

Sample Problems for CS 152 Final

Q1. Below is the specification of a control flow construct and its grammar:

Example:

```
If x < y then
    <otherstatements>
elseif a > b then
    <otherstatements>
else
    <otherstatements>
endif
```

Relevant Grammar Productions:

$\langle S \rangle \rightarrow \text{if } \langle \text{condt} \rangle \text{ then } \langle \text{otherstatements} \rangle \langle \text{rest} \rangle$

$\langle \text{rest} \rangle \rightarrow \text{elseif } \langle \text{condt} \rangle \text{ then } \langle \text{otherstatements} \rangle \langle \text{rest} \rangle$
| $\text{else } \langle \text{otherstatements} \rangle \text{ endif}$

$\langle \text{condt} \rangle \rightarrow \text{id rel op id}$

Provide the semantic rules that generate code for the above control flow construct. The generated code should be available in the $\langle S \rangle.\text{code}$ attribute. Assume that semantic rules specified by other productions not shown here will place the code generated for $\langle \text{otherstatements} \rangle$ in the attribute $\langle \text{otherstatements} \rangle.\text{code}$.

Solution:

```
 $\langle \text{condt} \rangle \rightarrow \text{id}_1 \text{ rel op id}_2 \{$   
     $\text{truelabel} = \text{newlabel}();$   
     $\langle \text{condt} \rangle.\text{falselabel} = \text{newlabel}();$   
     $\langle \text{condt} \rangle.\text{code} = \text{gen}(\text{"if" id}_1.\text{place "rel op" id}_2.\text{place "go to" truelabel})$   
         $\parallel \text{gen}(\text{"go to" } \langle \text{condt} \rangle.\text{falselabel}) \parallel \text{gen}(\text{truelabel}.);$   
}
```

```
 $\langle S \rangle \rightarrow \text{if } \langle \text{condt} \rangle \text{ then } \langle \text{otherstatements} \rangle$   
     $\{$   
         $\langle \text{rest} \rangle.\text{ifalselabel} = \langle \text{condt} \rangle.\text{falselabel};$   
         $\langle \text{rest} \rangle.\text{ixit} = \text{newlabel}();$   
         $\langle \text{rest} \rangle.\text{icode} = \langle \text{condt} \rangle.\text{code} \parallel \langle \text{otherstatements} \rangle.\text{code} \parallel$   
             $\text{gen}(\text{"go to" } \langle \text{rest} \rangle.\text{ixit})$   
    }  
 $\langle \text{rest} \rangle \{ \langle S \rangle.\text{code} = \langle \text{rest} \rangle.\text{icode} \}$ 
```

```

<rest1> → elseif <condt> then <otherstatements>
{
    <rest2>.icode = <rest1>.icode || gen(<rest1>.ifalselabel ":") ||
    <condt>.code || <otherstatements>.code || gen("go to" <rest1>.iexit);
    <rest2>.ifalselabel = <condt>.falselabel;
    <rest2>.iexit = <rest1>.iexit
}
<rest2> { <rest1>.scode = <rest2>.scode }

```

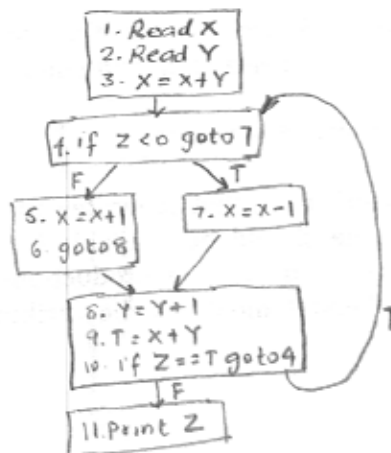
```

<rest1> → else <otherstatements> endif
{
    <rest1>.scode = <rest1>.icode || gen(<rest1>.ifalselabel ":")
    || <otherstatements>.code || gen(<rest1>.iexit ":")
}

```

Q2. Construct the control flow graph for the following intermediate code sequence.

1. Read X
2. Read Y
3. $X = X + Y$
- ✓ 4. If $Z < 0$ go to 7
- ✓ 5. $X = X + 1$
- ✓ 6. Go to 8
- ✓ 7. $X = X - 1$
- ✓ 8. $Y = Y + 1$
9. $T = X + Y$
10. If $Z == T$ go to 4
- ✓ 11. Print Z

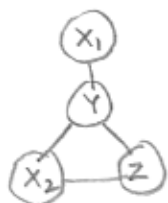
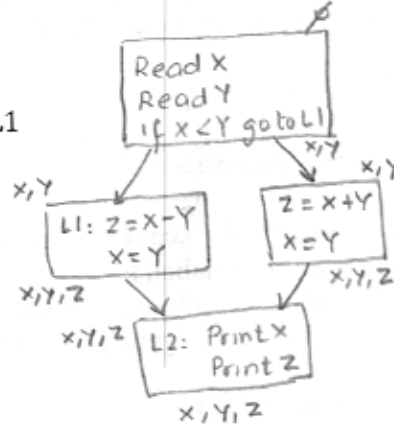


Q3. For the following code sequence: construct the control flow graph; perform liveness analysis for all variables; construct live ranges; and construct the interference graph. Color the interference graph using (a) 2 colors; and (b) 3 colors. Assume all variables are live at the end of the code sequence.

```

Read X
Read Y
If X < Y go to L1
Z = X + Y
X = Y
Go to L2
L1:
Z = X - Y
X = Y
L2:
Print X
Print Z

```



3 colors:
 $x_1, x_2, z; Y$
 $R_2 R_3 R_2 R_1$
 or
 R_3

2 colors:
 x_1 ; remove $Y; x_2, z$
 R_1 stays $R_2 R_1$
 or
 R_2 memory

Q4. For the program given below show the contents of the runtime stack (activation records, local variables, control links, and access links) for the following call sequence: $\text{Main} \rightarrow \text{F} \rightarrow \text{G} \rightarrow \text{H} \rightarrow \text{F}$

```

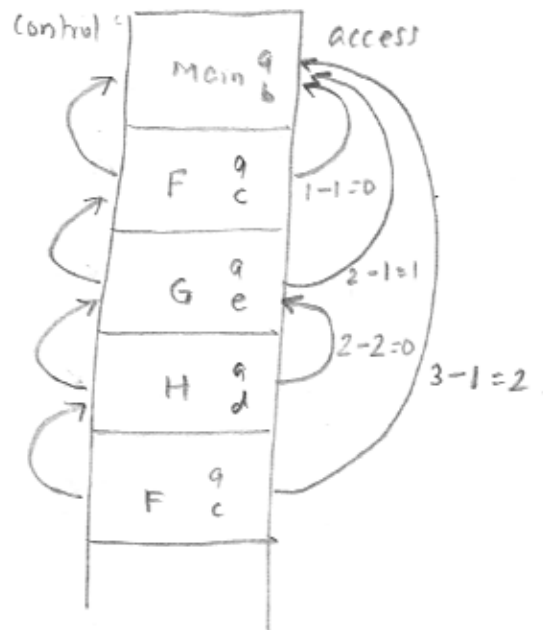
Main () {
    Int a, b;

    FO {
        Int a, c;
        Call GO;
    }

    GO {
        Int a, e;
        HO {
            Int a, d;
            Call FO;
        }
        Call HO;
    }

    Call FO
}

```



- Show how the access links are computed and setup; and
- Show how various non-local variables are accessed within each function using the access links.

eg. access variable b from H .

$$3-1=2 \quad (2, \text{offset})$$