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# Database Systems 2

## Lecture 18

### Normalisation

# *The Process of Normalization*

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**Formal technique for analyzing a relation based on its primary key and functional dependencies between its attributes.**

**Often executed as a series of steps. Each step corresponds to a specific normal form, which has known properties.**

**As normalization proceeds, relations become progressively more restricted (stronger) in format and also less vulnerable to update anomalies.**

# *Unnormalized Form (UNF)*

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A table that contains one or more repeating groups.

To create an unnormalized table:

- transform data from information source (e.g. paper based form) into table format with columns and rows.

# *First Normal Form (1NF)*

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A relation in which:

- the intersection of each row and column contains one and only one value, (i.e. is atomic)

and:

- One or more of the columns is designated as the primary key.

# *To get from UNF to 1NF*

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Identify repeating group(s) in unnormalized table which repeats for the key attribute(s).

Nominate an attribute (or group of attributes) to act as the key for the unnormalized table.

# *Get rid of repeating group*

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Remove repeating group by:

- entering appropriate data into the empty columns of rows containing repeating data ('flattening' the table).
- Note: some texts separate the repeating group out to it's own table at this stage, although technically it should be done when you go from 1NF to 2NF.

# *ClientRental UNF To 1NF By Flattening*

clientNo	cName	propertyNo	pAddress	rentStart	rentFinish	rent	ownerNo	oName
CR76	John Kay	PG4	6 Lawrence St	1-Jul-00	31-Aug-01	350	CO40	Tina Murphy
		PG16	5 Novar Dr	1-Sep-02	1-Sep-02	450	CO93	Tony Shaw
CR56	Aline Stewart	PG4	6 Lawrence St	1-Sep-99	10-Jun-00	350	CO40	Tina Murphy
		PG36	2 Manor Rd	10-Oct-00	1-Dec-01	375	CO93	Tony Shaw
		PG16	5 Novar Dr	1-Nov-02	10-Aug-03	450	CO93	Tony Shaw

clientNo	cName	propertyNo	pAddress	rentStart	rentFinish	rent	ownerNo	oName
CR76	John Kay	PG4	6 Lawrence St	1-Jul-00	31-Aug-01	350	CO40	Tina Murphy
<b>CR76</b>	<b>John Kay</b>	PG16	5 Novar Dr	1-Sep-02	1-Sep-02	450	CO93	Tony Shaw
CR56	Aline Stewart	PG4	6 Lawrence St	1-Sep-99	10-Jun-00	350	CO40	Tina Murphy
<b>CR56</b>	<b>Aline Stewart</b>	PG36	2 Manor Rd	10-Oct-00	1-Dec-01	375	CO93	Tony Shaw
<b>CR56</b>	<b>Aline Stewart</b>	PG16	5 Novar Dr	1-Nov-02	10-Aug-03	450	CO93	Tony Shaw

## *Now choose a primary key*

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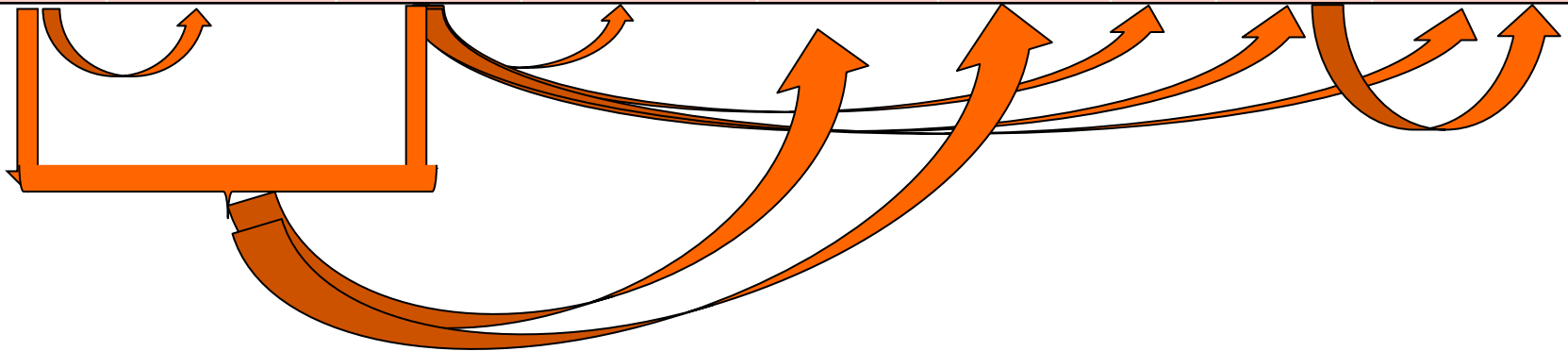
Might be fairly easy to spot.

