

Branched Rough Path Notes

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1 Trees

For higher order rough paths, it is useful to introduce tree notations to keep track of the algebraic and analytic conditions for the extra terms arising from the Taylor expansion.

Definition 1.1 Denote \mathcal{T} for the set of all rooted and labeled trees and \mathcal{F} for the set of the associated forests, i.e. \mathcal{F} is the free abelian monoid generated by \mathcal{T} and $f \in \mathcal{F}$ is a finite collection of rooted labeled trees (where we allow multiple copies of the same tree).

Some examples of rooted and labeled trees are

$$\bullet_a, \begin{array}{c} b \\ \downarrow \\ a \end{array}, \begin{array}{c} b \\ \swarrow \searrow \\ a \end{array}, \begin{array}{c} b \\ \downarrow \\ a \end{array}^c, \begin{array}{c} b \\ \downarrow \\ a \end{array}^d, \begin{array}{c} c \\ \downarrow \\ b \end{array}^d, \begin{array}{c} d \\ \downarrow \\ b \end{array}^c, \dots$$

We write $f = \prod_{i \in I} h_i$ as a commutative product where $h_i \in \mathcal{T}$ and I is a finite index set. We allow \mathcal{T} to contain the empty tree and we denote it by $\mathbf{1}$ since it is the multiplicative identity in \mathcal{F} . We introduce the rooting operation $[-]_a : \mathcal{F} \rightarrow \mathcal{T}$ where for $f \in \mathcal{F}$, we define $[f]_a$ for the tree obtained by joining all the trees in f to the root a , e.g.

$$[\mathbf{1}]_a = \bullet_a \text{ and } [\bullet_a \begin{array}{c} c \\ \downarrow \\ b \end{array}]_d = \begin{array}{c} d \\ \downarrow \\ a \end{array} \begin{array}{c} c \\ \downarrow \\ b \end{array}.$$

For $h \in \mathcal{T}$, denote f_h for the number of copies of h in f . Thus, we can write $f = \prod_{h \in \mathcal{T}} h^{f_h}$ where f_h is zero for all but finitely many terms. Namely, $h \in f$ if and only if $f_h \neq 0$. For $h \in \mathcal{T}$, the order of h : $|h|$ is simply the number of vertices of h and for $f \in \mathcal{F}$, $|f| = \sum_{h \in \mathcal{T}} f_h |h|$. We also write $\#f = \sum_{h \in \mathcal{T}} f_h$.

We denote $\langle \mathcal{F} \rangle$ for the free vector space generated by \mathcal{F} so $\langle \mathcal{F} \rangle$ is a unital commutative algebra with identity $\mathbf{1}$. We define the inner product of $\langle \mathcal{F} \rangle$ by setting

$$\langle f, f' \rangle = \mathbf{1}_{f=f'} \quad (1.1) \quad \boxed{\text{eq:inner-product}}$$

for all $f, f' \in \mathcal{F}$ and extending linearly.

We define the coproduct $\Delta : \langle \mathcal{F} \rangle \rightarrow \langle \mathcal{F} \rangle \otimes \langle \mathcal{F} \rangle$ by first defining Δ on \mathcal{F} inductively and then extending linearly. In particular, we set

- $\Delta \mathbf{1} = \mathbf{1} \otimes \mathbf{1}$.
- For $f \in \mathcal{F}$, $\Delta f = \prod_{h \in \mathcal{T}} (\Delta h)^{f_h}$.
- For $f \in \mathcal{F}$, $\Delta[f]_a = [f]_a \otimes \mathbf{1} + (\text{id} \otimes [-]_a) \Delta f$.

Here are some examples:

- $\Delta \cdot_a = \cdot_a \otimes \mathbf{1} + \mathbf{1} \otimes \cdot_a$.
- $\Delta {}^b\mathfrak{I}_a = {}^b\mathfrak{I}_a \otimes \mathbf{1} + \cdot_a \otimes \cdot_b + \mathbf{1} \otimes {}^b\mathfrak{I}_a$.
- $\Delta {}^b\mathfrak{V}_a^c = {}^b\mathfrak{V}_a^c \otimes \mathbf{1} + \cdot_a \cdot_b \otimes \cdot_c + \cdot_c \otimes {}^b\mathfrak{I}_a + \cdot_b \otimes {}^c\mathfrak{I}_a + \mathbf{1} \otimes {}^b\mathfrak{V}_a^c$.

