# Final Implementation

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### **Review**

- Tools
- Architecture

### **Tools**

- Python 2.7.1
- SimPy
- Matplotlib
- GitHub

### Input

- JSON
- Dictionary format:

 Congestion control algorithm can be specified for each flow

### Output

- Real time performance curves
  - Simulation run time
  - Data collecting interval
  - Objects to plot
    - link=1,2
    - flow
- Raw data files

### **Environment**

- Periodically pause simulation
- Collect statistics and update plots

### **Packet**

- Data packet (1 KB)
- Acknowledgement packet (64 B)
- Routing update packet (1 KB)
- FIN packet (64 B)
  - two-way handshake termination

#### **Flow**

- Sending flow:
  - Send data packets with congestion control
  - Variable start time and data amount
- Receiving flow:
  - Created on-the-fly
  - Acknowledges last data packet in sequence
  - Requests next packet in sequence

### Host

- Manage outgoing packets
- Coordinate incoming packets
  - Create receiving flows if necessary

#### Link

- Half-duplex
  - Packets transmitted in order of arrival
- Accepts packets from connected devices into buffers

#### Router

- Adjacent hosts
- Records minimum distances to all hosts from all outgoing links
- Immediately forwards data packets

### **Progress**

- Dynamic routing
- Congestion control
- Pass test cases 1 and 2

## **Dynamic Routing - Basics**

- Link weight:
  - queuing delay + propagation delay
- Bellman-Ford
- Routing update packets:
  - broadcast to adjacent routers
  - minimum distances to hosts

## **Dynamic Routing - Router Design**

- Distance table for each outgoing link
- Routing table selects link with minimum distance to host H
- Timestamp for each link indicates last arrival time of routing update packet

## **Dynamic Routing - Routing Update**

- Start of simulation
- Perform dynamic routing regularly
- Broadcast when routing table changed

## **Congestion Control**

- Go-back-N
- TCP Tahoe (with fast retransmit)
- FAST TCP

#### Go-back-N

#### Sender:

- Move window upon receiving ack of packet at the start of current window
- Send packets in the current window after timeout or receiving the previous batch

#### Receiver:

- Send ack of sequence number of the expected pkt
- ignore packets coming out of order

## **Tahoe (with Fast Retransmit)**

- Successful ack:
  - slide window;adjust window size
- Failed ack:
  - Only recognize dup ack at least 16ms apart
    - Protect flows from routing changes
  - o dup ack count = 3:
    - enter slow start
    - Resend lost packet

#### **FAST**

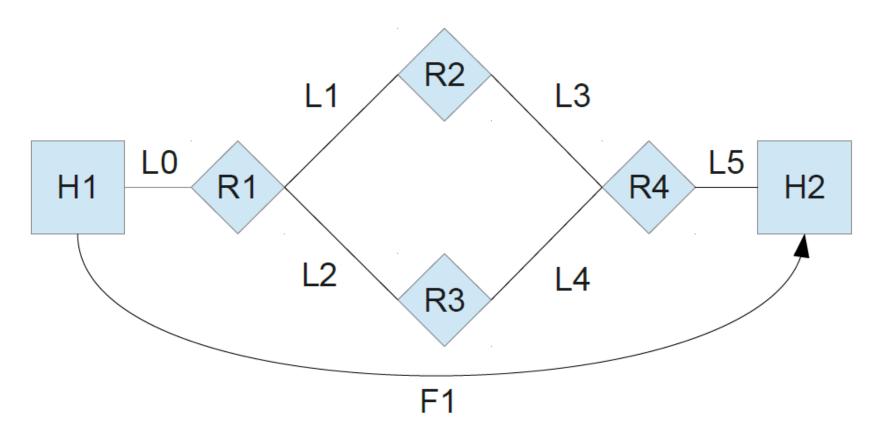
- alpha = 0.75 \* buffer size / # flow sharing link
- retransmit timeout = 3 \* rtt

```
FAST_monitor_incoming_pkts()
move window upon successful ack, update base rtt, rtt
```

