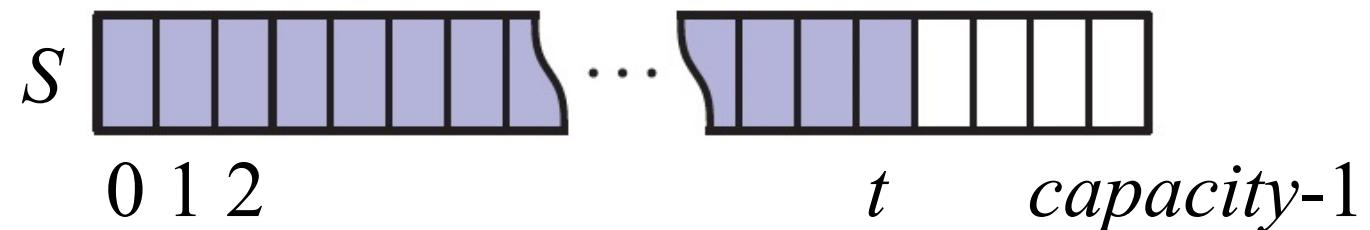
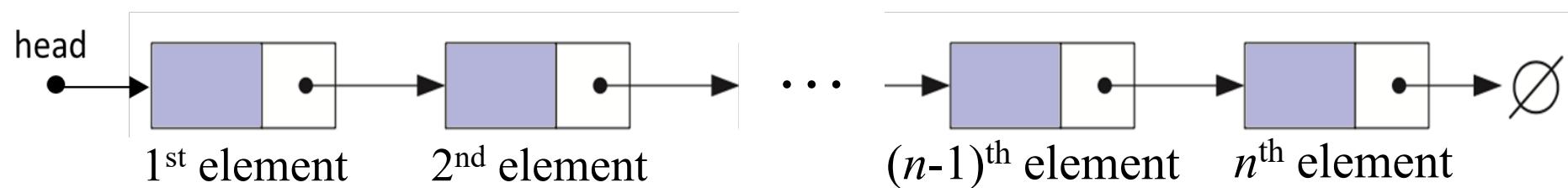


IMPLEMENTATION VIA A LINKED LIST

- Instead of implementing a stack using an array:



How about implementing it using a **linked list**:



THE STACK CLASS STRUCTURE

- Main Attributes:
 - A linked list to store the elements of the stack
 - A counter of the element in the stack

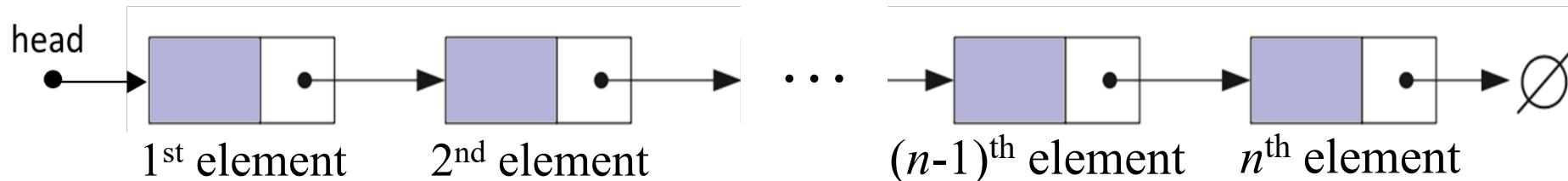
- Main Methods:
 - `top()`, `push(e)`, `pop()`, `size()` , `empty()`

CLASS DEFINITION

Compared to the array-based implementation, this part is **exactly the same**, except that “push” now doesn’t throw a “StackFull” exception, since the list size has no upper limit

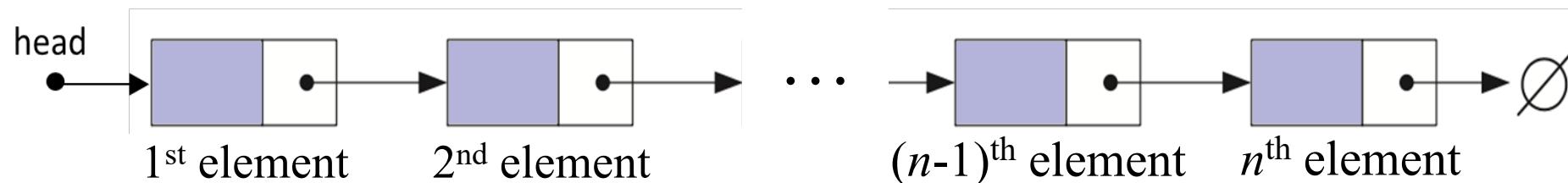
Compared to the array-based implementation, “S” is now defined as a linked list instead of an array, and we now have “n” instead of “t” and “capacity”

```
typedef string Elem;
class LinkedStack{
public :
    LinkedStack(); // constructor
    int size() const ;
    bool empty() const ;
    const E& top() const throw( StackEmpty );
    void push(const E& e);
    void pop() throw( StackEmpty );
private:
    SLinkedList<Elem> S; // linked list
    int n; // number of elements in stack
};
```



