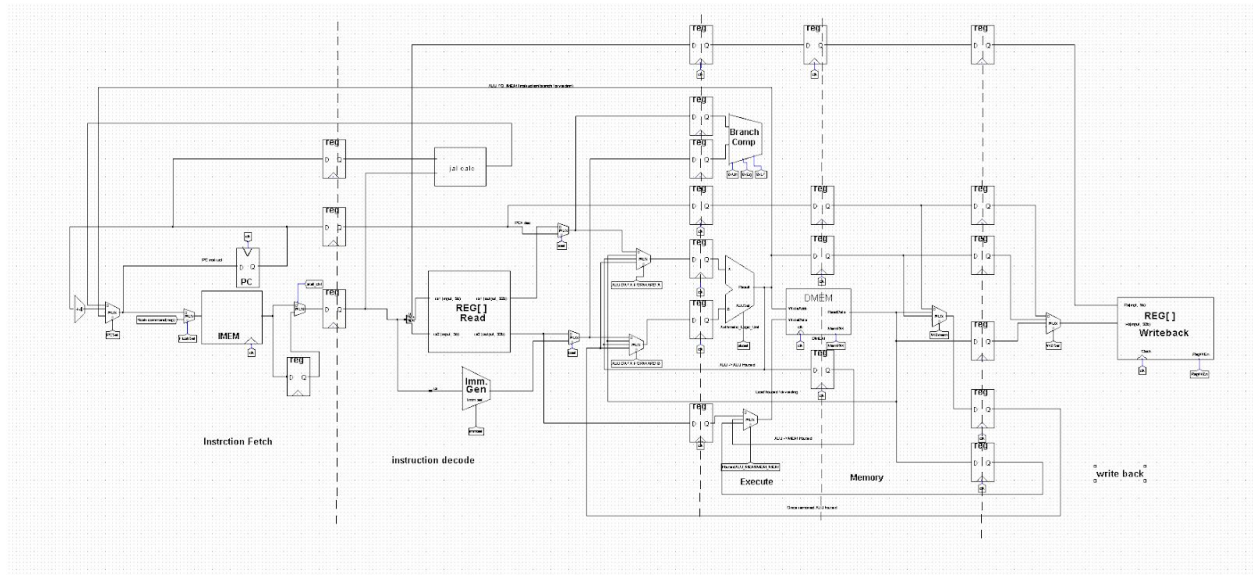


Final Report

EECS 151 ASIC Lab

Ryan Thornton and Kevin Jung

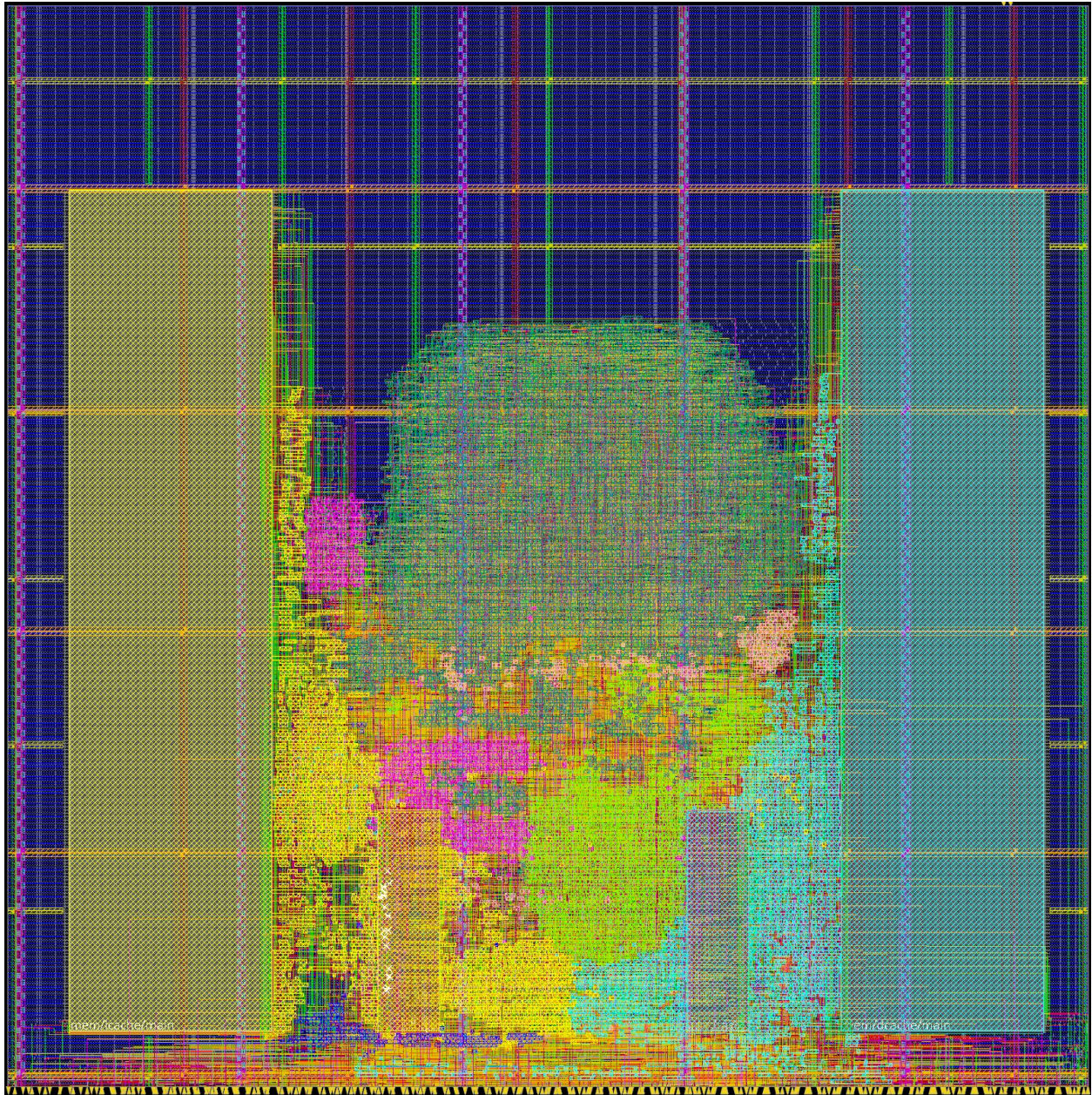
1) Final Pipeline Diagram:



2) Post-Synthesis Critical Path:

The post-synthesis critical path (with a target clock period of 0.6ns) is from the instruction IF/ID pipeline register through the JAL-specific adder through the PC selector mux to the icache request address buffer register. This route has a positive slack of 16ps.

3) Final Floorplan:



Yellow is instruction cache

Light blue is data cache

Red (barely visible, inside data cache) is memory arbiter

Dark blue at bottom is writeback stage

Pink is instruction fetch stage

Darker green is instruction decode stage

Light green is execute stage

Orange is memory stage

External facing modules (i\$ and d\$) are as close to the pins as possible to try to minimize delay

4) Post-PAR Critical Path:

The reported post-PAR critical path is from cpu/Bsel_pip/q_reg[4]/CLK to cpu/id_stage/haz/rd_log_ex/CLKGATE_RC_CG_HIER_INST6/RC_CGIC_INST/ENA.

This path travels through the ALU, into the data cache, through the data cache's stall logic, into the hazard controller in instruction decode. Interestingly, in this iteration, this post-route path fails the synthesis target clock of 600ps, with a 0.721ps negative slack. However, optimizing this path to meet the timing requirements has absolutely no impact on the clock speed that the CPU can be simulated at.

