Definition

Let S be a set. The power set (or powerset) of S, denoted $\mathcal{P}(S)$, is the set of S.

Example

- $\mathcal{P}(\emptyset) = \{\emptyset\}$
- For any non-empty S, $\mathcal{P}(S)$ contains at least two elements: \emptyset and S.
- $\mathcal{P}(\{\emptyset\}) = \{\emptyset, \{\emptyset\}\}$
- $\mathcal{P}(\{x,y\}) = \{\emptyset, \{x\}, \{y\}, \{x,y\}\}\$

Power set

Theorem

Let S be a set. Then $|\mathcal{P}(S)| = 2^{|S|}$

Proof.

Let $\operatorname{Map}(S,\{0,1\})$ be the set of maps from S to $\{0,1\}$. Since $|\operatorname{Map}(S,\{0,1\})| = 2^{|S|}$ (for S finite this can be proven; for S infinite this is the definition of $2^{|S|}$) the idea is to prove that $|\mathcal{P}(S)| = |\operatorname{Map}(S,\{0,1\})|$. Define $g:\operatorname{Map}(S,\{0,1\}) \to \mathcal{P}(S)$ by $g(f) = \{s \in S: f(s) = 1\}$ where f is a map $S \to \{0,1\}$.

Define $h: \mathcal{P}(S) \to \mathsf{Map}(S, \{0,1\})$ by $h(Y) = \chi_Y$, $Y \in \mathcal{P}(S)$ where χ_Y is the *characteristic map* (or the *indicator function*) of Y:

$$\chi_Y(s) = \begin{cases} 1, & s \in Y \\ 0, & s \notin Y \end{cases}$$

Thus, $Y = \{ s \in S : \chi_Y(s) = 1 \}.$

We claim that $g \circ h = \mathrm{id}_{\mathcal{P}(S)}$ and $h \circ g = \mathrm{id}_{\mathsf{Map}(S,\{0,1\})}$. Then g, h are both bijective and so $|\mathsf{Map}(S,\{0,1\})| = |\mathcal{P}(S)|$.

Power set

Proof (contd.)

Let
$$Y \in \mathcal{P}(S)$$
 (that is, $Y \subset S$). Then $(g \circ h)(Y) = g(h(Y))$

$$= \{ s \in S : h(Y)(s) = 1 \}$$
$$= \{ s \in S : \chi_Y(s) = 1 \}$$

$$= \{ s \in S : s \in Y \}$$
$$= Y$$

Thus,
$$g \circ h = id_{\mathcal{P}(S)}$$
.

Proof (contd.)

Let $f \in \mathsf{Map}(S, \{0,1\})$. We want to show that $(h \circ g)(f) = f$. That is, we want to show that for all $s \in S$, $((h \circ g)(f))(s) = f(s)$. We have

$$((h \circ g)(f))(s) = (h(g(f)))(s)$$

$$= \chi_{g(f)}(s)$$

$$= \begin{cases} 1, & s \in g(f) \\ 0, & s \notin g(f) \end{cases}$$

$$= \begin{cases} 1, & f(s) = 1 \\ 0, & f(s) = 0 \end{cases}$$

$$= f(s).$$

Thus, $h \circ g = id_{Map(S,\{0,1\})}$. We proved that g and h are inverses of each other.

Power set

If S is a countable set then we can think of $\mathcal{P}(S)$ as the set of bit strings. If S is finite, then we only need to consider bit strings of length |S|. Indeed, number each element of S as $s_0, s_1, s_2, s_3, \ldots$. Then $f: S \to \{0, 1\}$ is uniquely determined by the sequence a_0, a_1, a_2, \ldots where $a_i = f(s_i)$. Identifying f with the corresponding subset of S as in the proof of our previous theorem, we can think of a bit string as a subset of S. Thus, $a_i = 1$ if $s_i \in X$ and $a_i = 0$ if $s_i \notin X$.

