Week 7

Database Design

The Design Process¹

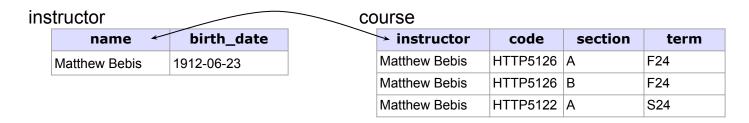
- 1. Determine the purpose of your database
 - What is the high level purpose of this database?
- 2. Find and organize the information required
 - What information should be stored to fulfill the high level purpose?
- 3. Divide the information into tables
 - What are the entities to store?
- 4. Turn information items into columns
 - What are the individual data pieces that make up those entities (tables)?
- 5. Specify primary keys
 - What is the unique field (column) for each entity (table)?
- 6. Set up the table relationships (and constraints)
 - O Do these tables relate? If so what is the connection?
- 7. Refine your design
 - Look for errors to inform improvements.
- 8. Apply the normalization rules
 - Review normalization rules and adjust as necessary.

Relationships Between Tables

- With tables, columns, and unique ids decided on... tables can be connected!
- Goals:
 - a. Identify where there are **relationships** between tables
 - b. Identify the relationship **type** between those tables

Finding Relationships

- Look for matching data in fields
 - This is often a field with the same name or data in different tables



- Determine if a table HAS A entity that is represented by another table
 - Example above: a course **HAS AN** instructor OR an instructor **HAS A** course to teach

- When a relationship is found, it will have a relationship type
- Relationships will fall into 1 of 3 categories
 - one-to-one relationship
 - one-to-many relationship b.
 - many-to-many relationship

one-to-one relationship

country				capital_city	
name	population	capital_city		city	population
Canada	38.93 million	Ottawa		Ottawa	1 million
Denmark	5.9 million	Copenhagen	←	Copenhagen	600 k

- 1 row in a table connects to exactly 1 row in a different table
 - Example above: A country **HAS A** capital city

one-to-many relationship

	_	Country		Oity		
	name	population		name	city	population
C	anada	38.93 million	←	Canada	Toronto	3 million
D	enmark	5.9 million		Canada	Vancouver	700 k

city

1 row in a table connects to many different rows in a different table OR

COUNTRY

- Many rows in a table connect to 1 row in a different table
 - Example above: A country **HAS MANY** cities, a city **HAS A** country

many-to-one relationship*same as one-to-many

name	city	population		name	population
Canada	Toronto	3 million	←	Canada	38.93 million
Canada	Vancouver	700 k		Denmark	5.9 million
		city	d	country	•

one-to-many relationship / Example

name	population			name	city	population
Canada	38.93 million	←		Canada	Toronto	3 million
Denmark	5.9 million	—		Canada	Vancouver	700 k
				Canada	Ottawa	1 million
				Denmark	Copenhagen	600 k
		L		Denmark	Odense	180 k

- Each ONE country HAS MANY cities
 - Canada has many cities
 - Denmark has many cities
- The MANY cities can HAVE just ONE country
 - Toronto, Vancouver, and Ottawa all belong to Canada
 - Copenhagen and Odense all belong to Denmark

Relationship Types

many-to-many relationship

student course

number	name
n01234567	Auston Matthews
n99999999	Vladimir Guerrero Jr.
n00000000	Natalie Spooner
n98765432	RJ Barrett



	instructor	code	section	term	class_list
_	Matthew Bebis	HTTP5126	Α	F24	Auston Matthews, Vladimir Guerrero Jr.
/	Matthew Bebis	HTTP5126	В	F24	Natalie Spooner, RJ Barrett
>	Sean Doyle	HTTP5122	Α	F24	Vladimir Guerrero Jr., RJ Barrett
7	Sean Doyle	HTTP5122	В	F24	Natalie Spooner, Auston Matthews