If it isn't, got through the following process.



# *ClientRental Functional Dependencies*

clientNo	cName	propertyNo	pAddress	rentStart	rentFinish	rent	ownerNo	oName
CR76	John Kay	PG4	6 Lawrence St	1-Jul-00	31-Aug-01	350	CO40	Tina Murphy
CR76	John Kay	PG16	5 Novar Dr	1-Sep-02	1-Sep-02	450	CO93	Tony Shaw
CR56	Aline Stewart	PG4	6 Lawrence St	1-Sep-99	10-Jun-00	350	CO40	Tina Murphy
CR56	Aline Stewart	PG36	2 Manor Rd	10-Oct-00	1-Dec-01	375	CO93	Tony Shaw
CR56	Aline Stewart	PG16	5 Novar Dr	1-Nov-02	10-Aug-03	450	CO93	Tony Shaw



What about rentStart and rentFinish?

# *ClientRental Functional Dependencies*

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clientNo	cName	propertyNo	pAddress	rentStart	rentFinish	rent	ownerNo	oName
CR76	John Kay	PG4	6 Lawrence St	1-Jul-00	31-Aug-01	350	CO40	Tina Murphy
CR76	John Kay	PG16	5 Novar Dr	1-Sep-02	1-Sep-02	450	CO93	Tony Shaw
CR56	Aline Stewart	PG4	6 Lawrence St	1-Sep-99	10-Jun-00	350	CO40	Tina Murphy
CR56	Aline Stewart	PG36	2 Manor Rd	10-Oct-00	1-Dec-01	375	CO93	Tony Shaw
CR56	Aline Stewart	PG16	5 Novar Dr	1-Nov-02	10-Aug-03	450	CO93	Tony Shaw

- clientNo → cName
- propertyNo → pAddress, rent, ownerNo, oName
- ownerNo → oName
- clientNo, propertyNo → cName, pAddress, rentStart, rentFinish, rent, ownerNo, oName

There are other possible dependencies:

- clientNo, rentStart → cName, propertyNo, pAddress, rentFinish, rent, ownerNo, oName
- propertyNo, rentStart → clientNo, cName, pAddress, rentFinish, rent, ownerNo, oName

# *ClientRental Functional Dependencies*

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clientNo	cName	propertyNo	pAddress	rentStart	rentFinish	rent	ownerNo	oName
CR76	John Kay	PG4	6 Lawrence St	1-Jul-00	31-Aug-01	350	CO40	Tina Murphy
CR76	John Kay	PG16	5 Novar Dr	1-Sep-02	1-Sep-02	450	CO93	Tony Shaw
CR56	Aline Stewart	PG4	6 Lawrence St	1-Sep-99	10-Jun-00	350	CO40	Tina Murphy
CR56	Aline Stewart	PG36	2 Manor Rd	10-Oct-00	1-Dec-01	375	CO93	Tony Shaw
CR56	Aline Stewart	PG16	5 Novar Dr	1-Nov-02	10-Aug-03	450	CO93	Tony Shaw

No one field can act as a primary key for the whole table:

What are the possible composite keys?

# *ClientRental Primary Key*

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Candidate keys are

- (clientNo, propertyNo)
- (clientNo, rentStart)
- (propertyNo, rentStart)

Choose (clientNo, propertyNo) as Primary key

<u>clientNo</u>	cName	<u>propertyNo</u>	pAddress	rentStart	rentFinish	rent	ownerNo	oName
CR76	John Kay	PG4	6 Lawrence St	1-Jul-00	31-Aug-01	350	CO40	Tina Murphy
CR76	John Kay	PG16	5 Novar Dr	1-Sep-02	1-Sep-02	450	CO93	Tony Shaw
CR56	Aline Stewart	PG4	6 Lawrence St	1-Sep-99	10-Jun-00	350	CO40	Tina Murphy
CR56	Aline Stewart	PG36	2 Manor Rd	10-Oct-00	1-Dec-01	375	CO93	Tony Shaw
CR56	Aline Stewart	PG16	5 Novar Dr	1-Nov-02	10-Aug-03	450	CO93	Tony Shaw