One advantage is that we can implement set operations on subsets of S using bitwise logical operations.

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Let S = \{1, 2, ..., 16\}, X = \{1, 3, 5, 6, 7, 9, 11, 12\}, Y = \{2, 3, 9, 10, 11, 13\}. Find X \cap Y, X \cup Y, X \triangle Y, X \setminus Y.
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X corresponds to the bit string 1010111010110000 and Y corresponds to the bit string 0110000011101000. Then $X \cap Y$ corresponds to 1010111010110000 \wedge 0110000011101000 = 0010000010100000 which is $\{3,9,11\}$,

$$X \cup Y$$
 corresponds to

 $1010111010110000 \lor 0110000011101000 = 1110111011111000$ which is $\{1,2,3,5,6,7,9,10,11,12,13\}$, $X \triangle Y$ corresponds to

 $\{1, 5, 6, 7, 12\}.$

Power set

This technique can be used to program set operations. For example, to check that X is a subset of Y we represent both by bit strings and perform

 \wedge . If the result is equal to X (which is easily checked) then $X \subset Y$.

Let S be a set and $\mathcal{P}(S)$. Then $|S| < |\mathcal{P}(S)|$. In particular, if S is countable infinite then $\mathcal{P}(S)$ is uncountable.

Corollary

Let n be a positive integer. Then $n < 2^n$.

Proof.

Define $g: S \to \mathcal{P}(S)$ by $g(s) = \{s\}$. This map is manifestly injective and so $|S| \leq |\mathcal{P}(S)|$.

Suppose that $|S| = |\mathcal{P}(S)|$. Then there is a bijection $f: S \to \mathcal{P}(S)$. Let $X = \{s \in S : s \notin f(s)\}$. Since $X \in \mathcal{P}(S)$ (as a subset of S), it is then equal to $f(s_0)$ for some $s_0 \in S$.

Now, by definition of X, $s_0 \in X$ if and only if $s_0 \notin f(s_0) = X$. This is clearly a contradiction.

We proved that a bijection $f: S \to \mathcal{P}(S)$ does not exist. Thus, $|S| \le |\mathcal{P}(S)|$ and $|S| \ne |\mathcal{P}(S)|$. Therefore, $|S| < |\mathcal{P}(S)|$.

Power set

Example

In particular, this gives an alternative proof that the set of all *infinite bit strings* is uncountable.

Indeed, the set of all infinite bit strings identifies with $\mathcal{P}(\mathbb{N})$ and it follows from our theorem that $|\mathbb{N}| < |\mathcal{P}(\mathbb{N})|$.

Counting subsets with a given number of elements.

Example

- How many ways are there to choose k objects from a collection of m indistinguishable objects (number of combinations)?
- How many subsets of cardinality k are in the set of m elements?
- What is the number of sequences of integers

$$1 < i_1 < i_2 < \cdots < i_k < m$$
?

Given a set S, denote $\mathcal{P}_k(S) \subset \mathcal{P}(S)$ the set of all subsets of S of cardinality $k \colon \mathcal{P}_k(S) = \{X \in \mathcal{P}(S) : |X| = k\} = \{X \subset S : |X| = k\}.$

Definition

Let m, k be non-negative integers. Define $\binom{m}{k}$ (reads "m choose k") as $|\mathcal{P}_k(S)|$ where |S| = m.

By definition $\binom{m}{k}$ is a non-negative integer.

Some properties of $\binom{m}{k}$

- $\binom{m}{k} = 0$ if k > m. Indeed, if |A| = m, all its subsets have cardinality $\leq m$
- $\binom{m}{0} = 1 = \binom{m}{m}$; $\binom{m}{1} = \binom{m}{m-1} = m$.
- $\binom{m}{k} = \binom{m}{m-k}$, $0 \le k \le m$.

Proof.

