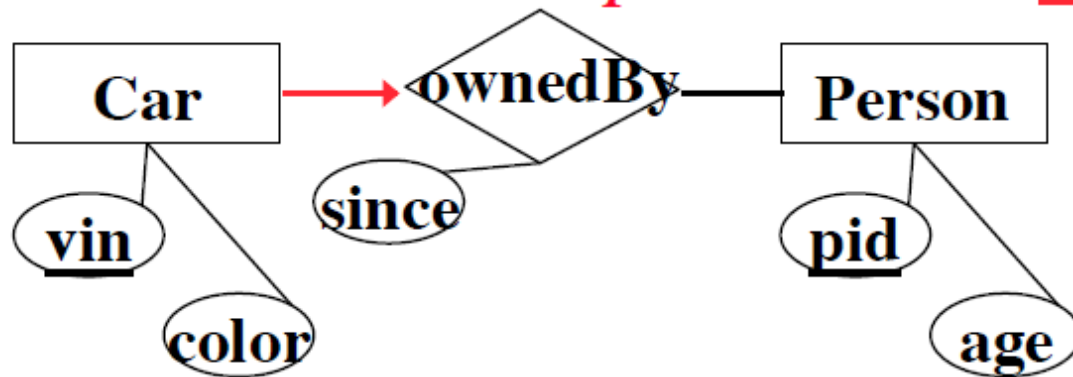


1-N Relationships Revisited e.g.



- **Merge** T_ownedBy(vin,pid,since) **into** T_Car(vin, color)

```
table T_Car
( vin char(25),
  color,
  pid int references T_Person,
  since date,
  primary key ( vin ) )
```

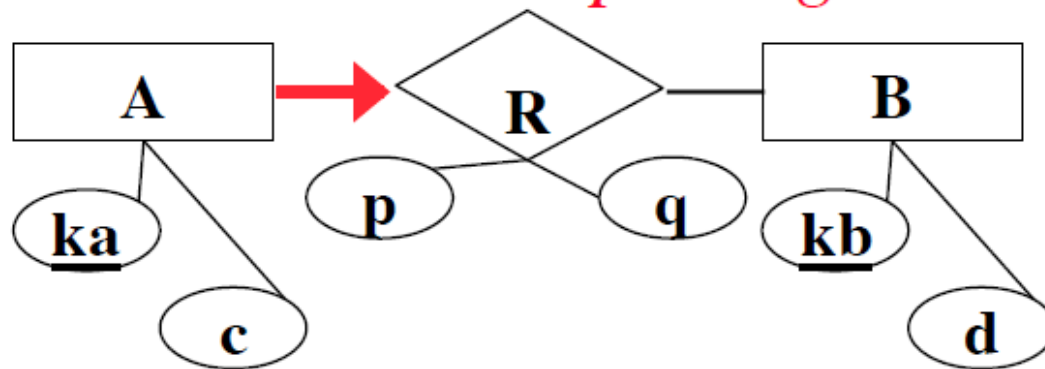
}

Describes a Car

}

*Describes the
Person related to
the car by the
ownedBy relation*

1-N Relationship merged: what if total?



