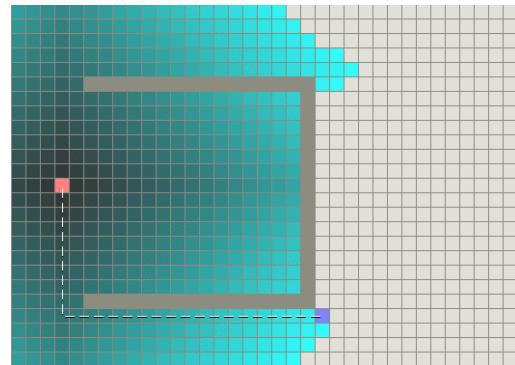
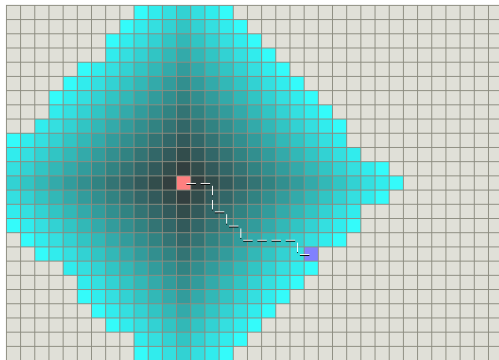
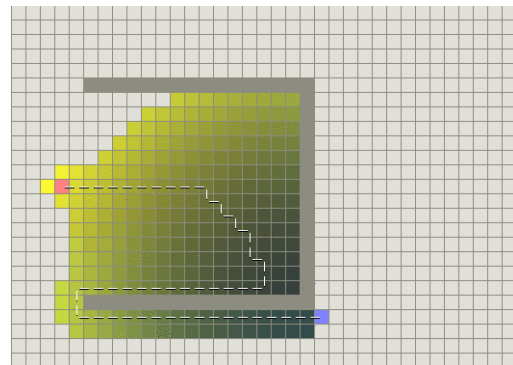
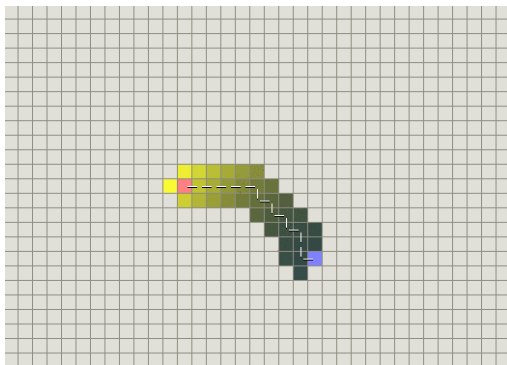


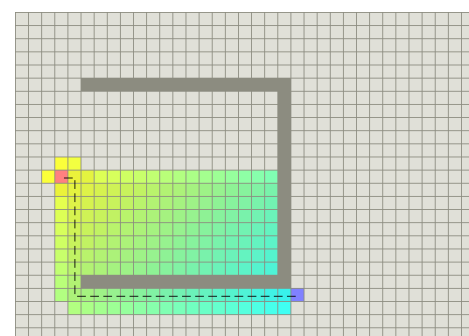
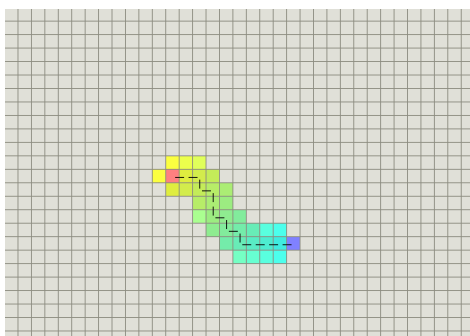
Dijkstra's Algorithm is guaranteed to find a shortest path from the starting point to the goal, as long as *none* of the edges have a negative cost:



The Greedy Best-First-Search algorithm works in a similar way, except that it has some estimate (called a *heuristic*) of how far from the goal any vertex is. It is *not* guaranteed to find a shortest path but runs much quicker:



A\* is like Dijkstra's Algorithm in that it can be used to find a shortest path, and is like Greedy Best-First-Search in that it can use a heuristic to guide itself. In the simple case, it is as fast as Greedy Best-First-Search, in the example with a concave obstacle, A\* finds a path as good as what Dijkstra's Algorithm found.



In short, A\* computes  $f(n) = g(n) + h(n)$ , where  $g(n)$  is the value calculated from starting point to current point, and  $h(n)$  is the heuristic value calculated from current point to the destination point.

- If  $h(n)$  is 0, then only  $g(n)$  plays a role, and A\* turns into Dijkstra's Algorithm.
- If  $h(n)$  is very high relative to  $g(n)$ , then A\* turns into Greedy Best-First-Search.
- If  $h(n) < \text{the cost of moving from } n \text{ to the goal}$ , then A\* is guaranteed to find a shortest path. The lower  $h(n)$  is, the more node A\* expands, making it slower.
- If  $h(n) > \text{the cost of moving from } n \text{ to the goal}$ , then A\* is not guaranteed to find a shortest path, but it can run faster.
- If  $h(n)$  is exactly equal to the cost of moving from  $n$  to the goal, then A\* will only follow the best path and never expand anything else, making it very fast.

### Comparison

Algorithm	Type	Purpose	Time Complexity	Space Complexity	Suitable for
Bellman-Ford	Single-source	Finding shortest paths with negative weights	$O(V * E)$ (with optimizations, $O(V^3)$ )	$O(V)$	Negative weight graphs
Floyd-Warshall	All-pairs	Finding shortest paths between all pairs	$O(V^3)$	$O(V^2)$	Dense graphs, with or without negative weights
Dijkstra	Single-source	Finding shortest paths in weighted graphs	$O((V + E) * \log(V))$	$O(V)$	Non-negative weighted graphs, single-source
A*	Single-source	Finding shortest path to a target	$O(b^d)$ (exponential, worst case)	$O(b^d)$ (for visited nodes)	Graphs with a well-defined heuristic, single-source

"v": number of vertices in the graph.

"E": number of edges in the graph.

"b": average number of edges from each node.

"d": number of nodes on the resulting path.