Consider the map $f: \mathcal{P}(S) \to \mathcal{P}(S)$, $X \mapsto S \setminus X$. Since $S \setminus (S \setminus X) = X$, we have $f \circ f = \mathrm{id}_{\mathcal{P}(S)}$. Thus, f is its own inverse and hence is bijective. Since S is finite, $|f(X)| = |S \setminus X| = m - |X|$. Thus, $f(\mathcal{P}_k(S)) \subset \mathcal{P}_{m-k}(S)$ and $f(\mathcal{P}_{m-k}(S)) \subset \mathcal{P}_k(S)$. Since $f|_{\mathcal{P}_k(S)}$ is injective and $f|_{\mathcal{P}_{m-k}(S)}$ is injective, it follows that $|\mathcal{P}_k(S)| \leq |\mathcal{P}_{m-k}(S)| \leq |\mathcal{P}_k(S)|$.

•
$$\binom{m}{0} + \binom{m}{1} + \dots + \binom{m}{m-1} + \binom{m}{m} = \sum_{k=0}^{m} \binom{m}{k} = 2^{m}$$

Proof.

If |S| = m then $\mathcal{P}(S)$ has cardinality 2^m . Since $\mathcal{P}(S) = \bigcup_{0 \le k \le m} \mathcal{P}_k(S)$ and $\mathcal{P}_k(S) \cap \mathcal{P}_l(S) = \emptyset$ if $k \ne l$ we have $2^m = |\mathcal{P}(S)| = |\mathcal{P}_0(S)| + \dots + |\mathcal{P}_m(S)| = \binom{m}{0} + \dots + \binom{m}{m}$.

Theorem (Pascal's rule)

Let
$$1 \le k \le m$$
. Then $\binom{m}{k} = \binom{m-1}{k} + \binom{m-1}{k-1}$.

Proof.

Denote
$$[r] = \{1, ..., r\}$$
. Let $A = \{X \in \mathcal{P}_k([m]) : m \in X\}$, $B = \{X \in \mathcal{P}_k([m]) : m \notin X\}$.

Clearly,
$$\mathcal{P}_k([m]) = A \cup B$$
 and $A \cap B = \emptyset$ so $|\mathcal{P}_k([m])| = |A| + |B|$.

Note that
$$B = \mathcal{P}_k([m-1])$$
. Therefore, $|B| = \binom{m-1}{k}$.

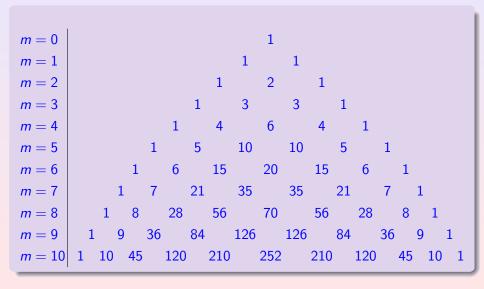
To find |A|, note that we have a map $g: A \to \mathcal{P}_{k-1}([m-1])$ given by $X \mapsto X \setminus \{m\}$. This map is bijective because the map

$$h: \mathcal{P}_{k-1}([m-1]) \to A$$
 defined by $Y \mapsto Y \cup \{m\}$ is easily seen to be its inverse. Thus, $|A| = \binom{m-1}{k-1}$.



Pascal's triangle

Pascal's rule allows us to compute the $\binom{m}{k}$ recursively.



Binomial identity

Theorem

Suppose that xy = yx. Then for any non-negative integer n

$$(x+y)^n = \sum_{k=0}^n {n \choose k} x^{n-k} y^k.$$

Proof.

We have $(x + y)^n = \sum_{S \subset \{1,...,n\}} a(S)$, where $a(S) = a(S)_1 \cdots a(S)_n$ with $a(S)_i = y$ if $i \in S$ and $a(S)_i = x$ if $i \notin S$. Since, clearly,

$$a(S) = x^{n-|S|}y^{|S|}$$
, it follows that

$$(x+y)^n = \sum_{k=0}^n |\mathcal{P}_k(\{1,\ldots,n\})| x^{n-k} y^k = \sum_{k=0}^n {n \choose k} x^{n-k} y^k.$$

Exercise

Use Pascal's rule and induction to prove the binomial identity.