### FAST()

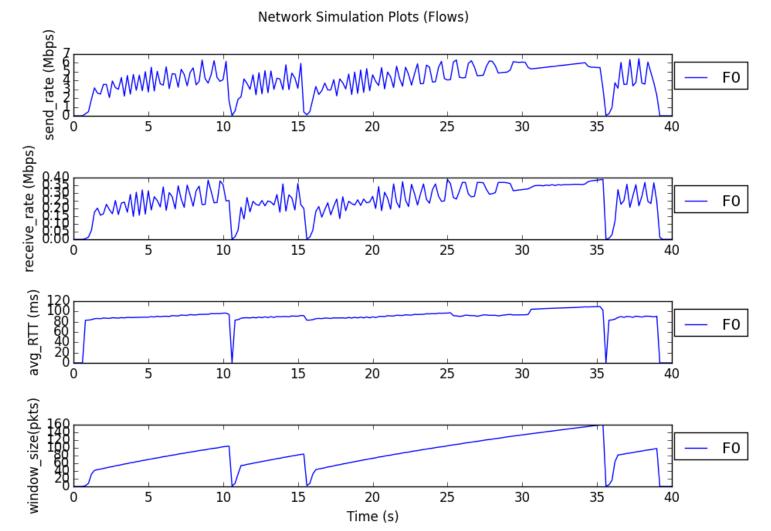
```
periodically adjust window size by delay
rtt = lastest rtt, base rtt = smallest rtt seen so far
```

