

## Homework 2

1a  $5n^3 + 2n^2 + 3n = O(n^3)$   
 $f(n) \leq c \cdot g(n)$

~~$5n^3 + 2n^2 + 3n \leq c \cdot n^3$~~

~~$5n^3 + 2n^2 + 3n \leq (5+2+3)n^3 = cn^3$~~ , for  $c=10$  when  $n \geq n_0=1$

1b.  $\sqrt{7n^2 + 2n - 8} = O(n)$

$\sqrt{7n^2 + 2n - 8} = O(n)$

$\sqrt{7n^2 + 2n - 8} \leq c \cdot n$

$(7n^2 + 2n - 8)^{1/2} \leq c \cdot n$

$\sqrt{7}n + \sqrt{2}n^{1/2} - \sqrt{8} \leq c \cdot n$

$\sqrt{7}n + \sqrt{2}n^{1/2} - \sqrt{8} \leq (\sqrt{7} + \sqrt{2})n = cn$   $c = \sqrt{7} + \sqrt{2}$

when  $n \geq n_0 = 1$

$\sqrt{7n^2 + 2n - 8} = \Omega(n)$

$\sqrt{7n^2 + 2n - 8} \geq c \cdot n$

~~$\sqrt{7n^2 + 2n - 8} \geq c \cdot n$~~ , for  $c = \sqrt{7}$

$\sqrt{7n^2 + 2n - 8} \geq c \cdot n$  for  $c = \sqrt{7}$

$\sqrt{2}n^{1/2} - \sqrt{8} \geq 0$

$n^{1/2} \geq 2$

$n \geq 4$

$\sqrt{7}n \leq \sqrt{7n^2 + 2n - 8} \leq (\sqrt{7} + \sqrt{2})n$  for  $n \geq 4$

1c  $d(n) = O(f(n))$   $e(n) = O(g(n))$   $d(n)e(n) = O(f(n)g(n))$

$d(n) \leq c_1 \cdot f(n)$   $n \geq n_0$   $e(n) \leq c_2 \cdot g(n)$   $n \geq m_0$   $d(n)e(n) \leq f(n)g(n) \cdot c$

$e(n) \cdot d(n)$

$\rightarrow c_1 \cdot c_2 \cdot f(n) \cdot g(n)$

$n \geq m_0$  or  $n$  will always be greater or

$n \geq n_0$  equal to the initial values therefore

$(d(n)e(n)) \leq f(n)g(n) \cdot c$

## Question 2:

def example1(lst):  $\overline{O(n^2)}$

$\Theta(1)$   $\left[ \begin{array}{l} n = \text{len}(lst) \quad ] \Theta(1) \\ \text{total} = 0 \quad ] \Theta(1) \end{array} \right.$

$\Theta(n^2)$   $\left[ \begin{array}{l} \text{for } j \text{ in range}(n): \\ \quad \Theta(n) \left[ \begin{array}{l} \text{for } k \text{ in range}(1+j): \\ \quad \text{total} += \text{lst}[k] \end{array} \right] \Theta(1) \end{array} \right.$

$\Theta(1)$   $\left[ \begin{array}{l} \text{return total} \end{array} \right] \Theta(1)$

def example3(n):  $\overline{O(\log(n))}$

$\Theta(1)$   $\left[ \begin{array}{l} i = 1 \quad ] \Theta(1) \\ \text{sum} = 0 \quad ] \Theta(1) \end{array} \right.$

$\Theta(\log n)$   $\left[ \begin{array}{l} \text{while } (i < n * n): \\ \quad i *= 2 \quad ] \Theta(1) \\ \quad \text{sum} += i \quad ] \Theta(1) \end{array} \right.$

$\Theta(1)$   $\left[ \begin{array}{l} \text{return sum} \end{array} \right] \Theta(1)$

def example2(lst):  $\overline{O(n)}$

$\Theta(1)$   $\left[ \begin{array}{l} n = \text{len}(lst) \quad ] \Theta(1) \\ \text{prefix} = 0 \quad ] \Theta(1) \\ \text{total} = 0 \quad ] \Theta(1) \end{array} \right.$

$\Theta(n)$   $\left[ \begin{array}{l} \text{for } j \text{ in range}(n) \\ \quad \text{prefix} += \text{lst}[j] \quad ] \Theta(1) \\ \quad \text{total} += \text{prefix} \quad ] \Theta(1) \end{array} \right.$

$\Theta(1)$   $\left[ \begin{array}{l} \text{return total} \end{array} \right] \Theta(1)$

def example4(n):  $\overline{O(n \log(n))}$

$\Theta(1)$   $\left[ \begin{array}{l} i = n \quad ] \Theta(1) \\ \text{sum} = 0 \quad ] \Theta(1) \end{array} \right.$

$\Theta(\log n)$   $\left[ \begin{array}{l} \text{while } (i > 1): \\ \quad \Theta(n) \left[ \begin{array}{l} \text{for } j \text{ in range}(i): \\ \quad \text{sum} += i * j \end{array} \right] \Theta(1) \\ \quad i //= 2 \end{array} \right.$

$\Theta(1)$   $\left[ \begin{array}{l} \text{return sum} \end{array} \right]$