# ***Client rental 1NF relation***

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Write down the ClientRental relation using standard notation

ClientRental (clientNo, propertyNo, cName, pAddress,  
rentStart, rentFinish, rent, ownerNo, oName)

**WE NOW HAVE A TABLE IN 1NF.**

# *Second Normal Form (2NF)*

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Based on concept of full functional dependency:

- $A_1, \dots, A_n$  and  $B$  are attributes of a relation,
- $B$  is fully dependent on  $A_1, \dots, A_n$  if  $B$  is functionally dependent on  $A_1, \dots, A_n$  but not on any proper subset of  $A_1, \dots, A_n$ .

2NF - A relation that is in 1NF and every non-primary-key attribute is fully functionally dependent on the whole primary key.

# *Getting from 1NF to 2NF*

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**Identify primary key for the 1NF relation.**

**Identify functional dependencies in the relation.**

**If partial dependencies exist on the primary key  
remove them by placing them in a new  
relation along with copy of their determinant.**

# *ClientRental Example: 1NF to 2NF*

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<u>clientNo</u>	<u>propertyNo</u>	cName	pAddress	rentStart	rentFinish	rent	owner No	oName
CR76	PG4	John Kay	6 Lawrence St	1-Jul-00	31-Aug-01	350	CO40	Tina Murphy
CR76	PG16	John Kay	5 Novar Dr	1-Sep-02	1-Sep-02	450	CO93	Tony Shaw
CR56	PG4	Aline Stewart	6 Lawrence St	1-Sep-99	10-Jun-00	350	CO40	Tina Murphy
CR56	PG36	Aline Stewart	2 Manor Rd	10-Oct-00	1-Dec-01	375	CO93	Tony Shaw
CR56	PG16	Aline Stewart	5 Novar Dr	1-Nov-02	10-Aug-03	450	CO93	Tony Shaw

Partial dependencies are:

- clientNo → cName
- propertyNo → pAddress, rent, ownerNo, oName



## ***ClientRental Example: clientNo → cName***

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Create new relation **Client**, with primary key clientNo

<u>clientNo</u>	cName
CR76	John Kay
CR56	Aline Stewart

Remove cName from the **ClientRental** relation

<u>clientNo</u>	<u>propertyNo</u>		pAddress	rentStart	rentFinish	rent	owner No	oName
CR76	PG4		6 Lawrence St	1-Jul-00	31-Aug-01	350	CO40	Tina Murphy
CR76	PG16		5 Novar Dr	1-Sep-02	1-Sep-02	450	CO93	Tony Shaw
CR56	PG4		6 Lawrence St	1-Sep-99	10-Jun-00	350	CO40	Tina Murphy
CR56	PG36		2 Manor Rd	10-Oct-00	1-Dec-01	375	CO93	Tony Shaw
CR56	PG16		5 Novar Dr	1-Nov-02	10-Aug-03	450	CO93	Tony Shaw

## *ClientRental Example:*

*propertyNo -> pAddress, rent, ownerNo, oName*

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Create new relation **PropertyOwner**, with primary key propertyNo

<u>propertyNo</u>	pAddress	rent	ownerNo	oName
PG4	6 Lawrence St	350	CO40	Tina Murphy
PG16	5 Novar Dr	450	CO93	Tony Shaw
PG36	2 Manor Rd	375	CO93	Tony Shaw

Remove attributes pAddress, rent, ownerNo, oName from the **ClientRental** relation

<u>clientNo</u>	<u>propertyNo</u>			rentStart	rentFinish			
CR76	PG4			1-Jul-00	31-Aug-01			
CR76	PG16			1-Sep-02	1-Sep-02			
CR56	PG4			1-Sep-99	10-Jun-00			
CR56	PG36			10-Oct-00	1-Dec-01			
CR56	PG16			1-Nov-02	10-Aug-03			

# *ClientRental Example: 2NF relations*

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Tidy up, and re-name the ClientRental relation “Rental”

