

# Chapter 1: roadmap

1.1 what *is* the Internet?

1.2 network edge

- end systems, access networks, links

1.3 network core

- packet switching, circuit switching, network structure

1.4 delay, loss, throughput in networks

1.5 protocol layers, service models

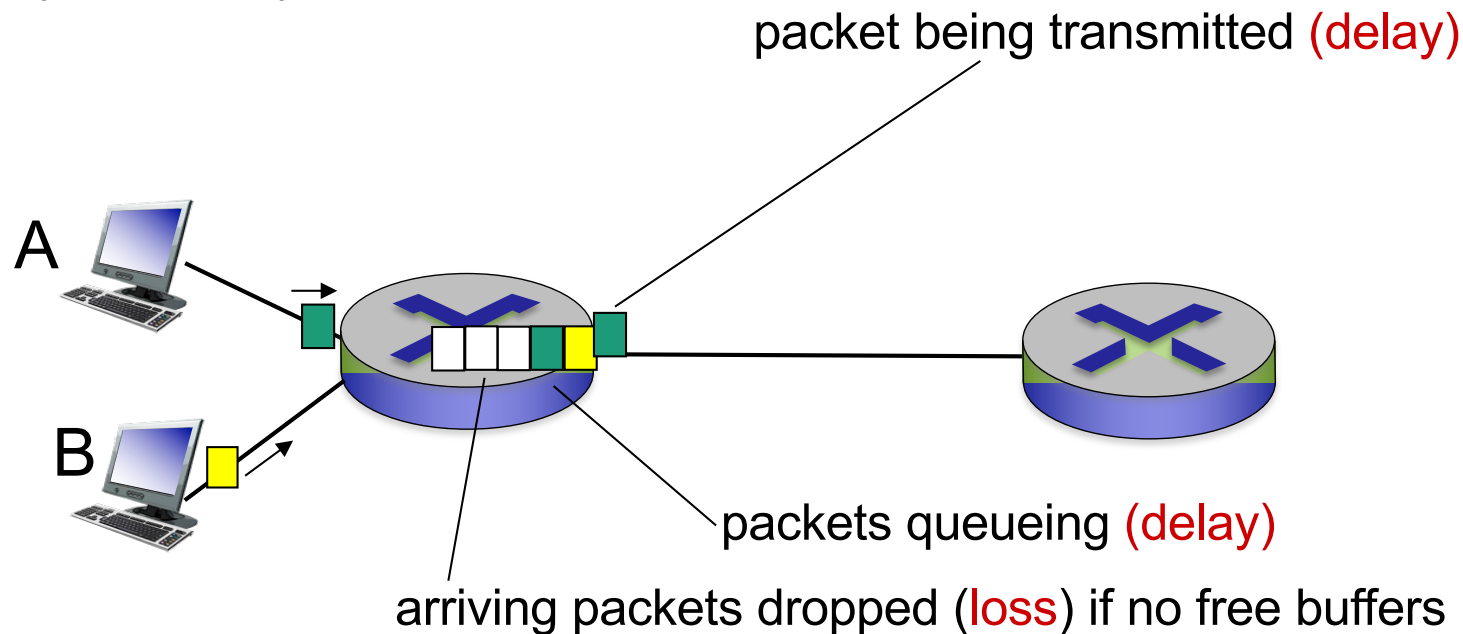
1.6 networks under attack: security

1.7 history

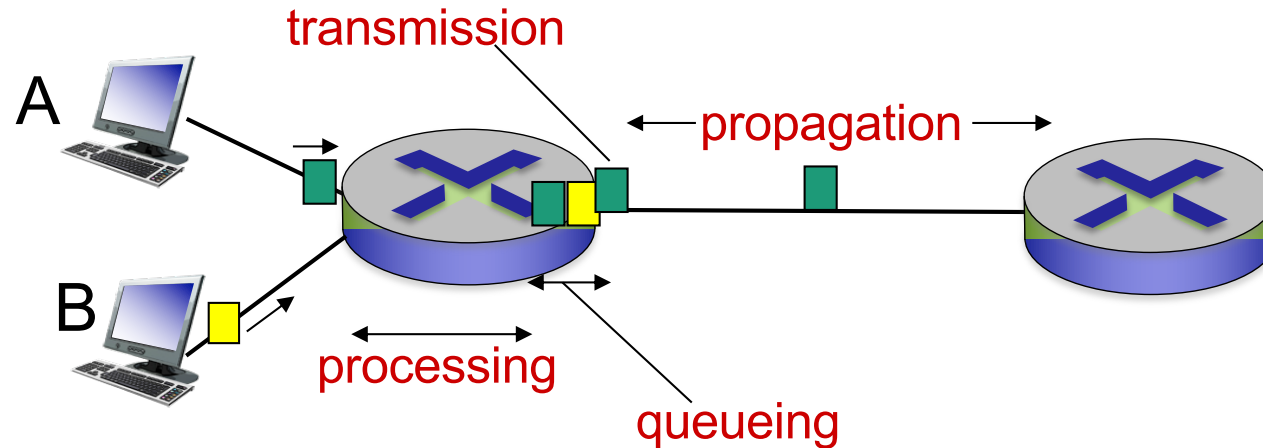
# How do loss and delay occur?

packets *queue* in router buffers

- packet arrival rate to link (temporarily) exceeds output link capacity
- packets queue, wait for turn



# Four sources of packet delay



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

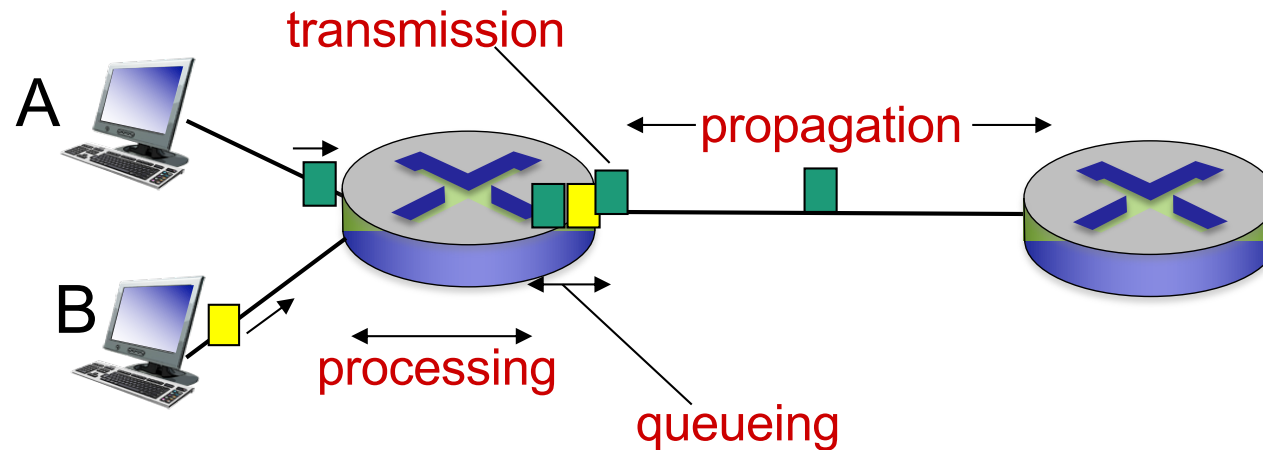
## $d_{\text{proc}}$ : nodal processing

- check bit errors
- determine output link
- typically < msec

## $d_{\text{queue}}$ : queueing delay

- time waiting at output link for transmission
- depends on congestion level of router

