



# Programming Languages

**Introduction to Computer**

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*(with most slides borrowed from Prof. Tian-Li Yu)*

# Outline

- Historical perspective
- Traditional programming concepts
- Procedural units
- Language translation process
- Object-oriented programming
- Programming concurrent activities

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# Programming Language Generations

**1<sup>st</sup>**

**2<sup>nd</sup>**

**3<sup>rd</sup>**

**4<sup>th</sup> ?**



**Machine  
instructions**

**Assembly**

**Fortran  
Cobol  
Basic  
C/C++  
Java**

**SQL  
SAS**

# From Machine Instructions to Assembly

1<sup>st</sup>  
Machine  
instructions

2<sup>nd</sup>  
Assembly

156C	→	LD R5, Price
166D	→	LD R6, ShippingCharge
5056	→	ADDI R0, R5, R6
306E	→	ST R0, TotalCost
C000	→	HTL

6E = 6C + 6D

TotalCost =  
Price + ShippingCharge

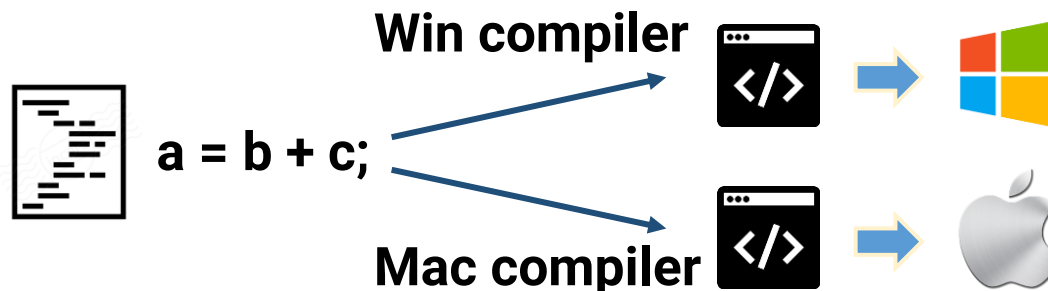
- **Mnemonic** names for op-codes
- **Program variables** or **identifiers**: descriptive names for memory locations, chosen by the programmer

# Assembly Language Characteristics

- **One-to-one** correspondence between machine instructions and assembly instructions
  - Programmer must think about the machine
- Inherently **machine-dependent**
- Converted to machine language by a program called an **assembler**

# 3<sup>rd</sup> Generation (High-level) Language

- Use high-level primitives
  - E.g., if-then, do-while
- Each primitive corresponds to a sequence of machine language instructions
- **Machine independent (mostly)**
- Converted to machine language by a program called a **compiler (or interpreter)**



# Programming Languages and Issues

- Natural v.s. Formal languages

- Formal language

- Use formal grammar

**Expression  $\rightarrow$  Term | Term + Expression | Term – Expression**

**Term  $\rightarrow$  Factor | Factor \* Term | Factor / Term**

**Factor  $\rightarrow$  x | y | z**

**x + y \* z**

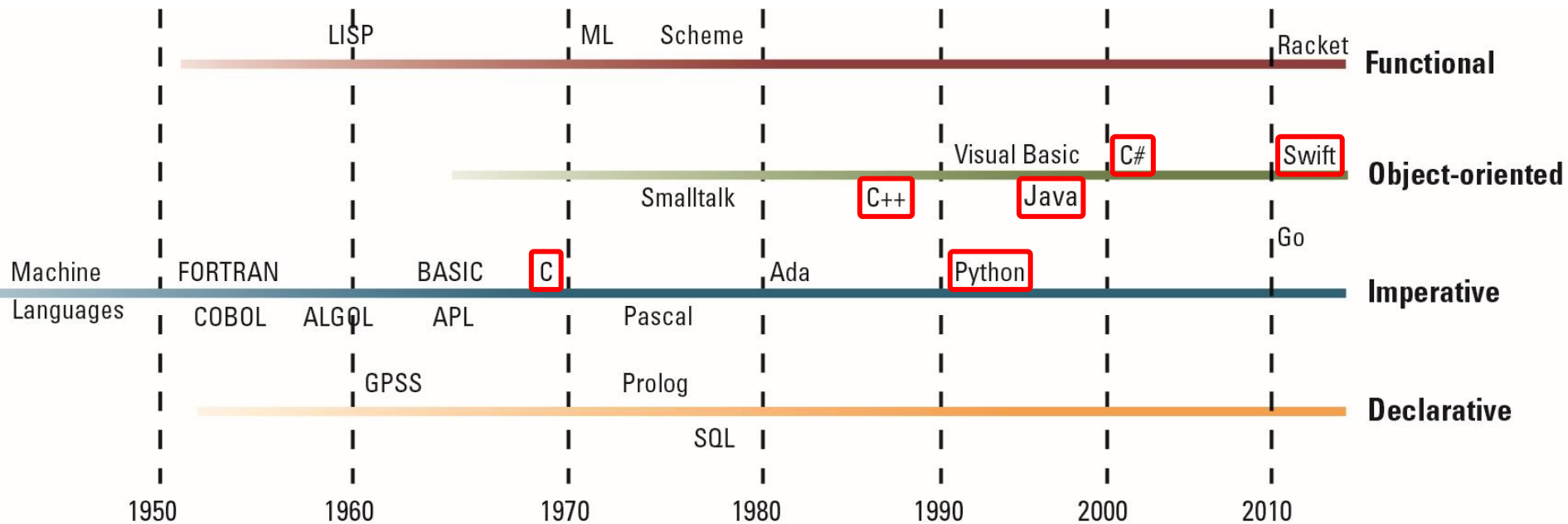
- Will be introduced later

- **Portability**

- Theoretically: same source code, different compilers
  - Reality: minor modifications

