Dragon Sensor Design Problem Solution



Problem

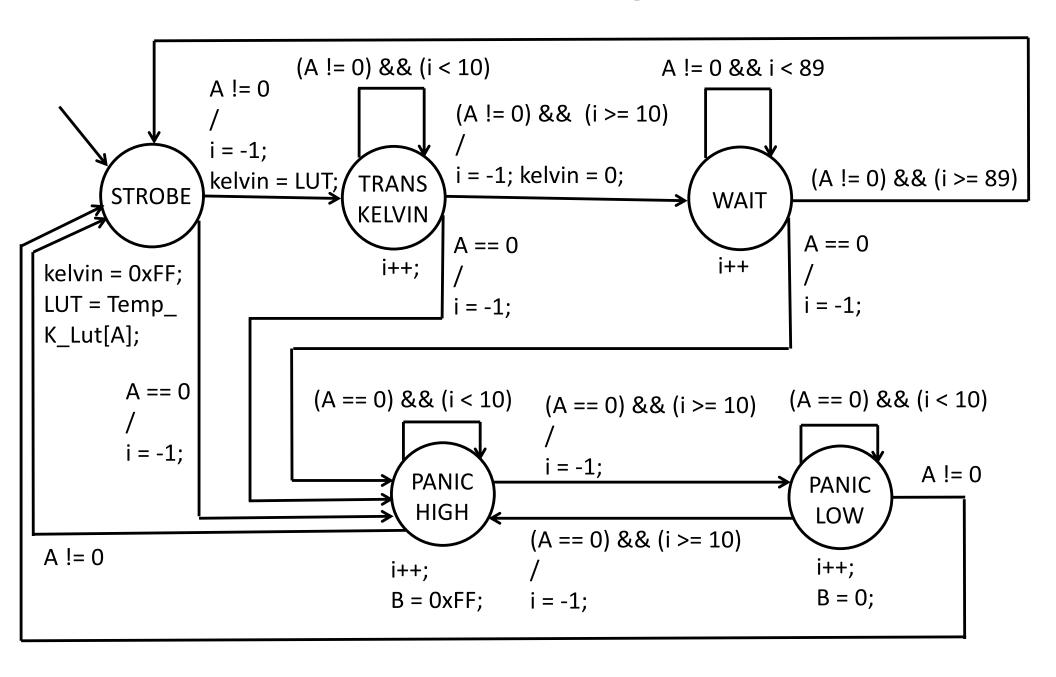
- A digital thermometer connected to A7..A0 reads the temperature inside the pressurized cabin of the SpaceX Dragon capsule (Above). The sensor samples at 1 second intervals.
- If pressure in the cabin is lost, then A == 0. (Space is very cold.)
- The software in the capsule uses a **lookup table "Temp_K_LUT"** to convert values of A into Kelvin temperatures. The values stored in the lookup table are 8-bit chars.
- The current temperature in the cabin (in Kelvin) needs to sent back to Earth, so ground control can monitor the current status of the capsule. A temperature is sent back to earth by strobing B = 0xff for 10 ms, and then placing the temperature from the lookup table on B for 100 ms.
- If cabin pressure is lost (A==0), the system should alternately set B=0xff, then B=0x00, repeatedly, at 100ms intervals, perhaps to warn other subsystems about the failure.

Solution #1 Single-Task

Overview

- 1. If (A != 0) // "Transmission" mode
 - 1.1 Strobe for 10ms; convert sensor A to Kelvin via LUT
 - 1.2 Transmit the kelvin value for 100ms
 - 1.3 Wait for the next 890ms to complete the 1s period
- 2. If (A == 0) // "PANIC" mode
 - 2.1 Alternate B=0xFF, B=0 for 100ms period
- 3. Swap between Transmission/PANIC mode
 - 3.1 Switch from Transmission to PANIC mode if A == 0
 - 3.2 Switch from PANIC to Transmission mode if A != 0

