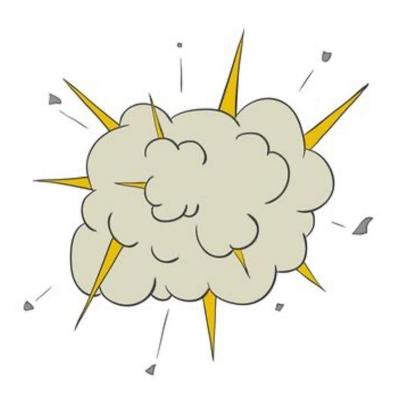
Attack Lab

Attack lab is released today!!



Agenda

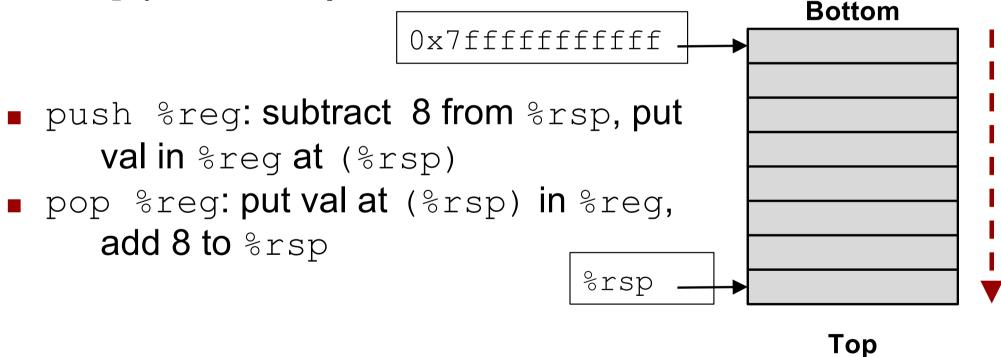
- Logistics
- Stack review
- Attack lab overview
 - Phases 1-3: Buffer overflow attacks
 - Phases 4-5: ROP attacks

Logistics

- Use remote Linux Server to complete this assignment ssh YourNetID@10.230.11.37 -p 4410
- Or use CentOS VM https://drive.google.com/file/d/1QLhvcloK5nrkv40PnfHBb1ZKcPlNibBG
 - Obtain Files from:
 http://DCLAP-V1111-CSD.ABUDHABI.NYU.EDU:15513/
 - Score Board:
 http://DCLAP-V1111-CSD.ABUDHABI.NYU.EDU:15513/scoreboard
- You need an active VPN connection to NYU/NYUAD network

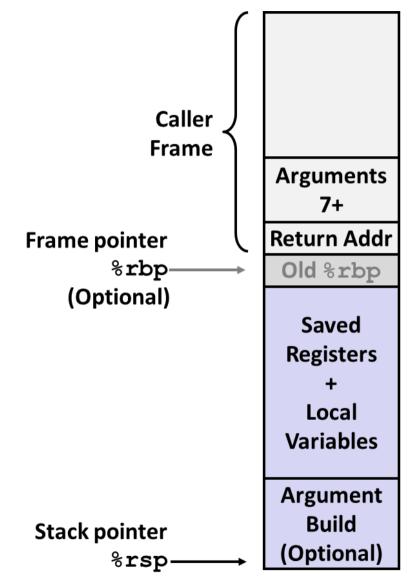
x86-64: The Stack

- Grows downward towards lower memory addresses
- %rsp points to top of stack



x86-64: Stack Frames

- Every function call has its own stack frame.
- Think of a frame as a workspace for each call.
 - Local variables
 - Callee & Caller-saved registers
 - Optional arguments for a function call



x86-64: Register Conventions

- Arguments passed in registers:
 - %rdi, %rsi, %rdx, %rcx, %r8, %r9
- Return value: %rax
- Callee-saved: %rbx, %r12, %r13, %r14, %rbp, %rsp
- Caller-saved: %rdi, %rsi, %rdx, %rcx, %r8, %r9, %rax, %r10, %r11
- Stack pointer: %rsp
- Instruction pointer: %rip

x86-64: Function Call Setup

Caller:

- Allocates stack frame large enough for saved registers, optional arguments
- Save any caller-saved registers in frame
- Save any optional arguments (in reverse order) in frame
- call foo: push %rip to stack, jump to label foo

Callee:

 Push any callee-saved registers, decrease %rsp to make room for new frame

x86-64: Function Call Return

Callee:

Increase %rsp, pop any callee-saved registers (in reverse order), execute ret: pop %rip

Attack Lab Overview: Phases 1-3

Overview

- Exploit x86-64 by overwriting the stack
- Overflow a buffer, overwrite return address
- Execute injected code

Key Advice

- Brush up on your x86-64 conventions!
- Use objdump –d to determine relevant offsets
- Use GDB to determine stack addresses