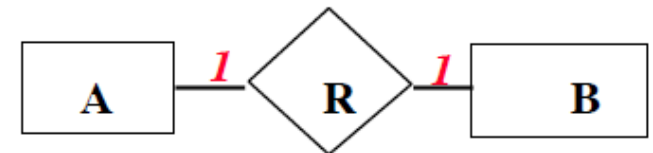
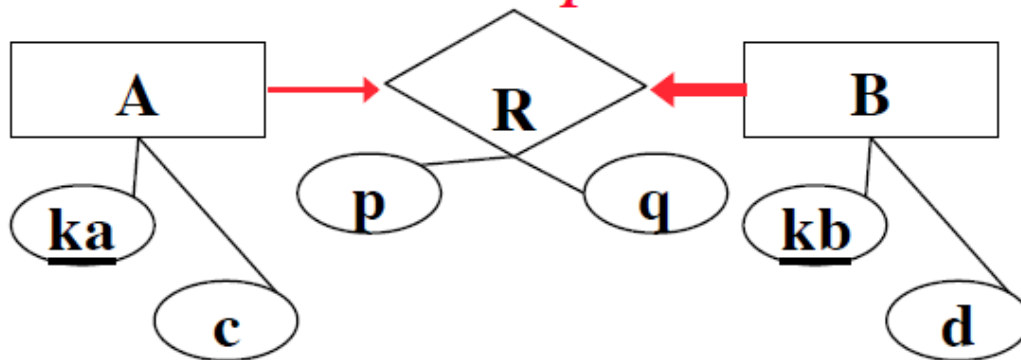
- **Merge** T_R(ka,kb,p,q) **into** T_A(ka, c)

```
table T_A
( ka
  c
  R_kb references T_B not null
  p
  q
  primary key ( ka )
```

*The “total”
constraint can
now be captured
in one table!*

Note that if B participated “total”, there was still no way to capture this in the the T_A table, merged with R. Need a general assertion.

1-1 Relationships Revisited



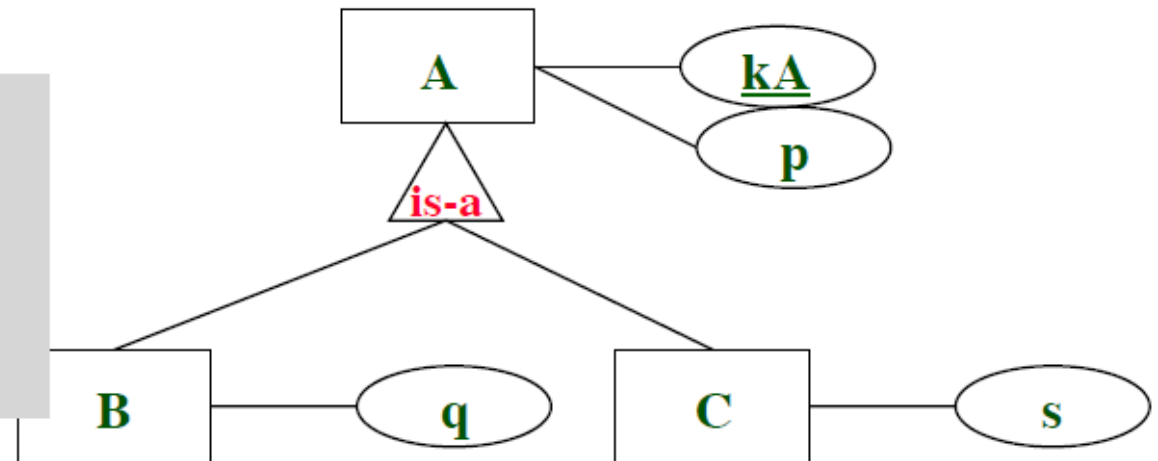
- *Can merge T_R into T_A **or** T_B ?*
- *Choose “total” end, so NOT NULL constraint enforces it*

```

table T_B
    ( kb
      d
    }   Describes a B
    R_ka references T_A not null
    p
    q
    primary key ( kb )
  
```

Mapping Subclasses - 2nd way (merge up)

```
table T_A(  
    kA primary key  
    p  
    q  
    s )
```