A formula for $\binom{m}{k}$

Theorem

For all
$$0 \le k \le m$$
 we have $\binom{m}{k} = \frac{m(m-1)\cdots(m-k+1)}{k!} = \frac{m!}{k!(m-k)!}$.

Inductive proof.

We use induction on m. For m=0 and m=1 the assertion is clear (we use the convention that an empty product is equal to 1).

For the inductive step, we have

$$\binom{m+1}{k} = \binom{m}{k} + \binom{m}{k-1}$$

$$= \frac{m!}{k!(m-k)!} + \frac{m!}{(k-1)!(m-k+1)!}$$

$$= \frac{m!}{k!(m-k+1)!} (m-k+1) + \frac{m!}{k!(m-k+1)!} k$$

$$= \frac{m!}{k!(m-k+1)!} (m-k+1+k)$$

$$= \frac{(m+1)m!}{k!(m+1-k)!}$$

$$= \frac{(m+1)!}{k!(m+1-k)!}$$

by Pascal's rule

by induction hypothesis

Bijective proof

Here we introduce another useful counting technique.

Theorem

Let A and B be finite sets. Suppose that there exists a surjective map $f: A \to B$ such that $|f^{-1}(\{b\})| = k$ for all $b \in B$. Then |A| = k|B|.

Proof.

If $g: A \to B$ is any map and $b \neq b' \in B$ then $g^{-1}(\{b\}) \cap g^{-1}(\{b'\}) = \emptyset$. If $g: A \to B$ is surjective then

$$A = \bigcup_{b \in B} g^{-1}(\{b\})$$

and so

$$|A| = \sum_{b \in B} |g^{-1}(\{b\})|.$$

Applying this to our map f we obtain $|A| = \sum_{b \in B} k = k|B|$.

Bijective proof.

We want to prove that $k!\binom{m}{k} = m(m-1)\cdots(m-k+1)$.

Denote $[r] = \{1, ..., r\}$. Let Inj([k], [m]) be the set of all injective maps $[k] \rightarrow [m]$.

The right hand side of our equality is | lnj([k], [m])|.

The idea is to construct a surjective map $Inj([k],[m]) \to \mathcal{P}_k([m])$ and to use Theorem from slide 17.

Define a map $F: \mathsf{Inj}([k],[m]) o \mathcal{P}([m])$ by

$$F(f) = \text{im } f = \{f(1), \dots, f(k)\}.$$

Since f is injective, $|\operatorname{im} f| = k$. Thus, F is actually a map $\operatorname{Inj}([k],[m]) \to \mathcal{P}_k([m])$.

We claim that $F: \text{Inj}([k], [m]) \to \mathcal{P}_k([m])$ is surjective. Indeed, let $X = \{i_1, \dots, i_k\} \in \mathcal{P}_k([m])$. Here we chose a way to enumerate elements

of X; for example, we can assume that $i_1 < i_2 < \cdots < i_k$. All the i_r are distinct (otherwise |X| < k). Define a map $f_X : [k] \to [m]$ by $f_X(r) = i_r$,

 $1 \le r \le k$. This map is injective and $F(f_X) = X$.

Bijective proof (cont.)

Now we claim that if $X \in \mathcal{P}_k([m])$ then

Indeed, F(f) = X if and only if f is a permutation of X (that is, a bijective map $[k] \to X$), and there are k! of them.