In reality, Innovus shows, with the command **report_timing -unconstrained -from mem/dcachelmem_req_valid -to mem_req_valid**:

```
#####
# Generated by:      Cadence Innovus 18.10-p002_1
# OS:                Linux x86_64(Host ID eda-6.EECS.Berkeley.EDU)
# Generated on:      Mon May 11 22:01:13 2020
# Design:            riscv_top
# Command:           report_timing -unconstrained -from
mem/dcachelmem_req_valid -to mem_req_valid
#####
```

Path 1:

```
View: PVT_0P63V_100C.setup_view
Startpoint: (R) mem/dcachelreq_address/q_reg[21]/CLK
Clock: (R) clk
Endpoint: (R) mem_req_valid
Clock:
```

	Capture	Launch
Src Latency:+	0.000	-85.256
Net Latency:+	0.000 (I)	89.700 (P)
Arrival:=	0.000	4.444

```
Launch Clock:= 4.444
Data Path:+ 903.100
```

```
#-----
-----
-
# Timing Point
Edge Cell Fanout Trans Flags Arc Delay Arrival
#
(ps) (ps) (ps)
```

```

#-----
-----
-
mem/dcache/req_address/q_reg[21]/CLK - CLK
R (arrival) 30 57.100 - 4.444
mem/dcache/req_address/q_reg[21]/Q - CLK->Q
F DFFHQx4_ASAP7_75t_SL 1 57.100 46.800 51.244
mem/dcache/req_address/FE_PHC3507_FE_RN_3/Y - A->Y
F HB2xp67_ASAP7_75t_SL 1 6.200 20.100 71.344
mem/dcache/req_address/FE_PHC1862_req_addr_21/Y - A->Y
F HB3xp67_ASAP7_75t_SRAM 3 12.100 83.500 154.844
mem/dcache/g8402/Y - B->Y
F XOR2xp5_ASAP7_75t_SL 1 52.100 18.500 173.344
mem/dcache/g8361/Y - E->Y
F OR5x1_ASAP7_75t_SL 2 27.800 33.100 206.444
mem/dcache/g8309/Y - A2->Y
R O2A101Ixp33_ASAP7_75t_SL 2 20.400 21.200 227.644
mem/dcache/g8307/Y - B->Y
F NAND2xp5_ASAP7_75t_SL 2 44.500 14.100 241.744
mem/dcache/g8306/Y - B->Y
R OAI21xp5_ASAP7_75t_SL 1 26.800 9.900 251.644
mem/dcache/FE_PHC5254_next_state/Y - A->Y
R HB2xp67_ASAP7_75t_SL 4 22.500 35.900 287.544
mem/dcache/g8305/Y - B->Y
F OAI21x1_ASAP7_75t_SL 3 43.700 17.300 304.844
mem/dcache/FE_OCPC1491_n_256/Y - A->Y
F HB1xp67_ASAP7_75t_SL 1 30.100 25.800 330.644
mem/dcache/FE_OCPC1488_n_256/Y - A->Y
R INVx1_ASAP7_75t_SL 2 24.700 11.300 341.944
mem/dcache/FE_OFC779_n_256/Y - A->Y
F INVx1_ASAP7_75t_SL 22 17.200 41.300 383.244
mem/dcache/g8295/Y - B->Y
R NAND2xp5_ASAP7_75t_SL 2 79.000 22.500 405.744
mem/dcache/mem_req_valid
cache_1 - - <<< mem_req_valid R
- 405.744
mem/arbiter/g544/Y - B->Y
R OR2x2_ASAP7_75t_SRAM 1 37.000 47.100 452.844
FE_OFC363_mem_req_valid/Y - A->Y
R BUFx24_ASAP7_75t_SL 1 40.300 454.700 907.544
mem_req_valid <<< mem_req_valid R
- 1 1369.700 438.700 907.544