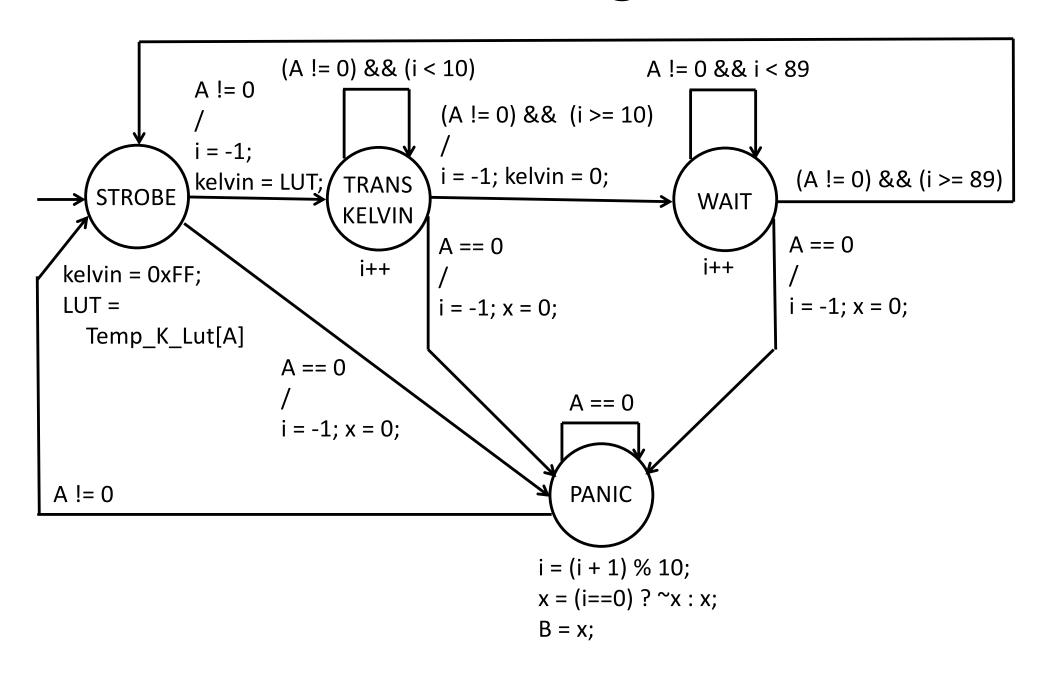
Solution #1, Single-Task



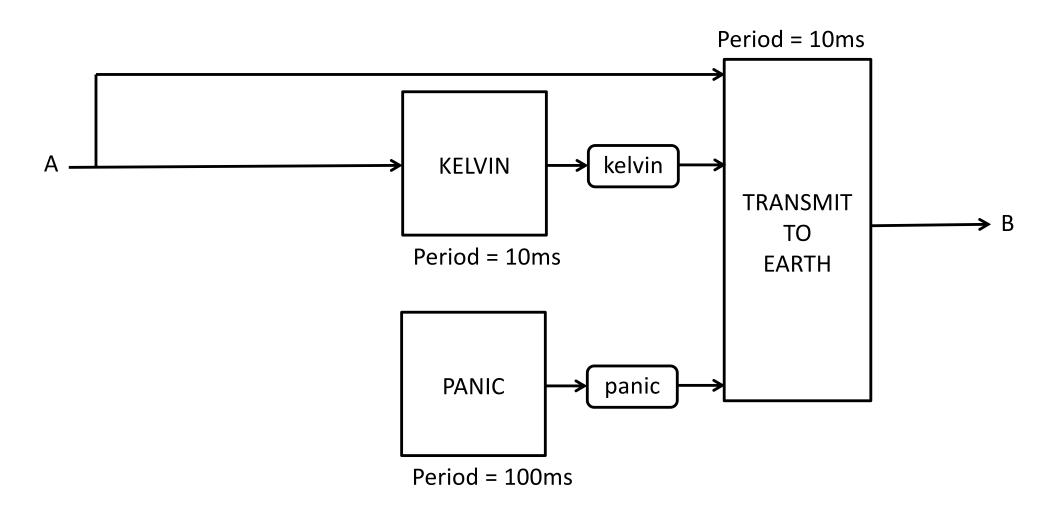
Solution #2 Single-Task (One "PANIC" mode state)

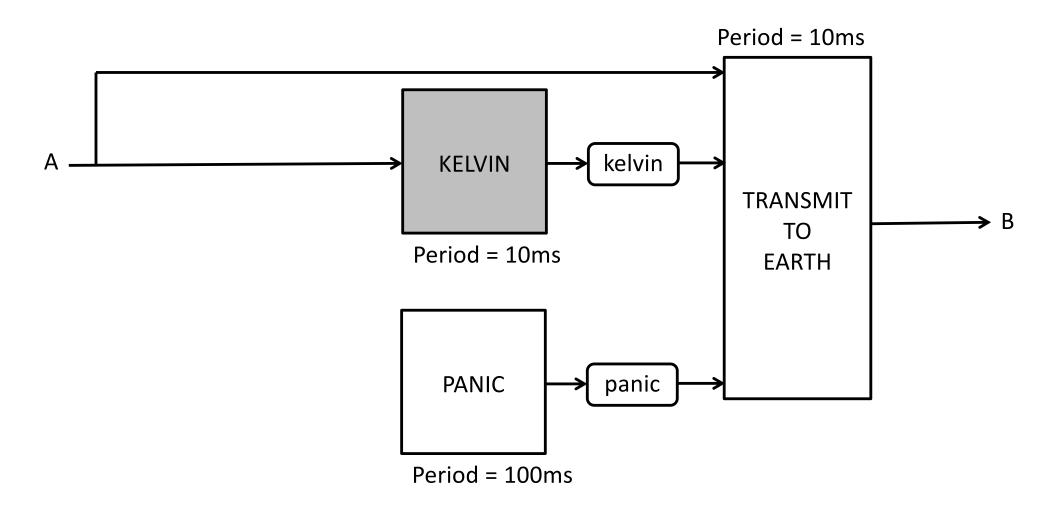
- Switching B from 0xFF to 0 and vice-versa is simply a bit inversion.
- Use a single state for "PANIC" mode
 - Use a variable x, initialized to zero
 - Every 100ms set $x = ^x$;
 - Output B = x
- Cleans up a bit of the mess

Solution #2, Single-Task

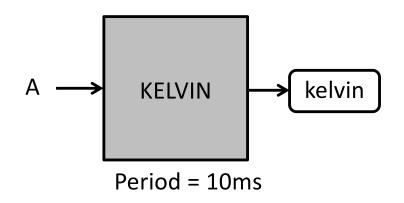


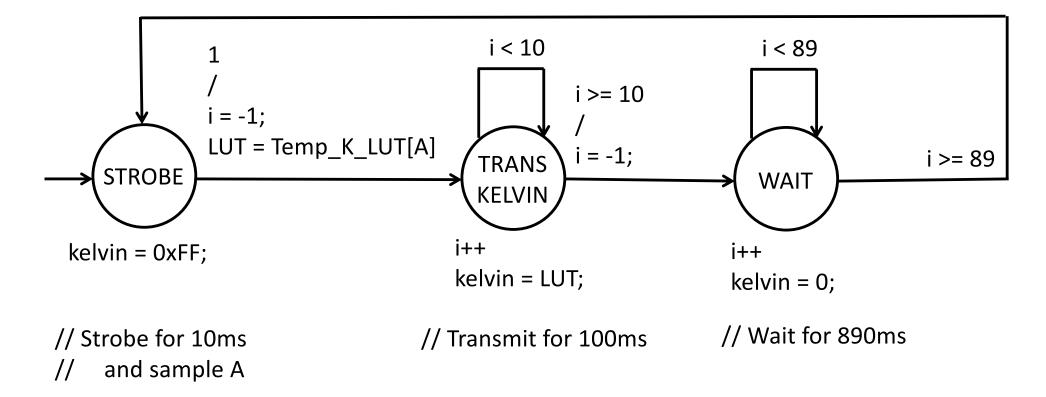
- Separate "Transmission" and "PANIC" mode into separate tasks
- A third task acts as a multiplexer
 - If (A == 0) Output "PANIC" mode result
 - If (A != 0) Output "Transmission" mode result

