

CS70 In Simpler Terms - Note 10

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1 RSA

- Encrypt message x : $E(X) = x^e \pmod{N}$, $N = pq$ (large primes)
- Decode $m = E(x)$: $D(m) = m^d \pmod{N}$, where d is the M.I. of e
 $(\pmod{()p-1})(q-1) \Rightarrow \gcd(e, (p-1)(q-1)) = 1$ ensures that d exists
since e is relatively prime with $(p-1)(q-1)$
- $D(E(x)) = x \pmod{N} \Rightarrow E$ is a bijection - $x^{ed} = x \pmod{N}$
Proof:
 - $ed = 1 \pmod{()p-1)(q-1)} = 1 + (p-1)(q-1)k$
 - $x^{ed} - x = x^{1+k(p-1)(q-1)} - x = x(x^{k(p-1)(q-1)} - 1)$
We are trying to show that the above statement $\equiv 0 \pmod{N}$, because then we will have proven that $x^{ed} \equiv x \pmod{N}$
 - We can do this by showing that x is divisibly by p and $q \Rightarrow$ divisible by pq
 - case 1: $p|x$
assuming that $x \not\equiv 0 \pmod{p}$ and using FLT
 $x^{p-1} \equiv 1 \pmod{p} \Rightarrow (x^{p-1})^{k(q-1)} \equiv 1^{k(q-1)} \pmod{p} \equiv 1 \pmod{p} \Rightarrow$
 $x^{k(p-1)(q-1)} - 1 \equiv 0 \pmod{p} \quad \square$
 - the proof for the case where $q|x$ is analagous

2 Polynomials

- degree d polynomial has at most d roots
- need $d + 1$ points to uniquely determine polynomial of degree d
- Lagrange interpolation - used to find a polynomial of lowest degree, given a set of points

Formula for calculating the final polynomial $p(x)$ of degree d

$$p(x) = \sum_{i=1}^{d+1} y_i \Delta_i(x), \text{ after constructing } d+1 \text{ polynomials } \Delta_i(x)$$

$$\text{Formula for finding the individual polynomials, } \Delta_i(x) = \frac{\prod_{j \neq i} (x - x_j)}{\prod_{j \neq i} (x_i - x_j)}$$

Example: Find the polynomial of lowest degree that passes through the points $(1, 1), (2, 2), (3, 4)$

- $d = 2, x_i = i$
- $i = 1, \Delta_1(x) = \frac{(x-2)(x-3)}{(1-2)(1-3)}$
 - * This makes sense because at $x = 2, x = 3, \Delta_1(x) = 0$, and at $x = 1, \Delta_1(x) = 1$, so $y_1\Delta(x) = 1(1) = 1$ which is correct
- $i = 2, \Delta_2(x) = \frac{(x-1)(x-3)}{(2-1)(2-3)}$
- $i = 3, \Delta_3(x) = \frac{(x-1)(x-2)}{(3-1)(3-2)}$

3 Finite Fields F_m or $GF(m)$

Restricting the fields to $(\text{mod } m)$

- x has inverse $(\text{mod } m)$ iff $\text{gcd}(m, x) = 1$ so m is prime
- $x \in \{1, \dots, m-1\}$ all have an inverse $(\text{mod } m) \Rightarrow$ bijection
- How many polynomials are there of degree at most 2 $(\text{mod } m)$?
 - There are 3 coefficients in a degree 2 polynomial
 - Each coefficient has m values to choose from $\{0, \dots, m-1\} \Rightarrow m^3$ polynomials

The chart below summarizes the result above:

Polynomials of degree $\leq d$ over F_m

Number of points given	Number of polynomials
$d + 1$	1
d	m
.	.
.	.
.	.
0	m^{d+1}