 $|F^{-1}(\{X\})| = k!$

- How many ways are there to select a soccer team out of twenty candidates (there are eleven players on the team)? This is the number of subsets of 11 elements in a set of 20 elements, that is, $\binom{20}{11} = \binom{20}{9} = 167,960$.
- What if we take into account that the team must have a goalkeeper, a captain (who will be a forward), two more forwards and four defenders? First we choose the captain (which can be done in 20 ways), then the goalkeeper (there are 19 choices), then two more forwards ($\binom{18}{2}$) ways), then four defenders ($\binom{16}{4}$) ways). We still need to select 11-1-1-2-4=11-8=3 players out of 16-4=12 remaining candidates, which gives $\binom{12}{3}$ possibilities. Thus, the total number is $20 \cdot 19 \cdot \binom{18}{2} \cdot \binom{16}{4} \cdot \binom{12}{3} = 23,279,256,000$.

We can also choose two forwards first, in $\binom{20}{2}$ ways, then four defenders, in $\binom{18}{4}$ ways, then the goalkeeper, in 14 ways and the captain, in 13 ways, and then the remaining 3 members of the team, in $\binom{12}{3}$ ways. The total is then $\binom{20}{2}\cdot\binom{18}{4}\cdot 14\cdot 13\cdot\binom{12}{3}=23,279,256,000$ as expected.

Another way to see this: we have 11 positions on the team and 20 candidates who can occupy these positions. If each position was numbered, we would be counting the number of injective maps from P (positions) to C (candidates), which is 20!/(20-11)! = 20!/9!. However, some of these positions need to be "identified": it does not matter to us in which order forwards or defenders are assigned, the only thing that matters is who is assigned to one of these positions. We have 4! permutations of defenders, 2! permutations of forwards and 3! permutations of the remaining members of the team. So, the answer is 20!/(9!2!3!4!). It is then easy to see that both our previous answers are equal to this number.

In how many permutations of the string *ABCDEFGH* the letters *A*, *C* and *D* appear in this order but not necessarily next to each other?

The positions in which 3 letters A, C and D may occur can be chosen in $\binom{8}{3} = 56$ ways. Once these letters are placed, there are 8 - 3 = 5 positions left and 5! ways to place the remaining 5 letters. So, the answer is $5!\binom{8}{3} = 6720$.

Example

What if A, C and D appear in this order but NOT next to each other?

There are (8-3+1)!=6! permutations of the original string containing ACD, so the answer is $5!\binom{8}{3}-6!=5!\binom{8}{3}-6)=5!\cdot 50=6000$. Another way of seeing this is that out of $\binom{8}{3}$ ways of selecting three places for letters A, C and D we are NOT interested in those which are adjacent, and there are 6 of them. So, there are $\binom{8}{3}-6$ places where A, C and D can be put, which should be multiplied by 5! permutations of the remaining 5 letters.

What if A, B, C must appear in that order, and also D, E, F?

There are $\binom{8}{3}$ ways of choosing places for A, B and C; once these are chosen, there are $\binom{5}{3}$ ways of choosing places for D, E, F. After these letters are placed, there are (8-6)=2! ways of placing the remaining two letters. So, the answer is $\binom{8}{3}\binom{5}{3}2!=1120$.

How many bit strings of length 10 contain at most three 1s? At least three 1s?

A bit string of length 10 represents subsets of a set of 10 elements. The number of 1s represents the number of elements in the corresponding subset. So, we are counting the number of subsets of a set of 10 elements containing at most 3 elements, that is $\binom{10}{0} + \binom{10}{1} + \binom{10}{2} + \binom{10}{3}$.

This equals to 1 + 10 + 45 + 120 = 176.

To answer the second question we count the number of subsets of a set of 10 elements containing at least 3 elements, that is, $\binom{10}{10} + \binom{10}{10} + \binom{10}{10}$

$$\binom{10}{3} + \binom{10}{4} + \dots + \binom{10}{10} = 2^{10} - \binom{10}{0} - \binom{10}{10} - \binom{10}{2} = 1024 - 1 - 10 - 45 = 968.$$

Question

How many bit strings of length 10 contain more zeros than ones?