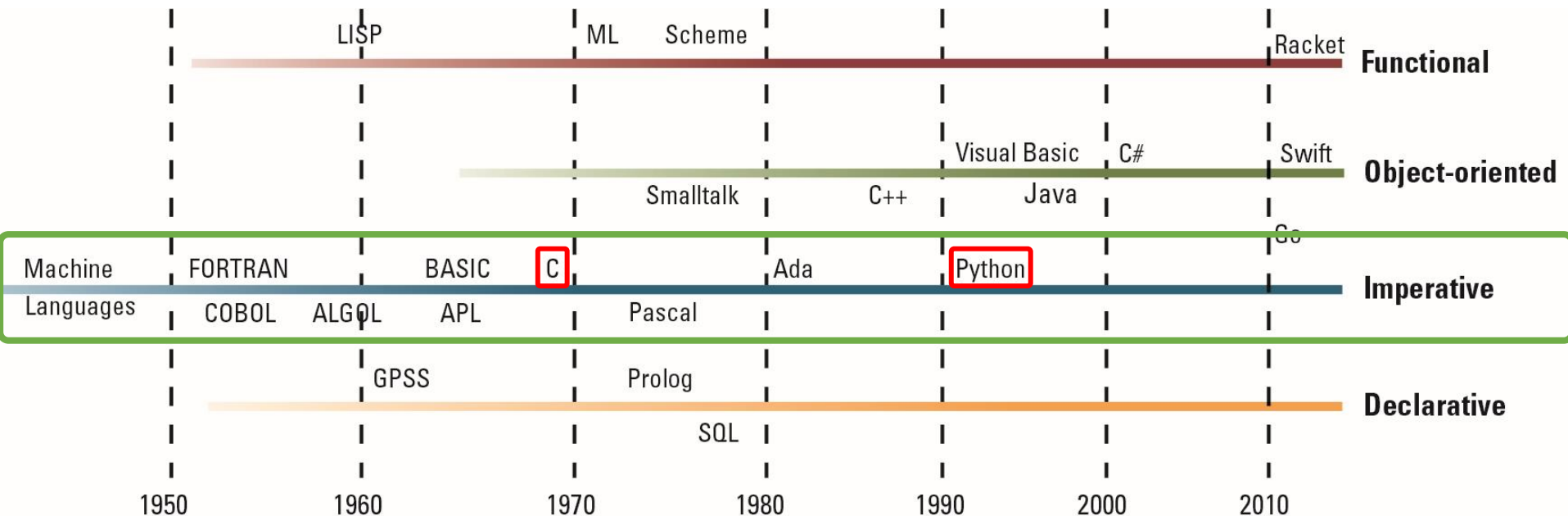


# Programming Language Paradigms



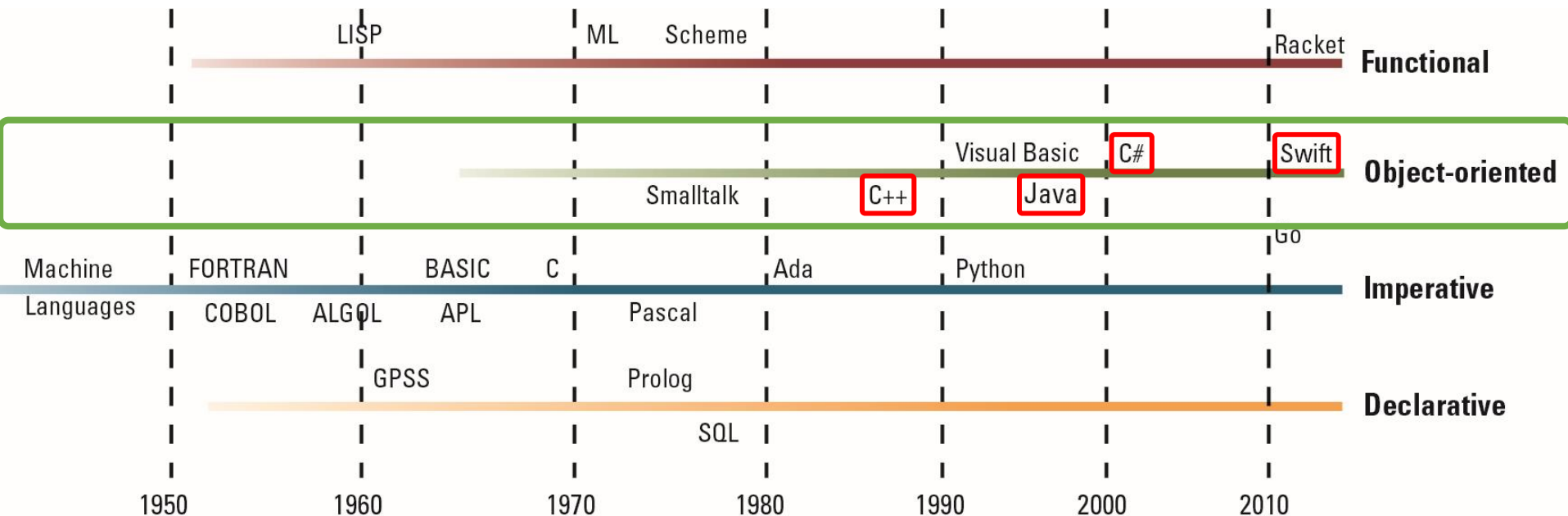
# Imperative Paradigms

- Procedural
- Approach a problem by **finding an algorithm** to solve the problem



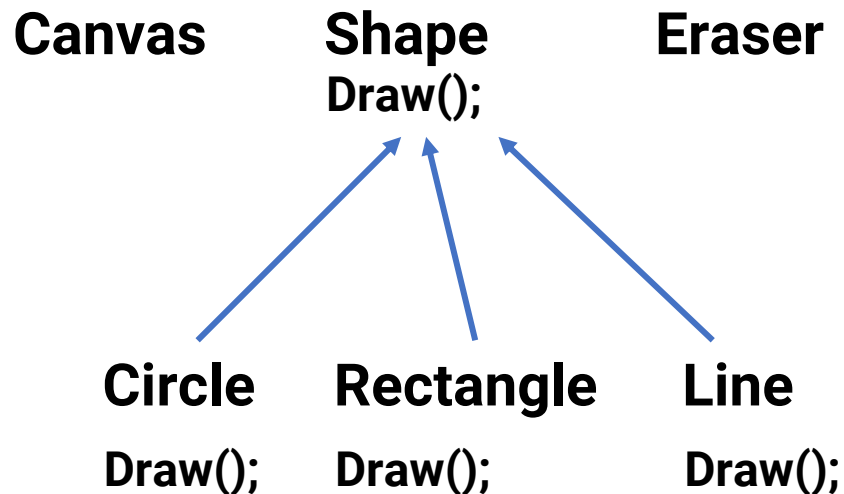
# Object-oriented Paradigms

- Implements **objects and their associated procedures** within the programming context to create software programs
- Information hiding, inheritance, polymorphism



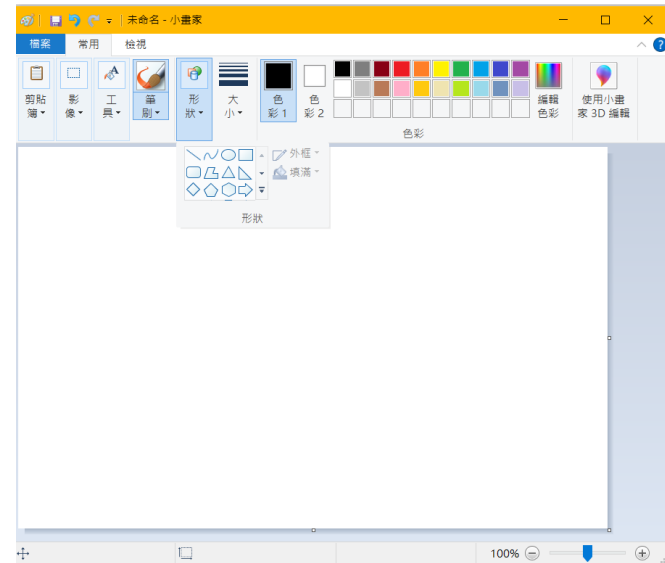
# Object-oriented Paradigms (cont.)

- Example: a painter program



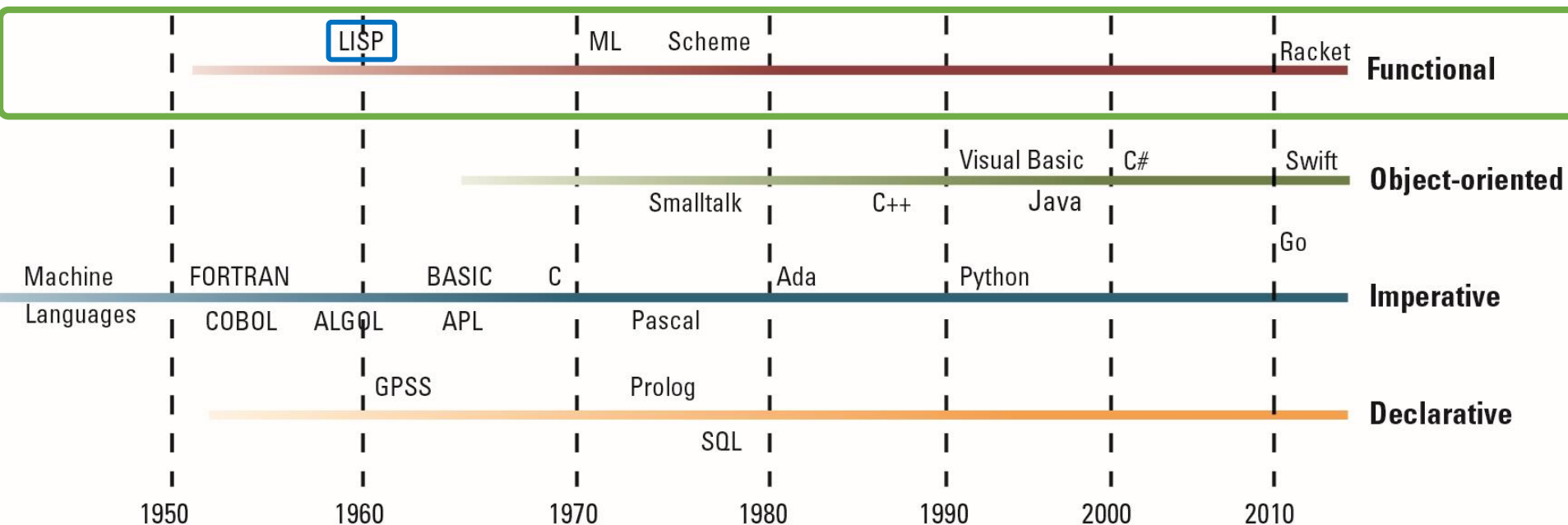
```

Shape* shapeList[10];
foreach shape in shapeList
    shape->Draw();
  