# Four sources of packet delay



$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

$d_{\text{trans}}$ : transmission delay:

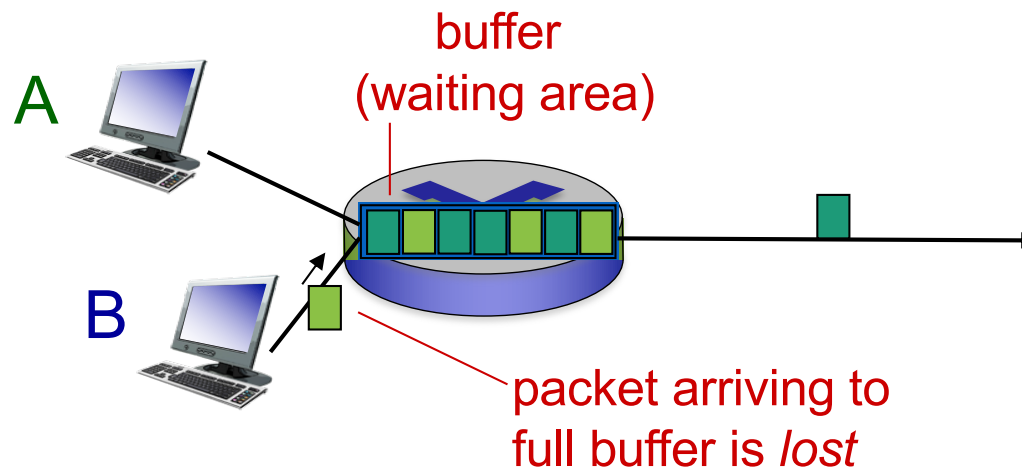
- $L$ : packet length (bits)
  - $R$ : link bandwidth (bps)
  - $d_{\text{trans}} = L/R$
- ←  $d_{\text{trans}}$  and  $d_{\text{prop}}$  →  
very different

$d_{\text{prop}}$ : propagation delay:

- $d$ : length of physical link
- $s$ : propagation speed ( $\sim 2 \times 10^8$  m/sec)
- $d_{\text{prop}} = d/s$

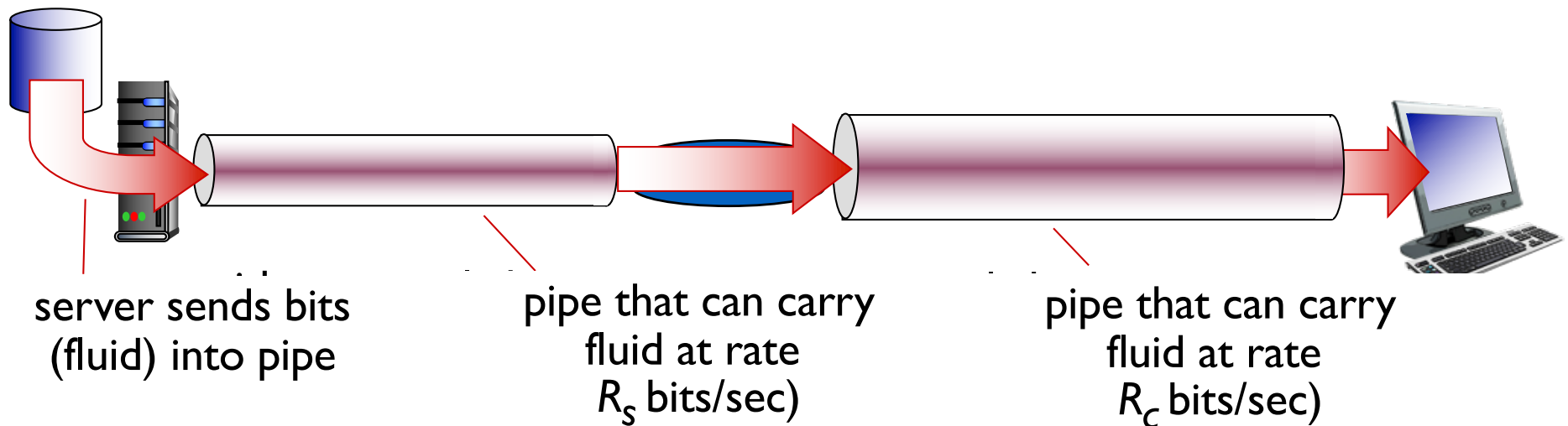
# Packet loss

- queue (aka buffer) in buffer has **limited** capacity
- packet dropped if arriving to **full** queue (aka lost)
- lost packet may be retransmitted
  - by previous node
  - by source end system
  - or not at all



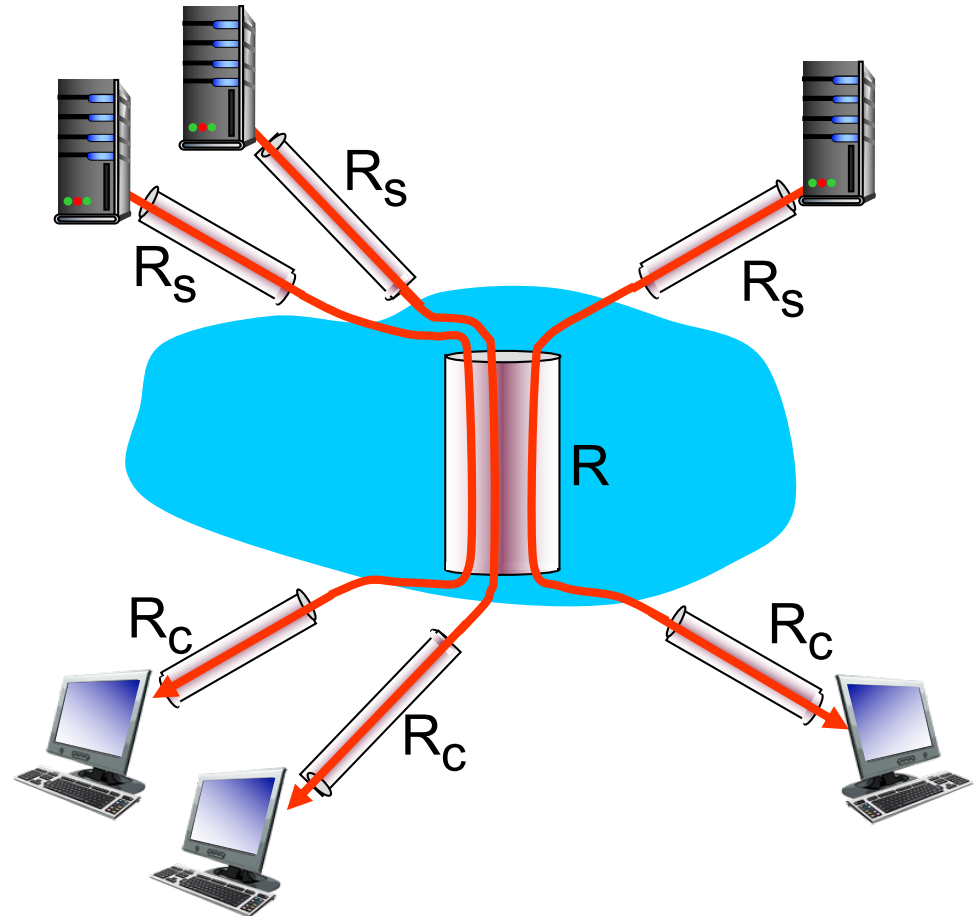
# Throughput

- *throughput*: rate (bits/time unit) at which bits transferred between sender/receiver
  - *Real-time*: rate at a given time point
  - *Average*: rate over longer period of time