2 Hopf Algebra

The algebra $\langle \mathcal{F} \rangle$ equipped with the coproduct Δ forms what is known as a Hopf algebra (in particular, it is the Connes-Kreimer Hopf algebra). We in this section provide a informal introduction to Hopf algebras.

Before we proceed with defining Hopf algebra, we first need to introduce *bialgebras*. Suppose we have an algebra \mathcal{H}^* acting on another algebra \mathcal{H}^* via the action

$$\langle \cdot, \cdot \rangle : \mathcal{H}^* \otimes \mathcal{H} \rightarrow \mathbb{R}.$$

Informally speaking, a bialgebra is then \mathcal{H} equipped with a *coproduct* $\Delta : \mathcal{H} \rightarrow \mathcal{H} \otimes \mathcal{H}$ which encodes this pairing and in particular, preserves the product structure of its *coalgebra* \mathcal{H}^* . More precisely, Δ is dual to the product operator ∇^* of \mathcal{H}^* in the sense that

$$\langle \nabla^*(f \otimes g), h \rangle = \langle f \otimes g, \Delta h \rangle$$

where $\nabla^* : \mathcal{H}^* \otimes \mathcal{H}^* \rightarrow \mathcal{H}^*$ is such that $\nabla^*(x, y) = xy$. Thus, in some sense, once we have the coproduct Δ , we can forget about \mathcal{H}^* and work solely with \mathcal{H} .

A bialgebra (\mathcal{H}, Δ) is *graded* if it has the decomposition

$$\mathcal{H} = \bigoplus_{n \in \mathbb{N}} \mathcal{H}_{(n)}$$

where $\mathcal{H}_{(n)}$ are vector spaces such that for all $n, m \in \mathbb{N}$,

$$\mathcal{H}_{(n)} \cdot \mathcal{H}_{(m)} \subseteq \mathcal{H}_{(n+m)}, \text{ and } \Delta \mathcal{H}_{(n)} \subseteq \bigoplus_{p+q=n} \mathcal{H}_{(p)} \otimes \mathcal{H}_{(q)}.$$

Denoting $\mathcal{F}_{(n)} = \{f \in \mathcal{F} : |f| = n\}$, taking $\mathcal{H}_{(n)} = \langle \mathcal{F}_{(n)} \rangle$, we observe that \mathcal{F} is a graded bialgebra.

While the coproduct preserves the product structure of \mathcal{H}^* , we would also like a map which preserves the units (invertible elements) of \mathcal{H}^* . This is achieved with the *antipole* map $S : \mathcal{H} \rightarrow \mathcal{H}$ which is required to satisfy

$$\nabla(S \otimes \text{id}) \Delta h = \nabla(\text{id} \otimes S) \Delta h = \langle \mathbf{1}, h \rangle \mathbf{1} \quad (2.1) \quad \boxed{\text{eq:antipole}}$$

for all $h \in \mathcal{H}$, where ∇ is the product operator of \mathcal{H} and

$$\langle \mathbf{1}, h \rangle = \begin{cases} 1 & \text{if } h = \mathbf{1}, \\ 0 & \text{otherwise.} \end{cases}$$

This condition is motivated by the fact that, if S is such that for any unit $f \in \mathcal{H}^*$, $S^*f = f^{-1}$ where $S^* : \mathcal{H}^* \rightarrow \mathcal{H}^*$ is the dual of S , then

$$\langle f \otimes f, (S \otimes \text{id})\Delta h \rangle = \langle \nabla^*(S^*f \otimes f), h \rangle = \langle S^*f \cdot f, h \rangle = \langle \mathbf{1}, h \rangle.$$

On the other hand,

$$\langle f \otimes f, \langle \mathbf{1}, h \rangle \mathbf{1} \otimes \mathbf{1} \rangle = \langle \mathbf{1}, h \rangle.$$

Hence, in order for S to preserve the unit structure of \mathcal{H}^* , it is necessary to require $\nabla(S \otimes \text{id})\Delta h = \nabla \langle \mathbf{1}, h \rangle \mathbf{1} \otimes \mathbf{1} = \langle \mathbf{1}, h \rangle$.