METHOD DEFINITION

```
LinkedStack::LinkedStack(): S(), n(0) { } // constructor, initializes S and sets n=0  
int LinkedStack::size() const { // returns the number of elements  
    return n;  $O(1)$   
}  
  
bool LinkedStack::empty() const { // returns true if stack is empty  
    return n == 0;  $O(1)$   
}  
const Elem& LinkedStack::top() const throw(StackEmpty) { // return top of stack  
    if (empty()) throw StackEmpty("Top of empty stack");  
    return S.front();  $O(1)$   
}
```



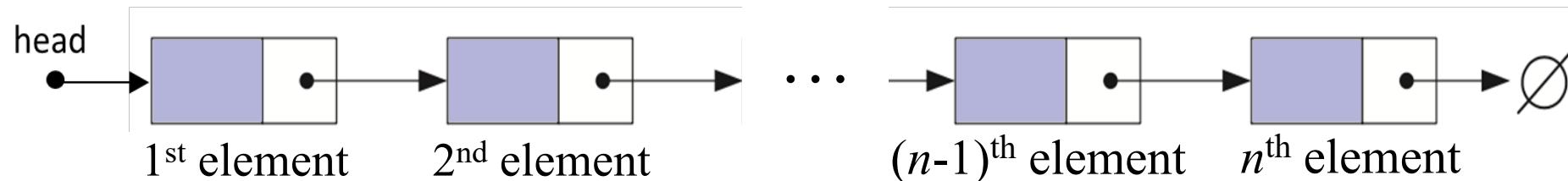
METHOD DEFINITION

```
void LinkedStack::push(const Elem& e) // push "e" onto stack
{
    ++n;
    S.addFront(e);
}

void LinkedStack::pop() throw(StackEmpty) // pop the stack
{
    if (empty()) throw StackEmpty("Pop from empty stack");
    --n;
    S.removeFront();
}
```

$O(1)$

$O(1)$



STACKS IN C++

- C++ provides a readily-available implementation of a stack (i.e., you do not have to implement it yourself).
- This comes as part of the “**Standard Template Library**” (STL) of C++, e.g., here is how to define a stack of integers:

```
#include <stack>  
using std::stack;  
stack<int> myStack;
```

This line defines myStack as a stack in which the elements are of type int; in this case we say the **base type** is int

ADT VS. STL

- **Abstract Data types (ADT)** is used to describe data structures without any implementation details
- **Standard Template Library (STL)**, which are readily-available data structures provided by C++
- **There might be some differences between the two:**
 - Example: In a stack **ADT**, we said earlier that the methods `pop()` and `top()` return **an error if the stack is empty**. While with stacks in C++'s **STL**, these two methods **do not throw an exception** if the stack is empty. Even though no exception is thrown, it may result in your program aborting!

STACK IMPLEMENTATION

- Practically, whenever you need to use a data structure such as a stack, you may use the available implementation in C++'s STL
 - Nevertheless, to better understand how data structures work, we will be required to implement them from scratch.
 - Moreover, you may want to handle exceptional cases an STL does not!

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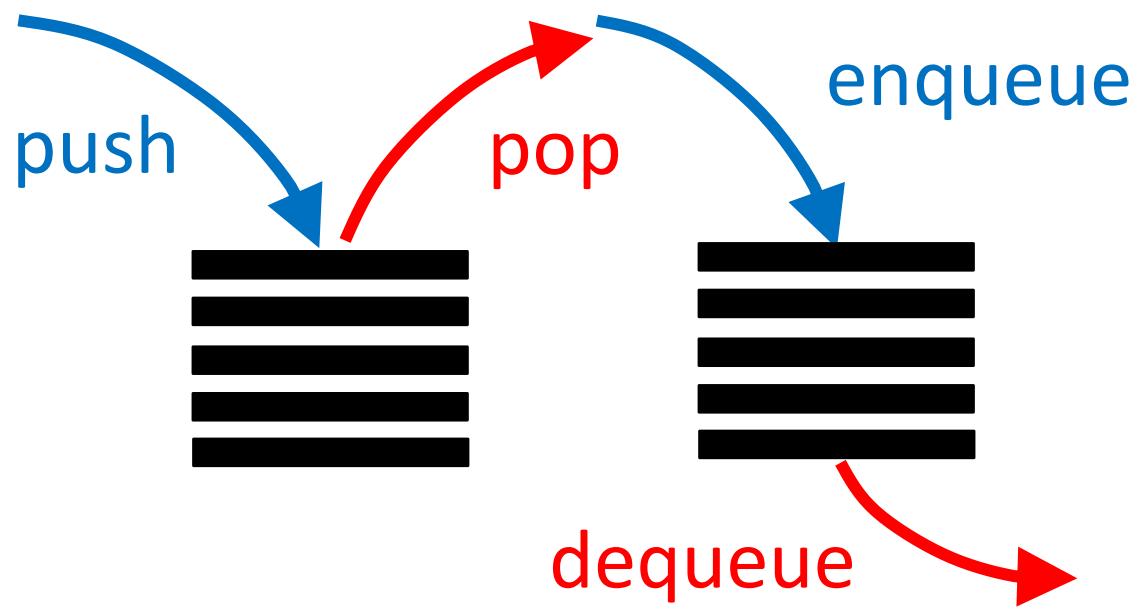
QUEUES



QUEUES

- A **queue** is a data structure that stores objects which can be added or removed according to the “*First-in First-out (FIFO)*” principle.