```

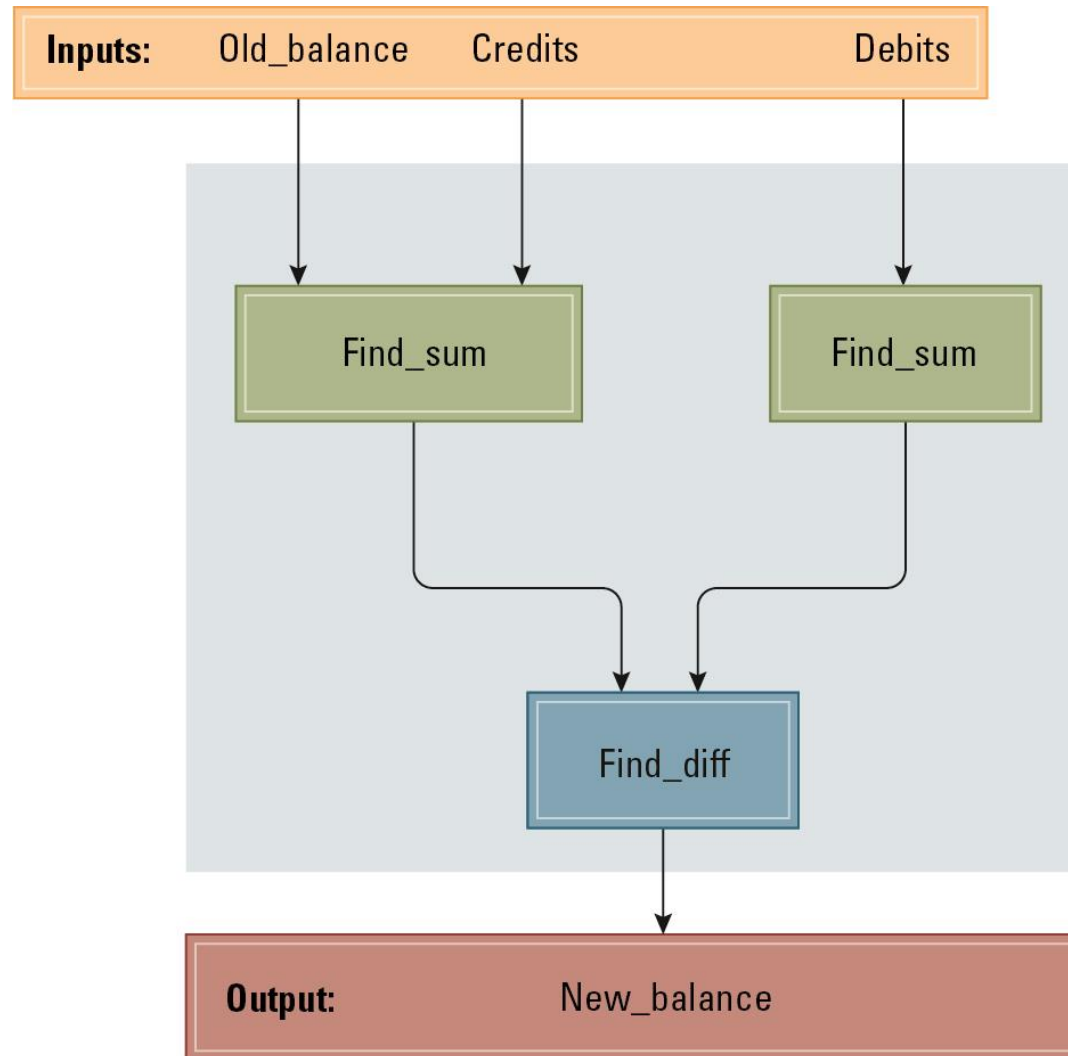


# Functional Paradigms

- Treat the entire program as a function
- A program consists of sub-problems that are handled by sub-functions



# Functional Paradigm (cont.)



# Functional v.s. Imperative

Temp\_balance  $\leftarrow$  Old\_balance + Credit

Total\_debits  $\leftarrow$  sum of all Debits

Balance  $\leftarrow$  Temp\_balance – Total\_debits

**LISP** (f x y)

**(Find\_diff (Find\_sum Old\_balance Credits) (Find\_sum Debits))**

Sum  $\leftarrow$  sum of all Numbers

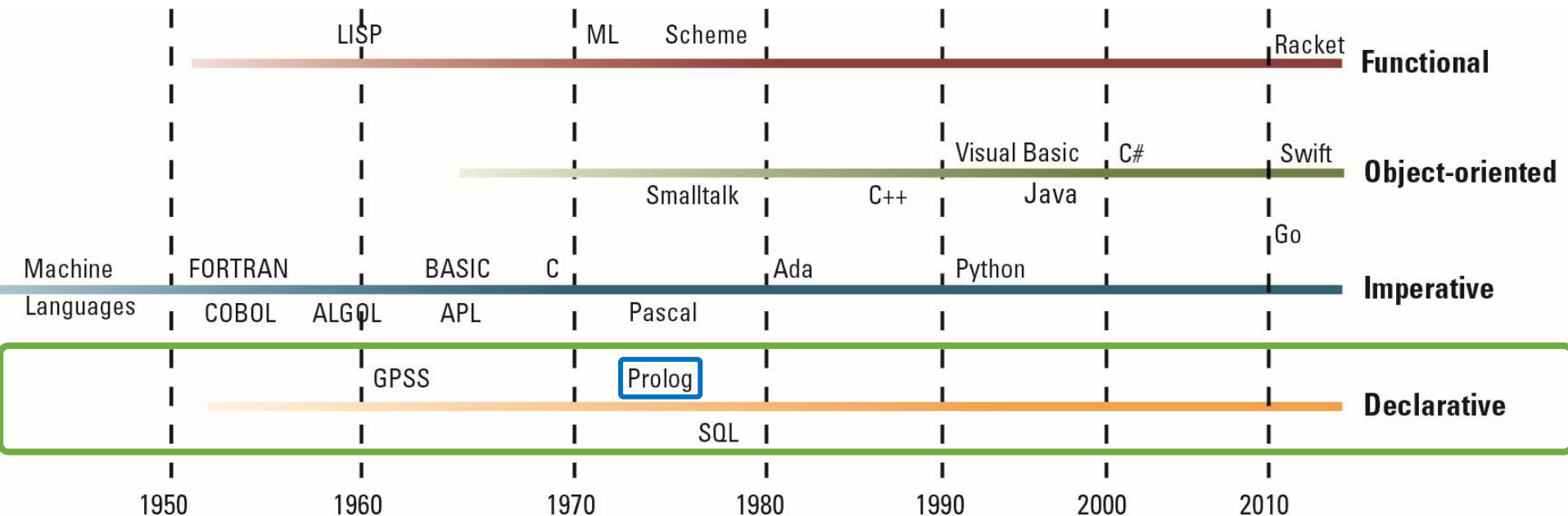
Count  $\leftarrow$  # of Numbers

Average  $\leftarrow$  Sum / Count

**(Find\_average (Find\_sum Numbers) (Find\_count Numbers))**

# Declarative Paradigms

- Implemented as a **general problem solver**
- Approach a problem by **finding a formal description of the problem**
  - E.g., define ***factorial*** by  $0! = 1$  and  $n! = n * (n-1)!$





