

Lecture 18

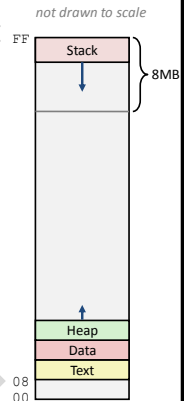
Buffer Overflow

CSPC 275
Introduction to Computer Systems

IA32 Linux Memory Layout

- Stack
 - Runtime stack (8MB limit)
 - e.g., local variables
- Heap
 - Dynamically allocated storage
 - When call `malloc()`, `calloc()`, `new()`
- Data
 - Statically allocated data
 - e.g., arrays & strings declared in code
- Text
 - Executable machine instructions
 - Read-only

Upper 2 hex digits
= 8 bits of address



Memory Allocation Example

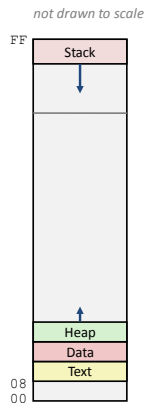
```
char big_array[1<<24]; /* 16 MB */
char huge_array[1<<28]; /* 256 MB */

int beyond;
char *p1, *p2, *p3, *p4;

int useless() { return 0; }

int main()
{
    p1 = malloc(1 << 28); /* 256 MB */
    p2 = malloc(1 << 8); /* 256 B */
    p3 = malloc(1 << 28); /* 256 MB */
    p4 = malloc(1 << 8); /* 256 B */
    /* Some print statements ... */
}
```

Where does everything go?

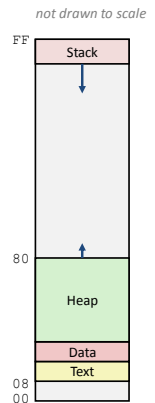


IA32 Example Addresses

address range $\sim 2^{32}$

\$esp	0xfffffbc0
p3	0x65586008
p1	0x55585008
p4	0x1904a110
p2	0x1904a008
&p2	0x18049760
&beyond	0x08049744
big_array	0x18049780
huge_array	0x08049760
main()	0x080483c6
useless()	0x08049744
malloc()	0x006be166

`malloc()` is dynamically linked
address determined at runtime



Internet Worm

- November, 1988
 - Internet Worm attacks thousands of Internet hosts.
 - How did it happen?
- It was based on *stack buffer overflow* exploits!
 - Many library functions do not check argument sizes.
 - Allows target buffers to overflow.

String Library Code

- Implementation of C Standard Library function `gets()`

```
/* Get string from stdin */
char *gets(char *dest)
{
    int c = getchar();
    char *p = dest;
    while (c != EOF && c != '\n') {
        *p++ = c;
        c = getchar();
    }
    *p = '\0';
    return dest;
}
```

- No way to specify limit on number of characters to read

String Library Code (con't)

- Similar problems with other library functions
 - strcpy, strcat**: Copy strings of arbitrary length
 - scanf, fscanf, sscanf**, when given **%s** conversion specification

Vulnerable Buffer Code

```
/* Echo Line */
void echo()
{
    char buf[4]; /* Way too small! */
    gets(buf);
    puts(buf);
}
```

```
void call_echo() {
    echo();
}
```

```
$ ./bufdemo
Type a string:123
```

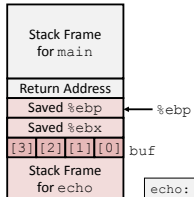
```
$ ./bufdemo
Type a string:1234567
Segmentation Fault
```

```
$ ./bufdemo
Type a string:12345678
Segmentation Fault
```

```
$ ./bufdemo
Type a string:123456789ABC
Segmentation Fault
```