Stack: Last in, first out **Queue:** First in, first out



QUEUES: ADT

- Formally, a queue is an abstract data type (ADT) that supports the following operations:
 - **enqueue(e)**: Insert element **e** at the rear of the queue.
 - **dequeue()**: Remove element at the front of the queue; an error occurs if the queue is empty.
 - **front()**: Return a reference to the **value of the front element** in the queue; an error occurs if the queue is empty.
 - **size()**: Return the **number of elements** in the queue.
 - **empty()**: Return **true if the queue is empty** and false otherwise.

EXAMPLE

- In this example, if we write a queue as follows:

(x,y,z)

it means that:

- z is the **back** of the queue
- x is at the **front**

Now, let's fill this table!

P.S., under “**Output**” write:

- “**error**” if there is an error
- The **value** returned (if the function returns a value)
- “**—**” if there is **no error** and the function returns **no value**

<i>Operation</i>	<i>Output</i>	<i>front ← Q ← rear</i>
enqueue(5)	—	(5)
enqueue(3)	—	(5,3)
front()	5	(5,3)
size()	2	(5,3)
dequeue()	—	(3)
enqueue(7)	—	(3,7)
dequeue()	—	(7)
front()	7	(7)
dequeue()	—	()
dequeue()	“error”	()
empty()	true	()

ARRAY-BASED QUEUE

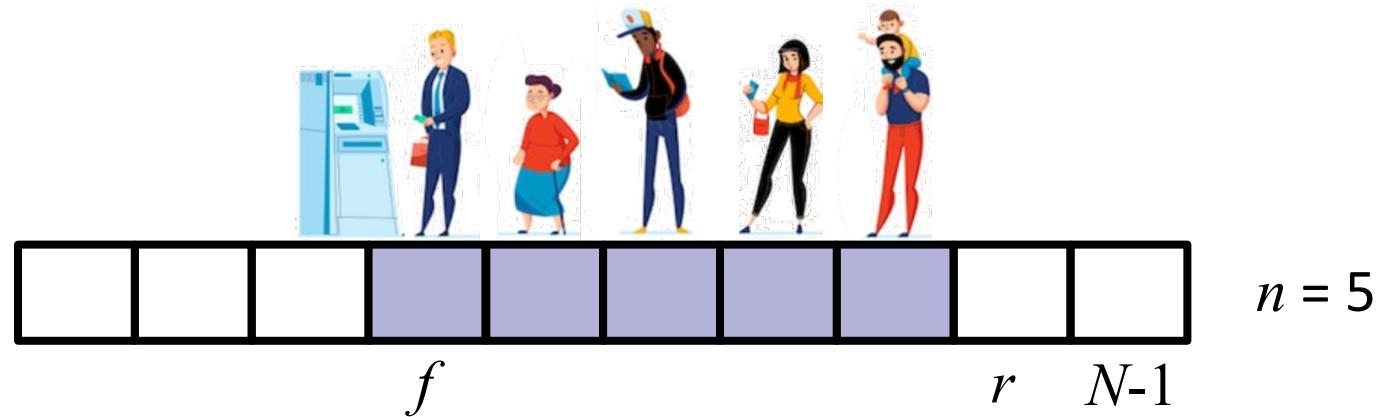
- In the array-based implementation of this queue, it helps to **imagine the ATM machine moving forward**, rather than the customers!



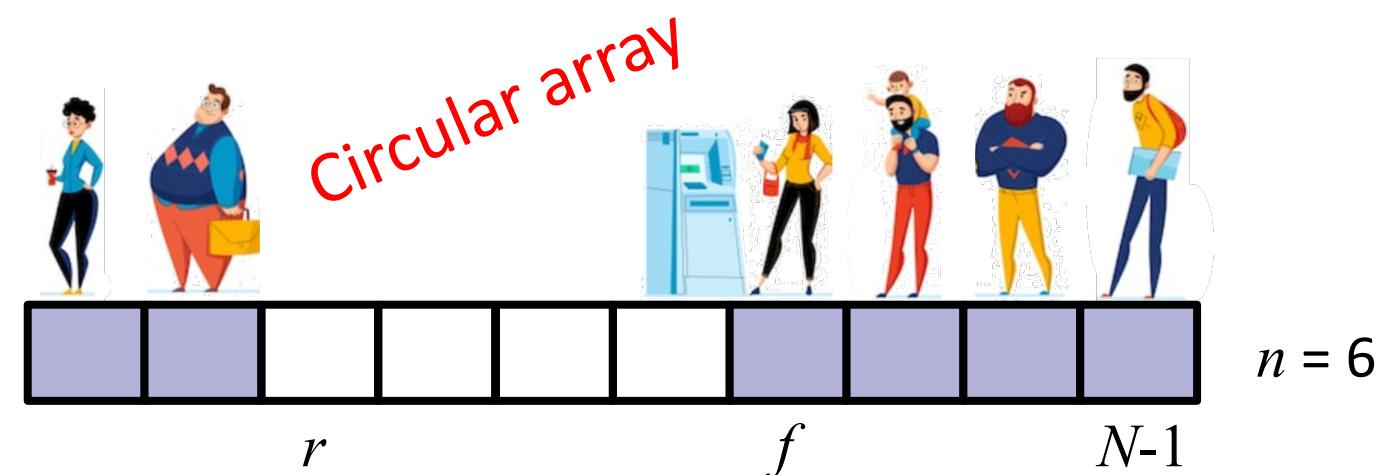
ARRAY-BASED QUEUE

f = index of front customer; r = index where a new joiner should stand;
 n = No. of elements; N = array size (i.e. capacity)