<u>clientNo</u>	<u>propertyNo</u>	rentStart	rentFinish
CR76	PG4	1-Jul-00	31-Aug-01
CR76	PG16	1-Sep-02	1-Sep-02
CR56	PG4	1-Sep-99	10-Jun-00
CR56	PG36	10-Oct-00	1-Dec-01
CR56	PG16	1-Nov-02	10-Aug-03

**Write down the 2NF relations:**

Client ( clientNo, cName)

PropertyOwner ( propertyNo, pAddress, rent, ownerNo, oName)

Rental ( clientNo, propertyNo, rentStart, rentFinish)

# *Third Normal Form (3NF)*

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Based on concept of transitive dependency:

- A, B and C are attributes of a relation such that  $A \rightarrow B$  and  $B \rightarrow C$ ,
- then C is transitively dependent on A through B. (Provided that A is not functionally dependent on B or C).

3NF - A relation that is in 1NF and 2NF and in which no non-primary-key attribute is transitively dependent on the primary key.

# *Getting from 2NF to 3NF*

---

Identify the primary key in the 2NF relation.

Identify functional dependencies in the relation.

If transitive dependencies exist on the primary key remove them by placing them in a new relation along with copy of their determinant.

# *ClientRental Example: 2NF to 3NF*

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Consider the relation:

PropertyOwner ( propertyNo, pAddress, rent, ownerNo, oName)

We have functional dependencies

- $\text{propertyNo} \rightarrow \text{ownerNo}$
- $\text{ownerNo} \rightarrow \text{oName}$

So oName is transitively dependent on  
propertyNo, the primary key

## *ClientRental Example:*

### *Remove Transitive Dependency On Primary Key*

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Create new relation **Owner**, with primary key ownerNo and attribute oName

<u>ownerNo</u>	oName
CO40	Tina Murphy
CO93	Tony Shaw

**Remove oName from PropertyOwner relation**

<u>propertyNo</u>	pAddress	rent	ownerNo	
PG4	6 Lawrence St	350	CO40	
PG16	5 Novar Dr	450	CO93	
PG36	2 Manor Rd	375	CO93	

# *ClientRental Example: 3NF Relations*

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Tidy up, and re-name PropertyOwner relation  
“PropertyForRent”

<u>propertyNo</u>	pAddress	rent	ownerNo*
PG4	6 Lawrence St	350	CO40
PG16	5 Novar Dr	450	CO93
PG36	2 Manor Rd	375	CO93

Write down the 3NF relations:

Client ( clientNo, cName)

Rental ( clientNo, propertyNo, rentStart, rentFinish)

PropertyOwner ( propertyNo, pAddress, rent, ownerNo)

Owner (ownerNo, oName)



# ***How To Show it in the Exam***

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We can represent normalisation:

- with text
- with a diagram
- with a table

### Zero Normal Form

- ClientRental (clientNo, propertyNo, cName, pAddress, rentStart, rentFinish, rent, ownerNo, oName)

### First Normal Form

ClientRental (clientNo, propertyNo, cName, pAddress, rentStart, rentFinish, rent, ownerNo, oName)

### Second Normal Form

Client ( clientNo, cName)

PropertyOwner ( propertyNo, pAddress, rent, ownerNo, oName)

Rental ( clientNo\*, propertyNo\*, rentStart, rentFinish)

### Third Normal Form

Client ( clientNo, cName)