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# Traditional Programming Concepts

- Variables and data types
- Data structure
- Constants and literals
- Assignments and operators
- Control
- Comments

# Variables and Data Types

- **Integer**: whole numbers
- **Floating-point (Real)**: numbers with fractions
- **Character**: symbols
- **Boolean**: true/false

## C/C++, Java

```
int a;  
float b;  
char c;  
bool d;
```

## FORTRAN

```
INTEGER a;  
REAL b;  
BYTE c;  
LOGICAL d;
```

# Data Structures

- Conceptual shape or **arrangement of data**
- A common data structure is the **array**
- **Homogeneous array**

**C/C++, Java**

```
int a[5][100];
```

**FORTRAN**

```
INTEGER a(5, 100);
```

- The starting index might differ in different programming languages

**Scores**


**Scores** (2, 4) in  
FORTRAN where  
indices start at one.

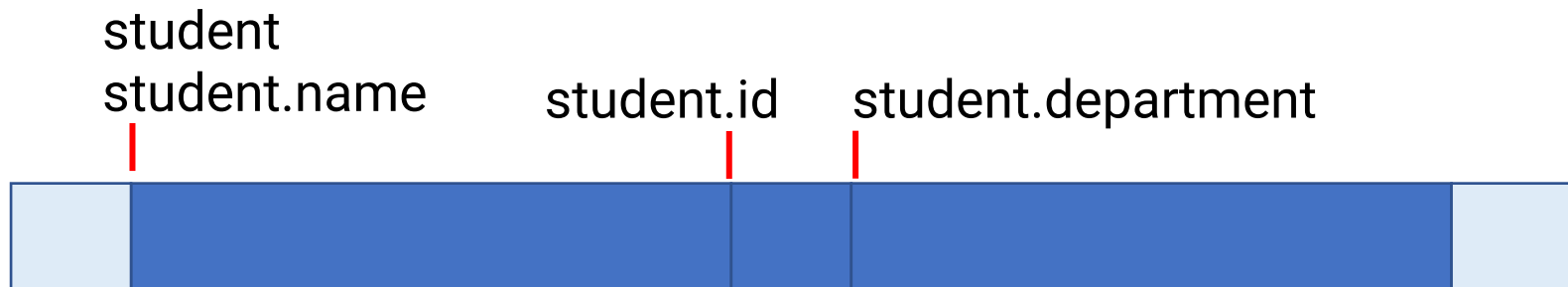
**Scores** [1][3] in C  
and its derivatives  
where indices start  
at zero.

# Data Structures (cont.)

- Conceptual shape or **arrangement of data**
- A common data structure is the **array**
- **Heterogeneous array**

**C/C++**

```
struct Student {  
    char name[30];  
    int id;  
    char department[30];  
};
```



# Literals and Constant

- **Literal**

- $a \leftarrow b + 100;$

- **Constant**

- `Const int a = 100;` (C/C++)
  - `final int a = 100;` (Java)
  - A constant cannot be a **l-value**
    - `const int a = 100;`  
`a = b + c;` ❌

# Assignment and Operators

- **Assignment**

- $a = b + c$ ; (C/C++/Java)

- **Operators**

- Operator precedence

- E.g.,  $\text{int } a = 3 + 4 * 5 \mid 6$ ;

- [https://en.cppreference.com/w/c/language/operator\\_precedence](https://en.cppreference.com/w/c/language/operator_precedence)

- Operator overloading

```
struct Complex {
    int real;
    int imag;
    Complex operator+ (Complex const& obj) {
        Complex res;
        res.real = real + obj.real;
        res.imag = imag + obj.imag;
        return res;
    }
};
```

```
Complex c1, c2;
c1.real = 3;    // c1 = 3 + 4i
c1.imag = 4;
c2.real = 5;    // c2 = 5 + 6i
c2.imag = 6;
Complex c3 = c1 + c2;
std::cout << c3.real << " + " << c3.imag << "i" << std::endl;
8 + 10i
```

# Control Statements

- **Old-fashion: goto**

- Not recommended

```
line #  goto 4
        print "passed."
        goto 7
4       if (grade < 60) goto 6
        goto 2
6       print "failed."
7       stop
```

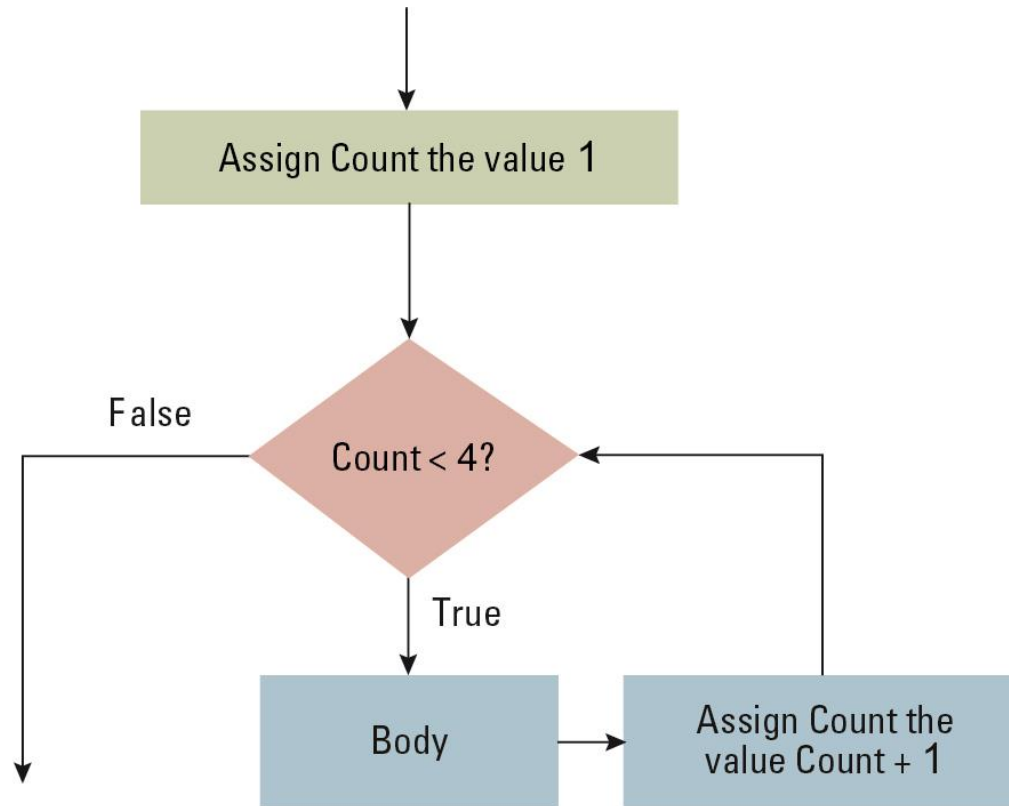
- Modern programming

- if / else if / else
- switch
- for
- while



# Control Statements (cont.)

- for



```
for (int Count = 1; Count < 4; Count++)  
    body;
```

# Comments

- Explanatory statements within a program
- Helpful when a human reads a program
- Ignored by the compiler

```
a = b + c;  // End-of-line comment.
```

```
/* Block comment */
```

```
a = b + c;
```

```
/**
```

```
 | Documentation comment.
```

```
*/
```