With the antipole motivated, a Hopf algebra is then simply a bialgebra equipped with an antipole.

Proposition 2.1 *A graded bialgebra satisfying $\mathcal{H}_{(0)} = \mathbb{R}$ is automatically a Hopf algebra.*

Thus, as $\langle \mathcal{F}_{(0)} \rangle = \mathbb{R}$, it follows $\langle \mathcal{F} \rangle$ is a Hopf algebra. The antipole S of $\langle \mathcal{F} \rangle$ can be computed explicitly by using the identity (2.1). We give some examples:

- $S\mathbf{1} = \mathbf{1}$.
 - $0 = \nabla(S \otimes \text{id})\Delta \cdot a = S \cdot a + \cdot a$. Thus, $S \cdot a = - \cdot a$.
 - $0 = \nabla(S \otimes \text{id})\Delta {}^b\mathbf{1}_a = S {}^b\mathbf{1}_a + (S \cdot a) \cdot b + {}^b\mathbf{1}_a$. Thus, $S {}^b\mathbf{1}_a = - {}^b\mathbf{1}_a + \cdot a \cdot b$.
 - $0 = \nabla(S \otimes \text{id})\Delta {}^b\mathbf{V}_a^c = S {}^b\mathbf{V}_a^c + (S \cdot a \cdot b) \cdot c + (S \cdot c) {}^b\mathbf{1}_a + (S \cdot b) {}^c\mathbf{1}_a + {}^b\mathbf{V}_a^c$.
- Thus, $S {}^b\mathbf{V}_a^c = - {}^b\mathbf{V}_a^c + \cdot c {}^b\mathbf{1}_a + \cdot b {}^c\mathbf{1}_a - \cdot a \cdot b \cdot c$.

Equipping $\langle \mathcal{F} \rangle$ with the inner product as defined by Equation (1.1), we can view $\langle \mathcal{F} \rangle$ acting on itself via the action $\langle f, \cdot \rangle$ for all $f \in \langle \mathcal{F} \rangle$. Thus, the underlying bialgebra of $(\langle \mathcal{F} \rangle, \Delta, S)$ can be viewed as $\mathcal{H} = \langle \mathcal{F} \rangle$ with itself as its own coalgebra $\mathcal{H}^* = \langle \mathcal{F} \rangle$ (with a different multiplication). We now give a description of the algebraic structure of $\mathcal{H}^* = \langle \mathcal{F} \rangle$.

By the property of the coproduct, if $\mathcal{H}^* = \langle \mathcal{F} \rangle$ has product $\nabla^*(f, g) = f * g$, then for any $f, g \in \mathcal{F}$ and $h \in \mathcal{F}$, we have the defining property for ∇^* :

$$\langle f * g, h \rangle = \langle \nabla^*(f, g), h \rangle = \langle f \otimes g, \Delta h \rangle.$$

Namely, if $\Delta h = \sum_{i=1}^k h_i \otimes h'_i$, then $h_i * h'_i = h$ (Check!). Moreover, we observe that $|f| + |g| = |f * g|$. $\langle \mathcal{F} \rangle$ equipped with $*$ forms an associative (*non-commutative*) algebra and we denote its coproduct by $\delta : \langle \mathcal{F} \rangle \rightarrow \langle \mathcal{F} \rangle \otimes \langle \mathcal{F} \rangle$ such that

$$\langle \delta f, h_1 \otimes h_2 \rangle = \langle f, h_1 h_2 \rangle.$$

Thus, if $\tau \in \mathcal{T}$, then, using Sweedler's notation,

$$1 = \langle \tau, \tau \cdot \mathbf{1} \rangle = \langle \delta \tau, \tau \otimes \mathbf{1} \rangle = \sum \langle \tau^{(1)}, \tau \rangle \langle \tau^{(2)}, \mathbf{1} \rangle.$$

Namely, there exists a unique component i of $\delta \tau$ for which $\tau^{(1)i} = \tau$ and $\tau^{(2)i} = \mathbf{1}$.