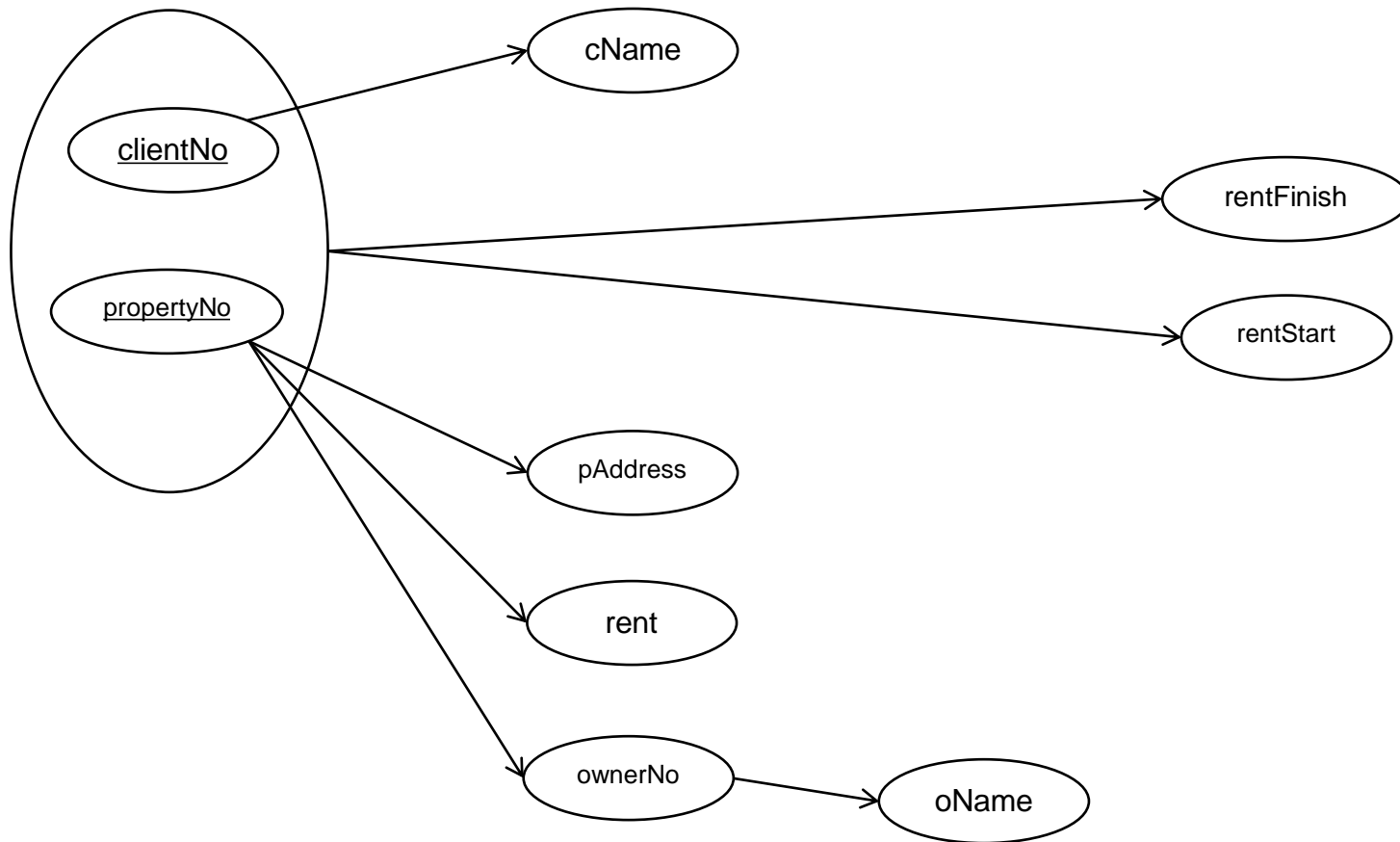
Rental ( clientNo, propertyNo, rentStart, rentFinish)

PropertyOwner ( propertyNo, pAddress, rent, ownerNo\*)

Owner (ownerNo, oName)