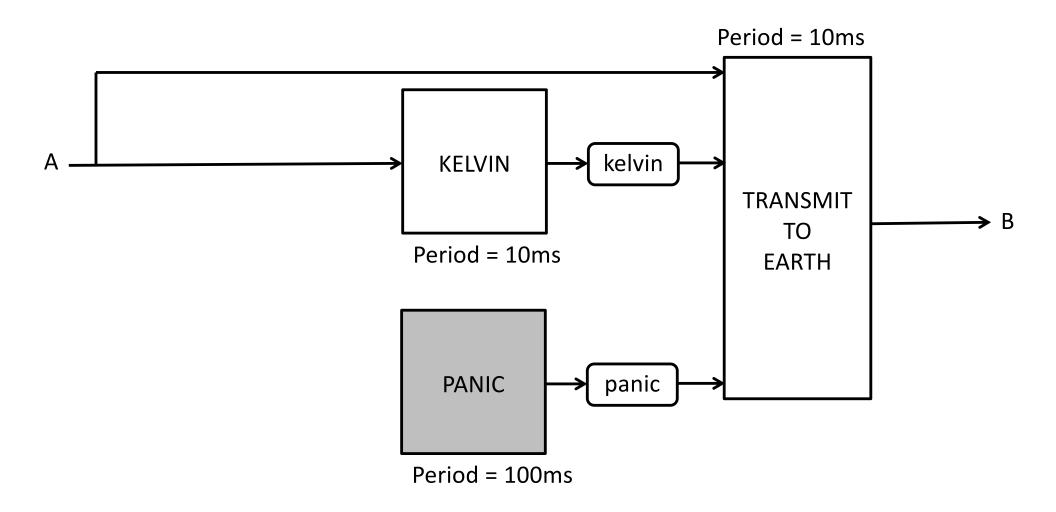




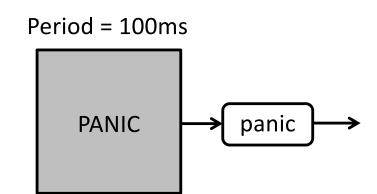
KELVIN Task



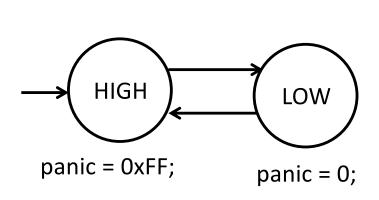




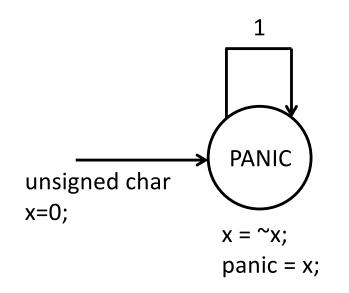
PANIC Task

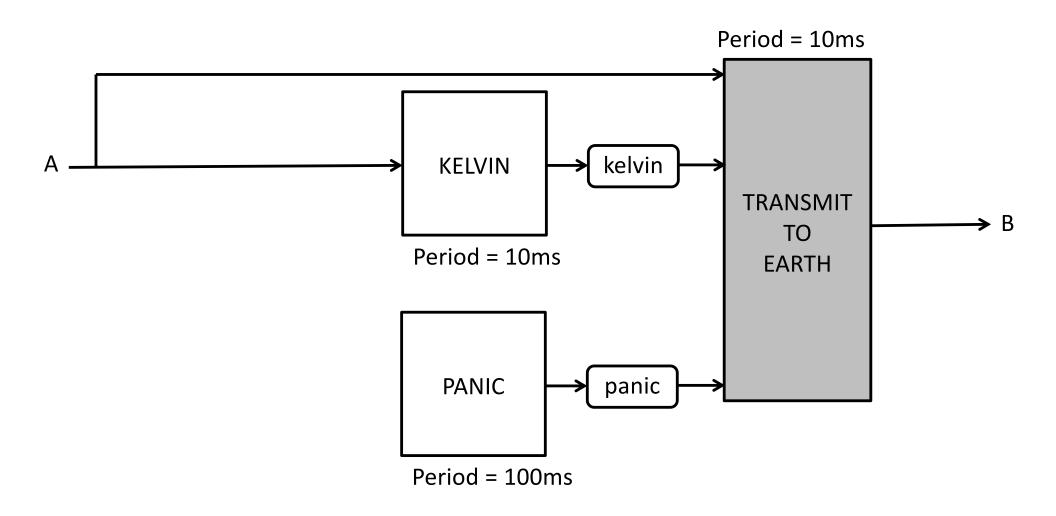


2-state Implementation

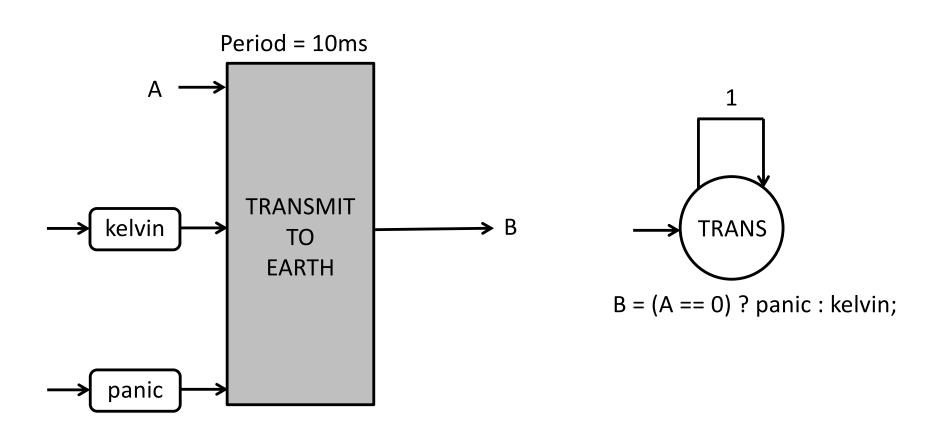


1-state Implementation

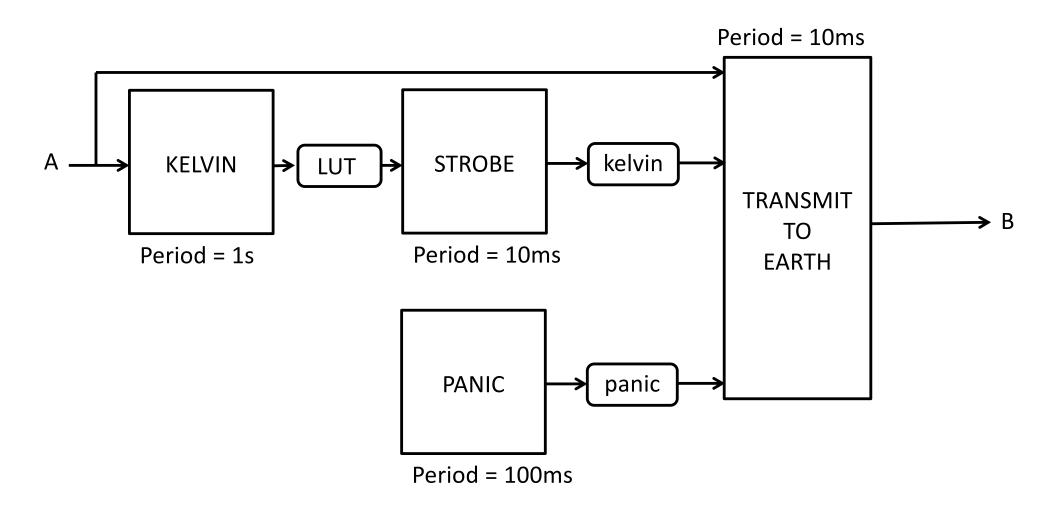