A potential state
of the queue:



Three customers
finished, and four
joined the queue:



QUEUE PSEUDO CODE

Algorithm size():

return n

Algorithm empty():

return ($n = 0$)

Algorithm front():

if empty() **then**

throw QueueEmpty exception

return $Q[f]$

Algorithm enqueue(e):

if size() = N **then**

throw QueueFull exception

$Q[r] \leftarrow e$

$r \leftarrow (r+1) \bmod N$

$n \leftarrow n+1$

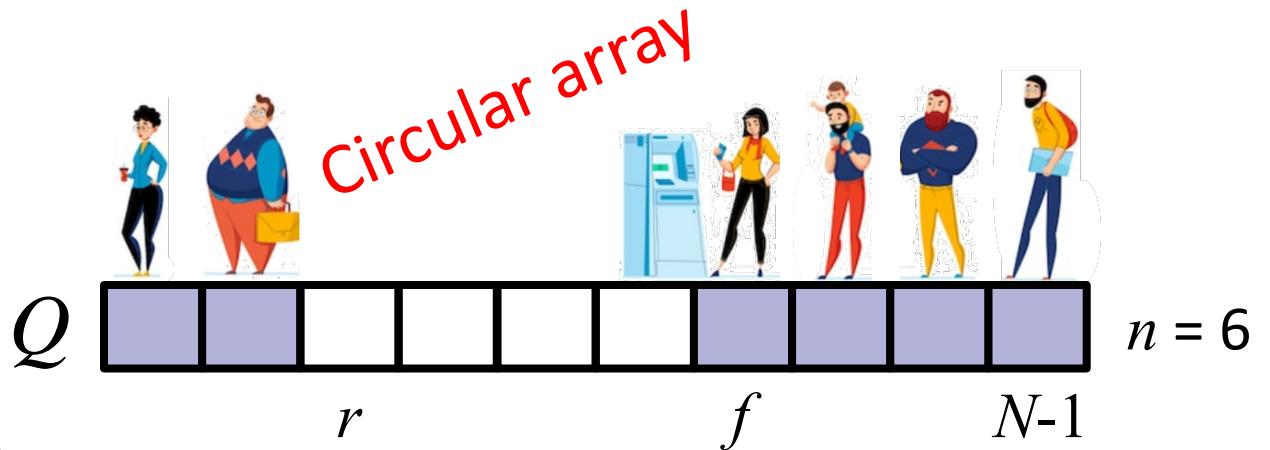
Algorithm dequeue():

if empty() **then**

throw QueueEmpty exception

$f \leftarrow (f+1) \bmod N$

$n \leftarrow n-1$



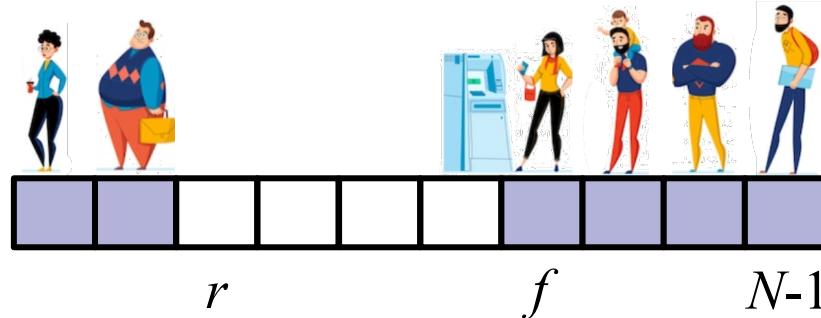
" $(r+1) \bmod N$ " is the remainder after dividing $(r+1)$ by N . That simply means:

"Increment r by 1, and if it falls out of the array, then go back to the beginning of the array."

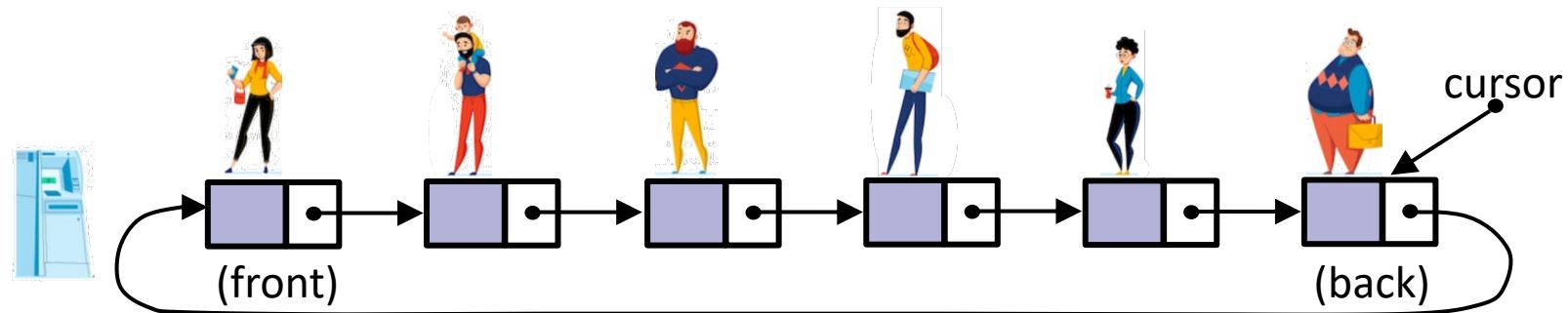
All methods are performed in $O(1)$ time, while the space usage is $O(N)$