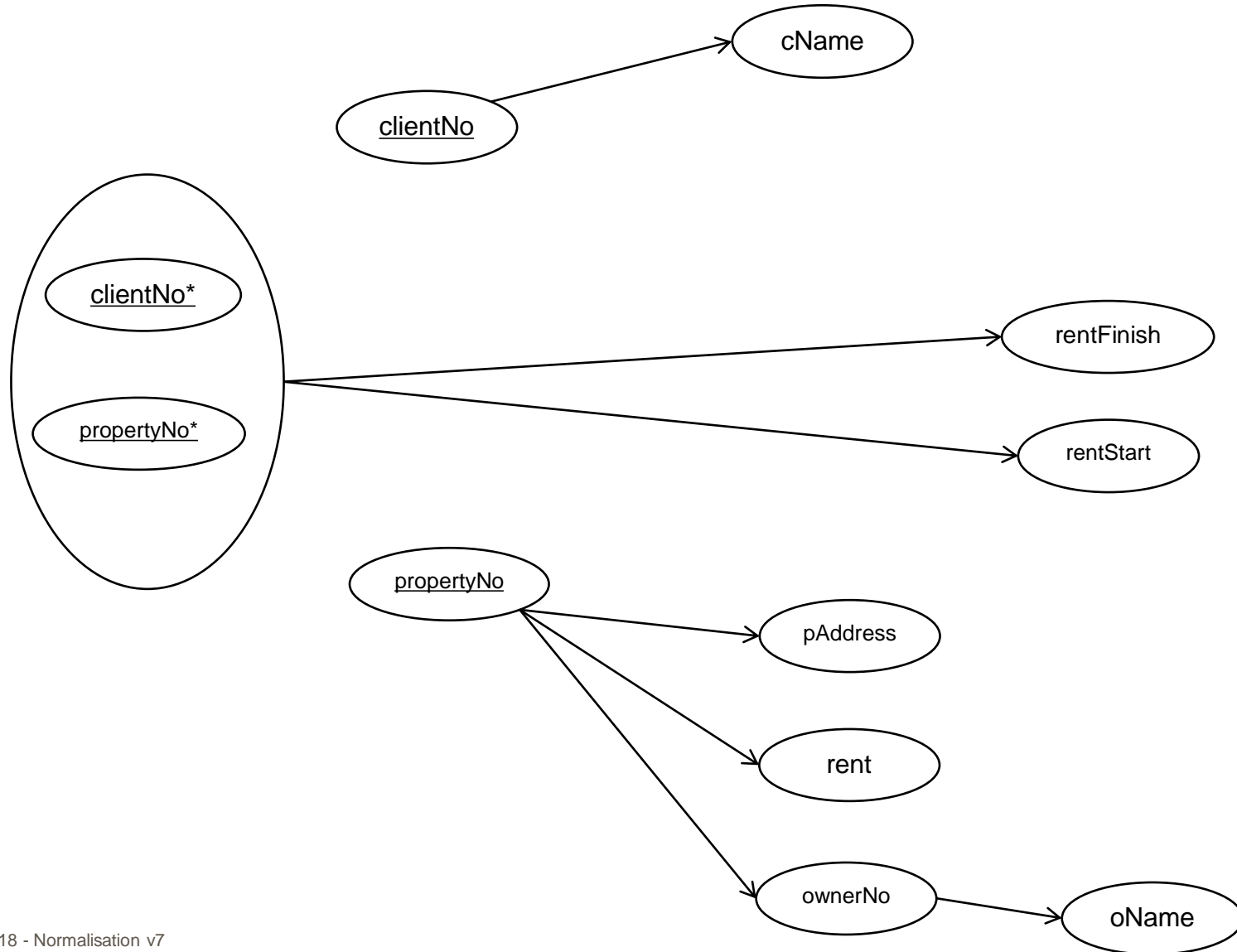
# 1NF

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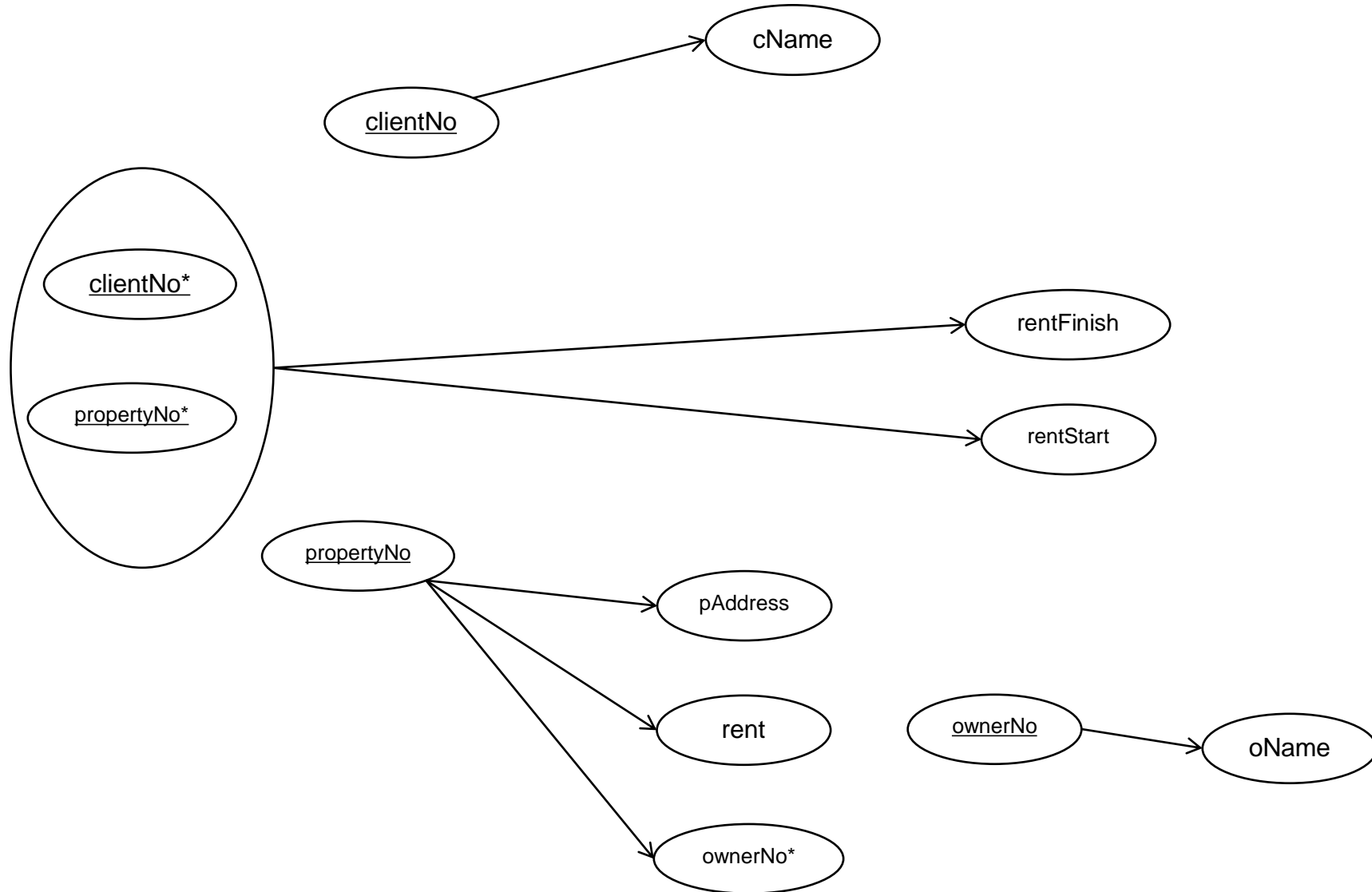
# 2NF

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# 3NF

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0NF	1NF	2NF	3NF	
clientNo propertyNo cName pAddress rentStart rentFinish Rent OwnerNo oName	<u>clientNo</u> <u>propertyNo</u> cName pAddress rentStart rentFinish Rent OwnerNo oName	<u>clientNo</u> * <u>propertyNo</u> * rentStart rentFinish	<u>clientNo</u> * <u>propertyNo</u> * rentStart rentFinish	Rental
		<u>clientNo</u> cName	<u>clientNo</u> cName	Client
		<u>propertyNo</u> pAddress Rent OwnerNo oName	<u>propertyNo</u> pAddress Rent OwnerNo*	PropertyOwner
			<u>OwnerNo</u> oName	Owner

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# Exam Questions

# *Boyce–Codd Normal Form (BCNF)*

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Based on functional dependencies that take into account all candidate keys in a relation, however BCNF also has additional constraints compared with general definition of 3NF.

BCNF - A relation is in BCNF if and only if every determinant in the relation is a candidate key for the whole relation.



# ***Boyce–Codd normal form (BCNF)***

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**For a relation with only one candidate key, 3NF and BCNF are equivalent.**

**To test whether a relation is in BCNF, we identify all the determinants and make sure that they are all candidate keys for the whole relation.**

**If they aren't, this can give rise to the following situation.**

# *Boyce–Codd normal form (BCNF)*

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It may be possible for there to be a functional dependency  $A \rightarrow B$  in a relation, such that:

B is part of a primary key attribute, and A is not a candidate key:

3NF will allow this. BCNF will not.

BCNF requires that you remove that dependency by separating out into a further table.

# *Boyce–Codd normal form (BCNF)*

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Violation of BCNF is quite rare.

Potential to violate BCNF may occur in a relation that:

- contains two (or more) composite candidate keys;
- the candidate keys overlap (ie. have at least one attribute in common).
- there is a functional dependency in a relation whose determinant is not a candidate key, but whose dependent is part of one of the candidate keys.