# Throughput: Internet scenario

- per-connection end-end throughput:  
 $\min(R_c, R_s, R/10)$
- in practice:  $R_c$  or  $R_s$  is often bottleneck



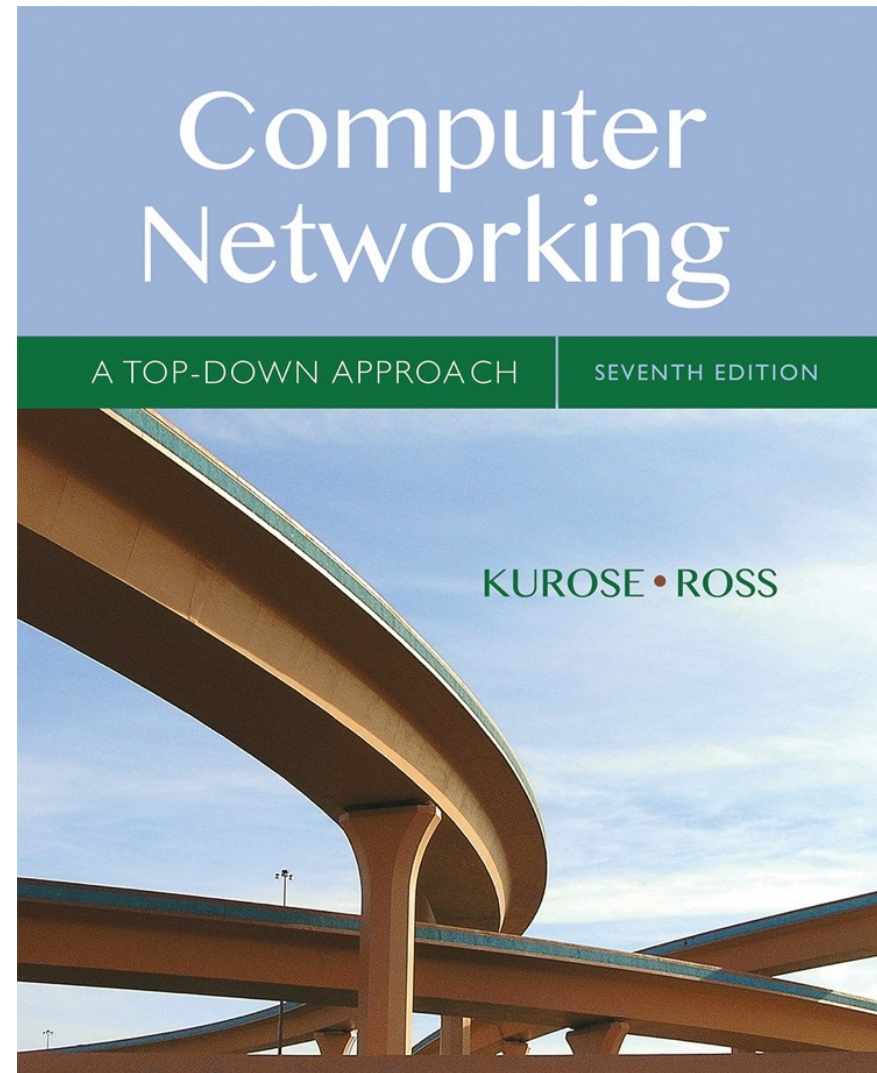
10 connections (fairly) share  
**bottleneck link**  $R$  bits/sec



# CS 352: Internet Technology (Section 3, 4, 5)

Computer Science  
Rutgers University

Fall 2019





# SPN Requests

- For all students who need a SPN
  - We are currently overwhelmed with SPN requests
  - We are constrained by classroom capacity, TA resources, class size, etc
  - We already issued a few SPNs to some students based on the department policy
  - We may be able to issue a few more if some students drop this course, but the chance is low.
  - For these who did not get SPNs so far, please try to register next semester
- All SPN questions should be sent to Prof. Zhang at [d.z@rutgers.edu](mailto:d.z@rutgers.edu)

# Class Website

The screenshot shows the Sakai LMS interface for the course CS 352: Internet Technology (Section 3, 4 and 5) at Rutgers University-Fall 2019. The page is titled "CS 352: INTERNET TECHNOLOGY (Section 3, 4 and 5)" and "Rutgers University-Fall 2019". The left sidebar contains navigation links: Home, Announcements, Chat Room, Email Archive, Mailtool, Messages, Gradebook, Calendar, Section Info, Site Members, Statistics, Assignments, Tests & Quizzes, Piazza, Site Info, and Help. The main content area is divided into sections: Information, Recitation, Sec 3 TA, Sec 4 TA, Sec 5 TA, Textbooks, and Grading. The right sidebar contains Recent Announcements, a Calendar for September 2019, and Message Center Notifications.

**CS 352: INTERNET TECHNOLOGY (Section 3, 4 and 5)**  
*Rutgers University-Fall 2019*

**Information**

**Instructor:** Desheng Zhang  
**Email:** d.z AT rutgers.edu  
**Office:** CoRE 307

**Lectures:** Mon and Wed, 5:00-6:20 pm  
**Classroom:** RWH 105  
**Office Hours:** Monday, 1:00-2:00pm

**Recitation:** Section 3 Tuesday 5:15 PM - 6:10 PM BUS PH-115;  
Section 4 Thursday 8:25 PM - 9:20 PM BUS SEC-203;  
Section 5 Wednesday 8:25 PM - 9:20 PM BUS SEC-202;

**Sec 3 TA:** Xiaoyang Xie  
(Email: xx88@cs.rutgers.edu)  
Office Hour: TBD.

**Sec 4 TA:** Abraham Gale  
(Email: ajg317@scarletmail.rutgers.edu)  
Office Hour: TBD.

**Sec 5 TA:** Behnam Babagholami Mohamadabadi  
(Email: bb510@cs.rutgers.edu)  
Office Hour: TBD.

**Textbooks:** James Kurose and Keith Ross, Computer Networking: A Top-Down Approach, 7th Edition.

**Grading:** In-class Quizzes: 10% of grade  
Homework: 10% of grade  
Midterm examination #1: 15% of grade  
Midterm examination #2: 15% of grade  
Programming assignment #1, #2: 20% of grade  
Final examination: 30% of grade