USING A CIRCULARLY-LINKED LIST

- Instead of implementing a queue using an “circular array”:



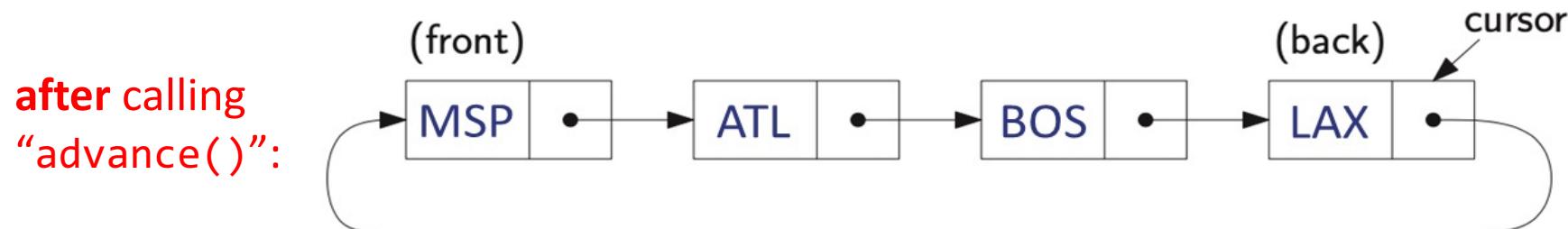
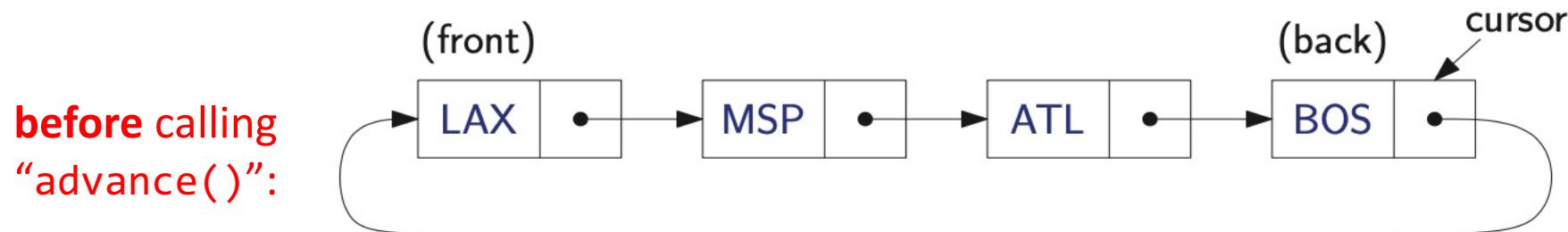
How about using a **circularly-linked list** instead?



Although the list can be used to implement a queue that rotates (i.e., where the ATM moves forward, just like with arrays); the ATM will always be on the left!

USING A CIRCULARLY-LINKED LIST

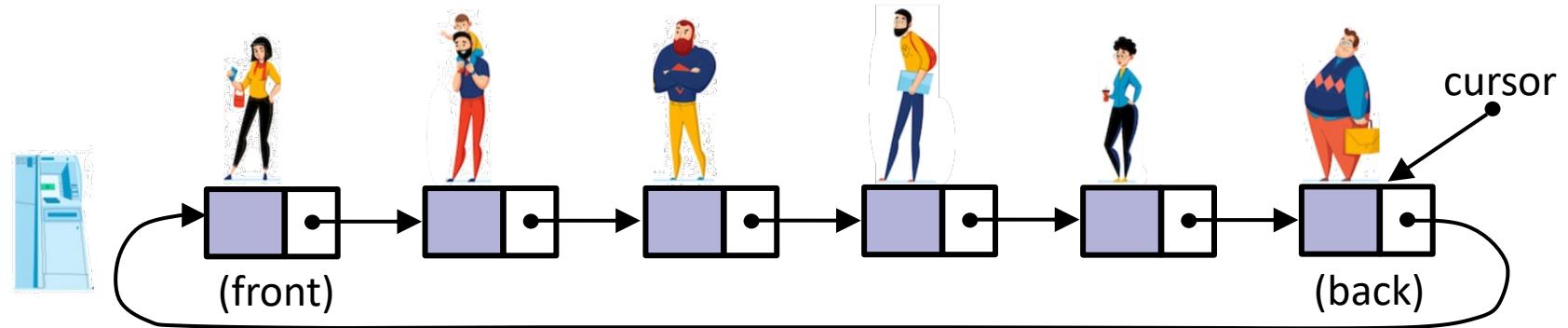
- Recall the main operations of a circularly-linked list:
 - `remove()`: deletes the node at the front of the list, i.e., the node right after the one that the cursor points to.
 - `add(e)`: puts `e` in a new node, which is added it to the front of the list.
 - `advance()`: moves the cursor one step forward in the list