- many rows in a table connect to many rows in a different table
 - o Example:
 - A student HAS MANY courses
 - o A course HAS MANY students

course

Relationship Types / Bridge Tables¹

many-to-many relationship

student

*listing items in 1 cell indicates need for another table, connected with either a one-to-many or many-to-many relationship

number	name	instructor	code	section	term	class_list
n01234567	Auston Matthews	Matthew Bebis	HTTP5126	A	F24	Auston Matthews, Vladimir Guerrero Jr.
n99999999	Vladimir Guerrero Jr.	Matthew Bebis	HTTP5126	В	F24	Natalie Spooner, RJ Barrett
n00000000	Natalie Spooner	Sean Doyle	HTTP5122	A	F24	Vladimir Guerrero Jr., RJ Barrett
n98765432	RJ Barrett	Sean Doyle	HTTP5122	В	F24	Natalie Spooner, Auston Matthews

student_number course_id n01234567 HTTP5126-A-F24 n99999999 HTTP5126-A-F24 HTTP5126-B-F24 n00000000 HTTP5126-B-F24 n01234567 n99999999 HTTP5122-A-F24 n98765432 HTTP5122-A-F24

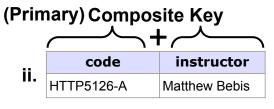
- many-to-many relationships require a bridging table to connect them
- Each record in one table can relate to multiple records in another table, and vice versa.
- With this, neither table directly relates to the other, instead both relate through the bridge

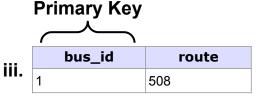
¹ Bridge Table in Database

Keys - Specify Primary Keys

- Every table should include a column that uniquely identifies each row
- The unique column is the **primary key** for the table
 - Primary key is referenced (by FK*) in creating the links (relationships) between tables
 - The primary key must be different in every row, there cannot be duplicated values
 - Primary keys should not change their value
- Primary Keys can be any column in your table, as long as they are unique
 - i. The key can be a column of data that is unique
 - ii. The key can also be a combination of columns that when combined make a unique value
 - iii. Typically creating an "id" column as the unique identifier is easiest

	Primary Key	
	number	name
ı.	n01234567	Auston Matthews





6.1 - Foreign Keys

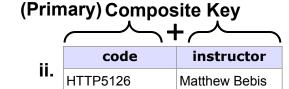
Keys - Specify Foreign Keys¹

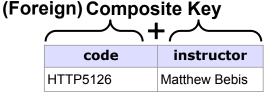
- Foreign Key: a column(s) in one table that refers to the Primary Key in another table, to formalize the relationship between tables
- Foreign Keys can reference any type of primary key from other tables
 - Unique Data PK*
 - ii. Composite PK*
 - iii. id PK*

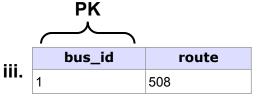
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- Coreign Key		
number	address	
n01234567	40 Bay St.	









ΕK

Keys - Specify Foreign Keys

one-to-one relationship

			1 11	
	name (PK)	population	capital_city	
country	Canada	38.93 million	Ottawa	
	Denmark	5.9 million	Copenhagen	
		•	•	