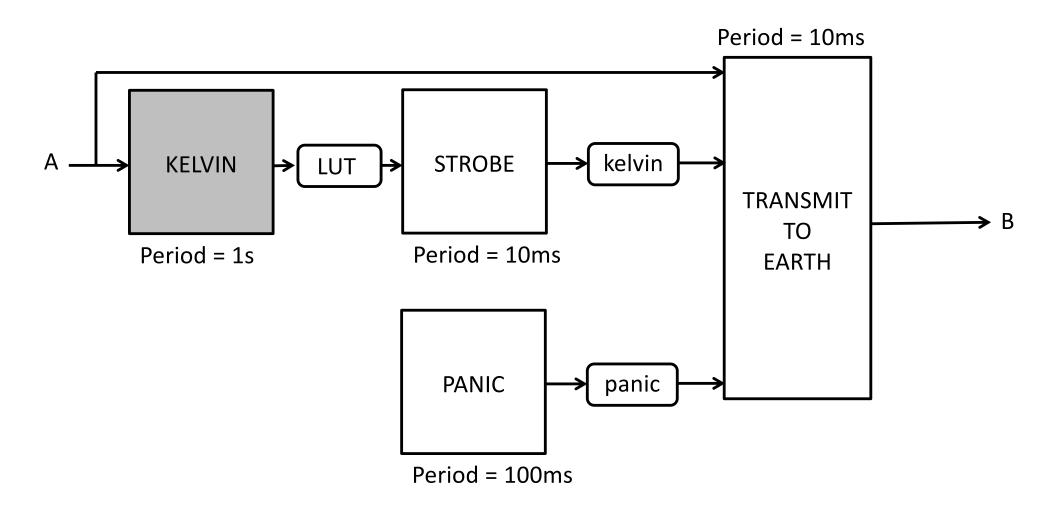


TRANSMIT TO EARTH Task

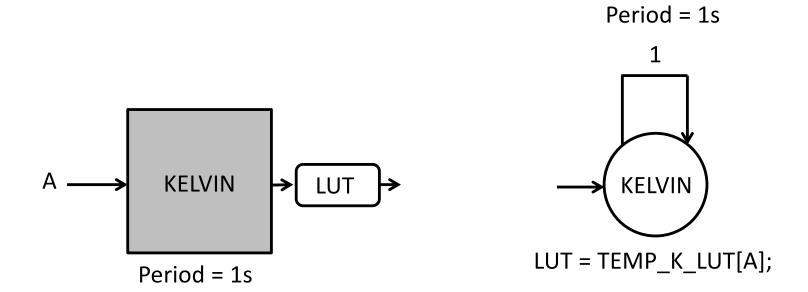


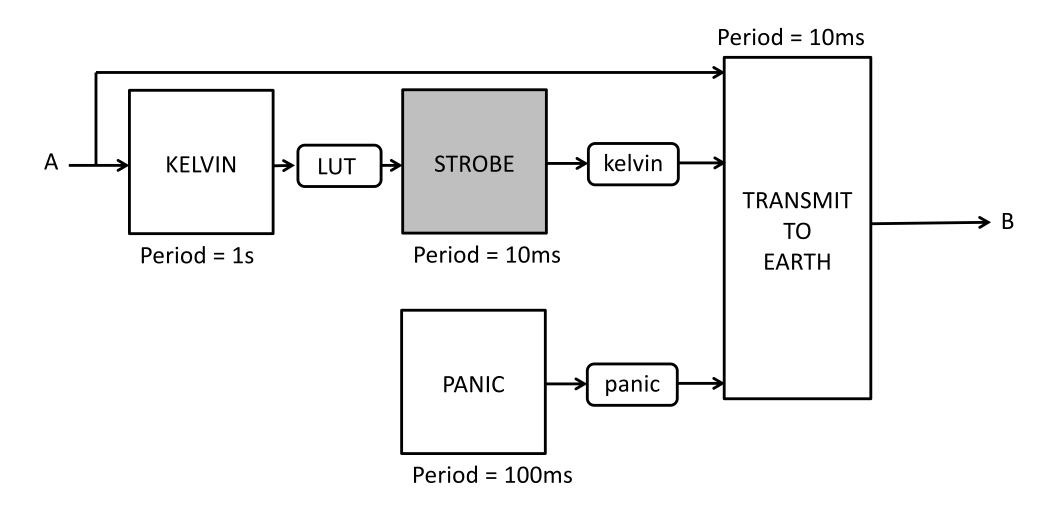
- Create a separate task with a 1 second period to perform the LUT-based conversion of A to kelvin
- Not much of an improvement over Solution #3, but leads to Solution #5 which is far more elegant