If a bit string of length 10 contains more 0s than 1s if it contains at least 6 zeros (or, which is the same, at most 4 ones). So, the answer is $\binom{10}{0} + \binom{10}{1} + \binom{10}{2} + \binom{10}{2} + \binom{10}{4} = 386$. Since

$$\begin{split} 2^{10} &= \binom{10}{0} + \binom{10}{1} + \binom{10}{2} + \binom{10}{3} + \binom{10}{4} + \binom{10}{5} + \binom{10}{6} + \binom{10}{7} + \binom{10}{8} + \binom{10}{9} + \binom{10}{10} \\ &= \binom{10}{0} + \binom{10}{10} + \binom{10}{1} + \binom{10}{9} + \binom{10}{2} + \binom{10}{8} + \binom{10}{3} + \binom{10}{7} + \binom{10}{4} + \binom{10}{6} + \binom{10}{5} \\ &= 2\binom{10}{0} + 2\binom{10}{1} + 2\binom{10}{2} + 2\binom{10}{3} + 2\binom{10}{4} + \binom{10}{5} \end{split}$$

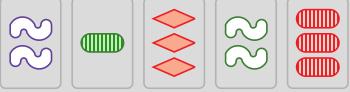
So

$$\begin{split} \binom{10}{0} + \binom{10}{1} + \binom{10}{2} + \binom{10}{3} + \binom{10}{3} + \binom{10}{4} &= 2^9 - \frac{1}{2} \binom{10}{5} \\ &= 2^9 - \frac{10 \cdot 9 \cdot 8 \cdot 7 \cdot 6}{2 \cdot 2 \cdot 3 \cdot 4 \cdot 5} = 2^9 - \frac{9 \cdot 8 \cdot 7 \cdot 6}{4!} = 2^9 - \binom{9}{4}. \end{split}$$

Example (The game of SET)

Each card has 4 attributes:

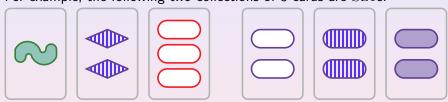
- Number of shapes on the card (1,2,3)
- The color of shapes on the card (red, green, purple)
- Shape (oval, diamond, squiggle)
- Shading (plain, striped, solid)



So, if $N = \{1, 2, 3\}$, $C = \{\text{red}, \text{green}, \text{purple}\}$, $S = \{\text{oval}, \text{diamond}, \text{squiggle}\}$ and $F = \{\text{plain}, \text{striped}, \text{solid}\}$ then each card can be represented as an element of $N \times C \times S \times F$. Alternatively, we can think of each card as an element of $\{0, 1, 2\}^4$ (each attribute has 3 possible values, encoded by 0, 1 and 2.

A SET is a collection of 3 cards, (a_1, a_2, a_3, a_4) , (b_1, b_2, b_3, b_4) and (c_1, c_2, c_3, c_4) such that for each $1 \le i \le 4$, either $a_i = b_i = c_i$ or $a_i \ne b_i \ne c_i$ (that is, each attribute must have the same value in all three cards, or different values in all three cards)

For example, the following two collections of 3 cards are Sets :



And the following two are not:













Question

What is the probability that 3 cards selected randomly form a Set ?

- How many cards are there in the deck?
 Each card has 4 properties, each property can take 3 values, so the total number is the cardinality of {0,1,2}⁴, that is 3⁴ = 81.
- How many ways are there to choose 3 cards from the deck? We are looking for the number of subsets of 3 elements in a set of 81 elements. This number is $\binom{81}{3} = 81 \cdot 80 \cdot 70/3! = 85\,320$.
- How many SETs (in the sense of the game) are there in the deck? Recall that once two cards X, Y are chosen, there is a unique card Z such that $\{X,Y,Z\}$ is a SET. Two cards can be selected in $\binom{81}{2}$ ways. However, this is not the end of the story: the same set $\{X,Y,Z\}$ is obtained in 3 different ways (from the pairs $\{X,Y\}$, $\{X,Z\}$ and $\{Y,Z\}$). So, the number of SETs is $\frac{1}{3}\binom{81}{2}=81\cdot80/(3\cdot2)=1,080$.