**Recent Announcements**  
(viewing announcements from the last 10 days)

**CS 352 Section 3.4.5: NO Recitation in the first week**  
(Desheng Zhang - Sep 2, 2019 5:21 pm)

**CALENDAR**  
OPTIONS PUBLISH (PRIVATE)

**September 2019**

Sun	Mon	Tue	Wed	Thu
1	2	3	4	
8	9	10	11	
15	16	17	18	
22	23	24	25	
29	30	1	2	

**MESSAGE CENTER NOTIFICATIONS**  
New Messages

- on Sakai
- also <https://www.cs.rutgers.edu/~dz220/CS352F2019.html>

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1.5 protocol layers, service models

1.6 networks under attack: security

1.7 history

# Protocol “layers”

*Networks are complex,  
with many “pieces”:*

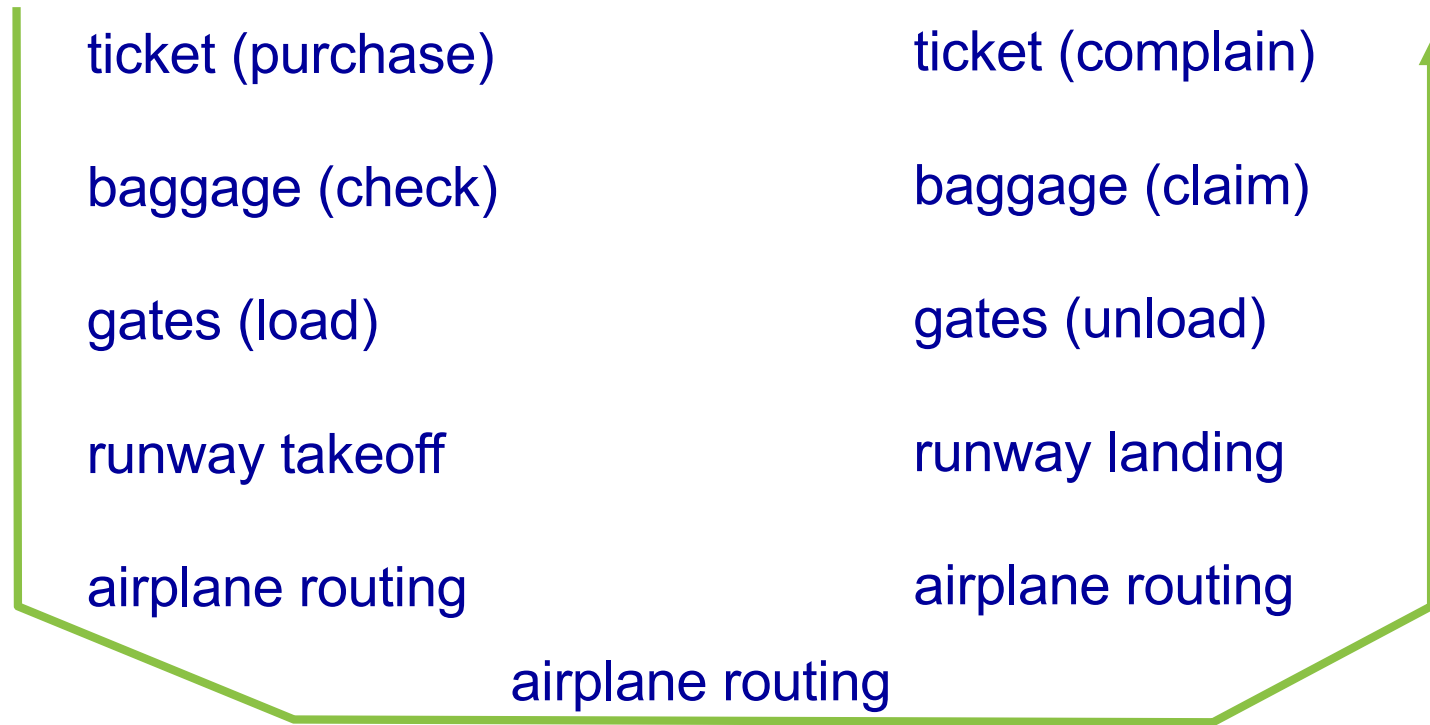
- hosts
- routers
- links of various media
- applications
- protocols
- hardware, software

*Question:*

is there any hope of  
*organizing* structure of  
network?

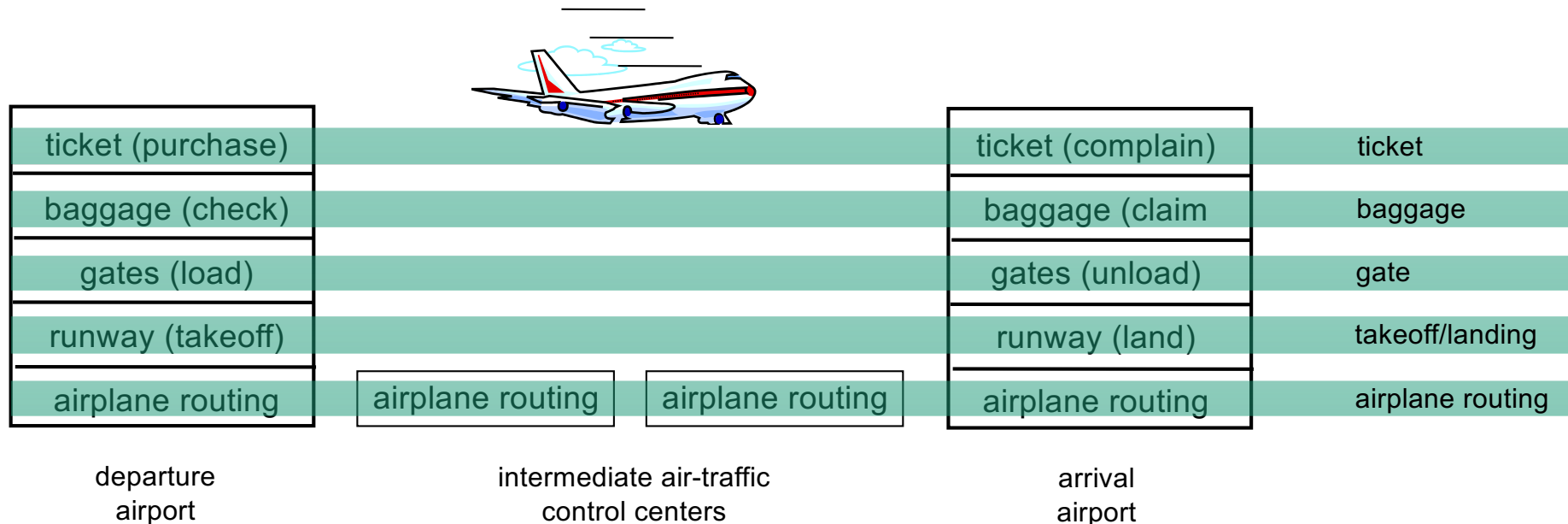
.... or at least our discussion  
of networks?