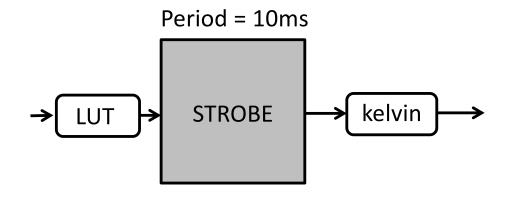


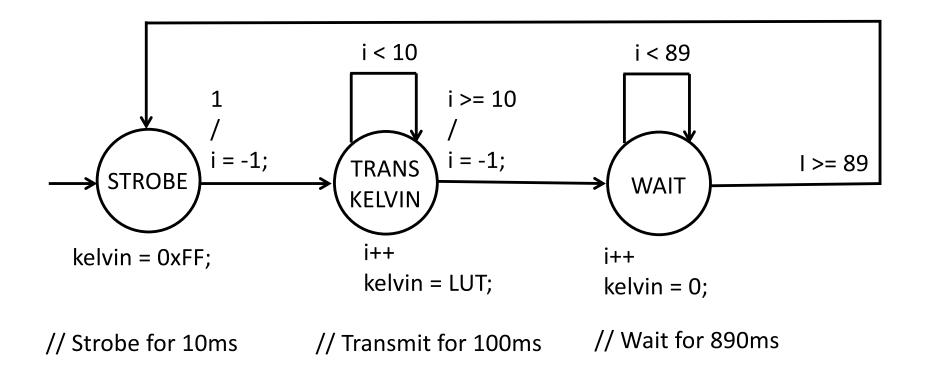
KELVIN Task

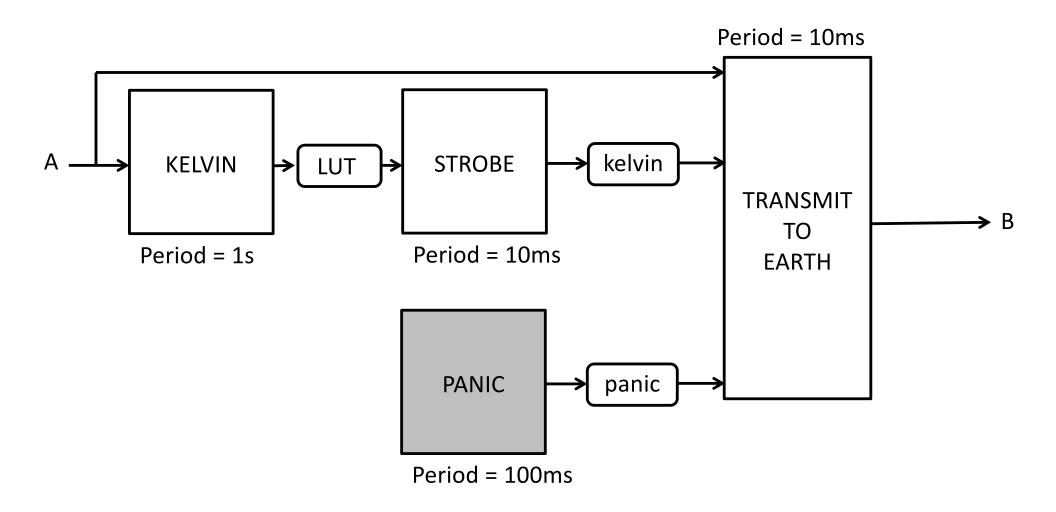




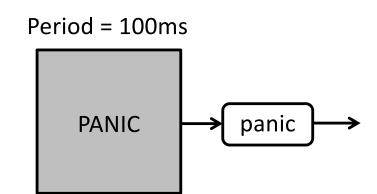
STROBE Task



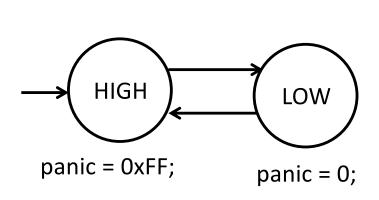




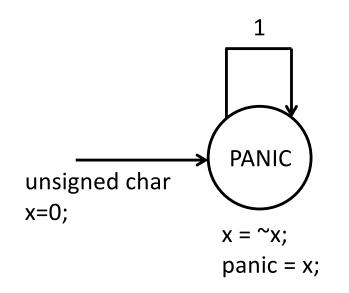
PANIC Task

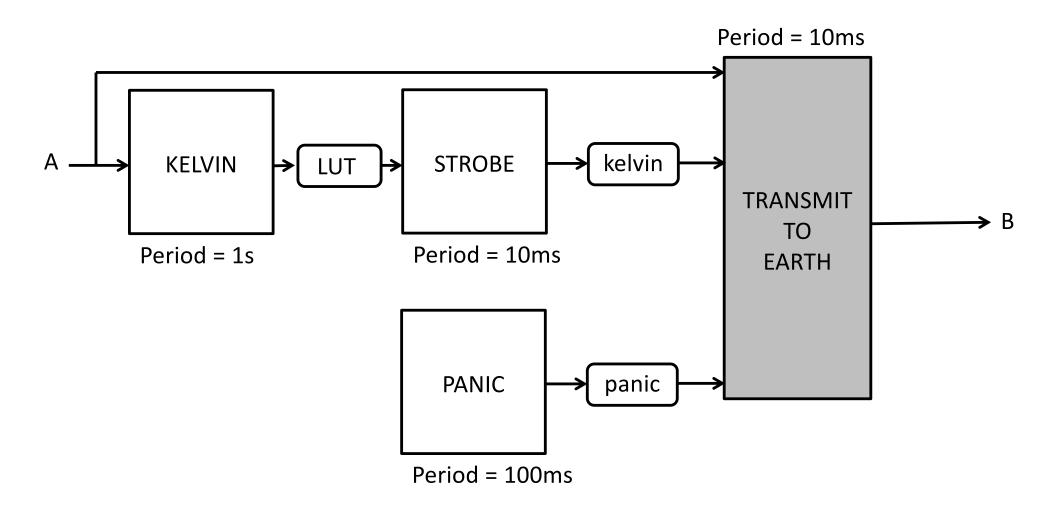


2-state Implementation

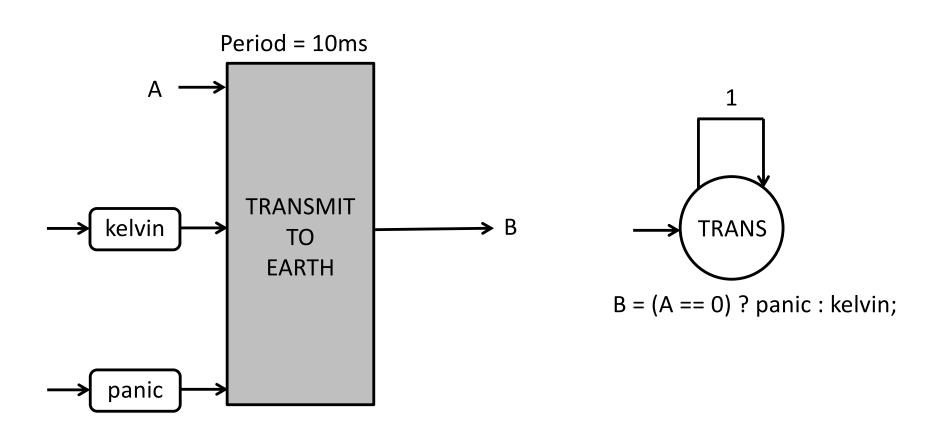


1-state Implementation





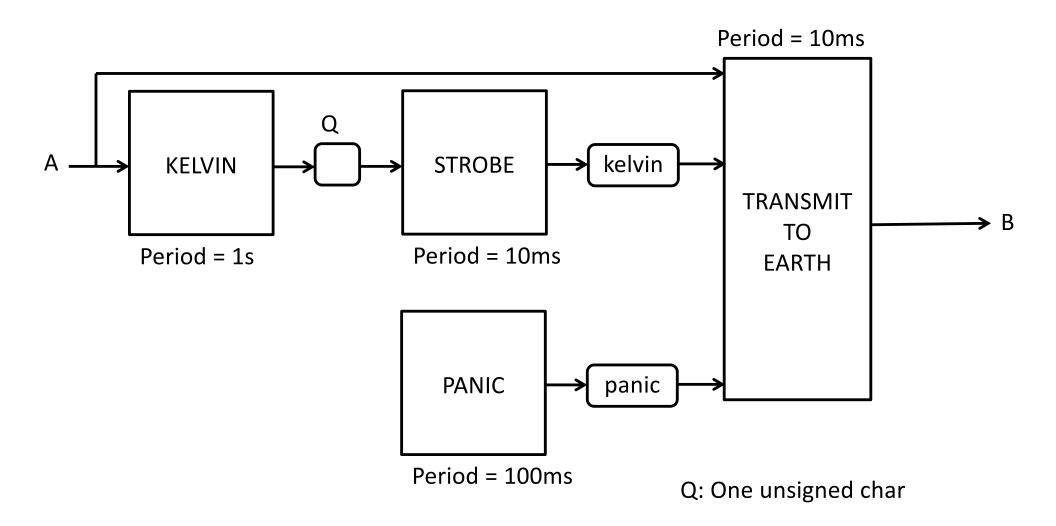