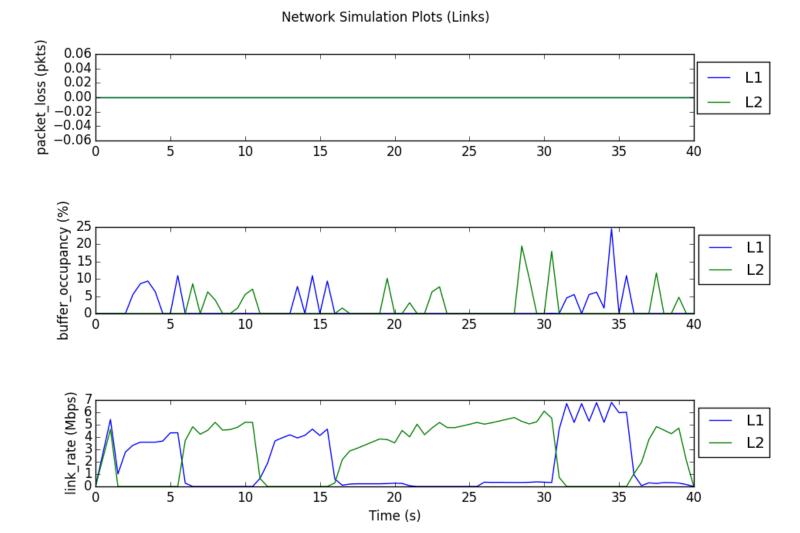
### **Test Case 1**

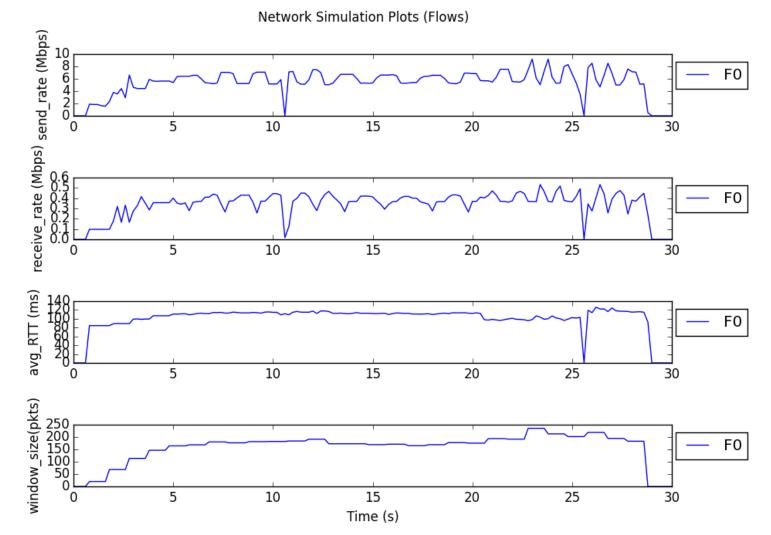


# **DEMO**

**Test Case 1** 

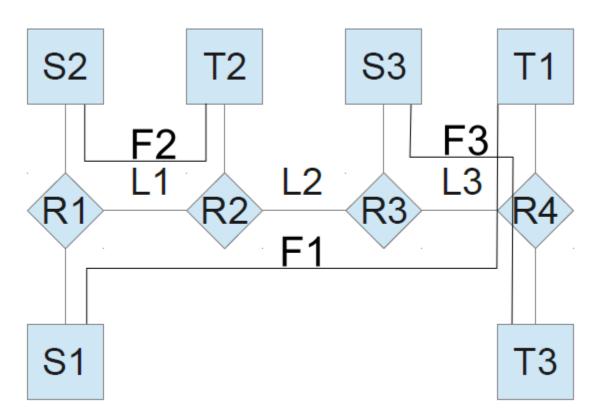






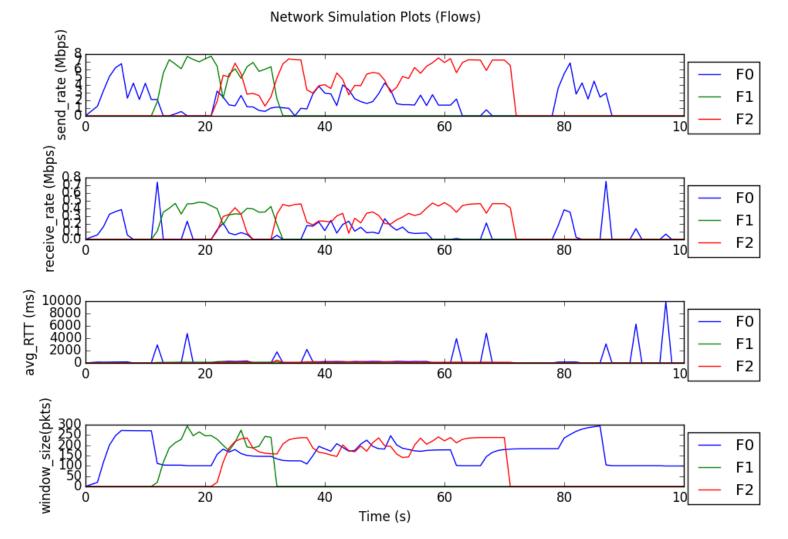
#### Network Simulation Plots (Links) packet\_loss (pkts) 0.06 0.04 L1 0.02 0.00 -0.02 -0.04 -0.06 10 15 20 25 5 30 buffer\_occupancy (%) 35 30 25 20 15 10 0 L1 15 10 30 link\_rate (Mbps) 15 Time (s) 30 5 10 20 25

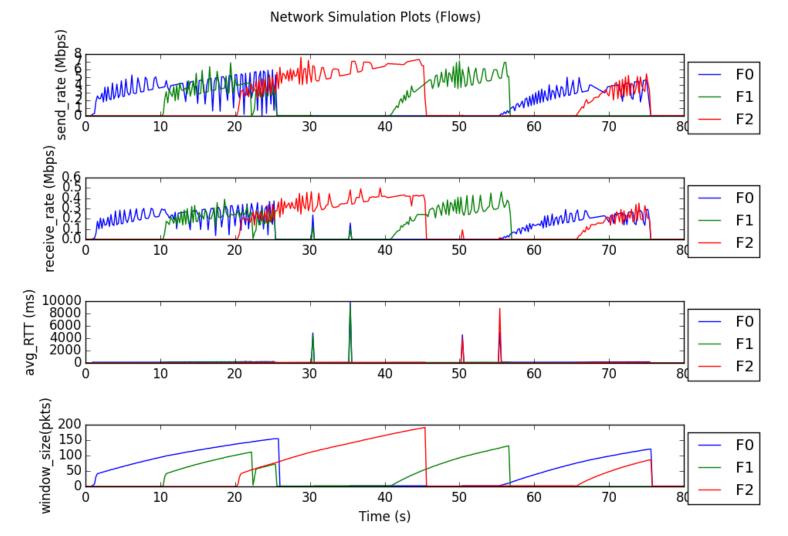
### **Test Case 2**



# **DEMO**

Test Case 2





### **Timeline and Division of Labor**

- Week 5 (Kexin): Environment, I/O, packets
- Week 6, 7 (Aman): Host, link, flow, router
- Week 8, 9 (Xander): congestion control, dynamic routing, \*additional goals
- Week 10 (Aditya): final presentation, report