	name (PK)	population
→	Ottawa	1 million
→	Copenhagen	600 k

capital city

one-to-many relationship

	name (PK)	population
country	Canada	38.93 million
	Denmark	5.9 million

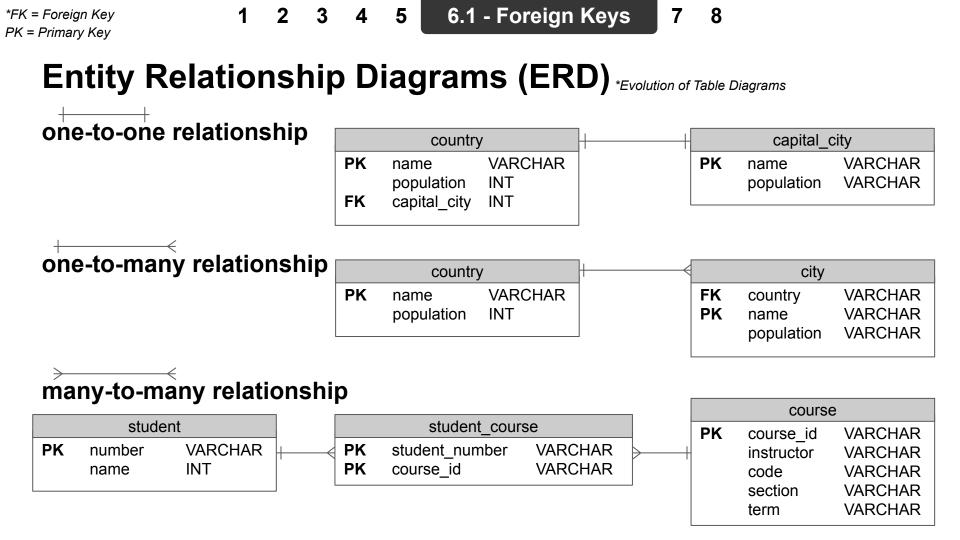
Canada Vancouver 700 k

many-to-many relationship Composite PK *That also acts as a FK

number (PK)	name
n01234567	Auston Matthews
n99999999	Vladimir Guerrero Jr.
n00000000	Natalie Spooner
n98765432	RJ Barrett

student_number	course_id
n01234567	1
n99999999	1
n00000000	2
n01234567	2

course_id (PK)	instructor	code	section	term
1	Matthew Bebis	HTTP5126	Α	F24
2	Matthew Bebis	HTTP5126	В	F24
3	Sean Doyle	HTTP5122	Α	F24
4	Sean Doyle	HTTP5122	В	F24

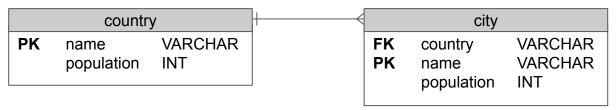


Multiple Relationship Types

one-to-one relationship

	country	/	+		capital_c	city
PK	name population	VARCHAR INT		PK	name population	VARCHAR INT
FK	capital_city	VARCHAR				

one-to-many relationship



Multiple Relationship Types

one-to-one relationship

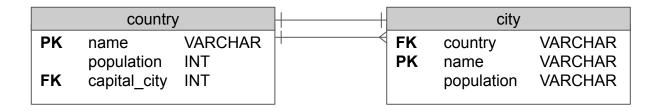
PK name VARCHAR FK country VARCHA	country	
population INT FK capital_city VARCHAR PK name VARCHAR population INT population INT	· ·	

one-to-many relationship

	countr	у]+	city	
FK capital_city VARCHAR population INT	population	INT		name	VARCHAR VARCHAR INT

Multiple Relationship Types

one-to-one relationship AND one-to-many relationship



• It is ok for 2 tables create multiple relationships with each other, as long as the tables pass the normalization rules

Table Constraints

- Constraints are used to specify rules and limit the data for columns in our tables
- This ensures the accuracy and reliability of the data in the table
- Violations of constraints abort queries
- Constraints can be column level or table level

Constraint Types¹

NOT NULL - Ensures that a column cannot have a NULL value

UNIQUE - Ensures that all values in a column are different

*PRIMARY and FOREIGN KEYS are also constraints

PRIMARY KEY - Both NOT NULL & UNIQUE

NOT NULL¹

	instructo	r
PK	instructor_id	INT
N	name	VARCHAR

UNIQUE²

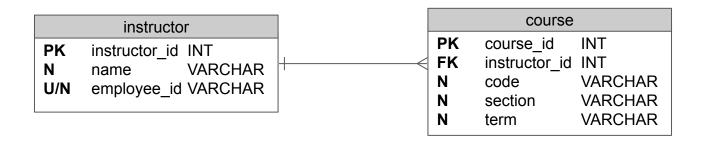
instructor				
PK	instructor_id	INT		
N	name	VARCHAR		
U	employee_id	VARCHAR		

NOT NULL & UNIQUE

	instructor
PK	instructor_id INT
N	name VARCHAR
U/N	employee_id VARCHAR

Entity Relationship Diagrams (ERD)