The tensor algebra $T(\mathbb{R}^d)$ can be identified in $\langle \mathcal{F} \rangle$ by identifying e_a by $\cdot a$, $e_b \otimes e_a$ by ${}^b\mathbf{1}_a$, $e_c \otimes e_b \otimes e_a$ by ${}^b\mathbf{1}_a^c$ and so on for $a, b, c = 1, \dots, d$. Namely, denoting $\iota : T(\mathbb{R}^d) \hookrightarrow \langle \mathcal{F} \rangle$ for this inclusion, $\iota(T(\mathbb{R}^d))$ corresponds to the set of all trees with a single branch.

Definition 2.2 Let \mathcal{H} be a Hopf algebra with coalgebra \mathcal{H}^* . Denote $\text{hom}(\mathcal{H}, \mathbb{R}) \subseteq \mathcal{H}^*$ for the space of characters (i.e. \mathbb{R} -valued algebra homomorphisms) of \mathcal{H} .

For all (by an abuse of notation) $f = \langle f, \cdot \rangle \in \text{hom}(\mathcal{H}, \mathbb{R})$, as f is multiplicative, it follows that

$$\langle f, h_1 h_2 \rangle = \langle f, h_1 \rangle \langle f, h_2 \rangle = \langle f \otimes f, h_1 \otimes h_2 \rangle.$$

On the other hand, denoting δ for the coproduct on the coalgebra \mathcal{H}^* (the cocoproduct?), we have

$$\langle \delta f, h_1 \otimes h_2 \rangle = \langle f, h_1 h_2 \rangle.$$

Thus, we can identify elements of $\text{hom}(\mathcal{H}, \mathbb{R})$ with $f \in \mathcal{H}^*$ which satisfies $\delta f = f \otimes f$. More precisely,

$$\text{hom}(\mathcal{H}, \mathbb{R}) \simeq \{f \in \mathcal{H}^* : \delta f = f \otimes f\} =: G(\mathcal{H}).$$

Definition 2.3 $G(\mathcal{H})$ forms a group and is known as the *Butcher group*.

Proposition 2.4 For all $f \in G(\mathcal{H})$, $f^{-1} = S^* f$ with S^* being the adjoint of the antipole.

3 Branched Rough Paths

Finally, we can now define branched rough paths. Defining $\mathcal{H}_{(n)}^* = \langle \mathcal{F}_{(n)} \rangle$ where the notation $*$ indicates we are working in the dual space, we define

$$G_N(\mathcal{H}) = \frac{G(\mathcal{H})}{G(\mathcal{H}) \cap \bigoplus_{n \geq N+1} \mathcal{H}_{(n)}^*}$$

where the fraction indicates the quotient group (it is not difficult to see that the denominator is indeed a normal subgroup). Namely, we identify the forests in $G(\mathcal{H})$ which has order greater than N with the identity element.

Definition 3.1 For a map $\mathbf{X} : [0, T] \rightarrow G_N(\mathcal{H})$, we denote $\mathbf{X}_{st} = \mathbf{X}_s^{-1} * \mathbf{X}_t$. \mathbf{X} is a γ -Hölder branched rough path if

$$\sup_{s < t \in [0, T]} \frac{|\langle \mathbf{X}_{st}, \tau \rangle|}{|t - s|^{\gamma|\tau|}} < \infty$$

for all $\tau \in \mathcal{F}_N$.

For a branched rough path \mathbf{X} , we define its path component X^a by

$$\delta X_{st}^a = \langle \mathbf{X}_{st}, \cdot^a \rangle.$$

We observe that the path components of a γ -Hölder branched rough path is γ -Hölder. Moreover, Chen's condition (that is $\mathbf{X}_{st} = \mathbf{X}_{su} * \mathbf{X}_{ut}$) is automatically satisfied by branching rough paths.