# BCNF Example

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Consider the relation:

ClientInterview (clientNo, interviewdate, interviewTime, staffNo, roomNo)

<u>clientNo</u>	<u>interviewDate</u>	interviewTime	staffNo	roomNo
CR76	13-May-02	10:30	SG5	G101
CR56	13-May-02	12:00	SG5	G101
CR74	13-May-02	12:00	SG37	G102
CR56	01-Jul-02	10:30	SG5	G102

A member of staff will be allocated a given room, during the day, while they are interviewing clients.

A room may be allocated to several members of staff throughout the day.

A client is only interviewed once on a given date, but maybe called back for other interviews.

The member of staff may be given a different room on a different day.

# BCNF Example: Functional dependencies

<u>clientNo</u>	<u>interviewDate</u>	<u>interviewTime</u>	<u>staffNo</u>	<u>roomNo</u>
CR76	13-May-02	10:30	SG5	G102
CR56	13-May-02	12:00	SG5	G102
CR74	13-May-02	12:00	SG5	G102
CR56	01-Jul-02	10:30	SG5	G102

There are three candidate keys for this relation. They all overlap on the interviewDate attribute.

Functional dependencies are

fd1: clientNo, interviewdate → interviewTime, staffNo, roomNo

fd2: staffNo, interviewdate, interviewTime → clientNo, roomNo

fd3: roomNo, interviewdate, interviewTime → staffNo, clientNo

fd4: staffNo, interviewdate → roomNo

# ***BCNF Example***

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ClientInterview relation is in 3NF

Is ClientInterview relation in BCNF?

fd1:                    the determinant, (clientNo, interviewdate), is  
                          the primary key

fd2 and fd3:   both determinants are candidate keys for the  
                          whole table

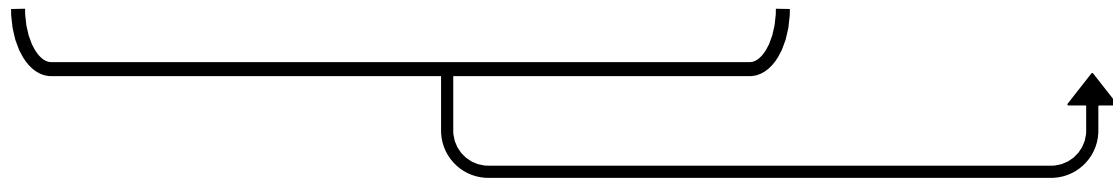
fd4:                    its determinant (staffNo, InterviewDate) is  
                          NOT a candidate key for the whole table, but  
                          it's dependent (room\_No) IS part of one of  
                          the other candidate keys.

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It is fd4 that will give rise to the possibility of update anomalies:

fd4: staffNo, interviewdate → roomNo

<u>clientNo</u>	<u>interviewDate</u>	interviewTime	staffNo	roomNo
CR76	13-May-02	10:30	SG5	G101
CR56	13-May-02	12:00	SG5	G101
CR74	13-May-02	12:00	SG37	G102
CR56	01-Jul-02	10:30	SG5	G102



Where is the redundancy?

# BCNF Example

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Create new relation **StaffRoom** with the attributes from fd4:

- Determinant attributes form the primary key of the new relation
- Include dependent attributes in the relation

Now can you see the duplication?

<u>interviewDate</u>	<u>staffNo</u>	roomNo
13-May-02	SG5	G101
13-May-02	SG5	G101
13-May-02	SG37	G102
01-Jul-02	SG5	G102

<u>interviewDate</u>	<u>staffNo</u>	roomNo
13-May-02	SG5	G101
13-May-02	SG37	G102
01-Jul-02	SG5	G102



# BCNF Example

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Remove the dependent attributes of fd4 from the ClientInterview relation and rename it **Interview**.

<u>clientNo</u>	<u>interviewDate</u>	interviewTime	staffNo
CR76	13-May-02	10:30	SG5
CR56	13-May-02	12:00	SG5
CR74	13-May-02	12:00	SG37
CR56	01-Jul-02	10:30	SG5

**Write down the BCNF relations:**

**StaffRoom (staffNo, interviewdate, roomNo)**

**Interview (clientNo, interviewdate, interviewTime, staffNo)**