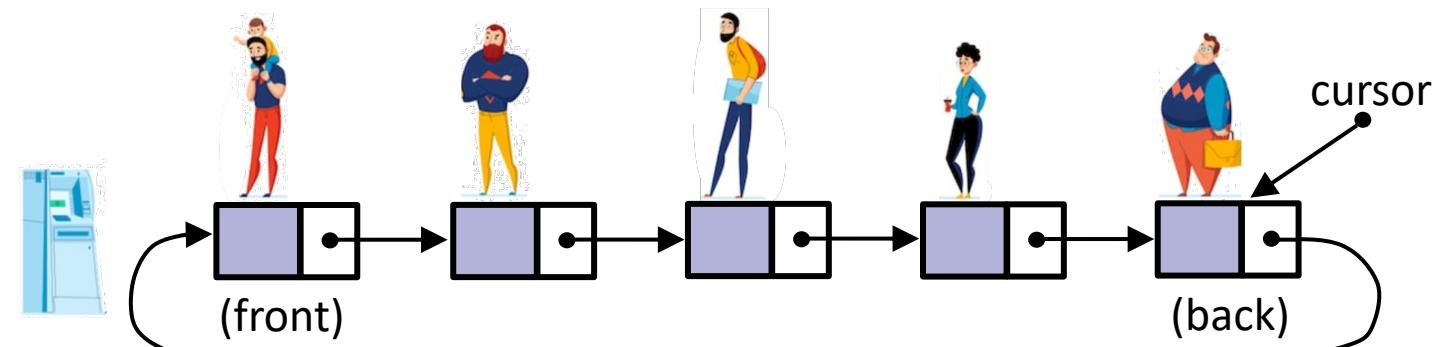
USING A CIRCULARLY-LINKED LIST

- To perform a **dequeue**, i.e., to remove the front customer:
 - call the method `remove()`, which deletes the node at the front, i.e., the node right after the one that the cursor points to

**before calling
“remove()”:**



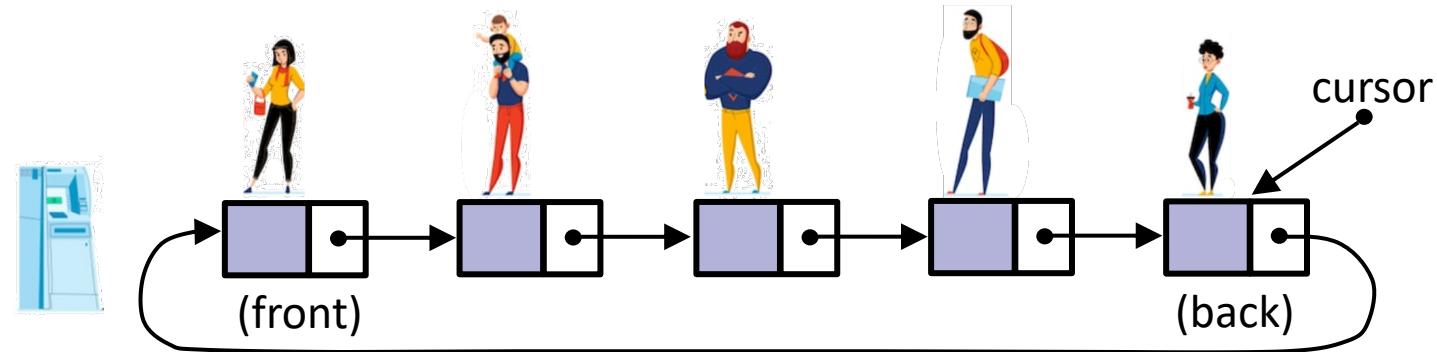
**after calling
“remove()”:**



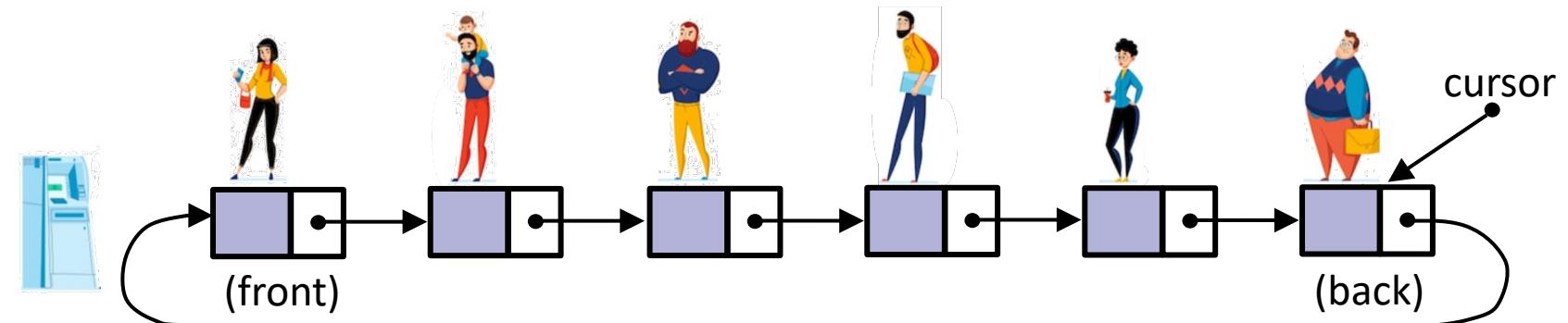
USING A CIRCULARLY-LINKED LIST

- To perform an **enqueue**, i.e., to add customer to the queue:
 - call **add(e)**, which puts **e** in a **new node at the front** of the list
 - then call **advance()**, which moves the **cursor one step forward**

before calling
add(e) and
advance():