- Used to represent a single databases architecture
- ERDs have varying layouts but will generally show this information:
 - o Entities (tables)
 - Fields (columns)
 - Data types
 - Relationships
 - Constraints



Refine your design with Guiding Principles¹

- **Reduce Redundancy, duplicate information is bad**
 - It wastes space and increases the likelihood of errors and inconsistencies
- **Correctness and Completeness**, no exceptions
 - Incorrect data leads to incorrect data driven decisions

A good database design is, therefore, one that:

- Divides your information into subject-based tables to reduce redundant data
- Creates relationships between tables to join information together as needed
- Helps support and ensure the accuracy and integrity of your information
- Accommodates data processing and reporting needs

Normalization Forms

Normalization is the process of organizing data in a database to eliminate redundancy and inconsistencies

Normalization can be achieved by applying each normal form step by step from the earliest normal forms to the latest

The normal forms are:

1NF - 1st Normal Form	4NF - 4 th N ormal F orm	DKNF - D omain- K ey N ormal
2NF - 2nd Normal Form	ETNF - E ssential T uple	Form
3NF - 3rd Normal Form	Normal Form	6NF - 6 th N ormal F orm
	5NF - 5 th N ormal F orm	

Normalized DBs

- There are differing opinions for when a database is considered normalized
- Some argue normalization occurs after 5NF
- It is best to assume a DB is normalized when it meets the standards of your use case
- 3NF is is generally accepted as normalized

Normal Forms

First Normal Form (1NF)

- Each attribute (column) has a unique name
- Domain of attributes must not change (same data type)
- Each row is uniquely identifiable (no duplicate rows)
- Each cell must have only a single value (atomicity)

Second Normal Form (2NF) – Including 1NF

Every column is dependent on the whole primary key

Third Normal Form (3NF) – Including 1NF, 2NF

All non-key columns are functionally dependent on the primary key

1 2 3 4 5 6 7 8 - Normalization

Normal Forms

Example Unnormalized Order Table

customer_name	product_name	supplier	supplier_location	price	brand	order_id	email
John	Laptop	XYZ Electronics	Toronto	800	Dell	1	john@gmail.com
Jane	Smartphone	ABC Gadgets	Montreal	Sa:600	Apple, Samsung	2	jane@gmail.com
John	Camera	XYZ Electronics	Toronto	300	Canon	3	john@gmail.com
Alice	Chair	XYZ Furniture	Vancouver	100	Herman Miller	4	alice@gmail.com

Apply 1st Normal Form

customer_name	product_name	supplier	supplier_location	price	brand	order_id	email
John	Laptop	XYZ Electronics	Toronto	800	Dell	1	john@gmail.com
Jane	Smartphone	ABC Gadgets	Montreal	Sa:600	Apple, Samsung	2	jane@gmail.com
John	Camera	XYZ Electronics	Toronto	300	Canon	3	john@gmail.com
Alice	Chair	XYZ Furniture	Vancouver	100	Herman Miller	4	alice@gmail.com

1st Normal Form

- Each column has a unique name
- Every cell in a column must hold the same data type
- ✓ Each row is uniquely identifiable (no duplicate rows)
- A cell may not contain a set of values or a nested record

Apply 1st Normal Form

Candidate	primary	key

customer_name	supplier	supplier_location	order _id	email	product_id
John	XYZ Electronics	Toronto	1	john@gmail.com	1
Jane	ABC Gadgets	Montreal	2	jane@gmail.com	2
Jane	ABC Gadgets	Montreal	2	jane@gmail.com	3
John	XYZ Electronics	Toronto	3	john@gmail.com	4
Alice	XYZ Furniture	Vancouver	4	alice@gmail.com	5

product_id	name	price	brand
1	Laptop	800	Dell
2	Smartphone	600	Apple
3	Smartphone	600	Samsung
4	Camera	300	Canon
5	Chair	100	Herman Miller

1st Normal Form

- Each column has a unique name
- Every cell in a column must hold the same data type
- ✓ Each row is uniquely identifiable (no duplicate rows)
- A cell may not contain a set of values or a nested record
 - If it does, it should be extracted to its own table and relate back with a relationship