TRANSMIT TO EARTH Task



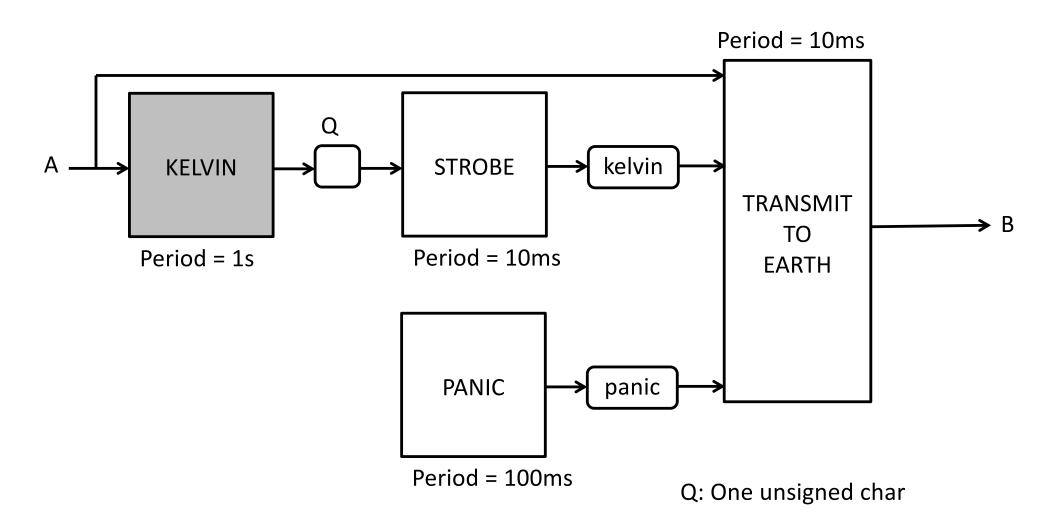
Solution #5 Concurrent SynchSM w/Shared Variables and a Queue

- Sample input A every 1 s and put the result into a queue, rather than a shared variable
- The STROBE task no longer needs to track waiting time (890ms in Solutions #1 - #4)
- The STROBE can simply poll the queue and wait until data is available to process. This simplifies the control logic significantly.

