

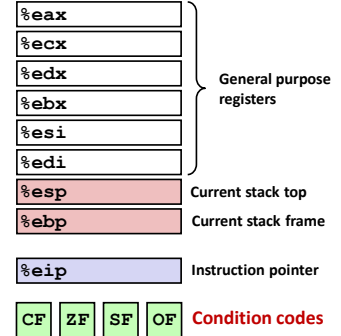
Lecture 13

Control

CSPC 275
Introduction to Computer Systems

Processor State

- Information about currently executing program
 - Temporary data (`%eax, ...`)
 - Location of runtime stack (`%ebp, %esp`)
 - Location of current code control point (`%eip, ...`)
 - Status of recent tests (`CF, ZF, SF, OF`)



Condition Codes (Implicit Setting)

- Single bit registers
 - CF** Carry Flag (for unsigned) **SF** Sign Flag (for signed)
 - ZF** Zero Flag **OF** Overflow Flag (for signed)
- Implicitly set by arithmetic operations
 - Example: `addl src, dest` \leftrightarrow `t = a + b`
 - CF set** if carry out from most significant bit (unsigned overflow)
 - ZF set** if `t == 0`
 - SF set** if `t < 0` (as signed)
 - OF set** if two's-complement (signed) overflow
(`a > 0 && b > 0 && t < 0`) || (`a < 0 && b < 0 && t >= 0`)
- Flags are *not* set by `leal`, `inc`, or `dec` instructions.

Condition Codes (Explicit Setting)

- Explicit Setting by Compare Instruction
 - `cmpl src2, src1`
 - (like computing `src1 - src2` without setting destination)
 - CF set** if carry out from most significant bit (used for unsigned comparisons)
 - ZF set** if `src1 == src2`
 - SF set** if `src1 - src2 < 0` (as signed)
 - OF set** if two's-complement (signed) overflow
(`src1 > 0 && src2 < 0 && (src1 - src2) < 0`) ||
(`src1 < 0 && src2 > 0 && (src1 - src2) > 0`)

Condition Codes (Explicit Setting)

- Explicit Setting by Test instruction
 - `testl src2, src1`
 - Sets condition codes based on value of `src1` & `src2`
 - Like computing `src1 & src2` without setting destination
 - Useful to have one of the operands be a *mask*
 - ZF set** when `src1 & src2 == 0`
 - SF set** when `src1 & src2 < 0`
- `testl %eax, %eax` (??)
- See whether `%eax` is negative, zero, or positive.

Jumping

- `jx` Instructions
 - Jump to different part of code depending on condition codes

| <code>jx</code> | Condition | Description |
|------------------|---------------------------------|---------------------------|
| <code>jmp</code> | 1 | Unconditional |
| <code>je</code> | <code>ZF</code> | Equal / Zero |
| <code>jne</code> | <code>~ZF</code> | Not Equal / Not Zero |
| <code>js</code> | <code>SF</code> | Negative |
| <code>jns</code> | <code>~SF</code> | Nonnegative |
| <code>jg</code> | <code>~(SF^OF) & ~ZF</code> | Greater (Signed) |
| <code>jge</code> | <code>~(SF^OF)</code> | Greater or Equal (Signed) |
| <code>jl</code> | <code>(SF^OF)</code> | Less (Signed) |
| <code>jle</code> | <code>(SF^OF) ZF</code> | Less or Equal (Signed) |
| <code>ja</code> | <code>~CF & ~ZF</code> | Above (unsigned) |
| <code>jb</code> | <code>CF</code> | Below (unsigned) |

Conditional Branch Example

C Code

```
int absdiff(int x, int y)
{
    int result;
    if (x > y)
        result = x - y;
    else
        result = y - x;
    return result;
}
```

Goto Version

```
int goto_ad(int x, int y)
{
    int result;
    if (x <= y) goto Else;
    result = x - y;
    goto Exit;
Else:
    result = y - x;
Exit:
    return result;
}
```

- C allows “goto” as means of transferring control
 - Closer to machine-level programming style
- Generally considered bad coding style

Conditional Branch Example (Cont.)

```
absdiff:
    pushl %ebp
    movl %esp, %ebp
    movl 8(%ebp), %edx
    movl 12(%ebp), %eax
    cmpl %eax, %edx
    jle .L6
    subl %eax, %edx
    movl %edx, %eax
    jmp .L7
.L6:
    subl %edx, %eax
.L7:
    popl %ebp
    ret
```

Labels and instructions are grouped into sections:

- Setup: `pushl %ebp`, `movl %esp, %ebp`
- Body1: `movl 8(%ebp), %edx`, `movl 12(%ebp), %eax`, `cmpl %eax, %edx`
- Body2a: `jle .L6`, `subl %eax, %edx`, `movl %edx, %eax`
- Body2b: `subl %edx, %eax`
- Finish: `popl %ebp`, `ret`

Note that `%eax` contains the return value.

“Do-While” Loop Example

C Code

```
int pcount_do(unsigned x)
{
    int result = 0;
    do {
        result += x & 0x1;
        x >>= 1;
    } while (x);
    return result;
}
```

Goto Version

```
int pcount_do(unsigned x)
{
    int result = 0;
loop:
    result += x & 0x1;
    x >>= 1;
    if (x)
        goto loop;
    return result;
}
```

“Do-While” Loop Compilation

Goto Version

```
int pcount_do(unsigned x) {
    int result = 0;
loop:
    result += x & 0x1;
    x >>= 1;
    if (x)
        goto loop;
    return result;
}
```

Registers:

| Register | Variable |
|-------------------|---------------------|
| <code>%edx</code> | <code>x</code> |
| <code>%ecx</code> | <code>result</code> |

```
movl $0, %ecx # result = 0
.L2:
movl %edx, %eax # t = x & 1
andl $1, %eax
addl %eax, %ecx # result += t
shrl %edx # x >>= 1
jne .L2 # If !0, goto loop
```

“While” Loop Example

C Code

```
int pcount_while(unsigned x) {
    int result = 0;
    while (x) {
        result += x & 0x1;
        x >>= 1;
    }
    return result;
}
```

Goto Version

```
int pcount_do(unsigned x) {
    int result = 0;
    if (!x) goto done;
loop:
    result += x & 0x1;
    x >>= 1;
    if (x)
        goto loop;
done:
    return result;
}
```

General “While” Translation

While version

```
while (Test)
    Body
```

Do-While Version

```
if (!Test)
    goto done;
do
    Body
while (Test);
done:
```

Goto Version

```
if (!Test)
    goto done;
loop:
    Body
    if (Test)
        goto loop;
done:
```

Practice Problems

- Read CSaPP Sec. 3.6.1-3.6.5 and try the following problems:
3.16, 3.17, 3.18, 3.22, 3.23