2 3 4 5 6 7 8 - Normalization

Apply 2nd Normal Form

Candidate primary key

customer_name	supplier	supplier_location	order_id	email	product_id
John	XYZ Electronics	Toronto	1	john@gmail.com	1
Jane	ABC Gadgets	Montreal	2	jane@gmail.com	2
Jane	ABC Gadgets	Montreal	2	jane@gmail.com	3
John	XYZ Electronics	Toronto	3	john@gmail.com	4
Alice	XYZ Furniture	Vancouver	4	alice@gmail.com	5

- ✓ 1NF
- **x** Every column must depend on the **whole primary key**, not just part of it
 - Fields that do not depend on the primary key should be extracted to a new or existing table in which they do depend on the primary key

customer_name	order_id	email	product_id
John	1	john@gmail.com	1
Jane	2	jane@gmail.com	2
Jane	2	jane@gmail.com	3
John	3	john@gmail.com	4
Alice	4	alice@gmail.com	5

supplier	supplier_location
XYZ Electronics	Toronto
ABC Gadgets	Montreal
ABC Gadgets	Montreal
XYZ Electronics	Toronto
XYZ Furniture	Vancouver

No primary key

- 1NF Each row is uniquely identifiable (no duplicate rows)
- Every column must depend on the **whole primary key**, not just part of it
 - Fields that do not depend on the primary key should be extracted to a new or existing table in which they do depend on the primary key

customer_name	order_id	email	product_id
John	1	john@gmail.com	1
Jane	2	jane@gmail.com	2
Jane	2	jane@gmail.com	3
John	3	john@gmail.com	4
Alice	4	alice@gmail.com	5

supplier_id (PK)	name	location
1	XYZ Electronics	Toronto
2	ABC Gadgets	Montreal
3	XYZ Furniture	Vancouver

product_id	name	price	brand	supplier_id
1	Laptop	800	Dell	1
2	Smartphone	600	Apple	2
3	Smartphone	600	Samsung	2
4	Camera	300	Canon	1
5	Chair	100	Herman Miller	3

- ✓ 1NF
- ✓ Every column must depend on the whole primary key, not just part of it
 - Fields that do not depend on the primary key should be extracted to a new or existing table in which they do depend on the primary key

customer_name	order_id	email	product_id
John	1	john@gmail.com	1
Jane	2	jane@gmail.com	2
Jane	2	jane@gmail.com	3
John	3	john@gmail.com	4
Alice	4	alice@gmail.com	5

^{*}Customer name depends on email

- ✓ 2NF
- Columns should not depend on non primary keys
 - If a column depends on another column that is not the primary key, it will violate 3NF
 - Columns that depend on non primary key columns should be extracted to their own table using the depended upon column as a primary key

order_id	email	product_id
1	john@gmail.com	1
2	jane@gmail.com	2
2	jane@gmail.com	3
3	john@gmail.com	4
4	alice@gmail.com	5

email	customer_name
john@gmail.com	John
jane@gmail.com	Jane
jane@gmail.com	Jane
john@gmail.com	John
alice@gmail.com	Alice

- 2NF 1NF Each row is uniquely identifiable (no duplicate rows)
- ✓ Columns should not depend on non primary keys
 - If a column depends on another column that is not the primary key, it will violate 3NF
 - Columns that depend on non primary key columns should be extracted to their own table using the depended upon column as a primary key

order_id	email	product_id
1	john@gmail.com	1
2	jane@gmail.com	2
2	jane@gmail.com	3
3	john@gmail.com	4
4	alice@gmail.com	5

email	customer_name
john@gmail.com	John
jane@gmail.com	Jane
alice@gmail.com	Alice

- ✓ 2NF 1NF Each row is uniquely identifiable (no duplicate rows)
- ✓ Columns should not depend on non primary keys
 - If a column depends on another column that is not the primary key, it will violate 3NF
 - Columns that depend on non primary key columns should be extracted to their own table using the depended upon column as a primary key