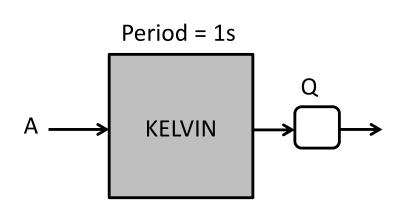
Solution #5 Concurrent SynchSM w/Queue

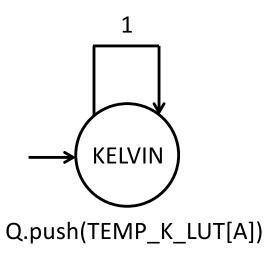


Solution #5 Concurrent SynchSM w/Queue

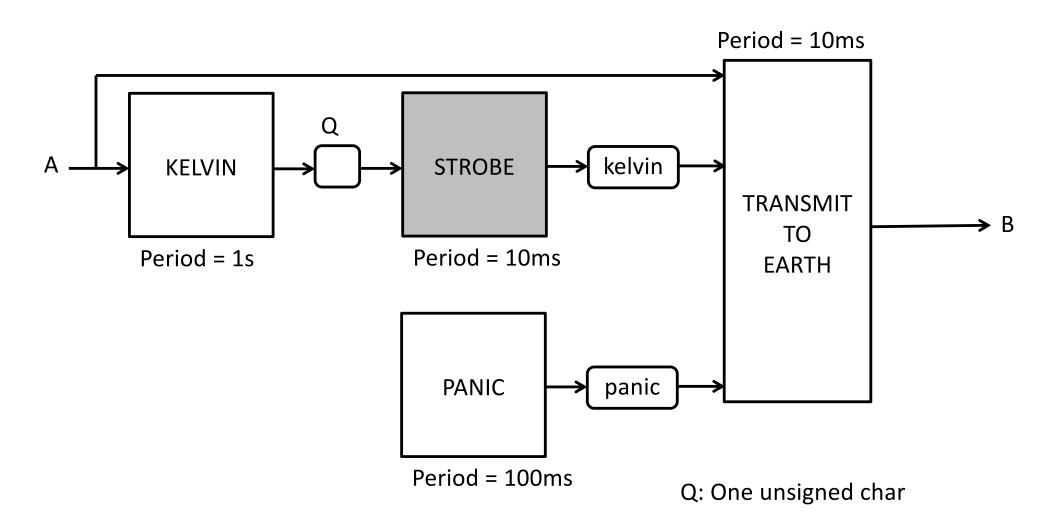


KELVIN Task

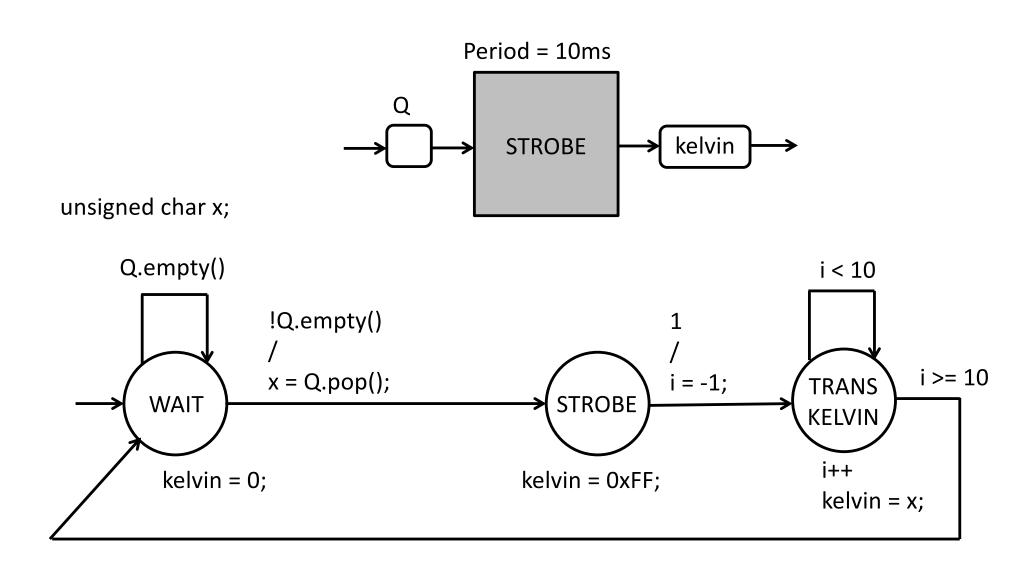




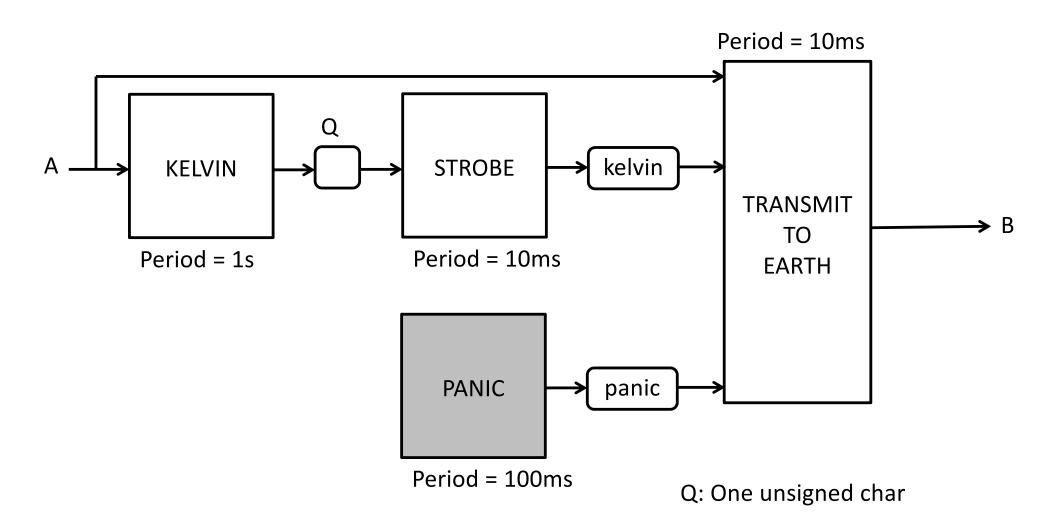
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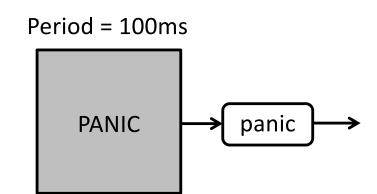
STROBE Task



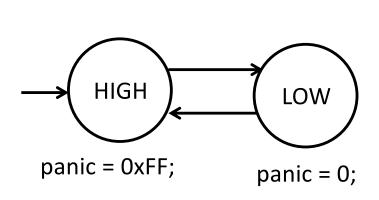
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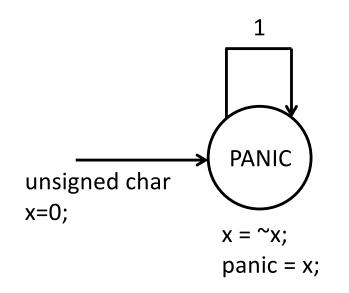
PANIC Task



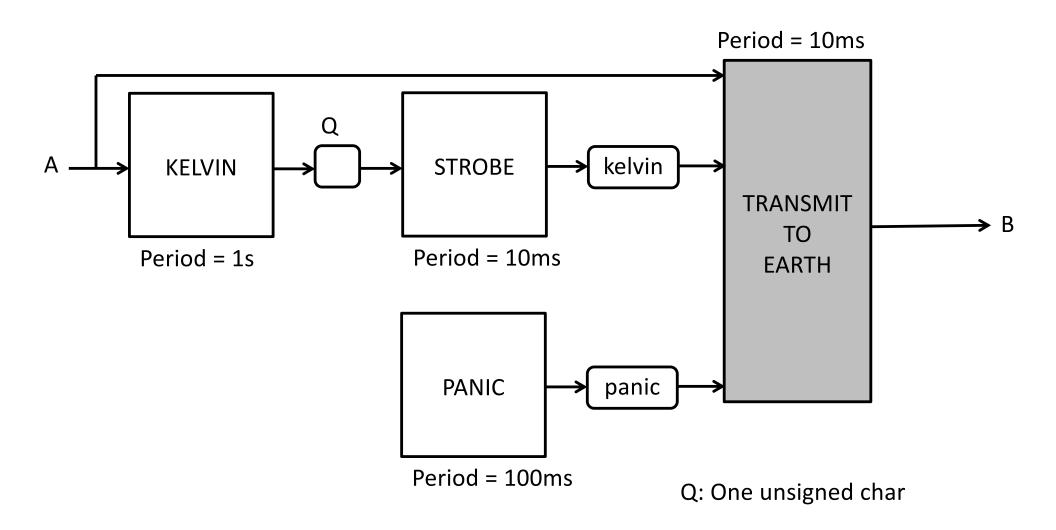
2-state Implementation



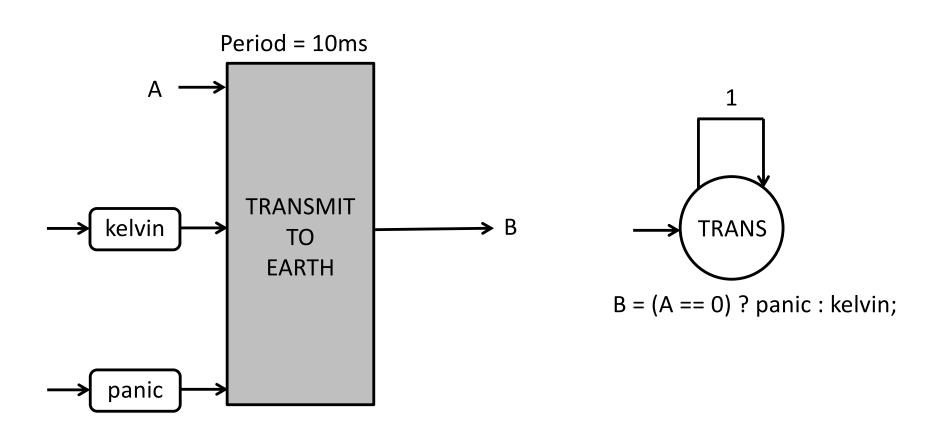
1-state Implementation



Solution #5 Concurrent SynchSM w/Queue



TRANSMIT TO EARTH Task



Conclusion

To get the simplest overall design:

- Concurrent task decomposition
 - Separate "Transmission" and "PANIC" functionalities
 - "STROBE" and "PANIC" tasks compute their output every tick;
 "TRANSMIT" selects which one to output to B
 - This is far simpler than switching between "STROBE" and "PANIC" modes of operations in a non-concurrent synchSM
- Uses both a queue and shared variables
 - Only one task writes to the queue in this case (unlike the Floating Beacon example)

And the Moral of the Story Is...

Queues make everything many things better

 Avoid "flashing" the value of a shared variable for a small number of ticks, as this requires complex cross-task synchronization