Buffer Overflow Stack

Before call to gets



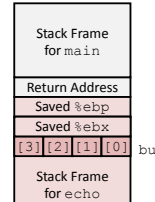
```
/* Echo Line */
void echo()
{
    char buf[4]; /* Way too small! */
    gets(buf);
    puts(buf);
}
```

```
echo:
    pushl %ebp          # Save %ebp on stack
    movl %esp, %ebp
    pushl %ebx          # Save %ebx
    subl $20, %esp      # Allocate stack space
    leal -8(%ebp), %ebx # Compute buf as %ebp-8
    movl %ebx, (%esp)   # Push buf on stack
    call gets           # Call gets
    . . .
```

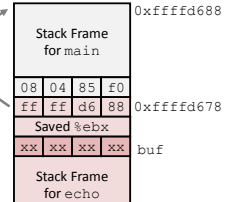
Buffer Overflow Stack Example

```
$ gdb bufdemo
(gdb) break echo
Breakpoint 1 at 0x80485c9
(gdb) run
Breakpoint 1, 0x80485c9 in echo ()
(gdb) print /x %ebp
$1 = 0xffffd678
(gdb) print /x *((unsigned *)$ebp
$2 = 0xffffd688
(gdb) print /x *((unsigned *)$ebp + 1)
$3 = 0x80485f0
```

Before call to gets



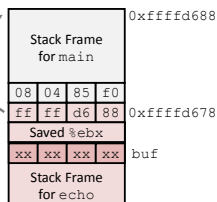
Before call to gets



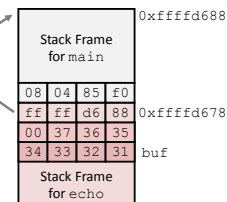
```
80485eb: e8 d5 ff ff ff call 80485c5 <echo>
80485f0: c9 leave
```

Buffer Overflow Example #1

Before call to gets



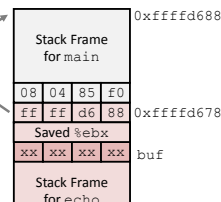
Input 1234567



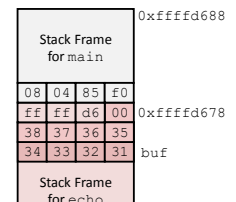
Overflow buf and corrupt %ebx

Buffer Overflow Example #2

Before call to gets



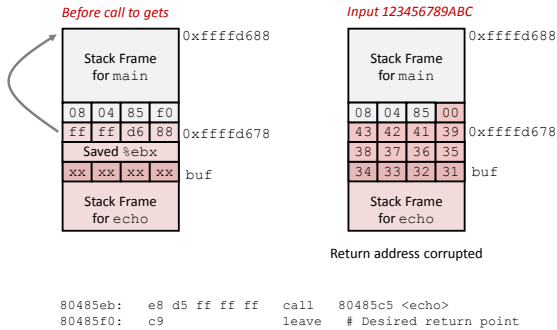
Input 12345678



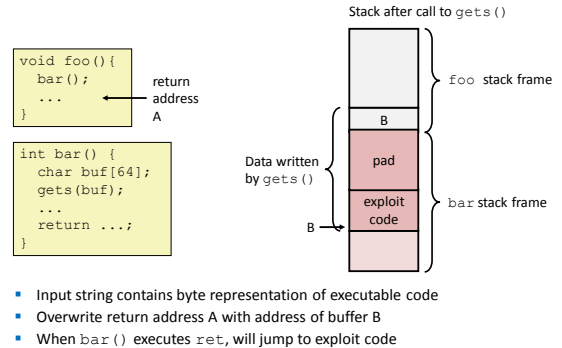
Base pointer corrupted

```
. . .
80485eb: e8 d5 ff ff ff call 80485c5 <echo>
80485f0: c9 leave # Set %ebp to corrupted value
80485f1: c3 ret
```

Buffer Overflow Example #3



Malicious Use of Buffer Overflow



Exploits Based on Buffer Overflows

- Buffer overflow bugs allow remote machines to execute arbitrary code on victim machines
- Internet worm
 - Early versions of the finger server (**fingerd**) used **gets ()** to read the argument sent by the client:
 - finger pyoon@lab.cs.trincoll.edu**
 - Worm attacked fingerd server by sending phony argument:
 - finger "exploit-code padding new-return-address"**
 - exploit code: executed a root shell on the victim machine with a direct TCP connection to the attacker.