Apply 3rd Normal Form
Candidate primary key

order_id email product_id

1 john@gmail.com 1

2 jane@gmail.com 2

2 jane@gmail.com 3

3 john@gmail.com 4

4 alice@gmail.com 5

email	customer_name		
john@gmail.com	John		
jane@gmail.com	Jane		
alice@gmail.com	Alice		

- 2NF Every column must depend on the whole primary key, not part of it
- Columns should not depend on non primary keys
 - If a column depends on another column that is not the primary key, it will violate 3NF
 - Columns that depend on non primary key columns should be extracted to their own table using the depended upon column as a primary key

order_id	product_id
1	1
2	2
2	3
3	4
4	5

order_id	email
1	john@gmail.com
2	jane@gmail.com
2	jane@gmail.com
3	john@gmail.com
4	alice@gmail.com

email	customer_name
john@gmail.com	John
jane@gmail.com	Jane
alice@gmail.com	Alice

- 1NF Each row is uniquely identifiable (no duplicate rows)
- ✓ 2NF Every column must depend on the whole primary key, not part of it
- ✓ Columns should not depend on non primary keys

order_id	product_id
1	1
2	2
2	3
3	4
4	5

order_id	email
1	john@gmail.com
2	jane@gmail.com
3	john@gmail.com
4	alice@gmail.com

email	customer_name
john@gmail.com	John
jane@gmail.com	Jane
alice@gmail.com	Alice

- ✓ 1NF Each row is uniquely identifiable (no duplicate rows)
- ✓ 2NF Every column must depend on the whole primary key, not part of it
- ✓ Columns should not depend on non primary keys

Unnormalized Data

customer_name	product_name	supplier	supplier_location	price	brand	order_id	email
John	Laptop	XYZ Electronics	Toronto	800	Dell	1	john@gmail.com
Jane	Smartphone	ABC Gadgets	Montreal	600	Samsung	2	jane@gmail.com
John	Camera	XYZ Electronics	Toronto	300	Canon	3	john@gmail.com
Alice	Chair	XYZ Furniture	Vancouver	100	Herman Miller	4	alice@gmail.com

Normalized Data

order_id (PK)	email (FK)
1	john@gmail.com
2	jane@gmail.com
3	john@gmail.com
4	alice@gmail.com

order_id (PK)	product_id (PK)
1	1
2	2
2	3
3	4
4	5

product_id	name	price	brand	supplier_id (FK)
1	Laptop	800	Dell	1
2	Smartphone	600	Apple	2
3	Smartphone	600	Samsung	2
4	Camera	300	Canon	1
5	Chair	100	Herman Miller	3

email (PK)	name
john@gmail.com	John
jane@gmail.com	Jane
alice@gmail.com	Alice

supplier_id (PK)	name	location
1	XYZ Electronics	Toronto
2	ABC Gadgets	Montreal
3	XYZ Furniture	Vancouver

Normalized Data ERD

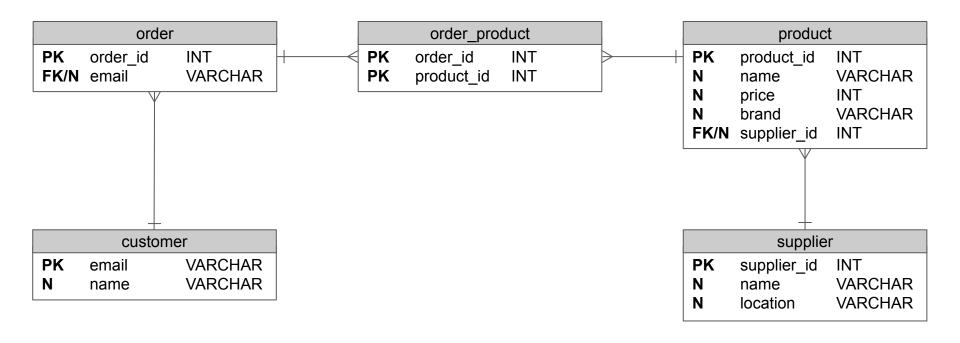


Table Relationships

- Determining the relationships between tables helps you ensure that you have the right tables and columns
- When a one-to-one or one-to-many relationship exists, the tables involved need to share a common column or columns
- When a many-to-many relationship exists, a third table is needed to represent the relationship.