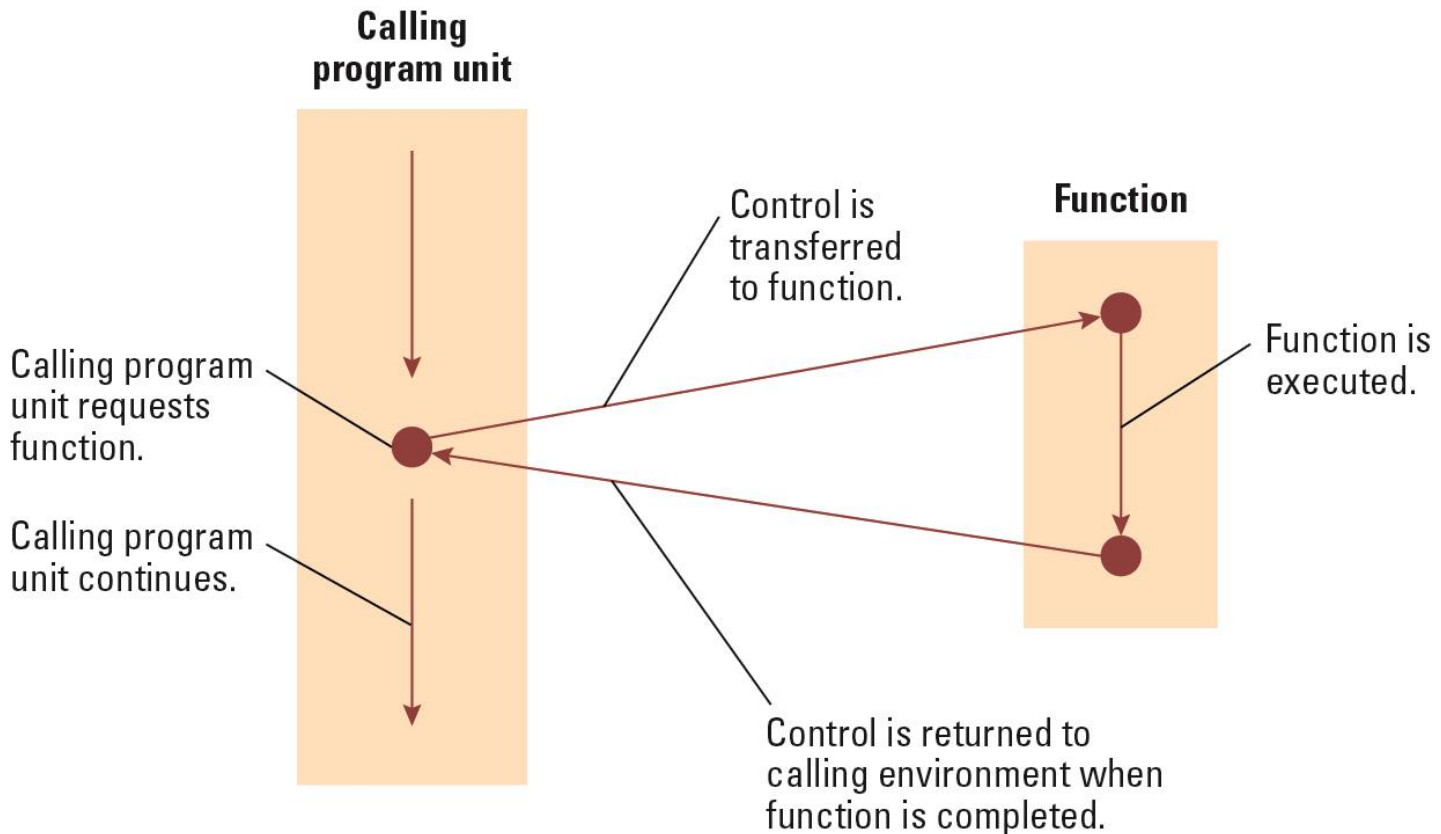
```
a = b + c;
```

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# Procedural Units

- Many terms for this concept:
  - Subprogram, subroutine, procedure, method, function



# Procedural Units (cont.)

- Terminology

Starting the header with the term “void” is the way that a C programmer specifies that the program unit returns no value. We will learn about return values shortly.

**type of the return value**

`void` `ProjectPopulation` (`float` `GrowthRate`)

The formal parameter list. Note that C, as with many programming languages, requires that the data type of each parameter be specified.

**formal parameters**

**function header**

```
{ int Year; // This declares a local variable named Year.
```

**local variable**

```
Population[0] = 100.0;
for (Year = 0; Year <= 10; Year++)
Population[Year+1] = Population[Year] + (Population[Year] * GrowthRate);
```

These statements describe how the populations are to be computed and stored in the global array named `Population`.

# Procedural Units (cont.)

- Terminology

The function header begins with the type of the data that will be returned.

**type of the return value**

```
float CylinderVolume (float Radius, float Height)
```

```
{ float Volume;
```

Declare a local variable named Volume.

```
Volume = 3.14 * Radius * Radius * Height;
```

```
return Volume;
```

Compute the volume of the cylinder.

**return value**

Terminate the function and return the value of the variable Volume.

```
}
```

# Procedural Units (cont.)

- Function's (procedure's) header

void Swap(int\*, int\*);      can be put in another header file

```
int a = 5;
int b = 3;
Swap(&a, &b);
std::cout << a << " " << b << std::endl;
```

```
void Swap(int* a, int* b)
{
    int temp = *a;
    *a = *b;
    *b = temp;
}
```

# Procedural Units (cont.)

- **Local variable** and **global variables**

```
// Global variable.
```


```
int var = 0;
```


```
int main()
```

```
{
```

```
    // Local variable.
```

```
    int var = 5;
```

```
    std::cout << var << std::endl;  5
```

```
    std::cout << ::var << std::endl;  0
```



# Procedural Units (cont.)

- **Formal parameters** and **actual parameters**

```
void Swap(int* a, int* b) a, b: formal parameters
{
    int temp = *a;
    *a = *b;
    *b = temp;
}

int main()
{
    int x = 5;
    int y = 3;
    Swap(&x, &y); x, y: actual parameters
    std::cout << x << " " << y << std::endl;
}
```

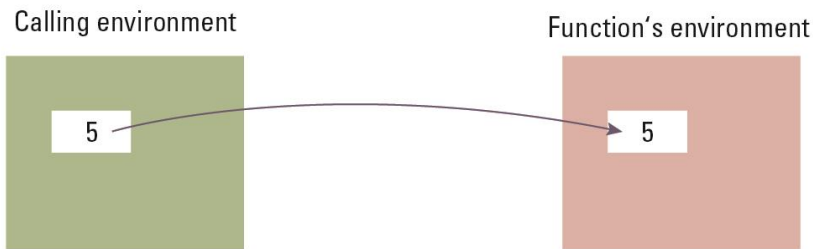
# Procedural Units (cont.)

- **Passing parameters**
  - Call by **value** (passed by value)
  - Call by **reference** (passed by reference)
  - Call by **address** (a variant of call-by-reference)

# Procedural Units (cont.)

- **Passing parameters**
  - Call by **value** (passed by value)

a. When the function is called, a copy of the data is given to the function



b. and the function manipulates its copy.



```
void Test(int v)
{
    v = 6;
}
```

```
int main()
{
    int val = 5;
    Test(val);
    std::cout << val << std::endl; 5
}
```

c. Thus, when the function has terminated, the calling environment has not been changed.