# Organization of air travel



- a series of steps

# Layering of airline functionality



*layers:* each layer implements a service

- via its own internal-layer actions
- relying on services provided by layer below

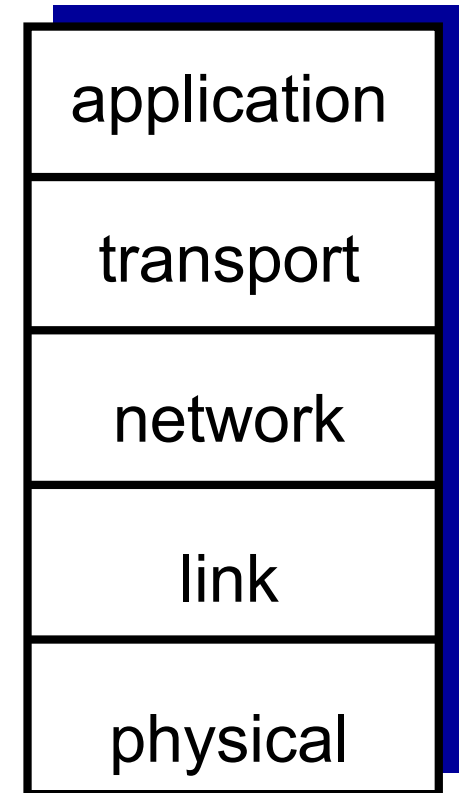
# Why layering?

dealing with complex systems:

- Layering structure allows **understanding** of relationship of complex system's pieces
- Layering eases **maintaining and updating** of system
  - change of implementation of layer's service transparent to rest of system
  - e.g., change in gate procedure doesn't affect rest of system
- layering considered harmful?

# Internet protocol stack

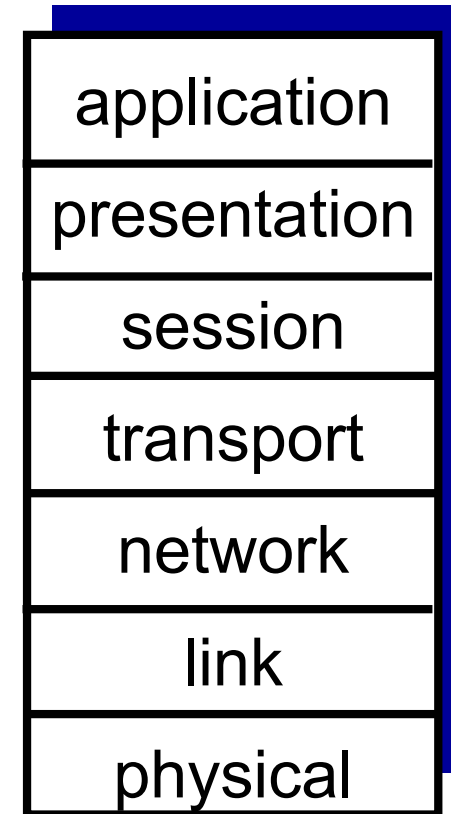
- *application*: supporting network applications
  - FTP, SMTP, HTTP
- *transport*: process-process data transfer
  - TCP, UDP
- *network*: routing of datagrams from source to destination
  - IP, routing protocols
- *link*: data transfer between neighboring network devices
  - Ethernet, 802.11 (WiFi)
- *physical*: bits “on the wire”





# ISO/OSI reference model

- *Open Systems Interconnections*
- *presentation*: allow applications to interpret meaning of data, e.g., encryption, compression, machine-specific conventions
- *session*: synchronization, recovery of data exchange
- Internet stack “missing” these layers!
  - these services, *if needed*, must be implemented in application



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# Network security

- **field of network security:**
  - how bad guys can **attack** computer networks
  - how we can **defend** networks against attacks
  - how to **design architectures** that are immune to attacks
- **Internet not originally designed with (much) security in mind**
  - *original vision:* “a group of **mutually trusting** users attached to a transparent network” 😊
  - security considerations in all layers!

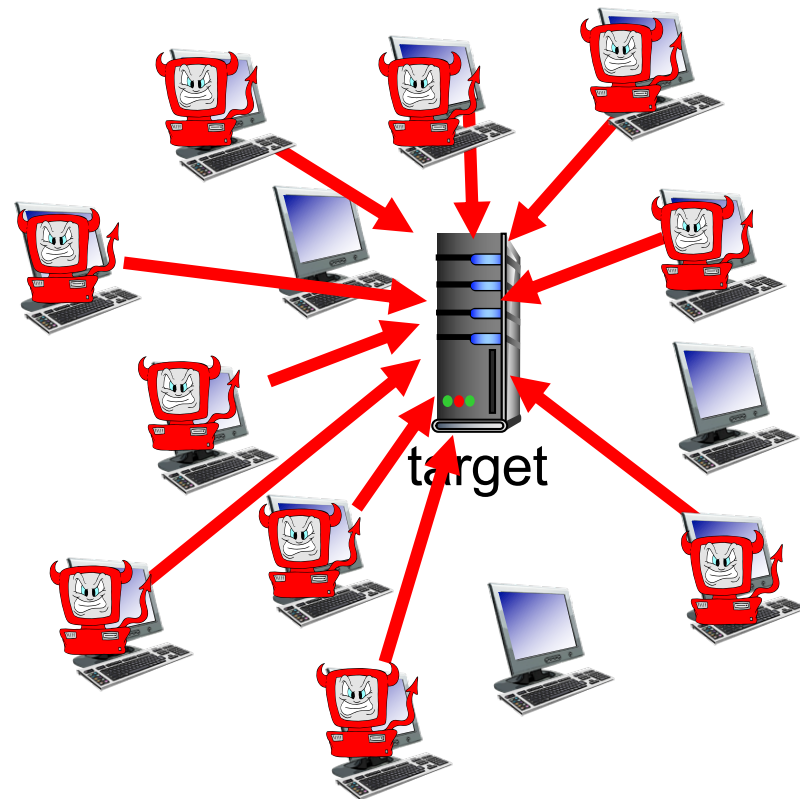
# Bad guys: put malware into hosts via Internet

- malware can get in host from:
  - *virus*: self-copying (self-replicating) infection by receiving/executing object (e.g., e-mail attachment)
  - *worm*: self-copying infection by passively receiving object that gets itself executed
- *spyware malware* can record
  - keystrokes
  - web sites visited
  - upload info to collection site

# Bad guys: attack server, network infrastructure

*Denial of Service (DoS):* attackers make **resources** (server, bandwidth) **unavailable** to legitimate traffic by overwhelming resource with bogus traffic

1. select target
2. break into hosts around the network (see botnet)
3. send packets to target from compromised hosts



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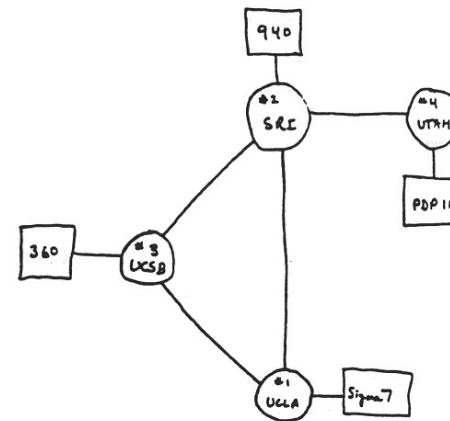
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# Internet history