Answer.

We divide the number of desirable outcomes by the total number of outcomes. This gives

$$\frac{1}{3} \cdot \frac{\binom{81}{2}}{\binom{81}{3}} = \frac{3!}{3 \cdot 2} \cdot \frac{81 \cdot 80}{81 \cdot 80 \cdot 79} = \frac{1}{79}.$$

Thus, the probability is slightly higher than the probability of picking a specific card (which is 1/81).

Question

How many Sets have d attributes different in all three cards, $1 \le d \le 4$?

There are 3^4 ways to select the first card. Once the first card is selected, the second card can be selected in $\binom{4}{d}2^d$ ways, since one value of each of d attributes is excluded and there are $\binom{4}{d}$ ways to select d attributes which are supposed to be different in all three cards. Once two cards are selected, the third card is determined uniquely. Thus, we have $3^4\binom{4}{d}2^d$ ordered triples which form a SET with the desired property. Since the order is irrelevant for our purposes, this number must be divided by 3!. So, the answer is $27 \cdot \binom{4}{d}2^{d-1}$.

The probability to encounter such a set is thus $81 \cdot \binom{4}{d} 2^{d-1} / \binom{81}{2} = \binom{4}{d} 2^d / 80$. The numbers are: 1/10, 3/10, 2/5 and 1/5 (so the most common SETs are those with 3 different attributes).

Hard questions

- When the game is played, 12 cards are put on the table at the same time. What is the probability that there will be no SETs among these 12 cards?
- ullet What is the minimal number of cards one needs to draw to ensure that there is a Set among them?

What else can we count using $\binom{m}{k}$?

There are $\binom{m}{k}$ ways to choose k indistinguishable objects from a collection of m objects or to fill k indistinguishable boxes so that no box contains more than one object (the number of *combinations* of k objects out of m objects is $\binom{m}{k}$).

There are $\binom{m}{k}$ sequences of integers $1 \le i_1 < \cdots < i_k \le m$.

Indeed, such a sequence identifies with a subset of k elements in $\{1, \ldots, m\}$.

Question

Find the number of sequences of integers $1 \le i_1 \le i_2 \le \cdots \le i_k \le m$.

Given
$$1 \le i_1 < \cdots < i_k \le m$$
, define a sequence (j_1, \ldots, j_k) by $j_r = i_r + r - 1$, $1 \le r \le k$.

Thus, $j_1 = i_1$, $j_2 = i_2 + 1$, $j_3 = i_3 + 2$ and so on.

Then
$$j_{r+1} - j_r = (i_{r+1} + r) - (i_r + r - 1) = i_{r+i} - i_r + 1 > 0$$
 because $i_{r+1} - i_r \ge 0$.

So,
$$1 \le j_1 < j_2 < \cdots < j_k \le m + k - 1$$
.

Conversely, from each such sequence we can obtain a sequence

$$1 \le i_1 \le \cdots \le i_k \le m$$
 where $i_r = j_r - r + 1$.

Thus, there is a bijection between the set of *weakly increasing* sequences of integers between 1 and m of length k and the set of *strictly increasing* sequences of integers between 1 and m + k - 1 of length k.

But the number of the latter sequences is $\binom{m+k-1}{k}$.

Permutations vs Combinations

r-permutations of a set of m elements: ordered arrangements of r elements of our sets.

r-combinations of a set of m elements: subsets containing r elements.



In the first case, the order is important. In the second case, the order does not matter.

Since there are r! ways to permute r different objects, the number of r-permutations is r! times the number of r-combinations, that is $r!\binom{m}{r} = m(m-1)\cdots(m-r+1)$.