

## Replacement Rules

The rules we will be looking at today are not so much inference rules (like the last 8 have been) where we go from a set of premises to a conclusion. The rules for today are known as replacement rules because we will be allowed to re-state a given line in our proof with another wff that is logically equivalent to the original line (remember our definition of equivalent).

9. De Morgan's rule (DM):

$$\sim(p \cdot q) :: (\sim p \vee \sim q)$$

$$\sim(p \vee q) :: (\sim p \cdot \sim q)$$

10. Commutativity (Com):

$$(p \vee q) :: (q \vee p)$$

$$(p \cdot q) :: (q \cdot p)$$

11. Associativity (Assoc):

$$[p \vee (q \vee r)] :: [(p \vee q) \vee r]$$

$$[p \cdot (q \cdot r)] :: [(p \cdot q) \cdot r]$$

12. Distribution (Dist):

$$[p \cdot (q \vee r)] :: [(p \cdot q) \vee (p \cdot r)]$$

$$[p \vee (q \cdot r)] :: [(p \vee q) \cdot (p \vee r)]$$

13. Double negation (DN):

$$p :: \sim\sim p$$

Some examples of how to use these replacement rules in proofs are provided below:

### Example 1:

1. $A \supset \sim(B \cdot C)$	
2. $A \cdot C$	/ $\sim B$
3. $A$	2, Simp
4. $\sim(B \cdot C)$	1, 3, MP
5. $\sim B \vee \sim C$	4, DM
6. $C$	2, Com, Simp
7. $\sim B$	5, 6, Com, DN, DS

### Example 2:

1. $D \cdot (E \vee F)$	
2. $\sim D \vee \sim F$	/ $D \cdot E$
3. $(D \cdot E) \vee (D \cdot F)$	1, Dist
4. $(D \cdot F) \vee (D \cdot E)$	3, Com
5. $\sim(D \cdot F)$	2, DM
6. $D \cdot E$	4, 5, DS