## *1961-1972: Early packet-switching principles*

- 1961: Kleinrock - **queueing theory** shows effectiveness of packet-switching
- 1964: Baran - packet-switching in **military nets**
- 1967: **ARPAnet** conceived by Advanced Research Projects Agency
- 1969: first ARPAnet node operational
- 1972:
  - ARPAnet public demo
  - NCP (Network Control Protocol) **first host-host protocol**
  - first e-mail program
  - ARPAnet has 15 nodes



THE ARPA NETWORK-67

# Internet history

## *1972-1980: Internetworking, new and proprietary nets*

- 1970: **ALOHAnet** satellite network in Hawaii
- 1974: Cerf and Kahn - **architecture for interconnecting networks**
- 1976: Ethernet at Xerox PARC
- late70' s: proprietary architectures: DECnet, SNA, XNA
- late 70' s: switching **fixed length packets** (ATM precursor)
- 1979: ARPAnet has 200 nodes

### Cerf and Kahn' s internetworking principles:

- minimalism, autonomy - no internal changes required to interconnect networks
- best effort service model
- stateless routers
- decentralized control

define today' s Internet  
architecture



# Internet history

## *1980-1990: new protocols, a proliferation of networks*

- 1983: deployment of TCP/IP
- 1982: smtp e-mail protocol defined
- 1983: DNS defined for name-to-IP-address translation
- 1985: ftp protocol defined
- 1988: TCP congestion control
- new national networks: CSnet, BITnet, NSFnet, Minitel
- 100,000 hosts connected to confederation of networks

# Internet history

## *1990, 2000 's: commercialization, the Web, new apps*

- early 1990' s: ARPAnet  
decommissioned
- 1991: NSF lifts restrictions on commercial use of NSFnet  
(decommissioned, 1995)
- early 1990s: Web
  - hypertext [Bush 1945, Nelson 1960' s]
  - HTML, HTTP: Berners-Lee
  - 1994: Mosaic, later Netscape
  - late 1990' s:  
commercialization of the Web

### late 1990' s – 2000' s:

- more killer apps: instant messaging, P2P file sharing
- network security
- est. 50 million host, 100 million+ users
- backbone links running at Gbps

# Internet history

## *2005-present*

- ~5B devices attached to Internet (2016)
  - smartphones and tablets
- aggressive deployment of broadband access
- increasing ubiquity of high-speed wireless access
- emergence of online social networks:
  - Facebook: ~ one billion users
- service providers (Google, Microsoft) create their own networks
  - bypass Internet, providing “instantaneous” access to search, video content, email, etc.
- e-commerce, universities, enterprises running their services in “cloud” (e.g., Amazon EC2)

# Introduction: summary

*covered a “ton” of material!*

- Internet overview
- what’s a protocol?
- network edge, core, access network
  - packet-switching versus circuit-switching
  - Internet structure
- performance: loss, delay, throughput
- layering, service models
- security
- history

*you now have:*

- context, overview, “feel” of networking
- more depth, detail *to follow!*

# Chapter 2

## Application Layer

# App-layer protocol defines

- ❑ Types of messages exchanged,
  - ❖ e.g., request, response
- ❑ Message format:
  - ❖ Syntax :what fields in messages
  - ❖ Semantics: meaning of information in fields
- ❑ Rules for when and how processes send & respond to messages

## Public-domain protocols:

- ❑ defined in RFCs
- ❑ allows for inter-operability
- ❑ e.g., HTTP, SMTP

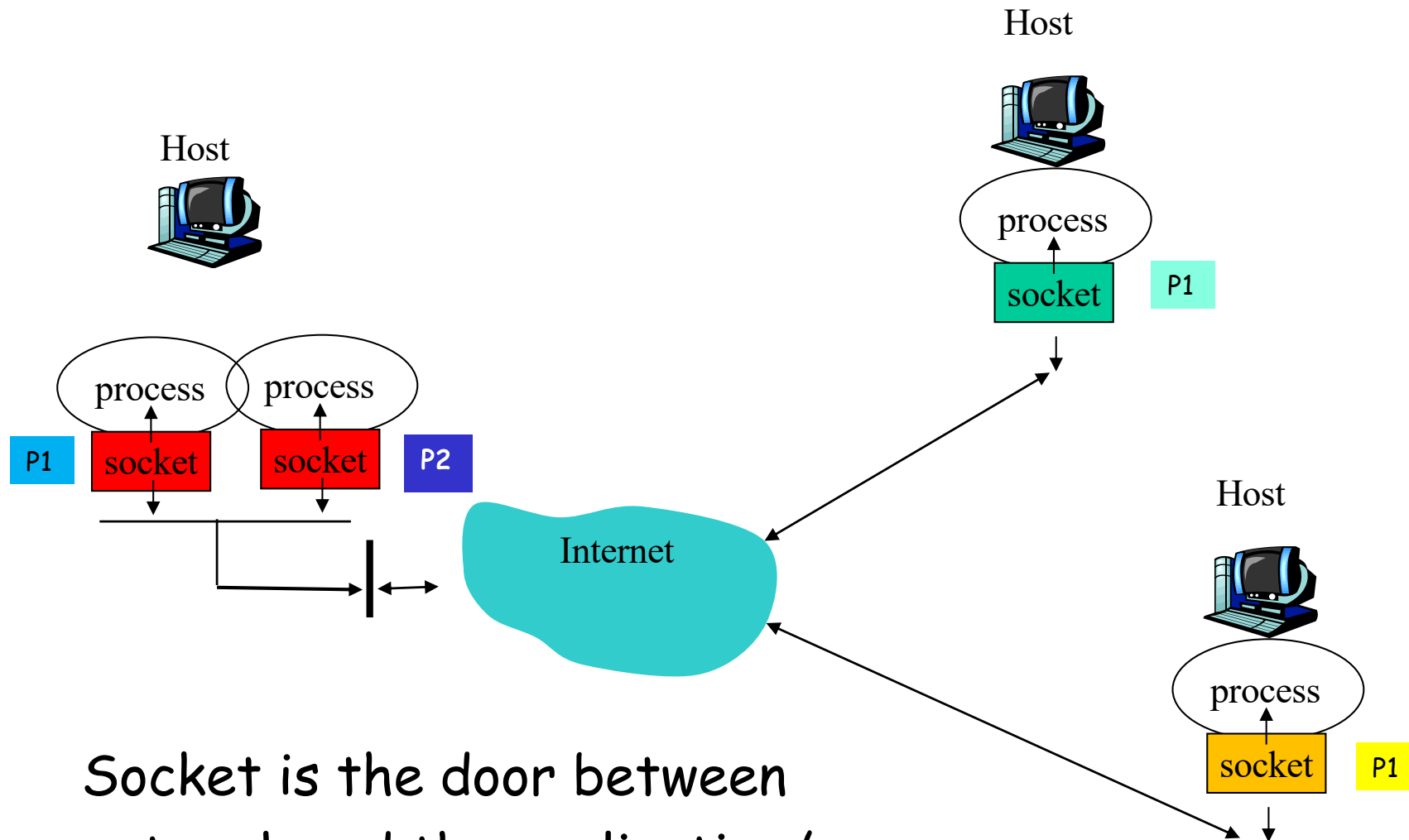
## Proprietary protocols:

- ❑ e.g., Skype, Hangout

# Network application

- ❑ Consider: Two **applications** on 2 different hosts connected by **a network**
- ❑ In order to communicate, need to **identify** the parties
- ❑ Phone network: **phone number** (10 digits)
- ❑ Computer network: **IP address**
  - ❖ IPv4 (32 bits) 128.6.24.78
  - ❖ IPv6 (128 bits)  
2001:4000:A000:C000:6000:B001:412A:8000
- ❑ In addition to find host address, we need one more.
- ❑ More than one program executing on a host
- ❑ Which Program to talk to ?
- ❑ We need another identity : port #

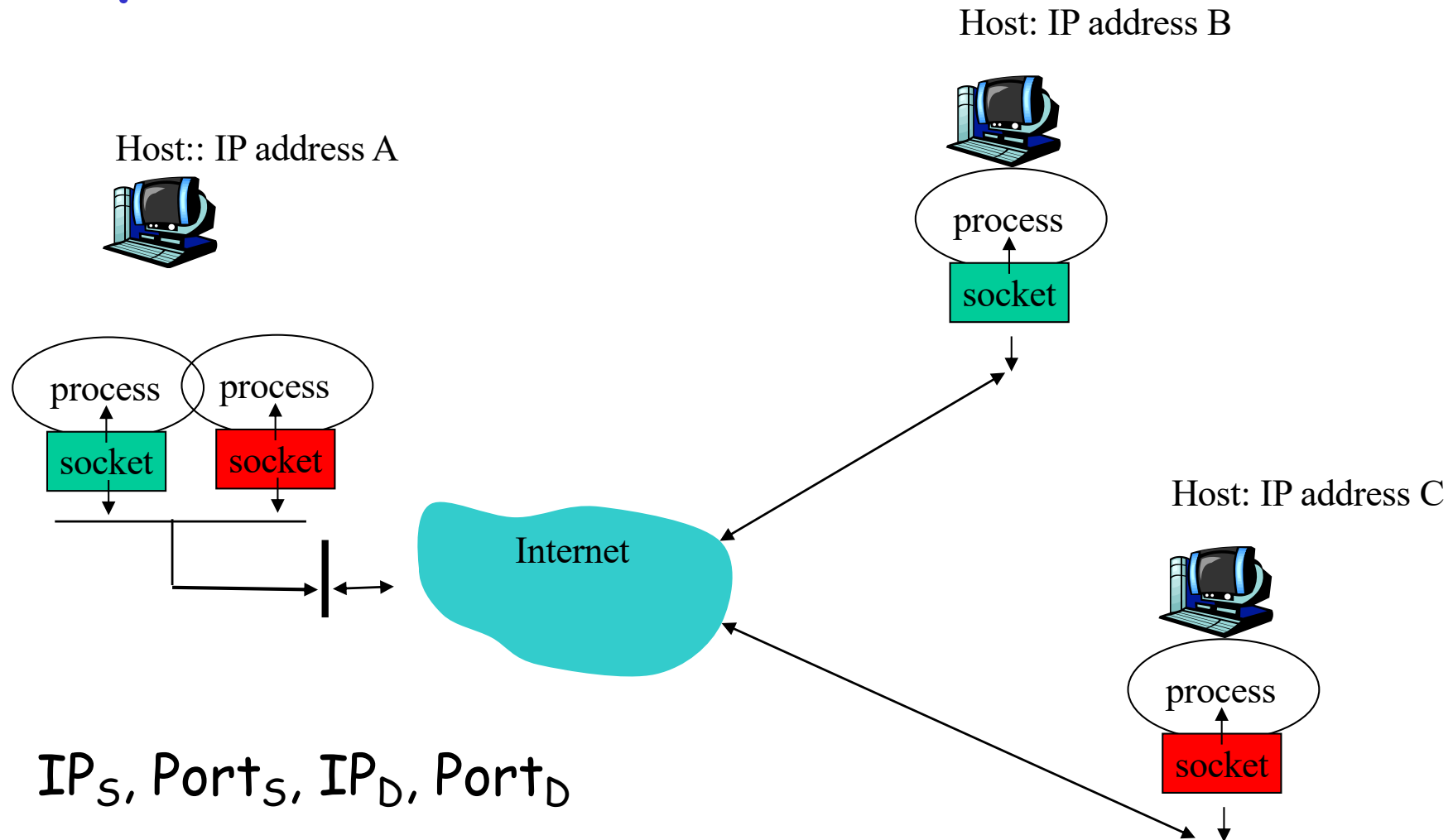
# IP address & port number



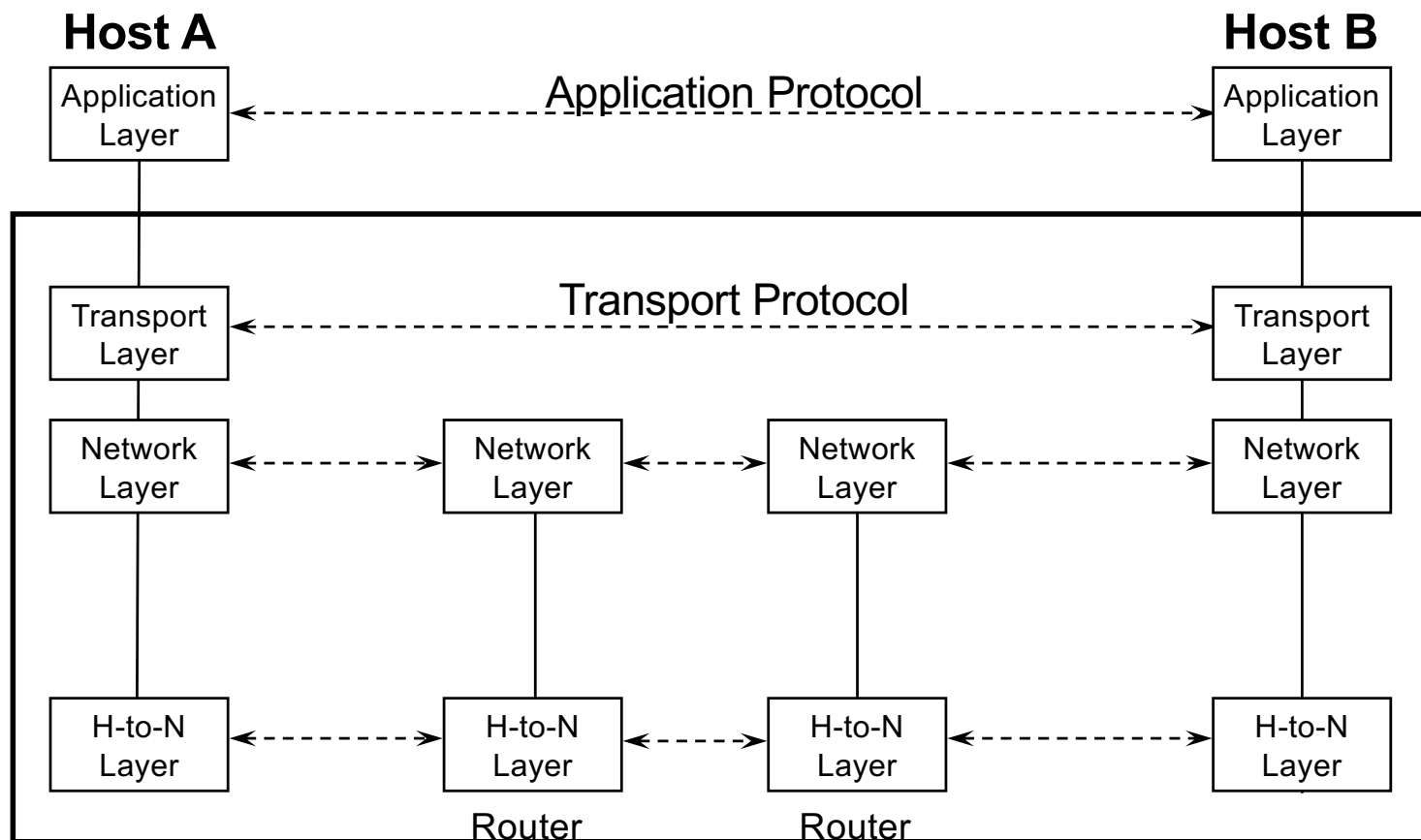
Socket is the door between network and the application/process