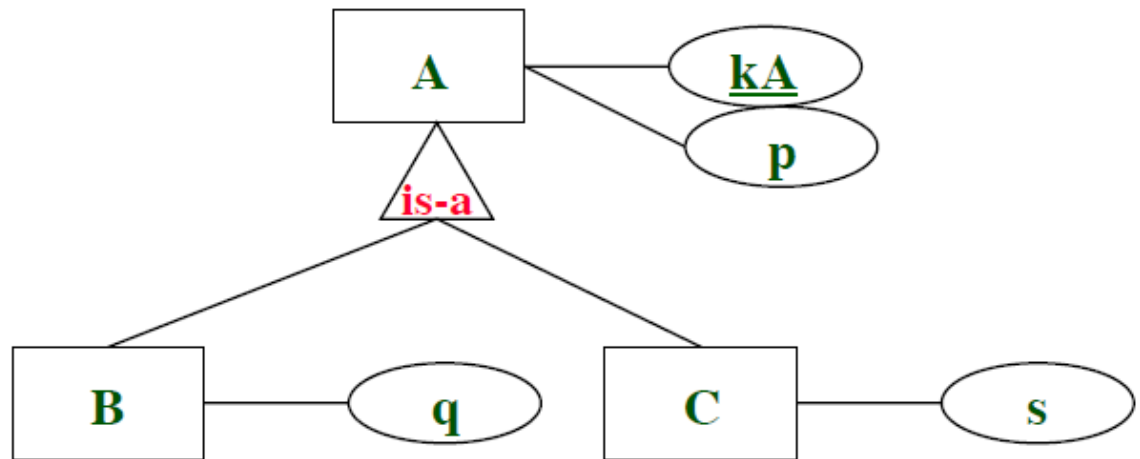
- Merge T_B (and/or T_C) into T_A: *note that merge conditions hold!*
- Note that objects in B but not in C will have null s; and objects in A but not B nor C will have null q and s.
- *Potentially serious violation of “know membership” rule:*
 - » **How can you tell if an object is in subclass B? Only if q value is not null. So this design only applicable when each subclass has some non-null attribute that can act as a discriminator.**
 - » ***Otherwise need a boolean flag column to indicate the subclass:**
inB? boolean,
inC? boolean,

What about total and disjoint constraints?

**Mapping Subclasses: 3rd way (collapse down - not a 'merge')*

```
table T_B(  
    kA primary key  
    p  
    q )
```

```
table T_C(  
    kA primary key  
    p  
    s )
```



- This design is used when subclasses **cover total** superclass
 - » otherwise, to represent membership in class A alone is “fake”: need to mark “non-Bs” in T_B and “non-Cs” in T_C (e.g., assume *q* and *s* are not null for B and C)
- and are **disjoint**
 - » otherwise A information, such as *kA*, *p* is repeated in both tables for some object <-- against no duplication principle (see later)

DB table design principles that motivate merging tables:

From database point of view: storage efficiency, making queries easier to state, making integrity checking faster, etc.

(i) Repetition of data is bad if it is **unnecessary**

e.g. `enrolledIn(studntId, courseId, studntAddress, grade)`

because

- » **need to verify consistency** (*do all occurrences of `studntId=4454` have the same address?*)
- » **need to maintain consistency** (*when `studntID 4454` changes address, all tuples need to be sought and changed*)
- » **wasted space**

The above example involved repeating non-key information (*studntAddress*). Note that key info cannot be repeated by definition of “key”.

DB Table design principles (cont'd)

- (ii) Having fewer tables is better because queries require fewer joins, which are**
 - » easier to write for the user
 - » less expensive to evaluate for the DBMS
- (iii) Having tables with too many nulls is undesirable**

As usual, there are trade-offs.

Merge rule is a consequence of DB design principles

If you have two tables of the form $T(\underline{K}, X)$ and $S(\underline{K2}, Y)$, Y not null, where $K2$ is a foreign key referencing $T(K)$ then you can merge S into T to get instead a single table $TS(\underline{K}, X, Y)$. Column Y will have NULLs for all keys in T but not S .

- This is good by principle (ii): avoid joins.
- By def'n of key, for every K' value there is only one Y , and for every K there is only one X , so when you combine them you get only 1 row, and principle “(i) *no repetition*” is preserved
- If $T('ann', 45)$ has no matching $S('ann', ...)$, we need to use null in final table: $T('ann', 45, \text{NULL})$. This may be a problem with “(iii) not too many nulls”, if there are very many columns in S , and very many T 's that are not in S .