```

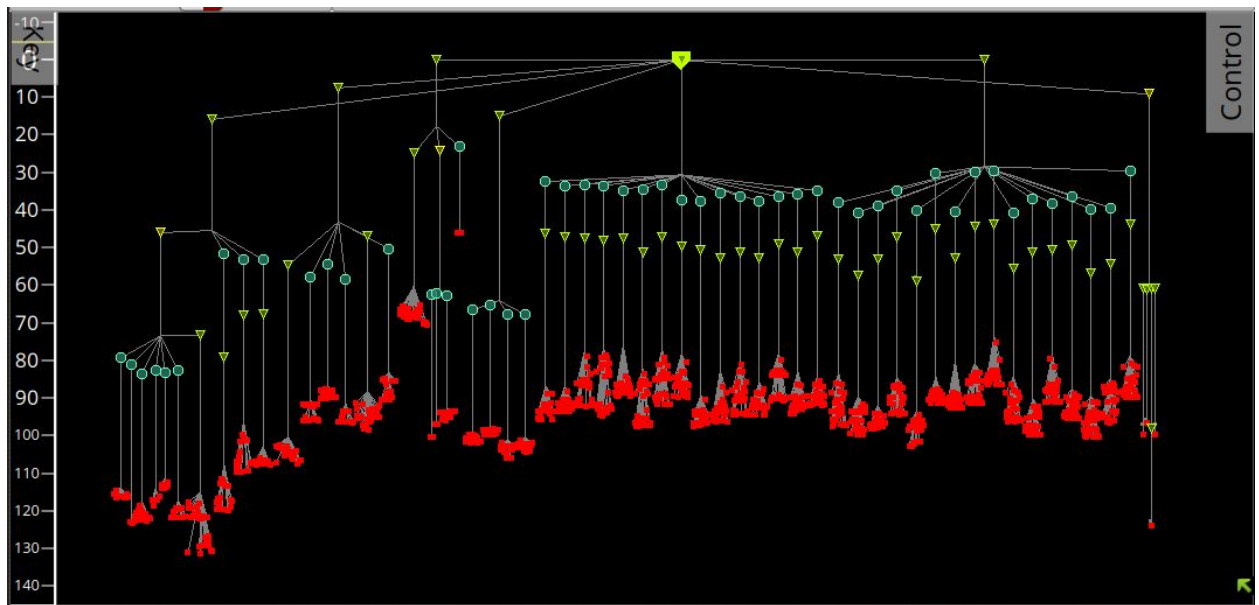
#-----

-

From this report, the final route from the memory arbiter to the off-chip ExtMemModel delays the signal by nearly the entire clock period. In an older PAR iteration with a target clock of 1ns, this single jump from mem/arbiter to mem_req_valid took 1.2ns.

This delay vastly limits the overall speed of the system. Despite completing a successful synthesis at 0.5ns, and having a PAR tapeout that met timing requirements at 0.6ns, the best simulated clock period is at 1.0ns. The overall delay was slightly reduced by forcing the cache logic and memory arbiter to be as close as possible to the output pins, but in the best case the delay from arbiter to external memory was near or larger than the target clock period.

5) Final Clock Tree:



insertion delay max: 82ps

skew max: 55ps

6) area used

The final design used 69192.202752 μm^2 .

The majority of the CPU area was from the ID stage ($6891\mu\text{m}^2$), as the register file is spacious and the hazard control logic is complex. The caches took up double the area of the CPU, with $46293.437952\mu\text{m}^2$ combined.

7) Cycles and Optimizations:

Bmark Short:

[PASSED]

/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/bmark_short_output/cachetest.out
() after 108252 simulation cycles

```

    [ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/bmark_short_output/final.out
() after 9004 simulation cycles
    [ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/bmark_short_output/fib.out
() after 7101 simulation cycles
    [ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/bmark_short_output/sum.out
() after 98190 simulation cycles
    [ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/bmark_short_output/replace.o
ut      () after 98271 simulation cycles

```

ASM:

```

    [ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/addi.out      ()
after 393 simulation cycles
    [ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/add.out      ()
after 694 simulation cycles
    [ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/andi.out     ()
after 331 simulation cycles
    [ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/and.out      ()
after 708 simulation cycles
    [ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/auipc.out     ()
after 144 simulation cycles
    [ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/beq.out      ()
after 488 simulation cycles
    [ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/bge.out      ()
after 536 simulation cycles
    [ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/bgeu.out     ()
after 561 simulation cycles
    [ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/blt.out      ()
after 488 simulation cycles
    [ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/bltu.out     ()
after 519 simulation cycles
    [ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/bne.out      ()
after 492 simulation cycles

