level array

```
int get_digit
(size_t index, size_t digit)
{
    return univ[index][digit];
}
```



Accesses looks similar in C, but address computations very different:

$\text{Mem}[\text{M} + 20 * \text{index} + 4 * \text{digit}]$

$\text{Mem}[\text{Mem}[\text{univ} + 8 * \text{index}] + 4 * \text{digit}]$

Summary

■ Arrays

- We have seen
 - One-dimensional arrays
 - Multi-dimensional arrays
- For each one, we have seen
 - How to declare it in C?
 - How is it allocated in memory?
 - Contiguous region of memory
 - Array of arrays (multi-level arrays)
 - How to calculate the address of an individual element?
 - Use index arithmetic to locate individual elements
 - We have seen how to do this in C and in Assembly