Buffer Overflows

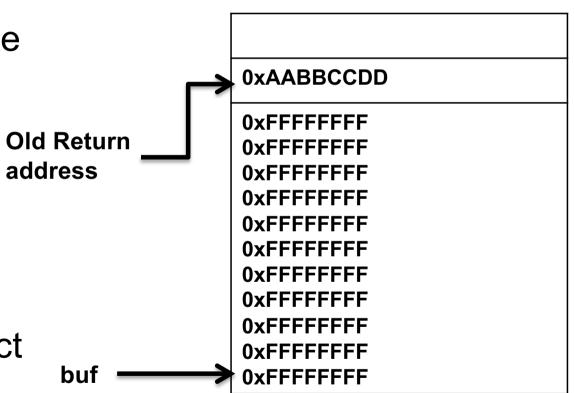
Exploit strcpy
 vulnerability to overwrite
 important info on stack
 Important info on stack

When this function returns, where will it begin executing?

Recall

ret:pop %rip

What if we want to inject new code to execute?



Attack Lab Overview: Phases 4-5

Overview

- Utilize return-oriented programming to execute arbitrary code
 - Useful when stack is non-executable or randomized
- Find gadgets, string together to form injected code

Key Advice

 Use mixture of pop & mov instructions + constants to perform specific task

ROP Example

Use ROP exploit to pop a value 0xBBBBBBB into %rbx and move it into %rax

```
void foo(char *input){
   char buf[32];
   ...
   strcpy (buf, input);
   return;
}
```

Gadgets:

address₁: mov %rbx, %rax; ret

address₂: pop %rbx; ret

Inspired by content created by Professor David Brumley

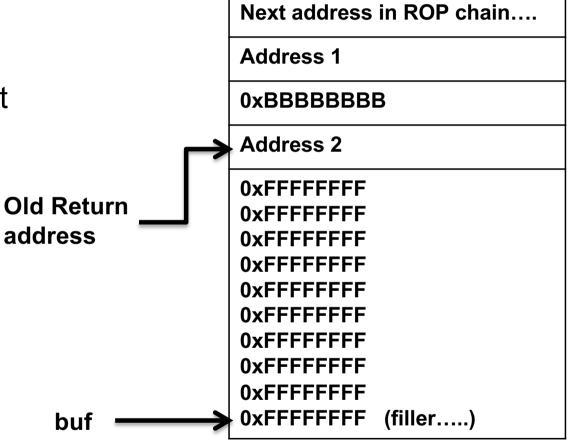
ROP Example: Solution

Gadgets:

Address 1: mov %rbx, %rax; ret

Address 2: pop %rbx; ret

```
void foo(char *input){
   char buf[32];
   ...
   strcpy (buf, input);
   return;
}
```



ROP Demonstration: Looking for Gadgets

How to identify useful gadgets in your code

Gadget Example #1

```
long ab_plus_c
  (long a, long b, long c)
{
   return a*b + c;
}
```

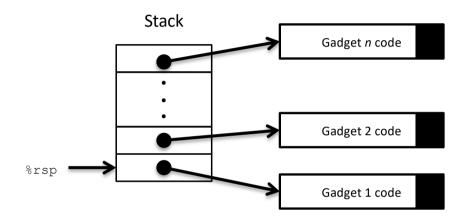
Use tail end of existing functions

Gadget Example #2

```
void setval(unsigned *p) {
    *p = 3347663060u;
}
```

Repurpose byte codes

ROP Execution



- Trigger with ret instruction
 - Will start executing Gadget 1
- Final ret in each gadget will start next one

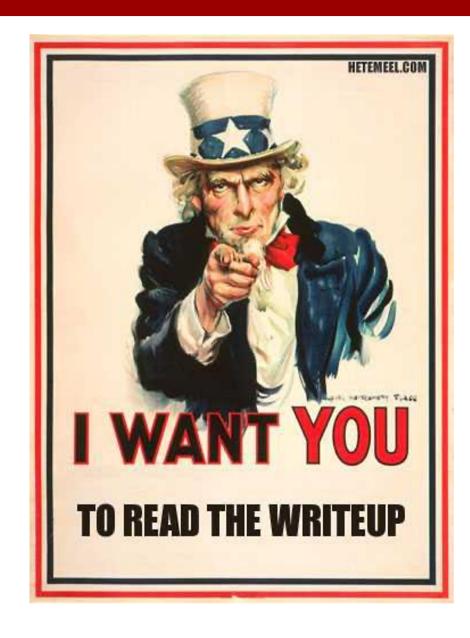
Tools

- objdump –d
 - View byte code and assembly instructions, determine stack offsets
- ./hex2raw
 - Pass raw ASCII strings to targets
- gdb
 - Step through execution, determine stack addresses
- gcc -c
 - Generate object file from assembly language file

More Tips

- Draw stack diagrams
- Be careful of byte ordering (little endian)

Also...



Questions?