```

[PASSED]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/jal.out ()
after 142 simulation cycles

[PASSED]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/jalr.out ()
after 242 simulation cycles

[PASSED]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/lb.out ()
after 399 simulation cycles

[PASSED]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/lbu.out ()
after 399 simulation cycles

[PASSED]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/lh.out ()
after 417 simulation cycles

[PASSED]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/lhu.out ()
after 424 simulation cycles

[PASSED]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/loi.out ()
after 156 simulation cycles

[PASSED]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/lw.out ()
after 427 simulation cycles

[PASSED]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/ori.out ()
after 344 simulation cycles

[PASSED]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/or.out ()
after 711 simulation cycles

[PASSED]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/sb.out ()
after 833 simulation cycles

[PASSED]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/sh.out ()
after 911 simulation cycles

[PASSED]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/simple.out ()
after 118 simulation cycles

[PASSED]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/slli.out ()
after 392 simulation cycles

[PASSED]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/sll.out ()
after 728 simulation cycles

[PASSED]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/slti.out ()
after 388 simulation cycles

```

[ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/sltiu.out      ()
after 388 simulation cycles
[ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/slt.out      ()
after 682 simulation cycles
[ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/sltu.out      ()
after 682 simulation cycles
[ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/srai.out      ()
after 413 simulation cycles
[ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/sra.out      ()
after 759 simulation cycles
[ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/srli.out      ()
after 407 simulation cycles
[ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/srl.out      ()
after 747 simulation cycles
[ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/sub.out      ()
after 680 simulation cycles
[ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/sw.out      ()
after 911 simulation cycles
[ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/xori.out      ()
after 346 simulation cycles
[ PASSED ]
/home/cc/eecs151/sp20/class/eecs151-aad/asic_proj/asm_output/xor.out      ()
after 710 simulation cycles

```

The major optimization made was a more aggressive pipeline to boost the clock speed. Since the pipeline already had a stage between the register file / hazard logic and the ALU, the only needed change was splitting up the instruction cache and the register file, and adding registers between the dcache and the writeback mux. One optimization issue was the writeback mux had to be replicated at the end of the MEM stage to provide outputs to the bypass logic in the correct stage, but since this wasn't the critical path it did not hinder performance. Initially, the five-stage pipeline had a combined write-back / ID stage with double-pumping on the register file. Since the bypassing logic was so complex, it was easier just to put the writeback mux in the MEM stage. In theory, the commit point of the processor is the fifth stage:

Instruction fetch (icache) | Instruction decode (rf read and hazard / alu logic) | execute (ALU) | MEM (dcache and writeback mux) | commit (register file has written, CSR is written to).

Adding the extra stages improved clock speed, but the CPI went down. In the initial system, JAL commands were resolved in one cycle (no bubbles needed) and dynamic branches and jumps were resolved in the second cycle (one bubble inserted). After splitting up IF and ID, one bubble must be inserted for a JAL command, and two bubbles must be inserted after a taken branch or a JALR command.

This optimization was not the most useful in the end, as the speed of the system was limited by the off-chip delay rather than the logic itself. At one point, we had a post-par report with no timing violations at a 6ns clock speed. If this interfaced with memory in a reasonable time, the overall speedup would be much greater. If we had more time, or discovered the cause of this limitation sooner, a potential fix would be to add a register between the memory arbiter and the output pin, and simply add a cycle of delay to external memory requests.

With additional time, it is likely the system would benefit from a 2-bit branch predictor. The delay of a small (64 entry) branch predictor with a tag, cached destination and two bits designating the branch prediction would be lower than that of the cache, so it wouldn't contribute to the critical path. Similarly, implementing a return stack that detects JALR x0 0(x2) would be able to speed up return calls to result in one or zero bubbles.

division of labor:

labor was pretty evenly distributed for the most part