Relationships will fall into 1 of 3 categories

one-to-one relationship

one-to-many relationship

many-to-many relationship

Bridging Tables

AKA Junction or Intersection Tables

- Depending on your databases need, bridging tables can store:
 - just the relationship information OR
 - other related information for each unique many-to-many combination
- Depending on the situation we can use different naming styles for the tables

number	course_id
n01234567	HTTP5126-A-F24
n99999999	HTTP5126-A-F24
n00000000	HTTP5126-B-F24
n01234567	HTTP5126-B-F24
n99999999	HTTP5122-A-F24
n98765432	HTTP5122-A-F24

student_number	course_id	date_enrolled
n01234567	HTTP5126-A-F24	2024-06-11
n99999999	HTTP5126-A-F24	2024-07-02
n00000000	HTTP5126-B-F24	2024-09-01
n01234567	HTTP5126-B-F24	2024-06-12
n99999999	HTTP5122-A-F24	2024-08-29
n98765432	HTTP5122-A-F24	2024-09-02

student course

enrollment

Why use Bridging Tables

- many-to-many relationships require a **bridging table** to relate them, if not...
- Repeated data is a data redundancy
- Compromises data integrity
- Inefficient querying

number	name	course_id
n01234567	Auston Matthews	HTTP5126-A-F24
n99999999	Vladimir Guerrero Jr.	HTTP5126-A-F24
n00000000	Natalie Spooner	HTTP5126-B-F24
n01234567	Auston Matthews	HTTP5126-B-F24
n99999999	Vladimir Guerrero Jr.	HTTP5122-A-F24
n98765432	RJ Barrett	HTTP5122-A-F24

instructor	code	section	term	student
Matthew Bebis	HTTP5126	Α	F24	Auston Matthews
Matthew Bebis	HTTP5126	Α	F24	Vladimir Guerrero Jr.
Matthew Bebis	HTTP5126	В	F24	Natalie Spooner
Matthew Bebis	HTTP5126	В	F24	RJ Barrett

Keys - Specify Foreign Keys

- Purpose of foreign keys is to create a relationship between two tables that:
 - Maintain Consistency
 - Ensure values in one table correspond to valid rows in another table, preventing data inconsistencies
 - Can Prevent Invalid Data
 - Blocks attempts at adding data with foreign key that does not have matching primary key
 - Can Enable Cascading Actions
 - Changes in the parent table automatically apply to the related rows in the child table

Table Constraints

- Rules and conditions applied to a columns to maintain data integrity, accuracy, and consistency
- Data must meet specific criteria
- Safeguards against data anomalies, errors, and inconsistencies
- Constraints can be column level or table level

NOT NULL - Ensures that a column cannot have a NULL value

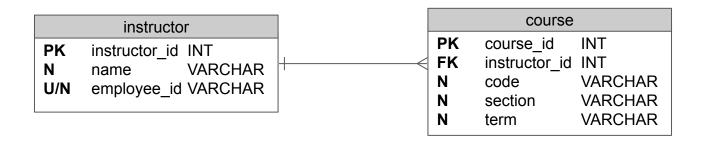
UNIQUE - Ensures that all values in a column are different

*PRIMARY and FOREIGN KEYS are also constraints

PRIMARY KEY - Both NOT NULL & UNIQUE

Entity Relationship Diagrams (ERD)

- Used to represent a single databases architecture
- ERDs have varying layouts but will generally show this information:
 - Entities (tables)
 - Fields (columns)
 - Data types
 - Relationships
 - Constraints



Normal Forms

First Normal Form (1NF)

- Each attribute (column) has a unique name
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Second Normal Form (2NF) – Including 1NF

Every column is dependent on the whole primary key

Third Normal Form (3NF) – Including 1NF, 2NF

All non-key columns are functionally dependent on the primary key