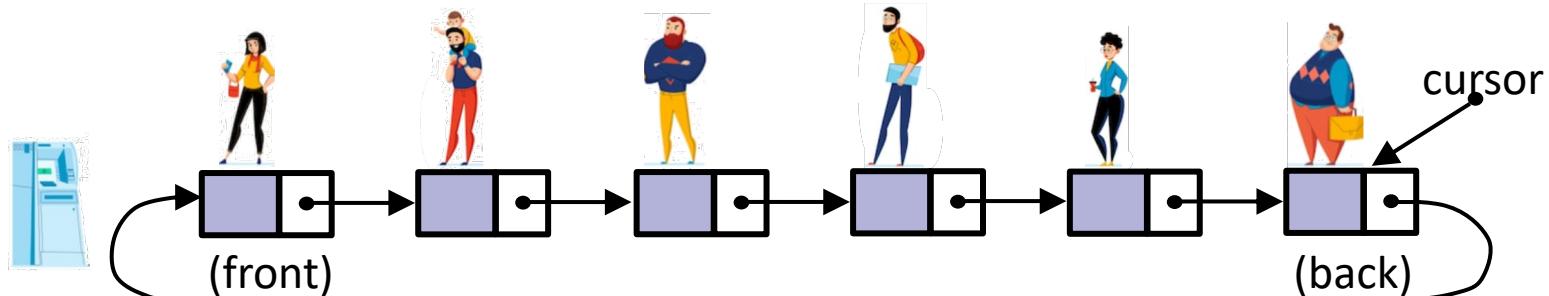


after calling
add(e) and
advance():



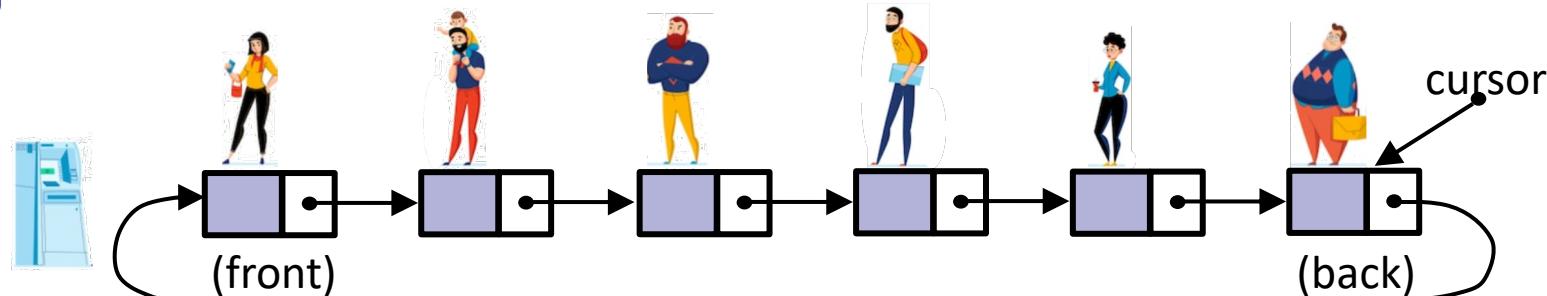
C++ IMPLEMENTATION

```
typedef string Elem; // queue element type
class LinkedQueue {
public:
    LinkedQueue(); // constructor
    int size() const; // number of items in the queue
    bool empty() const; // is the queue empty?
    const Elem& front() const throw(QueueEmpty); // the front element
    void enqueue(const Elem& e); // enqueue element at rear
    void dequeue() throw(QueueEmpty); // dequeue element at front
private:
    CircleList C; // circular list of elements
    int n; // number of elements
};
```



C++ IMPLEMENTATION

```
LinkedQueue::LinkedQueue() : C(), n(0) { } // constructor  
  
int LinkedQueue::size() const { return n; } // number of items in the queue  
  
bool LinkedQueue::empty() const { return n == 0; } // is the queue empty?  
  
const Elem& LinkedQueue::front() const throw(QueueEmpty) { // get front element  
    if (empty()) throw QueueEmpty("front of empty queue");  
    return C.front(); }  
  
void LinkedQueue::enqueue(const Elem& e) { // insert after cursor  
    C.add(e);  
    C.advance();  
    n++; }
```



C++ IMPLEMENTATION

```
LinkedQueue::LinkedQueue() : C(), n(0) { } // constructor

int LinkedQueue::size() const { return n; } // number of items in the queue

bool LinkedQueue::empty() const { return n == 0; } // is the queue empty?

const Elem& LinkedQueue::front() const throw(QueueEmpty) { // get front element
    if (empty()) throw QueueEmpty("front of empty queue");
    return C.front(); }

void LinkedQueue::enqueue(const Elem& e) { // insert after cursor
    C.add(e);
    C.advance();
    n++; }

void LinkedQueue::dequeue() throw(QueueEmpty) {
    if (empty()) throw QueueEmpty("dequeue of empty queue");
    C.remove();
    n--;
```

All methods are performed in O(1) time, and the space usage is O(n)

QUEUES: STL

- C++ provides a readily-available implementation of a queue (i.e., you do not have to implement it yourself).
- This comes as part of the “**Standard Template Library**” (STL) of C++, e.g., here is how to define a queue of floats:

```
#include <queue>
using std::queue;
queue<float> myQueue;
```