Avoiding Overflow Vulnerability

```

/* Echo Line */
void echo()
{
    char buf[4]; /* Way too small! */
    fgets(buf, 4, stdin);
    puts(buf);
}
  
```

- Use library routines that limit string lengths
 - fgets** instead of **gets**
 - strncpy** instead of **strcpy**
 - Don't use **scanf** with **%s** conversion specification
 - Use **fgets** to read the string
 - Or use **%ns** where **n** is a suitable integer

System-Level Protections

- Randomized stack offsets
 - At start of program, allocate random amount of space on stack
 - Makes it difficult for hacker to predict beginning of inserted code
- Nonexecutable code segments
 - In traditional x86, can mark region of memory as either "read-only" or "writeable"
 - Can execute anything readable
 - X86-64 added explicit "execute" permission

```

$ gdb bufdemo
(gdb) break echo

(gdb) run
(gdb) print /x $ebp
$1 = 0xfffffc638

(gdb) run
(gdb) print /x $ebp
$2 = 0xfffffb08

(gdb) run
(gdb) print /x $ebp
$3 = 0xffffc6a8
  
```

Stack Canaries

- Idea
 - Place special value ("canary") on stack just beyond buffer
 - Check for corruption before exiting function
- gcc Implementation
 - fstack-protector**
 - fstack-protector-all**

```

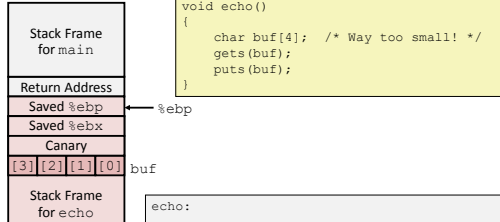
$ ./bufdemo-protected
Type a string:123
123
  
```

```

$ ./bufdemo-protected
Type a string:12345
*** stack smashing detected ***
  
```

Setting Up Canary

Before call to gets

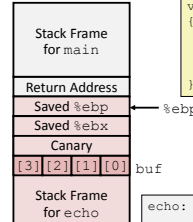


```
/* Echo Line */
void echo()
{
    char buf[4]; /* Way too small! */
    gets(buf);
    puts(buf);
}
```

```
echo:
    . . .
    movl    %gs:20, %eax    # Get canary
    movl    %eax, -8(%ebp)  # Put on stack
    xorl    %eax, %eax      # Erase canary
    . . .
```

Checking Canary

Before call to gets

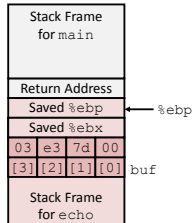


```
/* Echo Line */
void echo()
{
    char buf[4]; /* Way too small! */
    gets(buf);
    puts(buf);
}
```

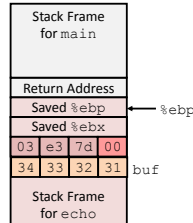
```
echo:
    . . .
    movl    -8(%ebp), %eax    # Retrieve from stack
    xorl    %gs:20, %eax      # Compare with Canary
    je      .L24              # Same: skip ahead
    call    __stack_chk_fail  # ERROR
.L24:
    . . .
```

Canary Example

Before call to gets



Input 1234



```
(gdb) break echo
(gdb) run
(gdb) stepi 3
(gdb) print /x *((unsigned *) $ebp - 2)
$1 = 0x3e37d00
```

Benign corruption!
(allows programmers to make
silent off-by-one errors)

Worms and Viruses

- Worm: A program that
 - Can run by itself
 - Can propagate a fully working version of itself to other computers
- Virus: Code that
 - Add itself to other programs
 - Cannot run independently
- Both are (usually) designed to spread among computers and to wreak havoc

Practice Problems

- Read CSaPP Sec. 3.12 and try the following problems:
3.43, 3.44 and 3.45