

# CS70 In Simpler Terms - Note 10

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## 1 RSA

- Encrypt message  $x$ :  $E(X) = x^e \pmod{N}$ ,  $N = pq$  (large primes)
- Decode  $m = E(x)$ :  $D(m) = m^d \pmod{N}$ , where  $d$  is the M.I. of  $e$   
 $(\pmod{()p-1})(q-1) \Rightarrow \gcd(e, (p-1)(q-1)) = 1$  ensures that  $d$  exists  
*since  $e$  is relatively prime with  $(p-q)(q-1)$*
- $D(E(x)) = x \pmod{N} \Rightarrow E$  is a bijection -  $x^{ed} = x \pmod{N}$   
Proof:
  - $ed = 1 \pmod{()p-1)(q-1)} = 1 + (p-1)(q-1)k$
  - $x^{ed} - x = x^{1+k(p-1)(q-1)} - x = x(x^{k(p-1)(q-1)} - 1)$   
We are trying to show that the above statement  $\equiv 0 \pmod{N}$ , because then we will have proven that  $x^{ed} \equiv x \pmod{N}$
  - We can do this by showing that  $x$  is divisibly by  $p$  and  $x$  is divisible by  $q \Rightarrow$  divisible by  $pq$ 
    - \* case 1:  $x$  is a multiple of  $p$ , in which case we are done
    - \* case 2:  $x$  is not a multiple of  $p$  so using FLT  
 $x^{p-1} \equiv 1 \pmod{p} \Rightarrow (x^{p-1})^{k(q-1)} \equiv 1^{k(q-1)} \pmod{p} \equiv 1 \pmod{p} \Rightarrow$   
 $x^{k(p-1)(q-1)} - 1 \equiv 0 \pmod{p} \quad \square$
  - the proof for the case where  $q|x$  is analagous

## 2 Polynomials

- degree  $d$  polynomial has at most  $d$  roots
- need  $d + 1$  points to uniquely determine polynomial of degree  $d$
- Lagrange interpolation - used to find a polynomial of lowest degree, given a set of points

*Formula for calculating the final polynomial  $p(x)$  of degree  $d$*

$$p(x) = \sum_{i=1}^{d+1} y_i \Delta_i(x), \text{ after constructing } d+1 \text{ polynomials } \Delta_i(x)$$

$$\text{Formula for finding the individual polynomials, } \Delta_i(x) = \frac{\prod_{j \neq i} (x - x_j)}{\prod_{j \neq i} (x_i - x_j)}$$

*Example:* Find the polynomial of lowest degree that passes through the points  $(1, 1), (2, 2), (3, 4)$

- $d = 2, x_i = i$
- $i = 1, \Delta_1(x) = \frac{(x-2)(x-3)}{(1-2)(1-3)}$ 
  - \* This makes sense because at  $x = 2, x = 3, \Delta_1(x) = 0$ , and at  $x = 1, \Delta_1(x) = 1$ , so  $y_1\Delta(x) = 1(1) = 1$  which is correct
- $i = 2, \Delta_2(x) = \frac{(x-1)(x-3)}{(2-1)(2-3)}$
- $i = 3, \Delta_3(x) = \frac{(x-1)(x-2)}{(3-1)(3-2)}$

### 3 Finite Fields $F_m$ or $GF(m)$

Restricting the fields to  $(\text{mod } m)$

- $x$  has inverse  $(\text{mod } m)$  iff  $\text{gcd}(m, x) = 1$  so  $m$  is prime
- $x \in \{1, \dots, m-1\}$  all have an inverse  $(\text{mod } m) \Rightarrow$  bijection
- How many polynomials are there of degree at most 2  $(\text{mod } m)$ ?
  - There are 3 coefficients in a degree 2 polynomial
  - Each coefficient has  $m$  values to choose from  $\{0, \dots, m-1\} \Rightarrow m^3$  polynomials

The chart below summarizes the result above:

**Polynomials of degree  $\leq d$  over  $F_m$**

Number of points given	Number of polynomials
$d + 1$	1
$d$	$m$
.	.
.	.
.	.
0	$m^{d+1}$