# Procedural Units (cont.)

- **Passing parameters**
  - Call by **address** (passed by address)

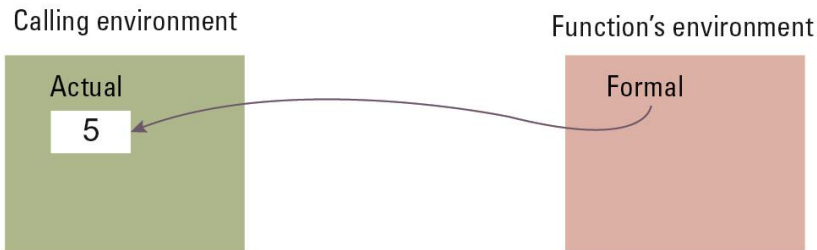
```
void Test(int *v)
{
    *v = 6;
}
```

```
int main()
{
    int val = 5;
    Test(&val);
    std::cout << val << std::endl; 6
}
```

# Procedural Units (cont.)

- **Passing parameters**
  - Call by **reference** (passed by reference)

a. When the function is called, the formal parameter becomes a reference to the actual parameter.



b. Thus, changes directed by the function are made to the actual parameter



c. and are, therefore, preserved after the function has terminated.



```
void Test(int &ref_v)
{
    ref_v = 6;
}
```

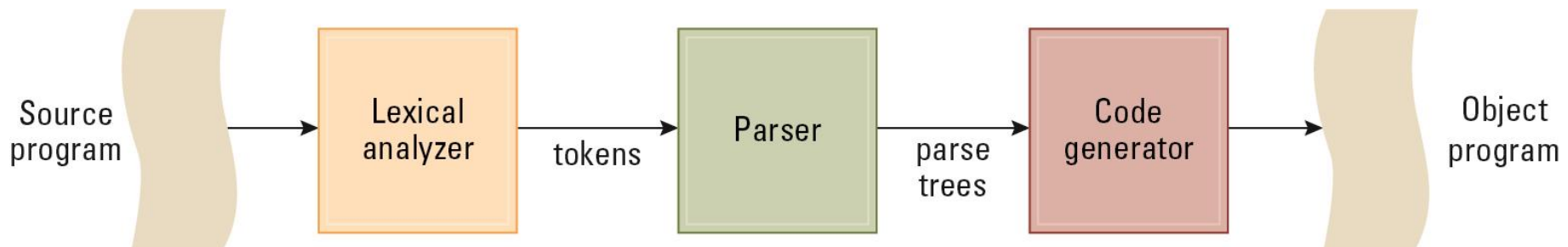
```
int main()
{
    int val = 5;
    Test(val);
    std::cout << val << std::endl; 6
}
```

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# Language Translation Process

- Converting a program written in a high-level language into a machine-executable form



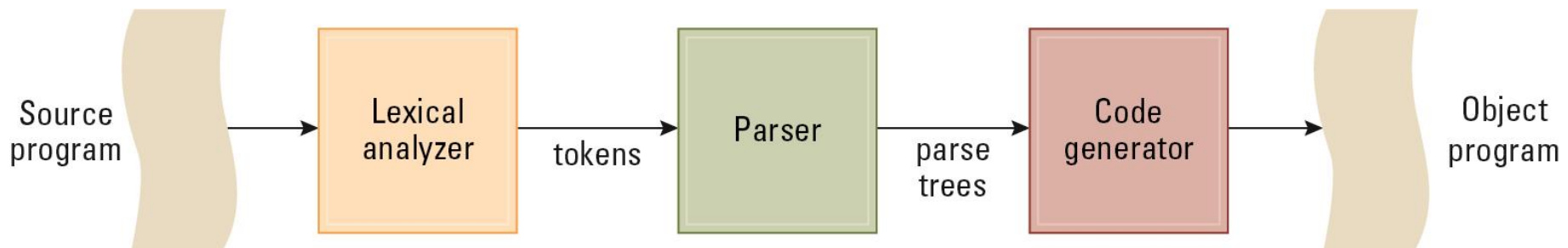
- **Lexical Analyzer:** recognize which strings of symbols represent a single entity, or token (identify tokens)

**y** **=** **x** **+** **1000**

**tokens**

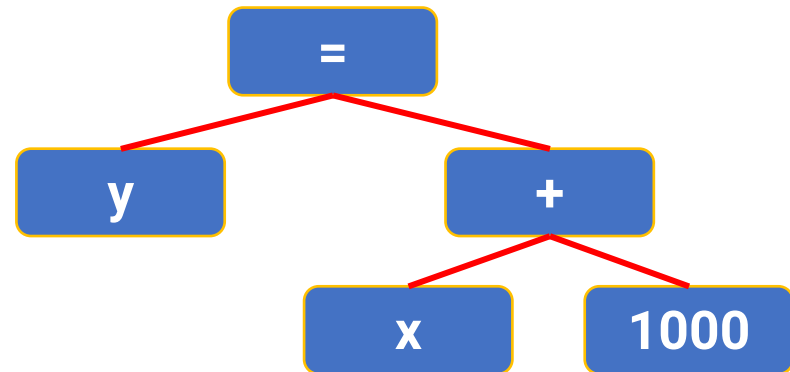
# Language Translation Process (cont.)

- Converting a program written in a high-level language into a machine-executable form



- Parser:** group tokens into statements, using syntax diagrams to make parse trees (identify syntax)

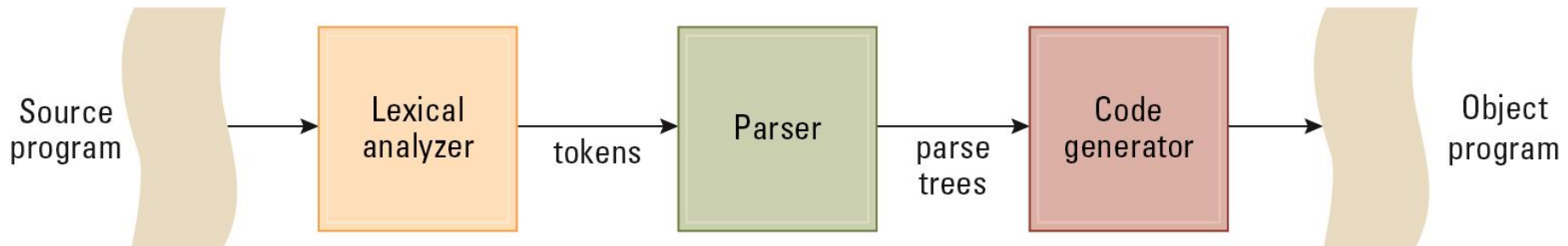
**y = x + 1000**  
tokens





# Language Translation Process (cont.)

- Converting a program written in a high-level language into a machine-executable form



- **Code Generator:** construct machine-language instructions to implement the statements
  - Link libraries

# Syntax Grammar for Algebra

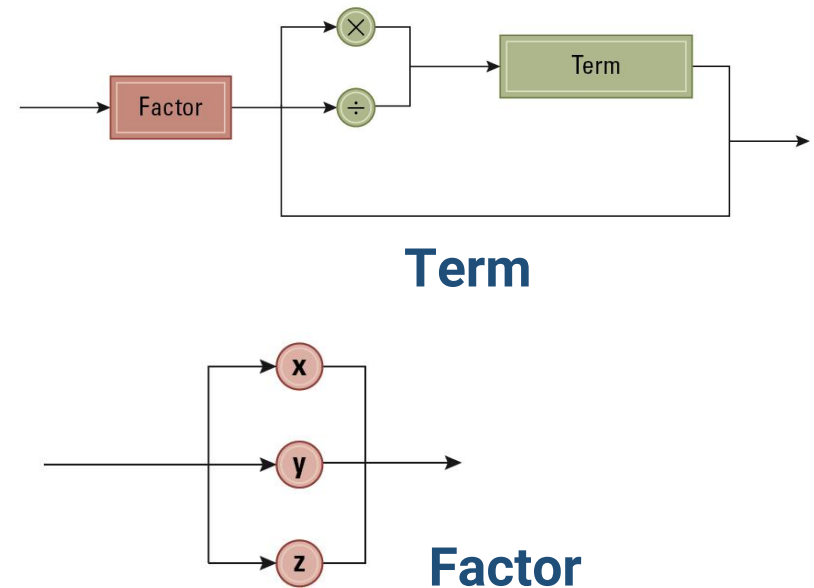
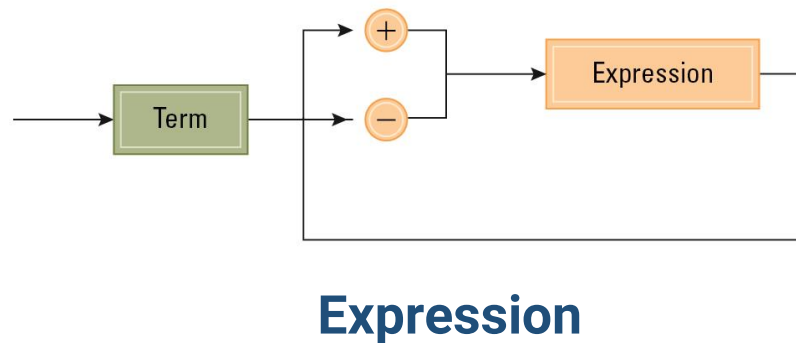
- A simple syntax grammar:

**Expression**  $\rightarrow$  **Term** | **Term** + **Expression**  
 | **Term** – **Expression**

**Term**  $\rightarrow$  **Factor** | **Factor** \* **Term** | **Factor** / **Term**

**Factor**  $\rightarrow$  **x** | **y** | **z**

**expression, term, factor: terminals**  
**x, y, z, +, -, \*, / : non-terminals**



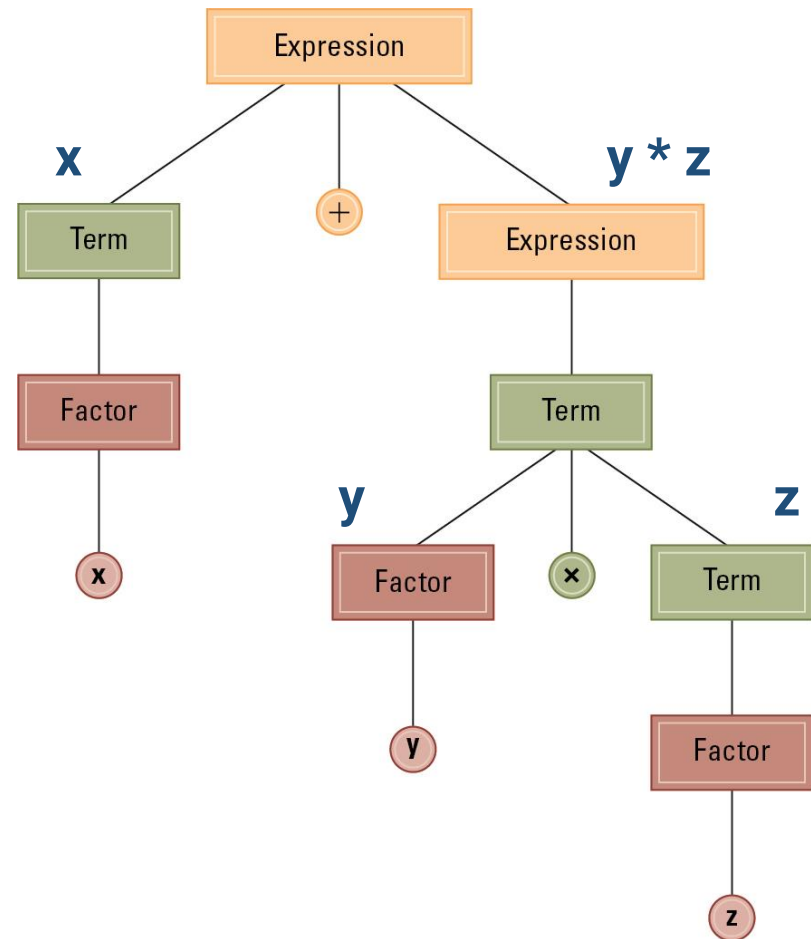
# Syntax Grammar for Algebra (cont.)

- Example: is  $x + y * z$  an expression?

$\text{Expression} \rightarrow \text{Term} \mid$   
 $\text{Term} + \text{Expression} \mid$   
 $\text{Term} - \text{Expression}$

$\text{Term} \rightarrow \text{Factor} \mid$   
 $\text{Factor} * \text{Term} \mid$   
 $\text{Factor} / \text{Term}$

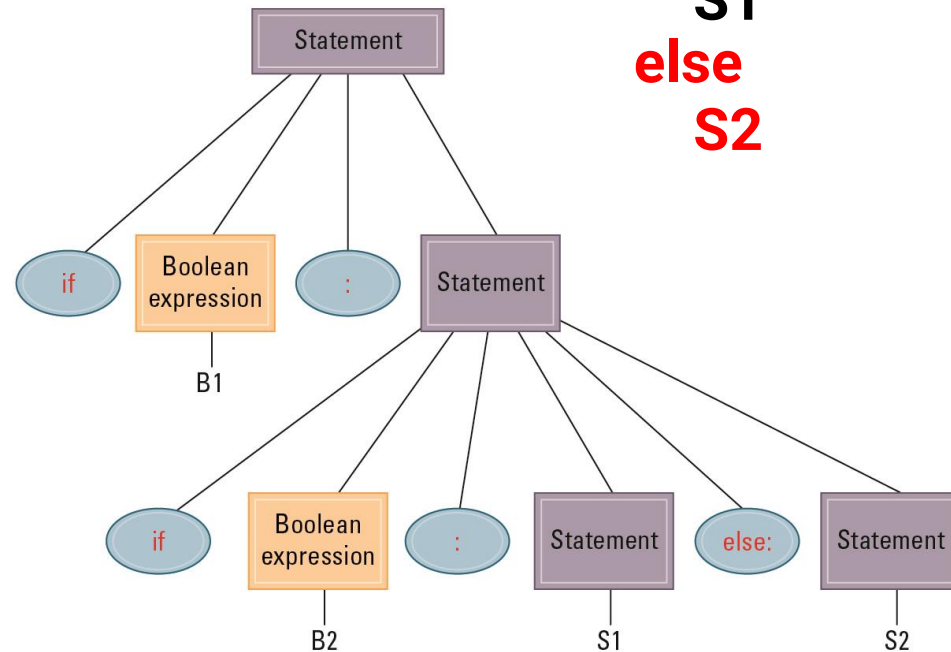
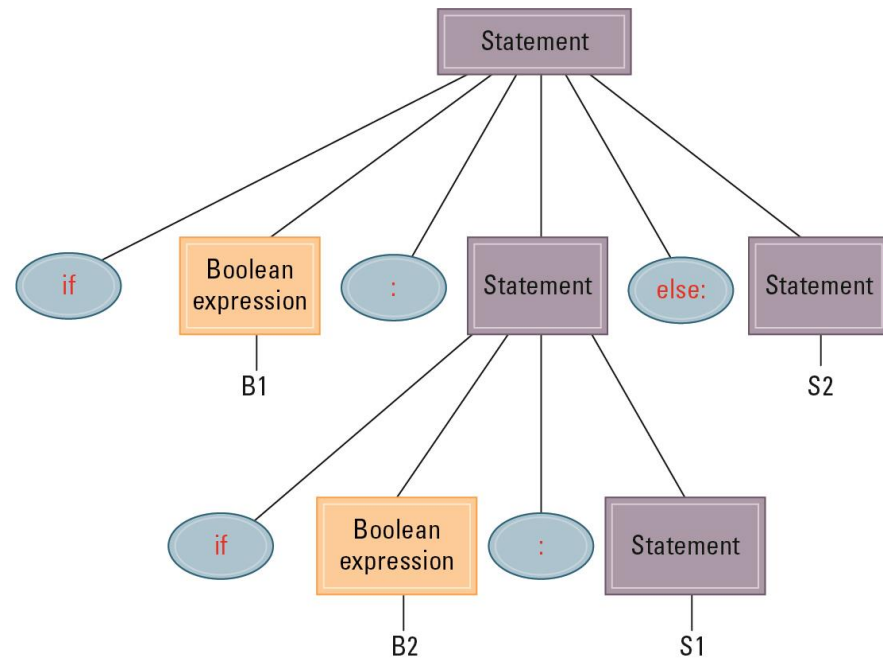
$\text{Factor} \rightarrow x \mid$   
 $y \mid$   
 $z$



# Ambiguity

- if B1 then if B2 then S1 else S2

if (B1)  
if (B2)  
S1  
else  
S2



# Code Generation

- **Coercion:** implicit conversion between data types
- **Strongly typed**
  - No coercion, data types must agree with each other
  - Handle type conversion by programmers
- **Code optimization**

**x = y + z;**

**w = x + z;**

 **w = y + (z << 1);**

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# Object-Oriented Programming

- **Object**
  - Active program unit containing both data and procedures
- **Class**
  - A template from which objects are constructed
- An object is called an **instance** of the class.