This line defines `myQueue` as a queue in which the elements are of type `float`; in this case, the “**base type**” is `float`

ADT VS. STL

- Reminder: Don't be confused between the following:
 - **Abstract Data types (ADT)**, which the textbook uses to describe data structures
 - **Standard Template Library (STL)**, which has data structures implemented in C++
- As we said earlier, there might be subtle differences between the two. For example:
 - The methods `enqueue()` and `dequeue()` in the queue **ADT** have different names in C++'s **STL**, which are `push()` and `pop()`, respectively.
 - Compared to the **ADT**, the queues in C++'s **STL** provides an additional method called `back()`, which returns a reference to the **element at the queue's rear**

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DEQUES



DOUBLE-ENDED QUEUE (DEQUE)

- It's a queue that supports **insertion** and **deletion** at both the **front** and the **back** of the queue!

A potential state
of the queue:



A customer can
skip the line!!



Last customer can
now leave before
reaching the ATM



DEQUE: ADT

- It's a queue that supports **insertion** and **deletion** at both the **front** and the **back** of the queue!
- Methods supported by DQ ADT:
 - `insertFront(e)`
 - `insertBack(e)`
 - `eraseFront()`
 - `eraseBack()`
 - `front()`
 - `back()`

<i>Operation</i>	<i>Output</i>	<i>D</i>
<code>insertFront(3)</code>	–	(3)
<code>insertFront(5)</code>	–	(5,3)
<code>front()</code>	5	(5,3)
<code>eraseFront()</code>	–	(3)
<code>insertBack(7)</code>	–	(3,7)
<code>back()</code>	7	(3,7)
<code>eraseFront()</code>	–	(7)
<code>eraseBack()</code>	–	()

DEQUES: STL

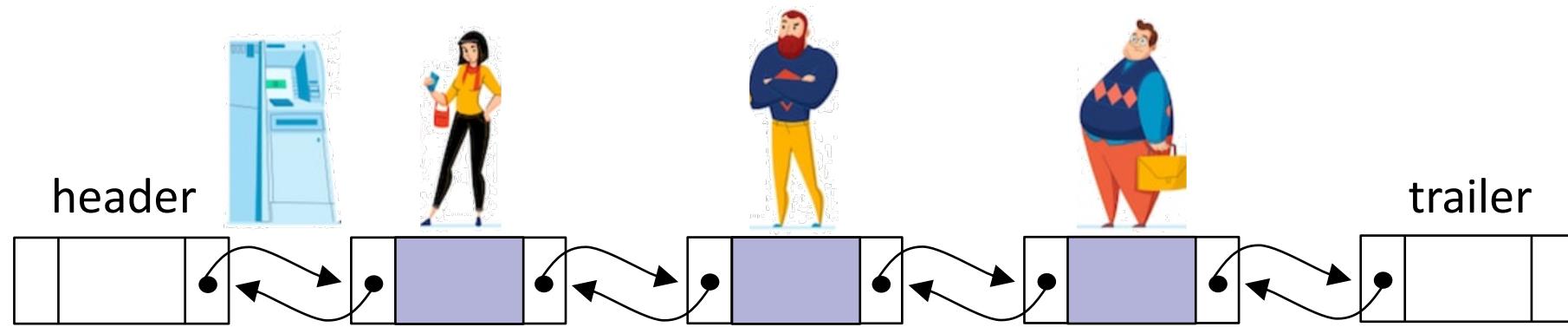
- C++ provides a readily-available implementation of a deque (i.e., you do not have to implement it yourself).
- This comes as part of the “**Standard Template Library**” (STL) of C++, e.g., here is how to define a queue of floats:

```
#include <deque>  
using std::deque;  
deque<string> myDeque;
```

This line defines deque as a deque in which the elements are of type string; in this case, the “**base type**” is string

USING A DOUBLY-LINKED LIST

- We will implement a deque using a **doubly-linked list**:



C++ IMPLEMENTATION

```
typedef string Elem; // deque element type
class LinkedDeque { // deque as doubly linked list
public:
    LinkedDeque(); // constructor
    int size() const; // number of items in the deque
    bool empty() const; // is the deque empty?
    const Elem& front() const throw(DequeEmpty); // the first element
    const Elem& back() const throw(DequeEmpty); // the last element
    void insertFront(const Elem& e); // insert new first element
    void insertBack(const Elem& e); // insert new last element
    void eraseFront() throw(DequeEmpty); // remove first element
    void eraseBack() throw(DequeEmpty); // remove last element
private: // member data
    DLinkedList D; // linked list of elements
    int n; // number of elements
};
```

C++ IMPLEMENTATION

```
void LinkedDeque::insertFront(const Elem& e) { D.addFront(e); n++; }
void LinkedDeque::insertBack(const Elem& e) { D.addBack(e); n++; }
void LinkedDeque::eraseFront() throw(DequeEmpty) {
    if (empty()) throw DequeEmpty("removeFront of empty deque");
    D.removeFront();
    n--;
}
void LinkedDeque::eraseBack() throw(DequeEmpty) {
    if (empty()) throw DequeEmpty("removeBack of empty deque");
    D.removeBack();
    n--;
}
```

All methods are performed in $O(1)$ time, and the space usage is $O(n)$