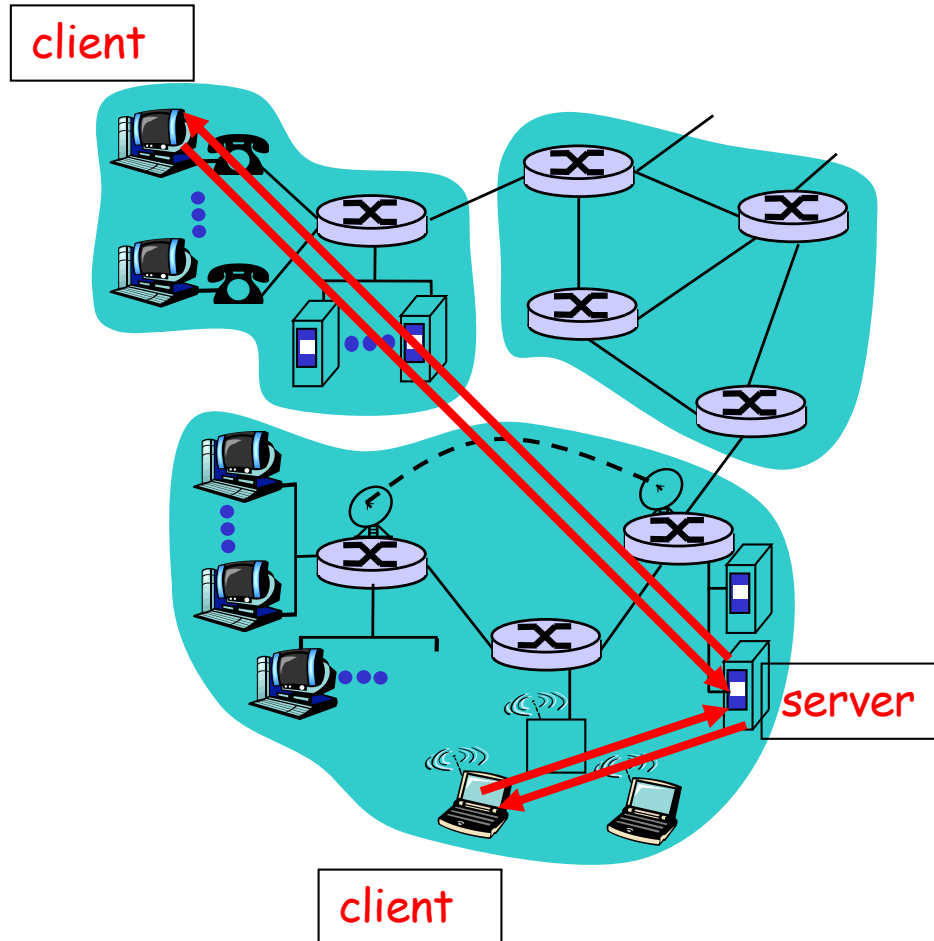
# A network connection is a 4-tuple



# Recall: Services provided by lower layers



# Client-server architecture



## server:

- ❖ always-on host
- ❖ permanent IP address
- ❖ server farms for scaling

## clients:

- ❖ communicate with server
- ❖ may not be always connected
- ❖ may have dynamic IP addresses
- ❖ do not communicate directly with each other

# DNS

Why?

For any networked application, we need to know the IP address of a host given its name

# Domain Name System (DNS)

## ❑ Problem statement:

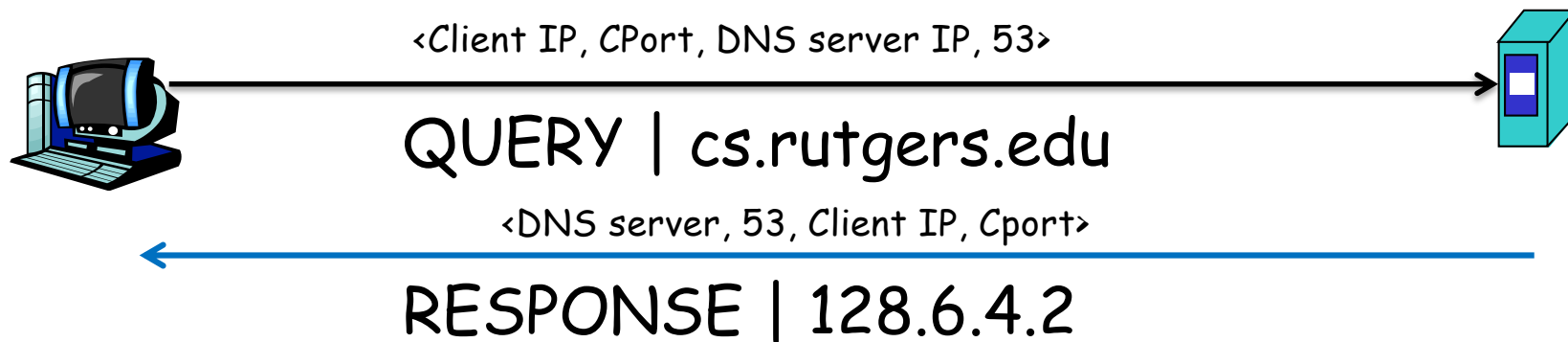
- ❖ Average brain can easily remember **7 digits** for a few names
- ❖ On average, IP addresses have **12 digits**
- ❖ We need an **easier** way to remember IP addresses

## ❑ Solution:

- ❖ Use **names** to refer to hosts
- ❖ Just as a contact or telephone **book** (white pages)
- ❖ Add a **service** (called DNS) to map between host names and binary IP addresses
- ❖ We call this ***Address Resolution***

# Simple DNS

DOMAIN NAME	IP ADDRESS
WWW.YAHOO.COM	98.138.253.109
cs.rutgers.edu	128.6.4.2
www.google.com	74.125.225.243
www.princeton.edu	128.112.132.86



- ❑ Simple but does not scale
- ❑ Every new host needs to be entered in this table
- ❑ Performance? Failure?

# DNS

## Centralize DNS?

- ❑ single point of failure
- ❑ traffic volume
- ❑ Distant centralized database
- ❑ maintenance

*doesn't scale!*