# Components of an Object

- **Instance variable (member variable)**
  - Variable within an object
  - Holds information within the object
- **Method (member function)**
  - Procedure within an object
  - Describes the actions that the object can perform
- **Constructor**
  - Special method used to initialize a new object when it is first constructed
- **Destructor v.s. garbage collection**



# Components of an Object (cont.)

- An example of Class

```
class LaserClass
{ int RemainingPower;
  LaserClass(InitialPower)
  { RemainingPower = InitialPower;
  }
  void turnRight()
  { ... }
  void turnLeft()
  { ... }
  void fire()
  { ... }
}
```

Constructor assigns a value to RemainingPower when an object is created.

# Object Integrity

- **Encapsulation**

- A way of restricting access to the internal components of an object
- Private, Public, and Protected

Components in the class are designated public or private depending on whether they should be accessible from other program units.

```
class LaserClass
{private int RemainingPower;
public LaserClass (InitialPower)
{RemainingPower = InitialPower;
}
public void turnRight ( )
{...}
public void turnLeft ( )
{...}
public void fire ( )
{...}
}
```

# Additional Object-oriented Concepts

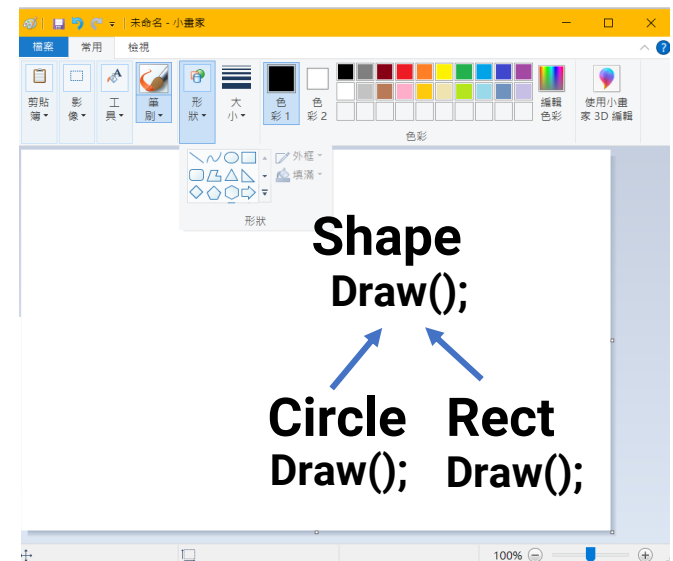
## • Inheritance

- Allows new classes to be defined in terms of previously defined classes

```
class Shape {
public:
    Shape(){}
    ~Shape(){}
    virtual void Draw() = 0;
};

class Circle : public Shape {
public:
    Circle(){}
    ~Circle(){}
    void Draw() { std::cout << "Draw Circle!" << std::endl; }
};

class Rect : public Shape {
public:
    Rect(){}
    ~Rect(){}
    void Draw() { std::cout << "Draw Rect!" << std::endl; }
};
```



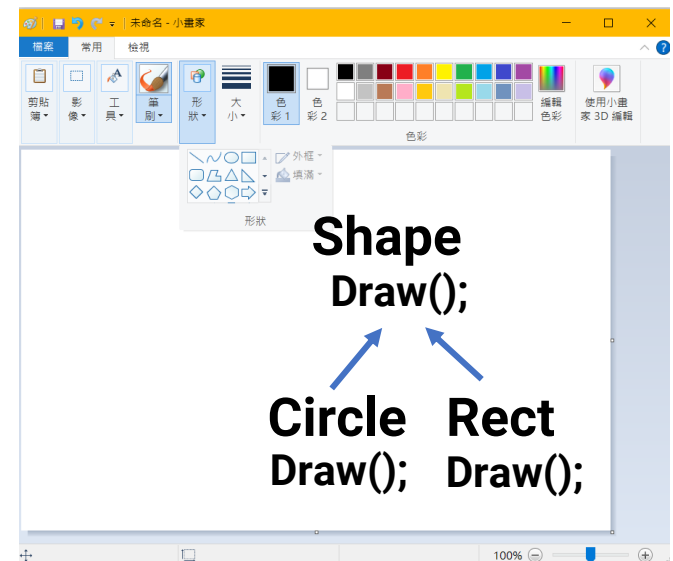
# Additional Object-oriented Concepts

- **Polymorphism**

- Allows method calls to be interpreted by the object that receives the call

```
Shape* shapeList[2];  
shapeList[0] = new Circle();  
shapeList[1] = new Rect();  
for (int i = 0; i < 2; ++i) {  
    shapeList[i]->Draw();  
}
```

```
Draw Circle!  
Draw Rect!
```

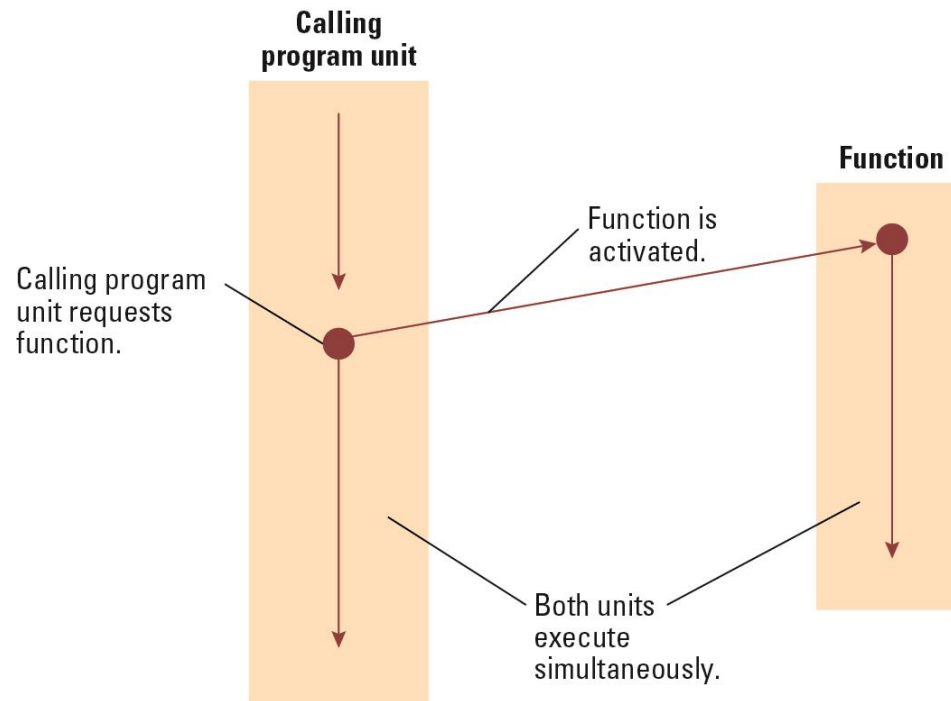


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# Programming Concurrent Activities

- Parallel (or concurrent) processing: simultaneous execution of multiple processes
  - True concurrent processing requires multiple CPUs
  - Can be simulated using time-sharing with a single CPU



**Any Questions?**